

# Competitive Programming I

## Computational Complexity and C++ STL

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# Computational Complexity

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# Computational Complexity

In competitions we usually care about:

- Runtime complexity
- Memory complexity

# Computational Complexity

Complexity is often discussed in Big-Oh notation, which characterizes a function according to their growth rate. We focus on the rate at large input values.

# Computational Complexity

Complexity is often discussed in Big-Oh notation, which characterizes a function according to their growth rate. We focus on the rate at large input values.

Essential classes of complexity worth knowing

- $O(1)$  — constant
- $O(\log N)$  — logarithmic
- $O(N)$  — linear
- $O(N \log N)$  — linearithmic
- $O(N^2)$  — quadratic
- $O(N^3)$  — cubic
- $O(2^N)$  — exponential
- $O(N!)$  — factorial

# Algorithm Analysis — Example 1

Consider the following problem:

We have a sorted vector  $V$  with  $N$  integers and want to find whether an integer  $x$  exists in  $V$  or not.

Solutions?

## Algorithm Analysis — Example 1

```
bool solve(vector<int> const& v, int x) {  
    for (auto const& val : v) {  
        if (val == x) {  
            return true;  
        }  
    }  
    return false;  
}
```

Runtime complexity?

(Extra) Memory complexity?

## Algorithm Analysis — Example 1

```
bool solve(vector<int> const& v, int x) {  
    for (auto const& val : v) {  
        if (val == x) {  
            return true;  
        }  
    }  
    return false;  
}
```

Runtime complexity?  $O(N)$

(Extra) Memory complexity?  $O(1)$



## Algorithm Analysis — Example 1

```
bool solve(vector<int> const& v, int x) {  
    return std::binary_search(v.begin(), v.end(), x);  
}
```

Runtime complexity?

(Extra) Memory complexity?

## Algorithm Analysis — Example 1

```
bool solve(vector<int> const& v, int x) {  
    return std::binary_search(v.begin(), v.end(), x);  
}
```

Runtime complexity?  $O(\log N)$

(Extra) Memory complexity?  $O(1)$

## Algorithm Analysis — Example 2

Consider the following problem:

Given a number  $N$  find the number of distinct factors of  $N$ . Examples:

- $N = 8$  has 4 distinct factors (1, 2, 4, 8)
- $N = 12$  has 6 distinct factors (1, 2, 3, 4, 6, 12).

Solutions?

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    vector<int> v;  
    for (int i = 1; i <= n; ++i) {  
        if (n % i == 0) {  
            v.push_back(i);  
        }  
    }  
    return v.size();  
}
```

Runtime complexity?

(Extra) Memory complexity?

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    vector<int> v;  
    for (int i = 1; i <= n; ++i) {  
        if (n % i == 0) {  
            v.push_back(i);  
        }  
    }  
    return v.size();  
}
```

Runtime complexity?  $O(N)$

(Extra) Memory complexity?  $O(N)$

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    size_t count = 0;  
    for (int i = 1; i <= n; ++i) {  
        if (n % i == 0) {  
            count += 1;  
        }  
    }  
    return count;  
}
```

Runtime complexity?

(Extra) Memory complexity?

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    size_t count = 0;  
    for (int i = 1; i <= n; ++i) {  
        if (n % i == 0) {  
            count += 1;  
        }  
    }  
    return count;  
}
```

Runtime complexity?  $O(N)$

(Extra) Memory complexity?  $O(1)$

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    size_t count = 0;  
    for (int i = 1; i*i <= n; ++i) {  
        if (n % i == 0) {  
            count += 2;  
        }  
    }  
    return count;  
}
```

Is this correct?



## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    size_t count = 0;  
    for (int i = 1; i*i <= n; ++i) {  
        if (n % i == 0) {  
            count += 2;  
        }  
    }  
    return count;  
}
```

Is this correct?

What if the number has an integer square root? For example for 25, this code would say there are 4 distinct factors, when in reality there are only 3 (1, 5, 25). What is the problem?

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    size_t count = 0;  
    int i = 1;  
    for (; i*i < n; ++i) {  
        if (n % i == 0) {  
            count += 2;  
        }  
    }  
    if (i*i == n) {  
        count += 1;  
    }  
    return count;  
}
```

Runtime complexity?

(Extra) Memory complexity?

## Algorithm Analysis — Example 2

```
size_t solve(int n) {  
    size_t count = 0;  
    int i = 1;  
    for (; i*i < n; ++i) {  
        if (n % i == 0) {  
            count += 2;  
        }  
    }  
    if (i*i == n) {  
        count += 1;  
    }  
    return count;  
}
```

Runtime complexity?  $O(\sqrt{N})$

(Extra) Memory complexity?  $O(1)$

We will go over more examples in future classes

# Rules of thumb — Time complexity

$n$	Worst AC Algorithm	Comment
$\leq [10..11]$	$O(n!), O(n^6)$	e.g., Enumerating permutations (Section 3.2)
$\leq [17..19]$	$O(2^n \times n^2)$	e.g., DP TSP (Section 3.5.2)
$\leq [18..22]$	$O(2^n \times n)$	e.g., DP with bitmask technique (Book 2)
$\leq [24..26]$	$O(2^n)$	e.g., try $2^n$ possibilities with $O(1)$ check each
$\leq 100$	$O(n^4)$	e.g., DP with 3 dimensions + $O(n)$ loop, ${}_nC_{k=4}$
$\leq 450$	$O(n^3)$	e.g., Floyd-Warshall (Section 4.5)
$\leq 1.5K$	$O(n^{2.5})$	e.g., Hopcroft-Karp (Book 2)
$\leq 2.5K$	$O(n^2 \log n)$	e.g., 2-nested loops + a tree-related DS (Section 2.3)
$\leq 10K$	$O(n^2)$	e.g., Bubble/Selection/Insertion Sort (Section 2.2)
$\leq 200K$	$O(n^{1.5})$	e.g., Square Root Decomposition (Book 2)
$\leq 4.5M$	$O(n \log n)$	e.g., Merge Sort (Section 2.2)
$\leq 10M$	$O(n \log \log n)$	e.g., Sieve of Eratosthenes (Book 2)
$\leq 100M$	$O(n), O(\log n), O(1)$	Most contest problem have $n \leq 1M$ (I/O bottleneck)

Table 1.4: Rule of Thumb for the ‘Worst AC Algorithm’ for various single-test-case input sizes  $n$ , assuming that a year 2020 CPU can compute  $100M$  operations in 1 second.

Taken from: Competitive Programming 4 — Book 1 — Section 1.3.3

# C++ STL

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# Arrays (Static and Dynamic)

- Contiguous block of objects (cache friendly)
- $O(1)$  for lookup by index
- $O(1)$  for insertion and deletions at the back
- $O(n)$  for insertions and deletions anywhere else
- Can be stored in the stack (static) or the heap (dynamic)
- Stack is preferable if you know the size of the array

*// Static size array saved on the stack*

*#include <array>*

*std::array<int, 3> arr;*

*arr[i]; // O(1)*

*// Dynamic size array saved on the heap*

*#include <vector>*

*std::vector<int> vec();*

*vec[i]; // O(1)*

*vec[i] = ...; // O(1)*

*vec.push\_back(int); // O(1) amortized*

*vec.pop\_back(); // O(1)*



- An array of arrays
- $O(1)$  insertion and deletion at the front and back
- $O(n)$  insertion and deletion anywhere else
- Small penalty in terms of random access (compared to vector/array)
- In C++ you can use `std::deque`

# Linked Lists

- List of dynamically allocated objects (not cache friendly)
- $O(1)$  insertions and deletions at the front and back
- $O(1)$  access, insertion, and deletion with an iterator to the desired position
- $O(n)$  if we need to lookup by index for access, insertion, or deletion
- Useful whenever there are many insertions or deletions in the middle or at the beginning, and we have quick access to the position where we want to insert/delete

```
#include <list>
```

```
std::list<int> l();  
l.push_back(int);           // O(1)  
l.pop_back();               // O(1)  
l.push_front(int);          // O(1)  
l.pop_front();              // O(1)  
l.insert(iterator, int);    // O(1)  
l.erase(iterator);          // O(1)
```

There is also `std::forward_list<T>` that can only move forward

# Stacks and Queues

- Stacks are used to implement last-in first-out (LIFO) mechanisms, which can be useful, for example, to simulate recursion
- Queues are used to implement first-in first-out (FIFO) mechanisms
- They can be easily implemented with arrays and deques or linked lists. Alternatively in the **STL** you also have:
- `std::stack` for a stack
- `std::queue` for a queue

# Self-balancing binary trees

- $O(\log n)$  for insertion, deletion and lookup
- Are often tricky to implement (luckily C++ does it for us)

## Self-balancing binary trees - C++

```
#include <set>
```

```
std::set<int> s();  
s.insert(int);           //  $O(\log n)$   
s.find(int);             //  $O(\log n)$   
s.lower_bound(int);      //  $O(\log n)$  smallest element  $\geq$  input  
s.upper_bound(int);      //  $O(\log n)$  smallest element  $>$  input  
s.erase(int);            //  $O(\log n)$ 
```

There is also `std::map<K, V>` for different keys and values. Also, you have `std::multiset<V>` and `std::multimap<K, V>` if you need to support repeated values in the tree.

# Heap Priority Queue

- The minimum (or maximum) element can be accessed or removed in  $O(1)$ , insertions are  $O(\log n)$
- You can use `std::priority_queue` in C++

- $O(1)$  for insertion, deletion and lookup on average
- $O(n)$  in the worst case
- A good hashing function is very important



```
#include <unordered_set>
```

```
std::unordered_set<int> s;  
s.insert(int);    // O(1) best, O(n) worst  
s.find(int);      // O(1) best, O(n) worst  
s.erase(int);     // O(1) best, O(n) worst
```

There is also `std::unordered_map<K, V>` for different keys and values.

## Summary - Usage guidelines

Use **static size arrays** (`std::array`) when:

- You want a fixed number of items
- You want a stack-allocated array (when supported by the programming language)

Use **dynamic size arrays** (`std::vector`) when:

- You want a re-sizable array
- You want an heap-allocated array
- You will mostly be inserting and removing from the end (or near the end)

# Summary - Usage guidelines

Use a `deque` (`std::deque`) when:

- You want a dynamic size array, but
- You want to insert/remove from both ends of the array (or near them)

Use `linked lists` (`std::list`, `std::forward_list`) when:

- You need to constantly insert and remove from the middle of array in known locations
- You don't need random access by index

# Summary - Usage guidelines

Use **priority queues** (`std::priority_queue`) when:

- You only ever need the largest or smallest element of the collection

Use **self-balancing binary trees** (`std::set`, `std::map`) when:

- You want to keep elements sorted
- You need to find if elements exists or not in a list
- You need to find the elements that are closest to another

Use **hash tables** (`std::unordered_set`, `std::unordered_map`) when:

- You need to see if elements exist or not in a list and you do not care about order

# Summary

- The right choice for a data structure depends on how many insertions, deletions, or lookups have to be performed
- There are other, more specific, data structures (segment trees, tries, suffix trees, ...) that are not covered by the **STL**. However, these should cover most standard cases.
- Explore  
<https://en.cppreference.com/w/cpp/container>

# Summary

Container	Access	Insert	Erase	Find
Array	$O(1)$ : Index	N/A	N/A	$O(n)$
Vector	$O(1)$ : Index	$O(1)$ : Back $O(n)$ : Other	$O(1)$ : Back $O(n)$ : Other	$O(n)$
Deque	$O(1)$ : Index	$O(1)$ : Back/Front $O(n)$ : Other	$O(1)$ : Back/Front $O(n)$ : Other	$O(n)$
List	$O(n)$ : Index	$O(1)$ : Iter/B/F $O(n)$ : Index/Other	$O(1)$ : Iter/B/F $O(n)$ : Index/Other	$O(n)$
Set/Map	$O(\log n)$ : Key $O(n)$ : Index	$O(\log n)$	$O(\log n)$	$O(\log n)$
Hash Set/Map	$O(1)$ : Avg $O(n)$ : Worst	$O(1)$ : Avg $O(n)$ : Worst	$O(1)$ : Avg $O(n)$ : Worst	$O(1)$ : Avg $O(n)$ : Worst

# Algorithms

The STL also includes many common simple algorithms, e.g.

- **shuffle** — shuffle elements in a sequence (e.g. vector)
- **sort** — sort elements in vectors, arrays, or deques
- **lower\_bound / upper\_bound** — find lower/upper bound on sorted sequences
- **binary\_search** — check if a value exists in  $O(\log n)$  on sorted sequences
- **merge** — merge two sorted sequences
- **max / min** — find max/min value in sequence
- **equal** — check if two sets of elements are equal

And many others, explore

<https://en.cppreference.com/w/cpp/header/algorithm>

Many algorithms in `std::algorithm` use iterators to define the range of values that they apply to.

We can think of iterators as a pointer to the value, e.g.

- `std::vector<T>::begin()` — returns an iterator that points to the first element of the vector
- `std::vector<T>::end()` — returns a special iterator that has no value, but represent the end of the object

If `begin() == end()` it means that the vector is empty.



## Using iterators

- `std::sort(v.begin(), v.end())` — sort the whole vector from beginning to end
- `auto it = set.find(value)` — returns an iterator that points to the position with the value. If the value does not exist, `it` will be equal to `set.end()`
- `++it` increments the iterator, `--it` decrements the iterator
- `std::next(it)` returns the following iterator, `std::prev(it)` returns the previous iterator

# Iterators

Accessing the value associated with an iterator can be done with the `*` operator

## Code

```
vector<int> v{4,3,2,1};  
auto it = v.begin();  
cout << *it << "\n";  
++it;  
cout << *it << "\n";  
cout << *std::next(it) << "\n";
```

## Result

```
4  
3  
2
```

# Iterators

Looping with iterators

## Code

```
vector<int> v{4,2,3,1};  
for (auto it = v.begin(); it != v.end(); ++it) {  
    cout << *it << "\n";  
}
```

## Result

4  
2  
3  
1

# Iterators

Looping with iterators (implicitly)

## Code

```
vector<int> v{4,2,3,1};  
for (auto value : v) {  
    cout << value << "\n";  
}
```

## Result

4  
2  
3  
1

# Iterators

For `map` and `unordered_map` iterators return a pair of (key, value)

## Code

```
map<int, char> m;  
m.emplace(1, 'a');  
m.emplace(2, 'b');  
m.emplace(3, 'c');  
  
for (auto it = m.begin(); it != m.end(); ++it) {  
    cout << it->first << " " << (*it).second << "\n";  
}
```

## Result

```
1 a  
2 b  
3 c
```

Check more about iterators

- <https://en.cppreference.com/w/cpp/iterator>
- [https://en.cppreference.com/w/cpp/container#Iterator\\_invalidation](https://en.cppreference.com/w/cpp/container#Iterator_invalidation)

# Algorithms with custom functions

Some algorithms make use of custom functions. Sort for example

## Code

```
struct Point {  
    int x, y;  
};  
  
vector<Point> v{{2,2}, {0,3}, {1,5}, {0, 2}};  
  
// Sort first by x, and in case of tie by y  
sort(v.begin(), v.end(), [](auto const& lhs, auto const& rhs) {  
    return lhs.x < rhs.x || (lhs.x == rhs.x && lhs.y < rhs.y);  
});  
  
for (auto const& p : v) {  
    cout << p.x << ", " << p.y << "\n";  
}
```

# Algorithms with custom functions

## Result

0, 2

0, 3

1, 5

2, 2



## Define custom operators

Another alternative is to define operators for the object

### Code

```
struct Point {  
    int x, y;  
    bool operator<(Point const& other) const {  
        return x > other.x || (x == other.x && y > other.y);  
    } };
```

```
vector<Point> v{{2,2}, {0,5}, {1,5}, {0, 2}};  
// Sort using the operator< function by default  
sort(v.begin(), v.end());
```

```
for (auto const& p : v) {  
    cout << p.x << ", " << p.y << "\n";  
}
```

# Define custom operators

## Result

2, 2

1, 5

0, 3

0, 2

# Conclusion

C++ is a powerful language where many things are possible  
However, it is also a big language with many advanced concepts  
Explore the documentation  
<https://en.cppreference.com/w/cpp>  
And ask us!