## q1. Ravishingly

- q2. a) Analysis of the first hash function  $h(M) = (\sum_{i=1}^{t} a_i) \mod n$ 
  - 1. Variable input Size: Yes, the function can take a sequence M of variable length t.
  - 2. Fixed Output Size: Yes, the Output is an integer modulo n, so it has a fixed size.
  - 3. Efficiency (Time-Space Complexity): Yes, the function is efficient with O(t) time complexity and constant space complexity.
  - 4. First and Second Pre-Image Resistance: No, it's easy to find another sequence M' such that h(M) = h(M') by just adding or subtracting multiples of n to any  $a_i$ .
  - 5. Strong Collision Resistant: No, for the same reason as above, it's easy to find two different sequences M and M' such that h(M) = h(M').
  - 6. Pseudo-randomness (Unpredictability of the Output): No, the output is directly related to the sum of the elements in *M*, making it predictable.
  - b) Analysis of the second hash function  $h_2(M) = \sum_{i=1}^t a_i^2 \pmod{n}$ .
  - 1. Variable Input Size: Yes, the function can take a sequence M of variable length t.
  - 2. Fixed Output Size: Yes, the output is an integer modulo n, so it has a fixed size.
  - 3. Efficiency (Time-Space Complexity): Yes, the function is efficient with O(t) time complexity and constant space complexity.
  - 4. First and Second Pre-image Resistance: no, it's still easy to find another sequence M'such that  $h_2(M) = h_2(M')$  by manipulating the squares of  $a_i$ .
  - 5. Strong Collision Resistance: no, for the same reason as above, it's easy to find two different sequences M and M' such that  $h_2(M) = h_2(M')$ .
  - 6. Pseudo-randomness (Unpredictability of the Output): No, the output is directly related to the sum of the squares of the elements in M, making it predictable.
  - c) Calculate the hash function of part(b) for M=(189,632,900,722,349) and n=989 to calculate  $h_2(M)$ , we first have to find the sum of the squares of the elements in M

$$\sum_{i=1}^{t} a_i^2 = 189^2 + 632^2 + 900^2 + 722^2 + 349^2$$

Then we take this sum modulo n

$$h2(M) = \left(\sum_{i=1}^{t} a_i^2\right) \bmod 989 = 229$$

So, the hash value is 229.

q3. 0x79 internationalization