

Red-Black Tree

[BST with following properties]

Properties of red black tree

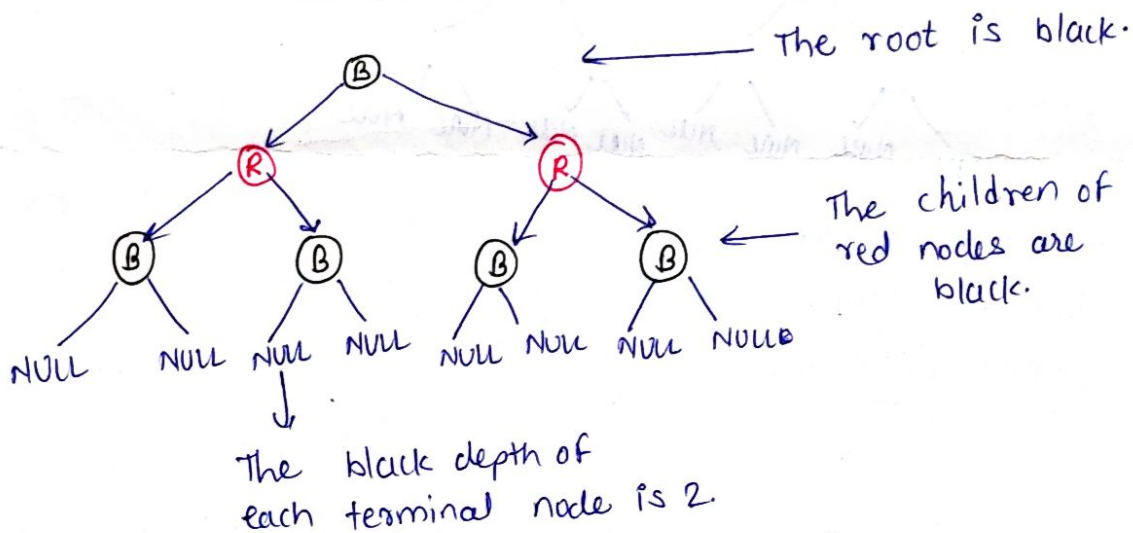
The nodes of a red-black tree are either red or black.

- ① The root of the tree is always black.
- ② A black node can have a black or a red child.
- ③ A red node cannot have a red child. It can only have a black child.
- ④ The black depth of a terminal node is the number of black nodes encountered while travelling from the terminal node to the root.
- ⑤ The black depth of a terminal node is always same.

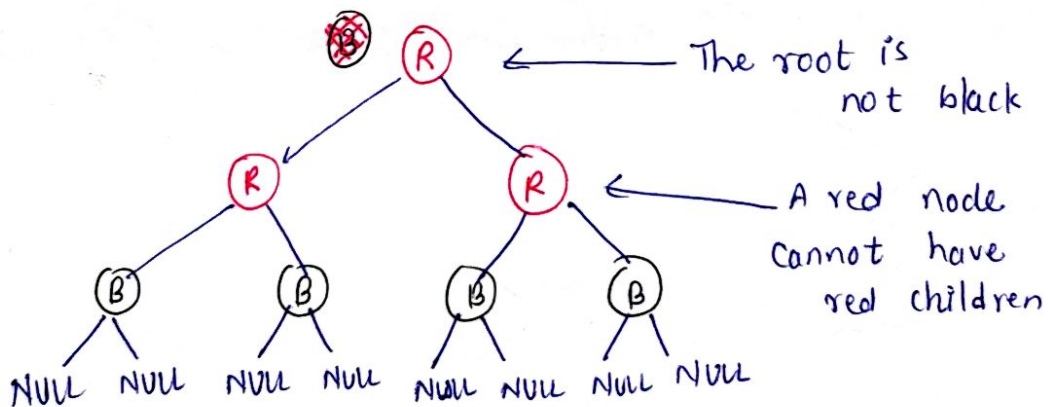
Black depth:- The number of black nodes from the terminal to the root is called the black depth of the node.

⑥ The leaves of a red-black tree would always be a NULL node.

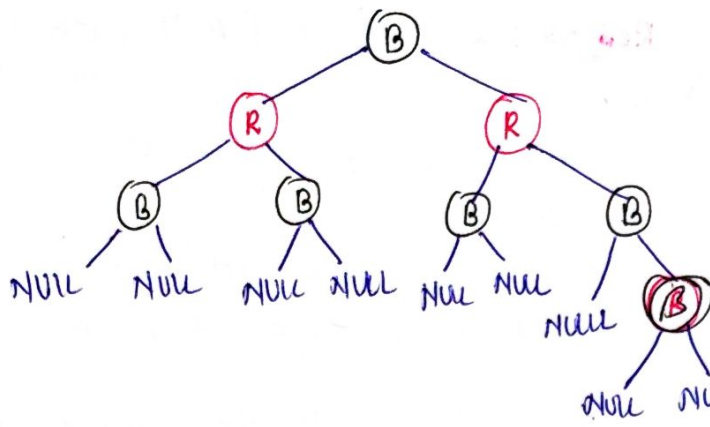
⑦ Each black or red node last in the hierarchy will have NULL nodes as children.



a) An example of a red-black tree



b) An example of a tree which is not red-black tree

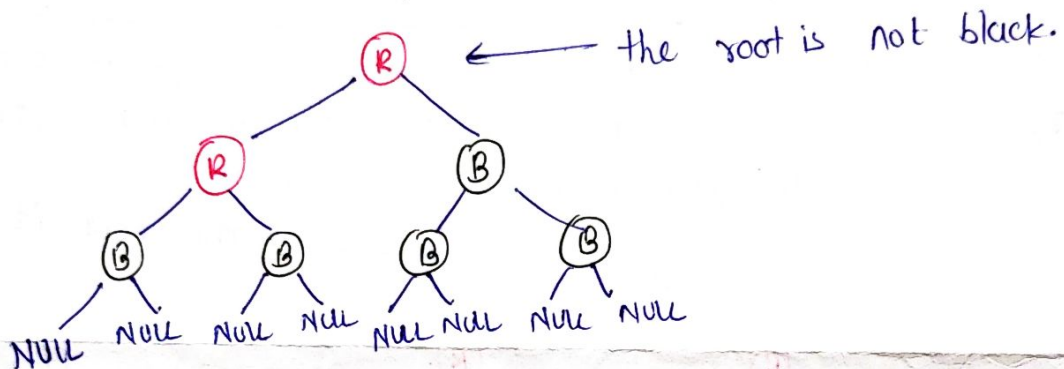


The black depth of this node & its sibling is 3, whereas for the other nodes it is 2.

c) An example of a tree which is not red-black tree

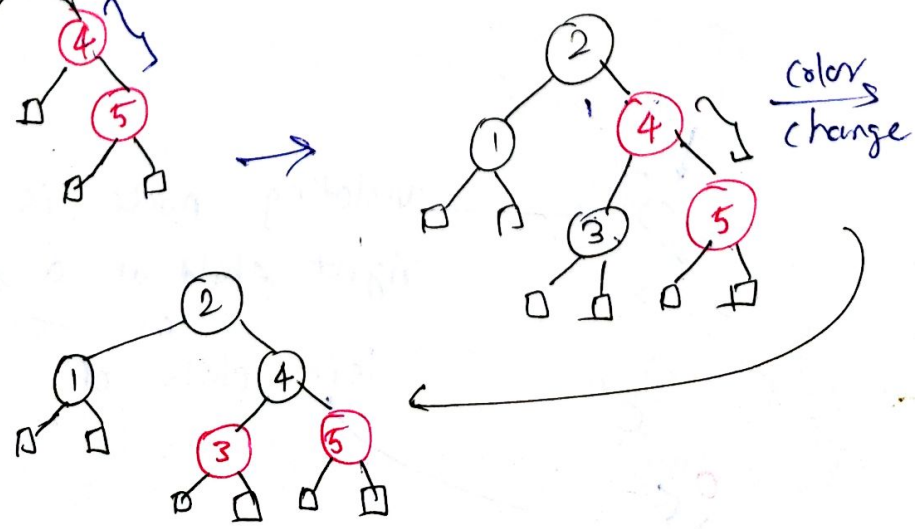
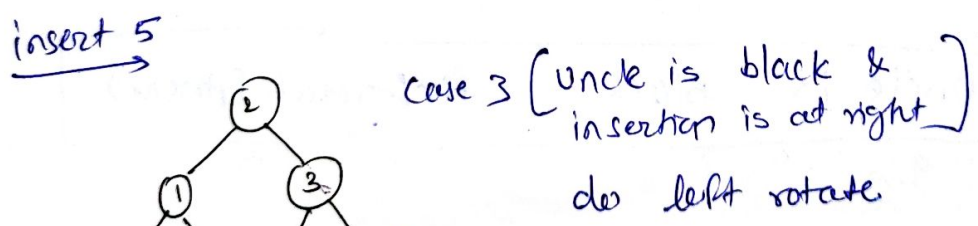
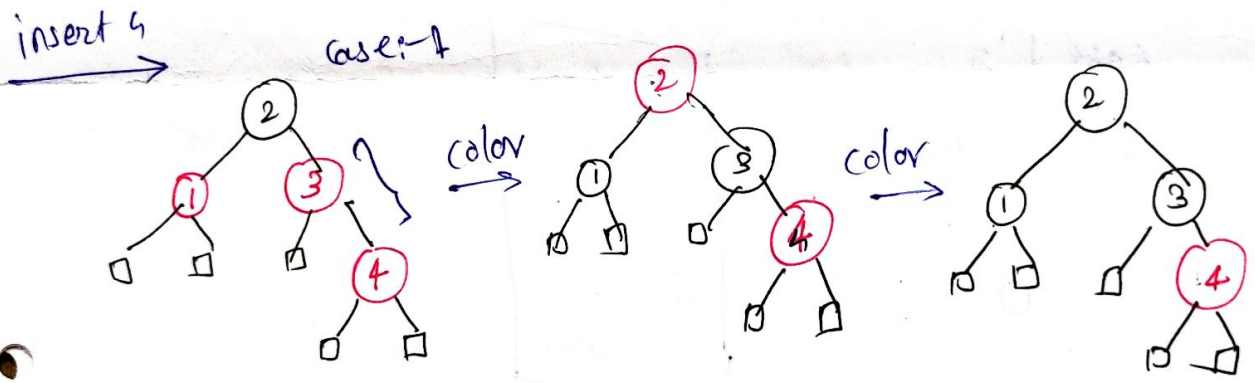
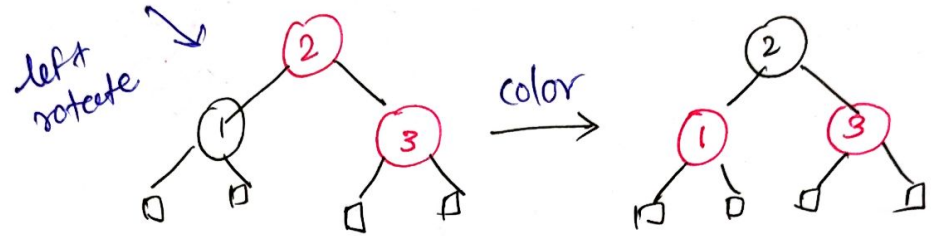
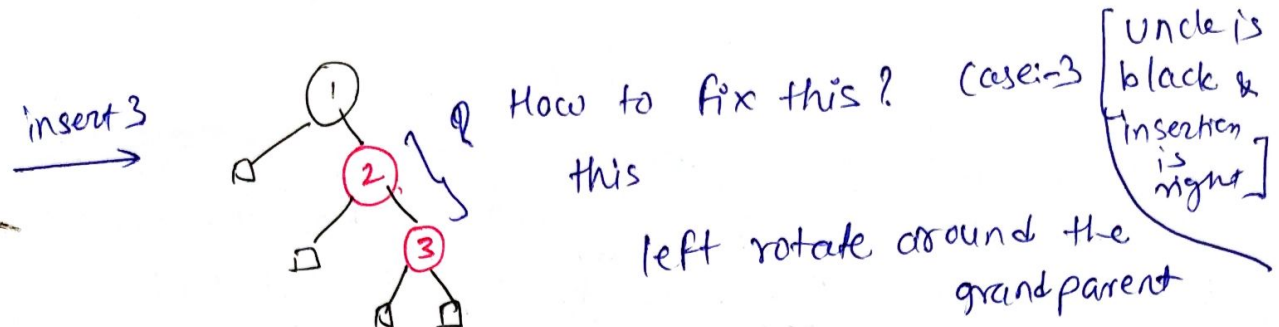
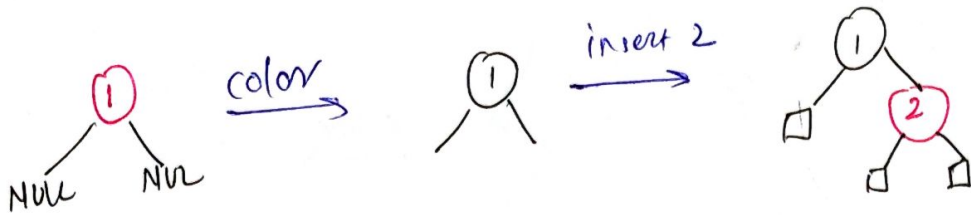
Double red problem

The case wherein the child of a red node is a red node is called the double red problem.

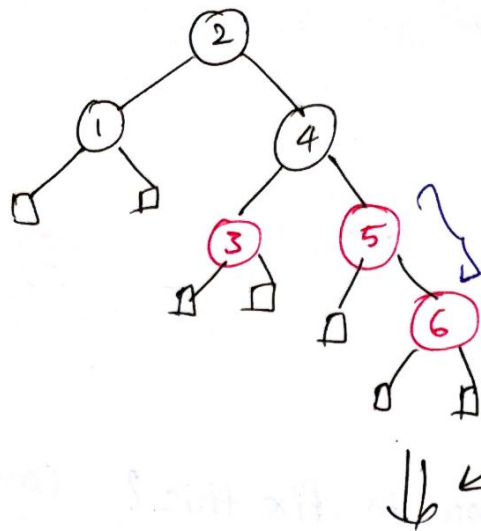


A red-black tree is a binary search tree with one extra bit of storage per node: its color which can be either RED or BLACK

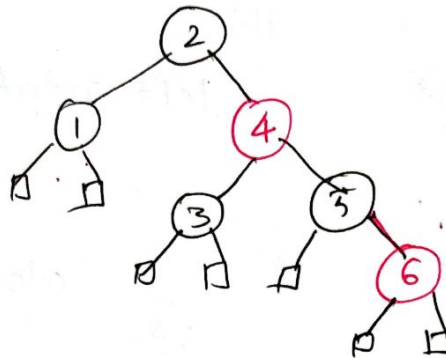
1 2 3 4 5 6



insert 6 →



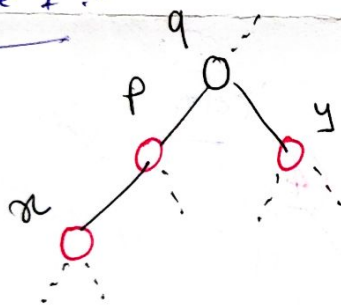
Uncle is red Case 1



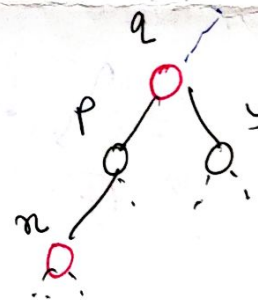
Black height of this tree is 3

Violations

Case 1 :-

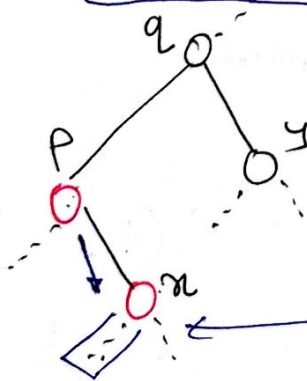


Uncle is Red
fix: color



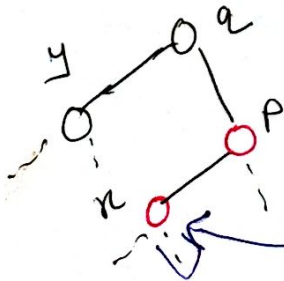
Case 2 :-

Uncle is Black fix: rotate (parent)

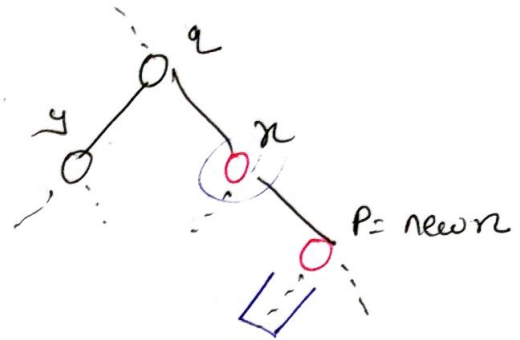
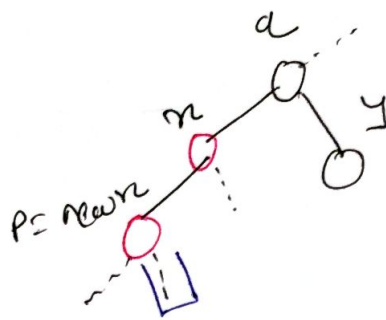


violating node is a right child of a left child or

left child of a right child



to fix this rotate around the parent node.



Convert to case 3.

