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Ans There exe two memory management techniques:

Contiguous and Non-Contiguous. In contiguous Technique,

emecuting process must be loaded entirely in main
memory. Contiguous Technique can be divided into:

- i) Fixed or static partitioning
- 2) Voxiable or dynamic postitioning

## · FIXED PARTITIONING:

This is the oldest and simplest technique used to put more than one processes in the main memory. In this partitioning, number of partitions in RAM are fixed but size of each partition may or may not be same. As it is configuous allocation, hence no spanning is allowed. Here partition are made before enecution or during system configure.

Block size = 4MB Free = 3MB & Internal Fragmentation.

P1 = 1MB

Block size = 8MB

P2 = 7MB

Block size = 6MB

P3 = 7MB

Block size = 16MB

P4 = 14MB

fixed size partition.

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As illustrated in the above figure, first process is only consuming 1MB out of 4MB in the main memory.

Hence, Internal pragmentation in livet block = 4-1 = 3MB

Sum of Internal pragmentation in every block = (4-1) + (8-7) + (14-14)

= 7 MB

suppose Ps process of 7MB comes. But this process connot be accommodated inspite of available free space because of contiguous allocation. Hence 7MB becomes part of Enternal Fragmentation

## 2) VARIABLE PARTITIONING:

It is a part of contiquous allocation technique. It is used to alleviate the problem faced by fixed partitioning. In contrast with fined partitioning, partitions are not made before the enecution or during system configure. Various jea tules associated with raviable partitions are made, during tritially partitions are made, during

- · Initially RAM is empty and partitions are made during the run-time according to processes need instead of partioning during system configure.
- . The size of partition will be equal to incoming process.

  The partition size varies according to the need of

the process so that the internal fragmentation can be avoided to ensure efficient utilization of RAM

on the number of incoming process and main memory's size

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	1	
	operating System	
	V V	
0 ,	P1 = 2 MB	Block size = 2MB
	P2 = 7 MB	BOCK Size = 7MB
	1	
	P3 = 1 MB	Block size = 1 MB
•		
	P4 = 5 MB	Block size = SMB
	Empty space	

postition size = process size

.. No internal Fragmentation

- · No internal fragmentation
- · No restriction on Degree of Multiprogramming
- · No limitation on size of proces &

· Easy to implement internal Fragmentation

- casy to implement . Internal Fragmentation Little os overhead. Enternal Fragmentation
  - · Limit Process size
  - Unitation on degree of multiprogram

Disadvantage of variable
Distinct implementation
Enternal Fragmentation

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FOR EDUCATIONAL USE

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QS	Process	AL	LDCAT	toN	M	ΑX			1	N EE	2		
		А	В	С	A	В	C		A	В	C		
	РО	0	l	0	7	5	3		7	4	3	_	
	P1	2	0	O	3	2	2	*1	1	2	2		
	P2	3	0	2	9	0	2		G	0	0	_	-
	P3	2	1	1.	-2	2-	2	;  ;	0	. 1	1		
	P4	0	0	2	4.	:3:	3		4	3	I		

TOTAL ALLOCATED: 7 2 5

TOTAL	resources			
A	В	C		
10	5	7 ;		

AVAILABLE
A B C
3 3 2

- ·: Need [Po] 4 Available, Po cannot be enecuted. >
  - : Avoilable = 5 3 2
- · · · Need [P2] { Available P2 can't be enecuted X
- : Need [P3] < Available P3 can be enecuted
  - : Available = 7 45
- .: Need [po] < Available po can be enecuted ~
  - : Available = 755
- .. Need of [P27 < Available, P2 can be enecuted V
  - .. Available = 10 5 7

NO MORE PROCESSE C

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:. System is in a sale state

:. sale sequence :-

$$P_1 \longrightarrow P_2 \longrightarrow P_4 \longrightarrow P_6 \longrightarrow P_2$$
  
 $P_1 \longrightarrow P_3 \longrightarrow P_4 \longrightarrow P_2 \longrightarrow P_0$