



Shri Vile Parle Kelavani Mandal's

DWARKADAS J. SANGHVI COLLEGE OF ENGINEERING

(Autonomous College Affiliated to the University of Mumbai)

NAAC Accredited with "A" Grade (CGPA : 3.18)



POA LAB EXPERIMENTS

Name	Junaid Girkar
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Division	A
Batch	A4
Branch	Computer Engineering
Semester	5
Subject	Processor Organization and Architecture
Subject Code	DJ19CEL502

**Experiment List**

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1	To implement Booth's multiplication algorithm.	3						
2	To study and implement Restoring division algorithm.	11						
3	To implement non restoring binary division method	21						
4	<p>Implement Sequential Memory Organization with given details: Processor can access one word at a time (Word accessible memory), L1 cache can store max 32 words, L2 cache can store 128 words, main memory has the capacity of 2048 bytes. Consider TL1=20 ns, TL2=60ns, TMM=120ns. Create output table containing</p> <table border="1"> <tr> <td>Processor Request (Sp. Word)</td><td>Location of Hit</td><td>Average Memory Access Time</td></tr> <tr> <td></td><td></td><td></td></tr> </table>	Processor Request (Sp. Word)	Location of Hit	Average Memory Access Time				29
Processor Request (Sp. Word)	Location of Hit	Average Memory Access Time						
5	Study various microprocessors and microcontrollers used in the market specifically targeting Aircrafts and Automated Vehicles.	37						
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8	Run the simulator using the given link https://ydoc2.github.io/8086-emulator-web/compile . Study and execute First 6 examples from the repository https://github.com/YJDoc2/8086-Emulator/tree/master/examples .	70						
9	Write Programs using Assembly Language: 1. To implement Macros for calculating Factorial of a number 2. To study and Implement DOS interrupts. Eg: calculate and display Sum of 2 user entered inputs using DOS interrupt	80						
10	Write a Program using ALP to Simulate Microcontroller interfacing with 7 segment display. Display your SAP ID using this tool.	86						



EXPERIMENT - 1

AIM: To demonstrate Booth's Algorithm for multiplication of 2 signed binary numbers.

Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	01-10-2021	01-10-2021	

THEORY:

Booth algorithm gives a procedure for multiplying binary integers in signed 2's complement representation in an efficient way, i.e., less number of additions/subtractions required. It operates on the fact that strings of 0's in the multiplier require no addition but just shifting and a string of 1's in the multiplier from bit weight 2^k to weight 2^m can be treated as $2^{(k+1)}$ to 2^m .

Booth's algorithm can be implemented by repeatedly adding (with ordinary unsigned binary addition) one of two predetermined values A and S to a product P , then performing a rightward arithmetic shift on P . Let m and r be the multiplicand and multiplier, respectively; and let x and y represent the number of bits in m and r .

1. Determine the values of A and S , and the initial value of P . All of these numbers should have a length equal to $(x + y + 1)$.
 1. A : Fill the most significant (leftmost) bits with the value of m . Fill the remaining $(y + 1)$ bits with zeros.
 2. S : Fill the most significant bits with the value of $(-m)$ in two's complement notation. Fill the remaining $(y + 1)$ bits with zeros.
 3. P : Fill the most significant x bits with zeros. To the right of this, append the value of r . Fill the least significant (rightmost) bit with a zero.
2. Determine the two least significant (rightmost) bits of P .
 1. If they are 01, find the value of $P + A$. Ignore any overflow.
 2. If they are 10, find the value of $P + S$. Ignore any overflow.
 3. If they are 00, do nothing. Use P directly in the next step.
 4. If they are 11, do nothing. Use P directly in the next step.



- Arithmetically shift the value obtained in the 2nd step by a single place to the right. Let P now equal this new value.
- Repeat steps 2 and 3 until they have been done y times.
- Drop the least significant (rightmost) bit from P . This is the product of m and r .

TIME COMPLEXITY: $O(n) * (\text{complexity_of_addition} + \text{complexity_of_shift})$

EXAMPLE: Multiplication of 7 and 3

Q_n	Q_{n+1}	$M = (0111)$ $M' + 1 = (1001)$ & Operation	AC	Q	Q_{n+1}	SC
1	0	Initial	0000	0011	0	4
		Subtract ($M' + 1$)	1001			
			1001			
		Perform Arithmetic Right Shift operations (ashr)	1100	1001	1	3
1	1	Perform Arithmetic Right Shift operations (ashr)	1110	0100	1	2
0	1	Addition ($A + M$)	0111			
			0101	0100		
		Perform Arithmetic right shift operation	0010	1010	0	1
0	0	Perform Arithmetic right shift operation	0001	0101	0	0

Final Answer = $(0001\ 0101)_2$
= $(21)_{10}$

CODE:

```
#include<stdio.h>
#include<stdlib.h>
#include<string.h>

void rightShift();

int main()
```



```
{
    printf("\n");
    printf("BOOTH's Algorithm\n");
    printf("\n");
    printf("Enter two numbers that are to be multiplied : "); //taking two numbers as inputs
    int a,b;
    scanf("%d %d",&a,&b);
    int ap=a,bp=b;
    if(ap<0) // Negative values check
        ap*=-1;
    if(bp<0) bp*=-1;
    if(bp>ap) //taking greater VALUE as multiplicand
    {
        ap=bp+ap-(bp=ap);
        a=b+a-(b=a);
    }
    int t1=ap,t2=bp;
    int ab[35]={};
    int bb[35]={};
    int i=0;
    while(t1>0)
    {
        ab[i]=t1%2;
        i++;
        t1/=2;
    }
    ab[i]=0;
    int j=0;
    while(t2>0)
    {
        bb[j]=t2%2;
        j++;
        t2/=2;
    }
    while(j<=i) //equating bits to the previous(ab) binary number(ab will either be larger or equal to
    bb).
        bb[j++]=0;
    int nb=i+1; //nb is number of bits
    i=0;j=0;
    while(i<nb/2) //converting VALUES to binary
    {
        ab[i]=ab[nb-i-1]+ab[i]-(ab[nb-i-1]=ab[i]);
        i++;
    }
    i=0;
    while(i<nb/2) { bb[i]=bb[nb-i-1]+bb[i]-(bb[nb-i-1]=bb[i]); i++; }
    int x[35]={0}; int y[35]={0}; i=0; if(a>=0)
    //taking actual binary numbers
```



```
{ //x is multiplicand and y is multiplier
    while(i<nb)
        x[i]=ab[i+++1];
}
else //2's compliment
{
    while(i<nb) { if(ab[i]==0) x[i]=1; else x[i]=0; i++; } i=1; x[nb-i]++; while(x[nb-i]==2) { x[nb-i]=0; i++;
x[nb-i]++; } } i=0; if(b>=0)
{
    while(i<nb)
        y[i]=bb[i+++1];
}
else //2's compliment
{
    while(i<nb) { if(bb[i]==0) y[i]=1; else y[i]=0; i++; } i=1; y[nb-i]++; while(y[nb-i]==2) { y[nb-i]=0; i++;
y[nb-i]++; } } printf("\n"); //output starts here printf("Multiplicand (Q) %d -> ",a);
i=0;
printf("Multiplicand (Q) -> ");
while(i<nb) printf("%d",x[i++]); printf("\nMultiplier (M) %d -> ",b);
i=0;
while(i<nb)
    printf("%d",y[i++]);
printf("\n");
i=0;
int ym[35]={0}; //calculating -M
if(b<0)
{
    while(i<nb)
        ym[i]=bb[i+++1];
}
else
{
    while(i<nb) { if(bb[i]==0) ym[i]=1; else ym[i]=0; i++; } i=1; ym[nb-i]++; while(ym[nb-i]==2) {
ym[nb-i]=0; i++; ym[nb-i]++; } } printf("we use -(M) i.e. %d -> ",-b);
i=0;
while(i<nb)
    printf("%d",ym[i++]);
printf("\n");
int q0=0;
int p[35]={0}; //p here is value that is stored in accumulator. initially set to zero.
int steps=nb;
printf("\n");
printf("n\t");
i=0;
while(i<nb)
{
    if(i*2==nb || i*2==nb-1)
```



```
printf("A");
else
printf(" ");
i++;
}
printf(" ");
i=0;
while(i<nb)
{
    if(i*2==nb || i*2==nb-1)
        printf("Q\t");
    else
        printf(" ");
    i++;
}
printf(" Q-1");
printf("\n");
j=0;

while(steps--) //counting down steps.
{
    printf("%d\t",j++);
    i=0;
    while(i<nb)
        printf("%d",p[i++]);
    printf(" ");
    i=0;
    while(i<nb)
        printf("%d",x[i++]);
    printf(" ");
    printf("%d\n",q0);
    if(x[nb-1]==0 && q0==0) //0-0 condition
    {
        q0=x[nb-1];
        rightShift(p,x,nb);
    }
    else if(x[nb-1]==0 && q0==1) //0-1 condition
    {
        printf(" A + M ");
        i=0;
        while(i<nb)
            printf("%d",y[i++]);
        i=0;
        while(i<nb)
        {
            p[nb-i-1]+=y[nb-i-1];
            if(p[nb-i-1]==2)
```



```
{
    p[nb-i-1]=0;
    if(nb-i-1!=0)
        p[nb-i-2]++;
}
if(p[nb-i-1]==3)
{
    p[nb-i-1]=1;
    if(nb-i-1!=0)
        p[nb-i-2]++;
}
i++;
}
printf("\n      ");
i=0;
while(i<nb)
    printf("%d",p[i++]);
printf("\n");
q0=x[nb-1];
rightShift(p,x,nb);
}
else if(x[nb-1]==1 && q0==0) //1-0 condition
{
    printf("  A - M ");
    i=0;
    while(i<nb)
        printf("%d",ym[i++]);
    i=0;
    while(i<nb)
    {
        p[nb-i-1]+=ym[nb-i-1];
        if(p[nb-i-1]==2)
        {
            p[nb-i-1]=0;
            if(nb-i-1!=0)
                p[nb-i-2]++;
        }
        if(p[nb-i-1]==3)
        {
            p[nb-i-1]=1;
            if(nb-i-1!=0)
                p[nb-i-2]++;
        }
        i++;
    }
    printf("\n      ");
    i=0;
```




```
while(i<nb)
    printf("%d",p[i++]);
printf("\n");
q0=x[nb-1];
rightShift(p,x,nb);
}
else if(x[nb-1]==1 && q0==1) //1-1 condition
{
    q0=x[nb-1];
    rightShift(p,x,nb);
}
}
printf("%d    ",j);
i=0;
while(i<nb)
    printf("%d",p[i++]);
printf(" ");
i=0;
while(i<nb)
    printf("%d",x[i++]);
printf(" ");
printf("%d\n",q0);
printf("\n");

printf("Final Product in signed binary number is : ");
i=0;
while(i<nb)
    printf("%d",p[i++]);
i=0;
printf(" ");
while(i<nb)
    printf("%d",x[i++]);
printf("\n\n");
return 0;
}

void rightShift(int p[],int x[],int nb)
{
    int i=0;
    while(nb-i-1)
    {
        x[nb-i-1]=x[nb-i-2];
        i++;
    }
    x[0]=p[nb-1];
    i=0;
    while(nb-i-1)
```



```
{  
    p[nb-i-1]=p[nb-i-2];  
    i++;  
}  
}
```

OUTPUT:

BOOTH's Algorithm

Enter two numbers that are to be multiplied : 20 35

Multiplicand (Q) -> 0100011

Multiplier (M) 20 -> 0010100

we use -(M) i.e. -20 -> 1101100

n	A	Q	Q-1
0	0000000	0100011	0
	A - M	1101100	
	1101100		
1	1110110	0010001	1
2	1111011	0001000	1
	A + M	0010100	
	0001111		
3	0000111	1000100	0
4	0000011	1100010	0
5	0000001	1110001	0
	A - M	1101100	
	1101101		
6	1110110	1111000	1
	A + M	0010100	
	0001010		
7	0000101	0111100	0

Final Product in signed binary number is : 0000101 0111100

...Program finished with exit code 0

Press ENTER to exit console.



EXPERIMENT - 2

AIM: To demonstrate the Restoring Division Algorithm for 2 unsigned binary numbers.

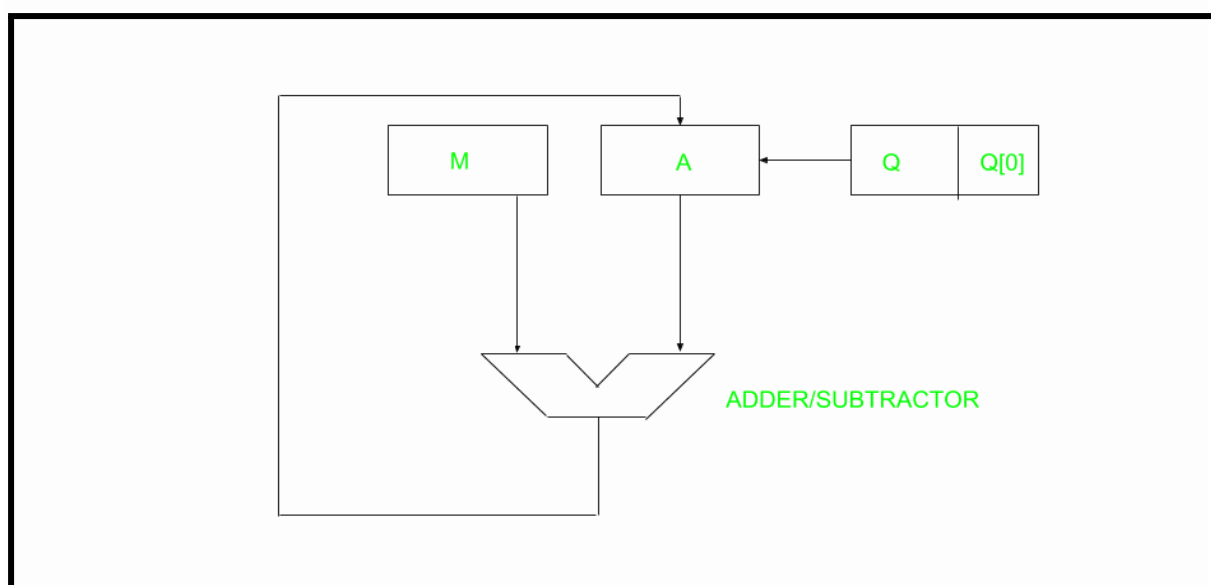
Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	08-10-2021	08-10-2021	

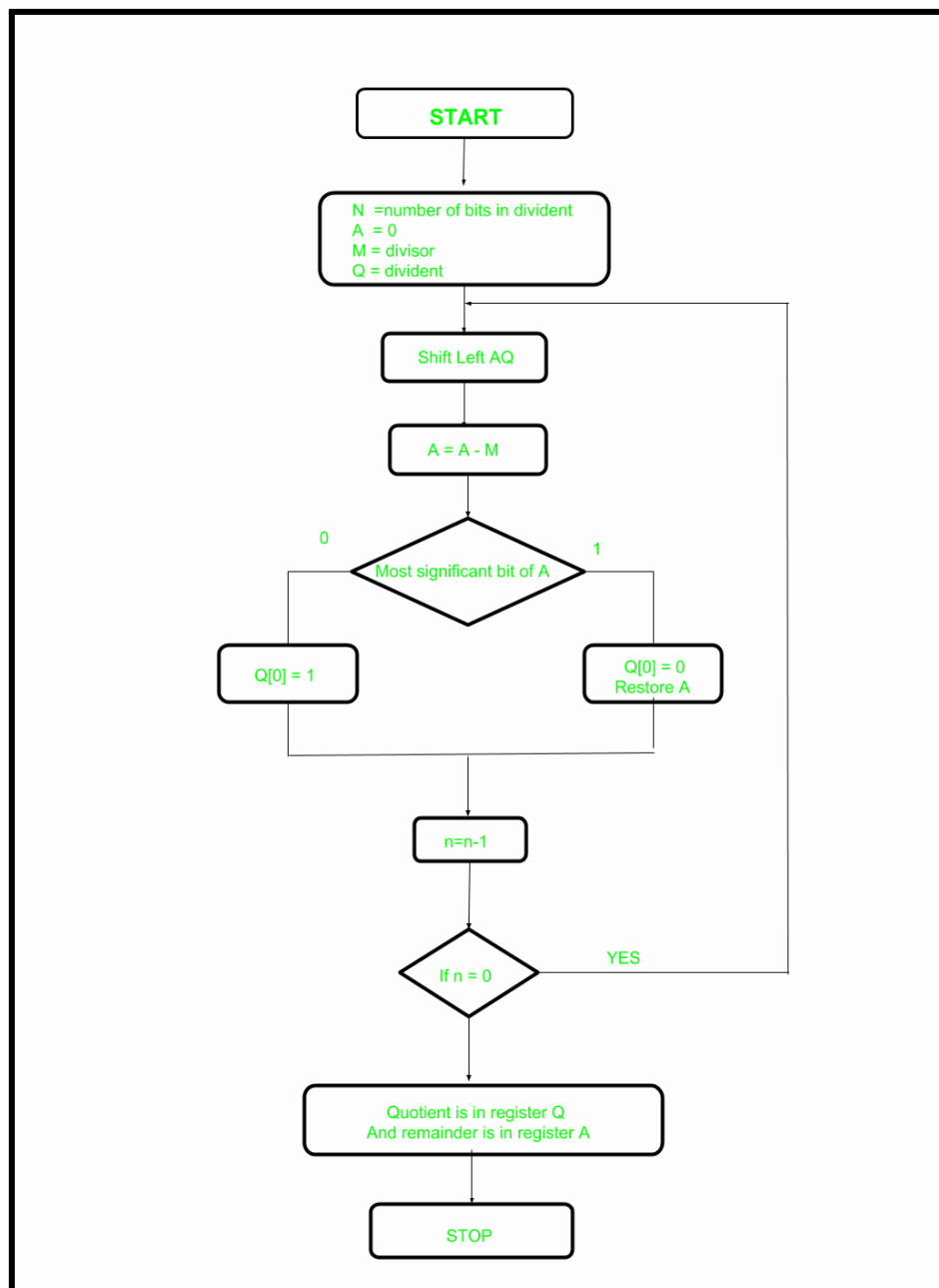
THEORY:

A division algorithm provides a quotient and a remainder when we divide two number. They are generally of two type **slow algorithm** and **fast algorithm**. Slow division algorithm are restoring, non-restoring, non-performing restoring, SRT algorithm and under fast comes Newton–Raphson and Goldschmidt.

In this article, will be performing restoring algorithm for unsigned integer. Restoring term is due to fact that value of register A is restored after each iteration.



Here, register Q contain quotient and register A contain remainder. Here, n-bit dividend is loaded in Q and divisor is loaded in M. Value of Register is initially kept 0 and this is the register whose value is restored during iteration due to which it is named Restoring.





Let's pick the step involved:

- **Step-1:** First the registers are initialized with corresponding values (Q = Dividend, M = Divisor, $A = 0$, n = number of bits in dividend)
- **Step-2:** Then the content of register A and Q is shifted right as if they are a single unit
- **Step-3:** Then content of register M is subtracted from A and result is stored in A
- **Step-4:** Then the most significant bit of the A is checked if it is 0 the least significant bit of Q is set to 1 otherwise if it is 1 the least significant bit of Q is set to 0 and value of register A is restored i.e the value of A before the subtraction with M
- **Step-5:** The value of counter n is decremented
- **Step-6:** If the value of n becomes zero we get of the loop otherwise we repeat fro step 2
- **Step-7:** Finally, the register Q contain the quotient and A contain remainder



Examples:

Perform Division Restoring Algorithm

Dividend = 11

Divisor = 3

n	M	A	Q	Operation
4	00011	00000	1011	initialize
	00011	00001	011_	shift left AQ
	00011	11110	011_	A=A-M
	00011	00001	0110	Q[0]=0 And restore A
3	00011	00010	110_	shift left AQ
	00011	11111	110_	A=A-M
	00011	00010	1100	Q[0]=0
2	00011	00101	100_	shift left AQ
	00011	00010	100_	A=A-M
	00011	00010	1001	Q[0]=1
1	00011	00101	001_	shift left AQ
	00011	00010	001_	A=A-M
	00011	00010	0011	Q[0]=1

Remember to restore the value of A, most significant bit of A is 1. As that register Q contains the quotient, i.e. 3 and register A contains remainder 2.

CODE:

```
#include<stdio.h>
#include<conio.h>
#include<math.h>
```



```
int getSize(int x)
{
    int c;
    if(x<=1)
        c = 2;
    else if(x < 4)
        c = 2;
    else if(x< 8)
        c = 3;
    else if(x< 16)
        c = 4;
    else if(x< 32)
        c = 5;
    else if(x< 64)
        c = 6;
    else if(x< 128)
        c = 7;
    else if(x< 256)
        c = 8;
    else if(x< 512)
        c = 9;
    return c;
}

int max(int x,int y)
{
    if(x< y)
        return(y);
    else
        return(x);
}

void main()
{
    int B,Q,Z,M,c,c1,e,f,g,h,i,j,x,y,ch,in,S,G,P;
    int a[24],b[12],b1[12],q[12],carry=0,count=0,option;
    long num;

    printf("\n\t\tPROGRAM FOR RESTORING DIVISION\t\t\n");
```



```
printf("\n\nENTER DIVIDEND\t: ");
scanf("%d",&Q);
y = getsize(Q);
printf("ENTER DIVISOR\t: ");
scanf("%d",&M);
x = getsize(M);
Z = max(x,y);
printf("\n\tTOTAL BITS CONSIDERED FOR RESULT => %d",2*Z+1);
printf("\n\tINITIALLY A IS RESET TO ZERO:");
for(i=0;i<=Z;i++)
printf("%d ",a[i]=0);
for(i=Z;i>=0;i--)
{
b1[i] = b[i] = M%2;
M = M/2;
b1[i] = 1-b1[i];
}
carry = 1;
for(i=Z;i>=0;i--)
{
c1 = b1[i]^carry;
carry = b1[i]&&carry;
b1[i]=c1;
}
for(i=2*Z;i>Z;i--)
{
a[i] = Q%2;
Q = Q/2;
}
printf("\n\n\tDivisor\t\t(M)\t: ");
for(i=0;i<=Z;i++)
printf("%d ",b[i]);
printf("\n\t2'C Divisor\t(-M)\t: ");
for(i=0;i<=Z;i++)
printf("%d ",b1[i]);
printf("\n\tDividend\t(Q)\t: ");
for(i=Z+1;i<=2*Z;i++)
printf("%d ",a[i]);
printf("\n\n\tBITS CONSIDERED:[ A ]\t [ M ]");
```




```
printf("\n\t\t\t");
for(i=0;i<=Z;i++)
printf("%d ",a[i]);
printf(" ");
for(i=Z+1;i<=2*Z;i++)
printf("%d ",a[i]);
count = Z;
do{
for(i=0;i<2*Z;i++)
a[i] = a[i+1];
printf("\n\nLeft Shift\t\t");
for(i=0;i<=Z;i++)
printf("%d ",a[i]);
printf(" ");
for(i=Z+1;i< 2*Z;i++)
printf("%d ",a[i]);
carry=0;
for(i=Z;i>=0;i--)
{
S=a[i]^(b1[i]^carry);
G=a[i]&&b1[i];
P=a[i]^b1[i];
carry=G||(P&&carry);
a[i]=S ;
}
printf("\nA< -A-M \t\t");
for(i=0;i<=Z;i++)
printf("%d ",a[i]);
printf(" ");
for(i=Z+1;i< 2*Z;i++)
printf("%d ",a[i]);
ch=a[0];
printf("\nBIT Q:%d",ch);
switch (ch)
{
case 0: a[2*Z]=1;
printf(" Q0< -1\t\t");
for(i=0;i<=Z;i++)
printf("%d ",a[i]);
```



```
printf(" ");
for(i=Z+1;i<=2*Z;i++)
printf("%d ",a[i]);
break;

case 1: a[2*Z]=0;
printf(" Q0< -0\t\t");
for(i=0;i<=Z;i++)
printf("%d ",a[i]);
printf(" ");
for(i=Z+1;i< 2*Z;i++)
printf("%d ",a[i]);
carry=0;
for(i=Z;i>=0;i--)
{
S=a[i]^(b[i]^carry);
G=a[i]&&b[i];
P=a[i]^b[i];
carry=G|(P&&carry);
a[i]=S ;
}
printf("\nA< -A+M");
printf("\t\t\t");
for(i=0;i<=Z;i++)
printf("%d ",a[i]);
printf(" ");
for(i=Z+1;i<=2*Z;i++)
printf("%d ",a[i]);
break;
}
count--;
}while(count!=0);
num=0;
printf("\n\n\t\tQUOTIENT IN BITS :");
for(i=Z+1;i<=2*Z;i++)
{
printf("%d ",a[i]);
num=num+pow(2,2*Z-i)*a[i];
}
```



```
printf("\n\t\tQUOTIENT IN DECIMAL :%ld",num);
num=0;
printf("\n\t\tREMAINDER IN BITS :");
for(i=0;i<=Z;i++)
{
printf("%d ",a[i]);
num=num+pow(2,Z-i)*a[i];
}
printf("\n\t\tREMAINDER IN DECIMAL :%ld",num);
getch();

getch();
}
```



OUTPUT:

```
PROGRAM FOR RESTORING DIVISION

ENTER DIVIDEND : 17
ENTER DIVISOR : 4

TOTAL BITS CONSIDERED FOR RESULT => 11
INITIALLY A IS RESET TO ZERO:0 0 0 0 0 0

Divisor      (M)      : 0 0 0 1 0 0
2'C Divisor  (-M)     : 1 1 1 1 0 0
Dividend     (Q)      : 1 0 0 0 1

BITS CONSIDERED:[ A ]      [ M ]
                0 0 0 0 0 0 1 0 0 0 1

Left Shift      0 0 0 0 0 1 0 0 0 1
A< -A-M         1 1 1 1 0 1 0 0 0 1
BIT Q:1 Q0< -0  1 1 1 1 0 1 0 0 0 1
A< -A+M         0 0 0 0 0 1 0 0 0 1 0

Left Shift      0 0 0 0 1 0 0 0 1 0
A< -A-M         1 1 1 1 1 0 0 0 1 0
BIT Q:1 Q0< -0  1 1 1 1 1 0 0 0 1 0
A< -A+M         0 0 0 0 1 0 0 0 1 0 0

Left Shift      0 0 0 1 0 0 0 1 0 0
A< -A-M         0 0 0 0 0 0 0 1 0 0
BIT Q:0 Q0< -1  0 0 0 0 0 0 0 1 0 0 1

Left Shift      0 0 0 0 0 0 1 0 0 1
A< -A-M         1 1 1 1 0 0 1 0 0 1
BIT Q:1 Q0< -0  1 1 1 1 0 0 1 0 0 1
A< -A+M         0 0 0 0 0 0 1 0 0 1 0

Left Shift      0 0 0 0 0 1 0 0 1 0
A< -A-M         1 1 1 1 0 1 0 0 1 0
BIT Q:1 Q0< -0  1 1 1 1 0 1 0 0 1 0
A< -A+M         0 0 0 0 0 1 0 0 1 0 0

QUOTIENT IN BITS :0 0 1 0 0
QUOTIENT IN DECIMAL :4
REMAINDER IN BITS :0 0 0 0 0 1
REMAINDER IN DECIMAL :1
```

CONCLUSION: From this experiment, we learn how to use the restoring division algorithm to divide unsigned bits. We understood that it is a type of slow algorithm along with non-restoring division. We also learn how to write the code for the same algorithm in C language and implement it successfully.



EXPERIMENT - 3

Aim: To study and implement non-restoring division algorithms.

Submission Sheet

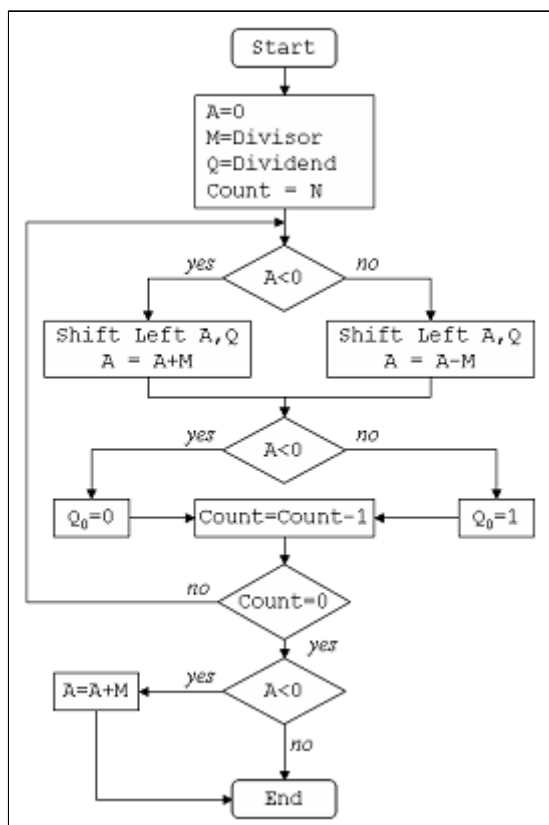
SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	30-10-2021	30-10-2021	

Theory

A division algorithm is an algorithm which, given two integers N and D, computes their quotient and/or remainder, the result of Euclidean division. Division algorithms fall into two main categories: slow division and fast division. Slow division algorithms produce one digit of the final quotient per iteration. Examples of slow division include restoring, non-performing restoring, non-restoring, and SRT division. Fast division methods start with a close approximation to the final quotient and produce twice as many digits of the final quotient on each iteration. Newton–Raphson and Goldschmidt algorithms fall into this category.

Non-Restoring division, it is less complex than the restoring one because simpler operations are involved i.e. addition and subtraction, also now the restoring step is performed. In the method, rely on the sign bit of the register which initially contains zero named as A.

Here is the flow chart given below.



Algorithm

- 1) Set the value of register A as 0 (N bits)
- 2) Set the value of register M as Divisor (N bits)
- 3) Set the value of register Q as Dividend (N bits)
- 4) Concatenate A with Q {A,Q}
- 5) Repeat the following "N" number of times (here N is no. of bits in divisor):
 - If the sign bit of A equals 0,
 - shift A and Q combined left by 1 bit and subtract M from A,
 - else shift A and Q combined left by 1 bit and add M to A
 - Now if sign bit of A equals 0, then set Q[0] as 1, else set Q[0] as 0
- 6) Finally if the sign bit of A equals 1 then add M to A.
- 7) Assign A as remainder and Q as quotient.

Example:

Dividend (A) = 101110, ie 46, and Divisor (B) = 010111, ie 23.



Initialization :

Set Register A = Dividend = 000000

Set Register Q = Dividend = 101110

(So AQ = 000000 101110 , Q₀ = LSB of Q = 0)

Set M = Divisor = 010111, M' = 2's complement of M = 101001

Set Count = 6, since 6 digits operation is being done here.

After this we start the algorithm, which I have shown in a table below :

In the table, SHL(AQ) denotes shift left AQ by one position leaving Q₀ blank.

Similarly, a square symbol in Q₀ position denote, it is to be calculated later

Action	A	Q	Count
Initial	000 000	101 110	6
A > 0 => SHL (AQ)	000 001	011 10□	
A = A-M	101 010	011 10□	
A < 0 => Q ₀ = 0	101 010	011 100	5
A < 0 => SHL (AQ)	010 100	111 00□	
A = A+M	101 011	111 00□	
A < 0 => Q ₀ = 0	101 011	111 000	4
A < 0 => SHL (AQ)	010 111	110 00□	
A = A+M	101 110	110 00□	
A < 0 => Q ₀ = 0	101 110	110 000	3
A < 0 => SHL (AQ)	011 101	100 00□	
A = A+M	110 100	100 00□	
A < 0 => Q ₀ = 0	110 100	100 000	2
A < 0 => SHL (AQ)	101 001	000 00□	
A = A+M	000 000	000 00□	
A < 0 => Q ₀ = 1	000 000	000 001	1
A > 0 => SHL (AQ)	000 000	000 01□	
A = A+M	101 001	000 01□	
A < 0 => Q ₀ = 1	101 001	000 010	0
Count has reached Zero, So final steps			
A < 0 => A = A+M	000 000	000 010	
	Reminder	Quotient	



CODE:

```
#include<stdio.h>
#include<malloc.h>
int *a,*q,*m,*mc,*c,n,d;
int powr(int x,int y)
{
    int s=1,i;
    for(i=0;i<y;i++)
        s=s*x;
    return s;
}
void print(int arr[],int n)
{
    int i;
    for(i=0;i<n;i++)
        printf("%d ",arr[i]);
}
void bin(int n, int arr[]){
    int r, i = 0;
    do{
        r = n % 2;
        n /= 2;
        arr[i] = r;
        i++;
    }while(n > 0);
}
void set(int array[], int x){
    int i,tmp[20]={0};
    for(i = x -1; i >=0; i--)
        tmp[x-1-i]=array[i];
    for(i=0;i<x;i++)
        array[i]=tmp[i];
}
int len(int x)
{
    int i=0;
    while(powr(2,i)<=x) i++;
    return ++i;
}
```




```
void addBinary(int a1[], int a2[])
{
    int bi[2]={0},ca[20]={0};
    int t=len(n),tmp=0;
    int *su=(int*)malloc(sizeof(int)*len(n));
    while(t-->0)
    {
        tmp=a1[t]+a2[t]+ca[t];
        bin(tmp,bi);
        su[t]=bi[0];
        ca[t-1]=bi[1];
        bi[0]=0;bi[1]=0;
    }
    for(t=0;t<len(n);t++)
        a1[t]=su[t];
    free(su);
}

void twoCom(int arr[]){
    int i;
    int *one=(int*)malloc(sizeof(int)*len(n));
    for(i=0;i<len(n)-1;i++)
        one[i]=0;
    one[i]=1;
    for(i = 0; i < len(n); i++){
        arr[i]=1-arr[i];
    }
    addBinary(arr, one);
    free(one);
}

void ls(int alen,int blen)
{
    int i=0;
    for(i=0;i<alen-1;i++)
        a[i]=a[i+1];
    a[i]=q[0];
    for(i=0;i<blen-1;i++)
        q[i]=q[i+1];
    q[i]=-1;
}
```



```

void printaq()
{
    print(a,len(n));
    printf("\t");
    print(q,len(n)-1);
    printf("\t");
    printf("\n");
}

int main()
{
    int i,cnt=0;
    printf("Enter The Numerator/Denominator: ");
    scanf("%d/%d",&n,&d);
    q=(int*)malloc(sizeof(int)*len(n)-1);
    bin (n,q);
    m=(int*)malloc(sizeof(int)*(len(n)));
    bin(d,m);
    a=(int*)malloc(sizeof(int)*(len(n)));
    for(i=0;i<len(n);i++)
        a[i]=0;
    mc=(int*)malloc(sizeof(int)*(len(n)));
    bin(d,mc);
    set(q,len(n)-1);
    set(m,len(n));
    set(mc,len(n));
    twoCom(mc);
    cnt=len(n)-1;
    printf("\t    A\t\t Q\t\t M\t    Count\n");
    printf("\t-----\t-----\t\t----- \n");
    while(cnt>0)
    {
        printf("\t");
        print(a,len(n));
        printf("\t");
        print(q,len(n)-1);
        printf("\t");
        print(m,len(n));
        printf("\t%d\n",cnt);
        if(a[0]==1)
        {

```



```
ls(len(n),len(n)-1);
printf("LSHIFT\t");
printaq();
addBinary(a,m);
printf("A=A+M\t");
printaq();
}
else
{
ls(len(n),len(n)-1);
printf("LSHIFT\t");
printaq();
addBinary(a,mc);
printf("A=A-M\t");
printaq();
}
if(a[0]==1)
{
q[len(n)-2]=0;
addBinary(a,m);
}
else
q[len(n)-2]=1;
printf("A=A+M\t");
printaq();
cnt-=1;
printf("\n");
}
return 0;
}
```



OUTPUT:

Enter The Numerator/Denominator: 10/3				
	A	Q	M	Count
	-----	-----	-----	-----
	0 0 0 0 0	1 0 1 0	0 0 0 1 1	4
LSHIFT	0 0 0 0 1	0 1 0 -1		
A=A-M	1 1 1 1 0	0 1 0 -1		
A=A+M	0 0 0 0 1	0 1 0 0		
	0 0 0 0 1	0 1 0 0	0 0 0 1 1	3
LSHIFT	0 0 0 1 0	1 0 0 -1		
A=A-M	1 1 1 1 1	1 0 0 -1		
A=A+M	0 0 0 1 0	1 0 0 0		
	0 0 0 1 0	1 0 0 0	0 0 0 1 1	2
LSHIFT	0 0 1 0 1	0 0 0 -1		
A=A-M	0 0 0 1 0	0 0 0 -1		
A=A+M	0 0 0 1 0	0 0 0 1		
	0 0 0 1 0	0 0 0 1	0 0 0 1 1	1
LSHIFT	0 0 1 0 0	0 0 1 -1		
A=A-M	0 0 0 0 1	0 0 1 -1		
A=A+M	0 0 0 0 1	0 0 1 1		

Conclusion :

The non-restorative division algorithm is an efficient way to perform binary division compared to traditional subtractive based algorithms by using the faster processed bit shift commands in the CPU registers. The algorithm is simple enough to be implemented in hardware in equipment like Arithmometers while also generalising to complex modern day systems. The algorithm serves as a good example in showing that considering lower-level system dependencies and physical limitations can be used to optimize algorithms. Non-restorative algorithm is more efficient than restorative algorithm as it uses simpler commands in terms of addition and subtraction, however it is slower than other algorithms.



EXPERIMENT - 4

Aim:

Implement Sequential memory organization with following details:

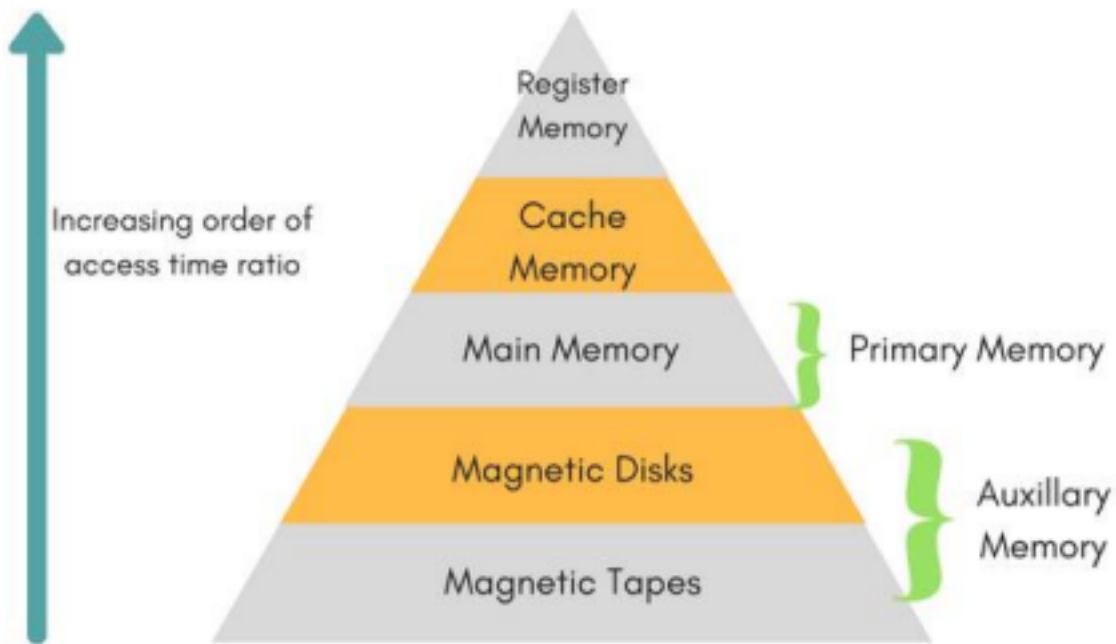
- Processor can access 1 word = 4 bytes.
- L1 cache can store max = 32 words.
- L2 cache can store max = 128 words.
- Main memory can store max = 2048 bytes.

Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	30-10-2021	30-10-2021	

Theory:

In Computer System Design, Memory Hierarchy is an enhancement to organize the memory such that it can minimize the access time. The Memory Hierarchy was developed based on a program behaviour known as locality of references. The figure below clearly demonstrates the different levels of memory hierarchy:



This Memory Hierarchy Design is divided into 2 main types:

- External Memory or Secondary Memory – Comprising of Magnetic Disk, Optical Disk, Magnetic Tape i.e., peripheral storage devices which are accessible by the processor via I/O Module.
- Internal Memory or Primary Memory – Comprising of Main Memory, Cache Memory & CPU registers. This is directly accessible by the processor.

Levels of memory:

- Level 1 or Register –

It is a type of memory in which data is stored and accepted that are immediately stored in CPU. Most commonly used register is accumulator, Program counter, address register etc.

- Level 2 or Cache memory –

It is the fastest memory which has faster access time where data is temporarily stored for faster access.

- Level 3 or Main Memory –

It is memory on which computer works currently. It is small in size and once power is off data no longer stays in this memory.



- Level 4 or Secondary Memory –

It is external memory which is not as fast as main memory but data stays permanently in this memory.

Types of Cache:

- Primary Cache –

A primary cache is always located on the processor chip. This cache is small and its access time is comparable to that of processor registers.

- Secondary Cache –

Secondary cache is placed between the primary cache and the rest of the memory. It is referred to as the level 2 (L2) cache. Often, the Level 2 cache is also housed on the processor chip.

Code:

```
memory.cpp
#include<bits/stdc++.h>
#include<time.h>
using namespace std;
int main() {
    vector<double> lc1(32,-1);
    vector<double> lc2(64,-1);

    double lc1Time = 20;
    double lc2Time = 60;
    double mainMemoryTime = 120;
    double totalHits = 0;
    double lc1Hits = 0;
    double lc2Hits = 0;
    double count = 0;
    int fr1 = -1;
    int re1 = 0;
    int fr2 = -1;
    int re2 = 0;
    char key = 100;
    srand(time(0));
    cout<<"Sequential Memory Organization:"<<endl<<endl;
    cout<<" "<<"Requirement"<<setw(22)<<"Location of Hit"<<" "<<setw(15)<<"Avg.
Access Time"<<endl;
    while(key--) {
        count++;
```



```
double avgTime = 0;
double input = rand()%512;
bool f1 = false;
bool f2 = false;
for(double i=0; i<32; i++) {
    if(lc1[i]==input) {
        lc1Hits++;
        f1 = true;
        f2 = true;
        avgTime += (lc1Hits/(count))*lc1Time + (1-
(lc1Hits/(count)))*(lc2Hits/(count-lc1Hits))*(lc1Time+lc2Time) + (1-
(lc1Hits/(count)))*(1-(lc2Hits/(count
lc1Hits)))*1*(lc1Time+lc2Time+mainMemoryTime);

    cout<<setw(12)<<input<<" "<<setw(20)<<"L1 Cache"<<"
"<<setw(8)<<avgTime<<endl;
    }
}
if(!f1) {
    double b = input/2;
    for(double i=0; i<64; i++) {
        if(lc2[i]==b) {
            lc2Hits++;
            f2 = true;
            // found in 2
            if(re1==fr1) {
                re1 = (re1+1)%32;
            }
            fr1 = (fr1+1)%32;
            lc1[fr1] = input;
            avgTime += (lc1Hits/(count))*lc1Time + (1-
(lc1Hits/(count)))*(lc2Hits/(count-lc1Hits))*(lc1Time+lc2Time) + (1-
(lc1Hits/(count)))*(1-(lc2Hits/(count
lc1Hits)))*1*(lc1Time+lc2Time+mainMemoryTime);
            cout<<setw(12)<<input<<" "<<setw(20)<<"L2 Cache"<<" "<<setw(8)<<avgTime<<endl;
        }
    }
}
if(!f2) {
    double b = input/2;
    if(re2==fr2) {
        re2 = (re2+1)%64;
    }
    fr2 = (fr2+1)%64;
    lc2[fr2] = b;
    if(re1==fr1) {
        re1 = (re1+1)%32;
```




```

}

fr1 = (fr1+1)%32;
lc1[fr1] = input;
avgTime += (lc1Hits/(count))*lc1Time + (1-
(lc1Hits/(count)))*(lc2Hits/(count-lc1Hits))*(lc1Time+lc2Time) + (1-
(lc1Hits/(count)))*(1-(lc2Hits/(count
lc1Hits)))*1*(lc1Time+lc2Time+mainMemoryTime);
cout<<setw(12)<<input<<" "<<setw(20)<<"Main Memory"<<"
"<<setw(8)<<avgTime<<endl;
}
}
cout<<endl;
cout<<"L1 Hit Ratio (H1) = "<<lc1Hits<<"/"<<count<<" =
"<<(double)lc1Hits/count<<endl;
cout<<"L1 Access Time (T1) = "<<lc1Time<<" ns"<<endl<<endl;
cout<<"L2 Hit Ratio (H2) = "<<lc2Hits<<"/"<<(count-lc1Hits)<<" =
"<<(double)lc2Hits/(count-lc1Hits)<<endl;
cout<<"L2 Access Time (T2) = "<<lc2Time<<" ns"<<endl<<endl;
cout<<"Main Memory Hit Ratio (Hm) = "<<1<<endl;
cout<<"Main memory Access Time = "<<mainMemoryTime<<" ns"<<endl<<endl;
cout<<"Average Access Time = [H1*T1] + [(1-H1)*H2*(T1+T2)] + [(1-H1)*(1-
H2)*Hm*(T1+T2+Tm)]"<<endl;
double finalAns = ((double)lc1Hits/count)*lc1Time + (1-
((double)lc1Hits/count))*((double)lc2Hits/(100-lc1Hits))*(lc1Time+lc2Time) + (1-
((double)lc1Hits/count))*(1-((double)lc2Hits/(count
lc1Hits)))*1*(lc1Time+lc2Time+mainMemoryTime);
cout<<"Average Access Time = "<<finalAns<<" ns"<<endl;

return 0;
}

```

Output:

```

Sequential Memory Organization:
Requirement Location of Hit Avg. Access Time
353 Main Memory 200
372 Main Memory 200
89 Main Memory 200
189 Main Memory 200
265 Main Memory 200
423 Main Memory 200
379 Main Memory 200
107 Main Memory 200
8 Main Memory 200
217 Main Memory 200

```



155 Main Memory 200
294 Main Memory 200
192 Main Memory 200
47 Main Memory 200
348 Main Memory 200
217 L1 Cache 188.75
264 Main Memory 189.412
375 Main Memory 190
100 Main Memory 190.526
351 Main Memory 191
476 Main Memory 191.429
12 Main Memory 191.818
178 Main Memory 192.174
138 Main Memory 192.5
338 Main Memory 192.8
250 Main Memory 193.077
306 Main Memory 193.333
363 Main Memory 193.571
208 Main Memory 193.793
99 Main Memory 194
143 Main Memory 194.194
205 Main Memory 194.375
381 Main Memory 194.545
34 Main Memory 194.706
329 Main Memory 194.857
259 Main Memory 195
245 Main Memory 195.135
429 Main Memory 195.263
7 Main Memory 195.385
228 Main Memory 195.5
65 Main Memory 195.61
320 Main Memory 195.714
46 Main Memory 195.814
19 Main Memory 195.909
262 Main Memory 196
56 Main Memory 196.087
178 L1 Cache 192.34
132 Main Memory 192.5

252 Main Memory 192.653
189 L2 Cache 190.4
301 Main Memory 190.588
315 Main Memory 190.769
418 Main Memory 190.943
163 Main Memory 191.111
302 Main Memory 191.273
87 Main Memory 191.429



49 Main Memory 191.579
402 Main Memory 191.724
467 Main Memory 191.864
80 Main Memory 192
462 Main Memory 192.131
101 Main Memory 192.258
410 Main Memory 192.381
125 Main Memory 192.5
3 Main Memory 192.615
465 Main Memory 192.727
116 Main Memory 192.836
217 L2 Cache 191.176
368 Main Memory 191.304
71 Main Memory 191.429
27 Main Memory 191.549
262 L1 Cache 189.167
190 Main Memory 189.315
283 Main Memory 189.459
166 Main Memory 189.6
146 Main Memory 189.737
469 Main Memory 189.87
229 Main Memory 190
347 Main Memory 190.127
122 Main Memory 190.25
335 Main Memory 190.37
296 Main Memory 190.488
452 Main Memory 190.602
492 Main Memory 190.714
44 Main Memory 190.824
305 Main Memory 190.93
261 Main Memory 191.034
81 Main Memory 191.136
455 Main Memory 191.236
66 Main Memory 191.333
364 Main Memory 191.429
139 Main Memory 191.522
8 Main Memory 191.613
240 Main Memory 191.702
74 Main Memory 191.789
379 Main Memory 191.875
476 Main Memory 191.959
458 Main Memory 192.041
1 Main Memory 192.121
455 L1 Cache 190.4
L1 Hit Ratio (H1) = $4/100 = 0.04$
L1 Access Time (T1) = 20 ns



L2 Hit Ratio (H_2) = $2/96 = 0.0208333$

L2 Access Time (T_2) = 60 ns

Main Memory Hit Ratio (H_m) = 1

Main memory Access Time = 120 ns

Average Access Time = $[H_1 * T_1] + [(1-H_1) * H_2 * (T_1 + T_2)] + [(1-H_1) * (1-H_2) * H_m * (T_1 + T_2 + T_m)]$

Average Access Time = 190.4 ns

Conclusion:

We learnt about sequential memory organization and implemented it in C++ programming language successfully on the given problem statement in a 2-Level Cache Memory architecture for 100 different number words required by the processor.



EXPERIMENT - 5

AIM: To simulate a direct-mapped cache

Submission Sheet

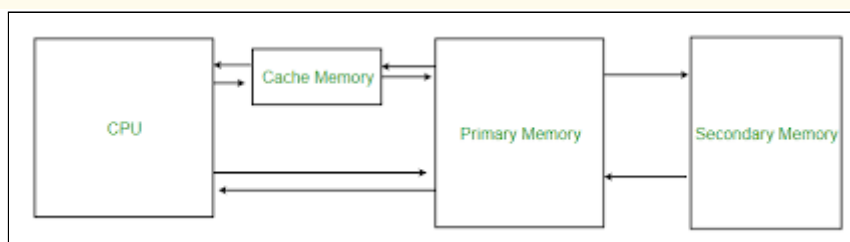
SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	17-12-2021	17-12-2021	

THEORY:

Cache memory, also called cache, is a supplementary memory system that temporarily stores frequently used instructions and data for quicker processing by the central processing unit (CPU) of a computer. The cache augments, and is an extension of, a computer's main memory. Both main memory and cache are internal random-access memories (RAMs) that use semiconductor-based transistor circuits. Cache holds a copy of only the most frequently used information or program codes stored in the main memory. The smaller capacity of the cache reduces the time required to locate data within it and provide it to the CPU for processing.

When a computer's CPU accesses its internal memory, it first checks to see if the information it needs is stored in the cache. If it is, the cache returns the data to the CPU. If the information is not in the cache, the CPU retrieves it from the main memory. Disk cache memory operates similarly, but the cache is used to hold data that have recently been written on, or retrieved from, a magnetic disk or other external storage device.

Hit ratio = $\text{hit} / (\text{hit} + \text{miss}) = \text{no. of hits} / \text{total accesses}$.



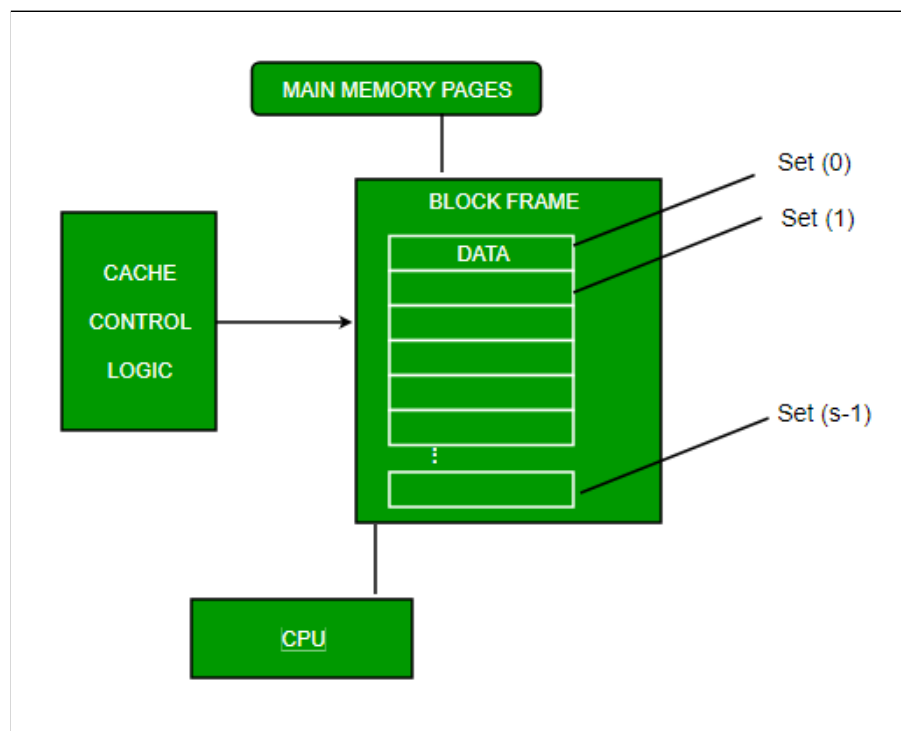
DIRECT CACHE MAPPING:

The simplest technique, known as direct mapping, maps each block of main memory into only one possible cache line. or

In Direct mapping, assign each memory block to a specific line in the cache. If a line is previously taken up by a memory block when a new block needs to be loaded, the old block is trashed. An address space is split into two parts: index field and a tag field. The cache is used to store the tag field whereas the rest is stored in the main memory. Direct mapping's performance is directly proportional to the Hit ratio.

```
i = j modulo m
where
i=cache line number
j= main memory block number
m=number of lines in the cache
```

For purposes of cache access, each main memory address can be viewed as consisting of three fields. The least significant w bits identify a unique word or byte within a block of main memory. In most contemporary machines, the address is at the byte level. The remaining s bits specify one of the 2^s blocks of main memory. The cache logic interprets these s bits as a tag of $s-r$ bits (most significant portion) and a line field of r bits. This latter field identifies one of the $m=2^r$ lines of the cache.



8 Bit:



CODE:

```
#include <stdio.h>

int tag[8];

int main( )
{
    int addr;
    int i, j, t;
    int hits, accesses;
    FILE *fp;

    fp = fopen("trace.txt", "r");
    hits = 0;
    accesses = 0;
    while (fscanf(fp, "%x", &addr) > 0) {
        /* simulate a direct-mapped cache with 8 words */
        accesses += 1;
        printf("%3d: 0x%08x ", accesses, addr);
        i = (addr >> 2) & 7;
        t = addr | 0x1f;
        if (tag[i] == t) {
            hits += 1;
            printf("Hit%d ", i);
        } else {
            /* allocate entry */
            printf("Miss ");
            tag[i] = t;
        }
        for (i = 0; i < 8; i++)
            printf("0x%08x ", tag[i]);
        printf("\n");
    }
    printf("Hits = %d, Accesses = %d, Hit ratio = %f\n", hits, accesses,
((float)hits)/accesses);
    close(fp);
}
```

OUTPUT:

```
1: 0x80000000 Miss 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
2: 0x80000004 Miss 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
```




```

72: 0x80000014 Miss 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000003f
73: 0x80000018 Hit6 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000003f
74: 0x8000000c Hit3 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000003f
75: 0x00000058 Miss 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x0000005f 0x0000003f
76: 0x80000010 Hit4 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x0000005f 0x0000003f
77: 0x80000014 Hit5 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x0000005f 0x0000003f
78: 0x80000018 Miss 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000003f
79: 0x8000000c Hit3 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000003f
80: 0x0000005c Miss 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
81: 0x80000010 Hit4 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
82: 0x80000014 Hit5 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
83: 0x80000018 Hit6 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
84: 0x8000000c Hit3 0x0000005f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
85: 0x00000060 Miss 0x0000007f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
86: 0x80000010 Hit4 0x0000007f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
87: 0x80000014 Hit5 0x0000007f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
88: 0x80000018 Hit6 0x0000007f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
89: 0x8000000c Hit3 0x0000007f 0x0000005f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
90: 0x00000064 Miss 0x0000007f 0x0000007f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
91: 0x80000010 Hit4 0x0000007f 0x0000007f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
92: 0x80000014 Hit5 0x0000007f 0x0000007f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
93: 0x80000018 Hit6 0x0000007f 0x0000007f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
94: 0x8000000c Hit3 0x0000007f 0x0000007f 0x0000005f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
95: 0x00000068 Miss 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
96: 0x80000010 Hit4 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
97: 0x80000014 Hit5 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
98: 0x80000018 Hit6 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
99: 0x8000000c Hit3 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
100: 0x0000006c Miss 0x0000007f 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
101: 0x80000010 Hit4 0x0000007f 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
102: 0x80000014 Hit5 0x0000007f 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
103: 0x80000018 Hit6 0x0000007f 0x0000007f 0x0000007f 0x0000007f 0x8000001f 0x8000001f 0x8000001f 0x0000005f
Hits = 68, Accesses = 103, Hit ratio = 0.660194

```

16 Bit:

CODE:

```

#include <stdio.h>

int tag[16];

int main()
{
    int addr;
    int i, j, t;
    int hits, accesses;
    FILE *fp;

    fp = fopen("trace.txt", "r");
    hits = 0;
    accesses = 0;

```



```

while (fscanf(fp, "%x", &addr) > 0) {
    /* simulate a direct-mapped cache with 8 words */
    accesses += 1;
    printf("%3d: 0x%08x ", accesses, addr);
    i = (addr >> 2) & 15;
    t = addr | 0x1f;
    if (tag[i] == t) {
        hits += 1;
        printf("Hit%d ", i);
    } else {
        /* allocate entry */
        printf("Miss ");
        tag[i] = t;
    }
    for (i = 0; i < 16; i++)
        printf("0x%08x ", tag[i]);
    printf("\n");
}
printf("Hits = %d, Accesses = %d, Hit ratio = %f\n", hits, accesses,
((float)hits)/accesses);
close(fp);
}

```

OUTPUT:

```

1: 0x80000000 Miss 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
2: 0x80000004 Miss 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
3: 0x80000008 Miss 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
4: 0x8000000c Miss 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
5: 0x00000020 Miss 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
6: 0x80000010 Miss 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
7: 0x80000014 Miss 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
8: 0x80000018 Miss 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
9: 0x8000000c Hit3 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
10: 0x00000024 Miss 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
11: 0x80000010 Hit4 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
12: 0x80000014 Hit5 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
13: 0x80000018 Hit6 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
14: 0x8000000c Hit3 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x8000001f 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000

```

[illegible]

[illegible]



32 Bit:

CODE:

```
#include <stdio.h>
#include <stdio.h>
int tag[32];
int main()
{
    int addr;
    int i, j, t;
    int hits, accesses;
    FILE *fp;

    fp = fopen("trace.txt", "r");
    hits = 0;
    accesses = 0;
    while (fscanf(fp, "%x", &addr) > 0) {
        /* simulate a direct-mapped cache with 8 words */
        accesses += 1;
        printf("%3d: 0x%08x ", accesses, addr);
        i = (addr >> 2) & 31;
        t = addr | 0x1f;
        if (tag[i] == t) {
            hits += 1;
            printf("Hit%d ", i);
        } else {
            /* allocate entry */
            printf("Miss ");
            tag[i] = t;
        }
        for (i = 0; i < 32; i++)
            printf("0x%08x ", tag[i]);
        printf("\n");
    }
    printf("Hits = %d, Accesses = %d, Hit ratio = %f\n", hits, accesses,
        ((float)hits)/accesses);
    close(fp);
}
```



OUTPUT:

[illegible]



Hits = 76, Accesses = 103, Hit ratio = 0.737864

Comparison Table

Number of Bits	Hits	Accesses	Hit Ratio
8	68	103	0.660194
16	72	103	0.699029
32	76	103	0.737864

CONCLUSION:

Cache memory is the closest to the CPU, it reduces the access time from Main Memory.

There are 3 ways to map cache to Main memory: Direct mapping, Associative mapping, and Set-Associative mapping. In Direct mapping, each memory block is assigned to a specific line in the cache. It is a simple way to map blocks but it has a lower cache hit ratio, as there is only one cache line available in a set. Every time a new memory is referenced to the same set, the cache line is replaced, this is the drawback of Direct Mapping.

As seen from the comparison table above which is derived from our program output, as we go on increasing the number of bits in cache memory, the hit ratio also increases.



EXPERIMENT - 6

Aim

To implement fully associative mapping and set associative mapping.

Submission Sheet

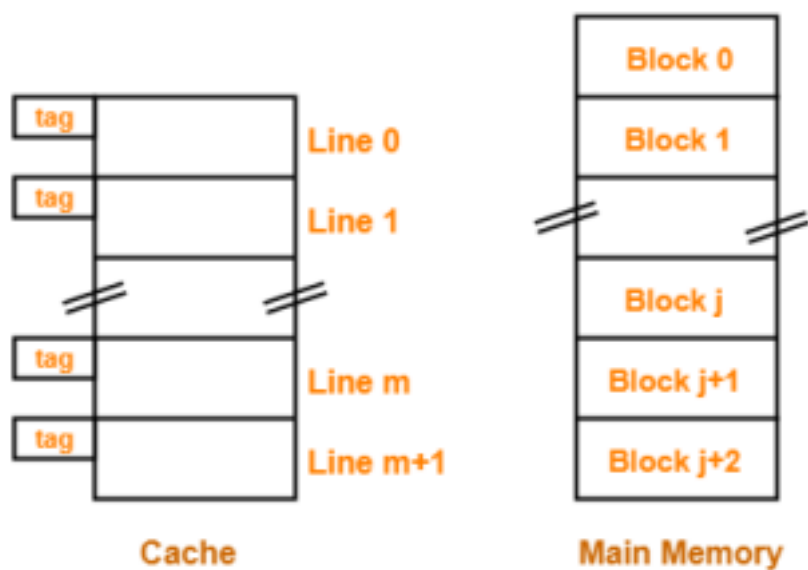
SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	6-11-2021	6-11-2021	

Theory

A CPU cache is a hardware used by the central processing unit (CPU) of a computer to reduce the average cost (time) of accessing data from the main memory. A cache is a smaller, faster memory, located closer to a processor core, which stores copies of the data from frequently used main memory locations. Cache memory is costlier than main memory or disk memory but economical than CPU registers. It is an extremely fast memory type that acts as a buffer between RAM and the CPU, and holds frequently requested data and instructions so that they are immediately available to the CPU when needed.

Fully associative mapping

Fully Associative Mapping is a cache mapping technique that allows to map a block of main memory to any freely available cache line. In this type of mapping, the associative memory is used to store content and addresses of the memory word. Any block can go into any line of the cache. This means that the word id bits are used to identify which word in the block is needed, but the tag becomes all of the remaining bits. This enables the placement of any word at any place in the cache memory. It is considered to be the fastest and the most flexible mapping form.



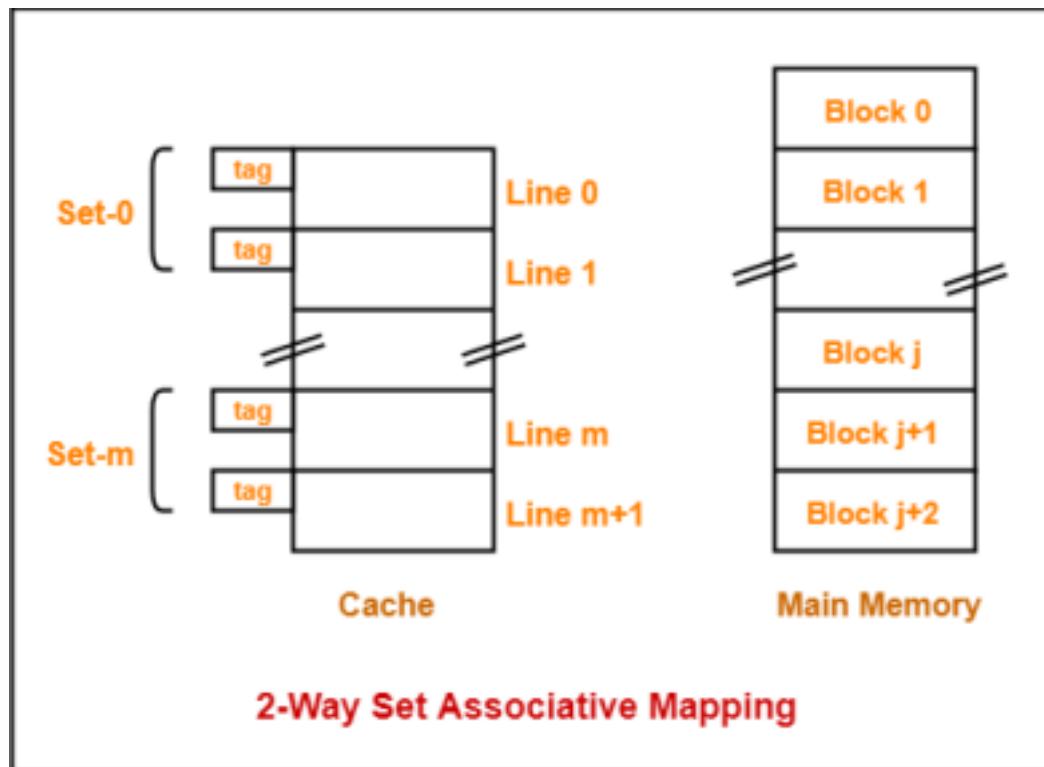
Fully Associative Mapping



Set associative mapping

Set-associative cache is a trade-off between direct-mapped cache and fully associative cache. This form of mapping is an enhanced form of direct mapping where the drawbacks of direct

mapping is removed. Set associative addresses the problem of possible thrashing in the direct mapping method. It does this by saying that instead of having exactly one line that a block can map to in the cache, we will group a few lines together creating a *set*. Then a block in memory can map to any one of the lines of a specific sea. Set-associative mapping allows that each word that is present in the cache can have two or more words in the main memory for the same index address. Set associative cache mapping combines the best of direct and associative cache mapping techniques.



Prelab code for generating trace file

```
main: addu $t0,$0,$0
      addiu $t1,$0,80
      addu $t2,$0,$0
loop: lw $t3,array($t0)
      addu $t2,$t2,$t3
```




```
addiu $t0,$t0,4
bne $t0,$t1,loop
*done: beq $0,$0,done
array: .word 1,2,3,4,5,6,7,8,9,10
       .word 11,12,13,14,15,16,17,18,19,20
```

Fully Associative Mapping – 8 bit

Code:

```
#include <stdio.h>
int tag[8];
int mru[8] = {7,6,5,4,3,2,1,0};
void mruUpdate(int index)
{
    int i;
    // find index in mru
    for (i = 0; i < 8; i++)
        if (mru[i] == index)
            break;
    // move earlier refs one later
    while (i > 0) {
        mru[i] = mru[i-1];
        i--;
    }
    mru[0] = index;
}

int main( )
{
    int addr;
    int i, j, t;
    int hits, accesses;
    FILE *fp;
    fp = fopen("trace.txt", "r");
    hits = 0;
    accesses = 0;
    while (fscanf(fp, "%x", &addr) > 0) {
        /* simulate fully associative cache with 8 words */
        accesses += 1;
        printf("%3d: 0x%08x ", accesses, addr);
        for (i = 0; i < 8; i++) {
            if (tag[i] == addr) {
                hits += 1;
                printf("Hit%d ", i);
                mruUpdate(i);
            }
        }
    }
}
```



```

break;
}
}
if (i == 8) {
/* allocate entry */
printf("Miss ");
i = mru[7];
tag[i] = addr;
mruUpdate(i);
}
for (i = 0; i < 8; i++)
printf("0x%08x ", tag[i]);
for (i = 0; i < 8; i++)

printf("%d ", mru[i]);
printf("\n");
}
printf("Hits = %d, Accesses = %d, Hit ratio = %f\n", hits, accesses,
((float)hits)/accesses);
pclose(fp);
}

```

Output:

```

79: 0x000000c Hit3 0x00000054 0x00000058 0x0000004c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 3 7 6 5 1 0 4 2
80: 0x0000005c Miss 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 2 3 7 6 5 1 0 4
81: 0x00000010 Hit5 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 5 2 3 7 6 1 0 4
82: 0x00000014 Hit6 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 6 5 2 3 7 1 0 4
83: 0x00000018 Hit7 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 7 6 5 2 3 1 0 4
84: 0x0000000c Hit3 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 3 7 6 5 2 1 0 4
85: 0x00000060 Miss 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 4 3 7 6 5 2 1 0
86: 0x00000010 Hit5 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 5 4 3 7 6 2 1 0
87: 0x00000014 Hit6 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 6 5 4 3 7 2 1 0
88: 0x00000018 Hit7 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 7 6 5 4 3 2 1 0
89: 0x0000000c Hit3 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 3 7 6 5 4 2 1 0
90: 0x00000054 Miss 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 0 3 7 6 5 4 2 1
91: 0x00000010 Hit5 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 5 0 3 7 6 4 2 1
92: 0x00000014 Hit6 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 6 5 0 3 7 4 2 1
93: 0x00000018 Hit7 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 7 6 5 0 3 4 2 1
94: 0x0000000c Hit3 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 3 7 6 5 0 4 2 1
95: 0x00000068 Miss 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 1 3 7 6 5 0 4 2
96: 0x00000010 Hit5 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 5 1 3 7 6 0 4 2
97: 0x00000014 Hit6 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 6 5 1 3 7 0 4 2
98: 0x00000018 Hit7 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 7 6 5 1 3 0 4 2
99: 0x0000000c Hit3 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 3 7 6 5 1 0 4 2
100: 0x0000006c Miss 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 2 3 7 6 5 1 0 4
101: 0x00000010 Hit5 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 5 2 3 7 6 1 0 4
102: 0x00000014 Hit6 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 6 5 2 3 7 1 0 4
103: 0x00000018 Hit7 0x00000054 0x00000058 0x0000005c 0x0000000c 0x00000050 0x00000010 0x00000014 0x00000018 7 6 5 2 3 1 0 4
Hits = 76, Accesses = 103, Hit ratio = 0.737864

```

Fully Associative Mapping – 16 bit**Code:**

```

#include <stdio.h>
int tag[16];
int mru[16] = {15,14,13,12,11,10,9,8,7,6,5,4,3,2,1,0};
void mruUpdate(int index)

```



```
{
int i;
// find index in mru
for (i = 0; i < 16; i++)
if (mru[i] == index)
break;
// move earlier refs one later
while (i > 0) {
mru[i] = mru[i-1];

i--;
}
mru[0] = index;
}
int main( )
{
int addr;
int i, j, t;
int hits, accesses;
FILE *fp;
fp = fopen("trace.txt", "r");
hits = 0;
accesses = 0;
while (fscanf(fp, "%x", &addr) > 0) {
/* simulate fully associative cache with 16 words */
accesses += 1;
printf("%3d: 0x%08x ", accesses, addr);
for (i = 0; i < 16; i++) {
if (tag[i] == addr) {
hits += 1;
printf("Hit%d ", i);
mruUpdate(i);
break;
}
}
if (i == 16) {
/* allocate entry */
printf("Miss ");
i = mru[15];
tag[i] = addr;
mruUpdate(i);
}

for (i = 0; i < 16; i++)
printf("0x%08x ", tag[i]);
printf("\n ");
}
```



```
for (i = 0; i < 16; i++)
printf("%d ", mru[i]);
printf("\n");
}
printf("Hits = %d, Accesses = %d, Hit ratio = %f\n", hits, accesses,
((float)hits)/accesses);
close(fp);
}
```

**Output: Fully Associative Mapping – 16 bit**

```

96: 0x00000010 Hit5 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
5 13 3 7 6 12 11 10 9 8 4 2 1 0 15 14
97: 0x00000014 Hit5 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
6 5 13 3 7 12 11 10 9 8 4 2 1 0 15 14
98: 0x00000018 Hit7 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
7 6 5 13 3 12 11 10 9 8 4 2 1 0 15 14
99: 0x00000000 Hit1 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
3 7 6 5 13 12 11 10 9 8 4 2 1 0 15 14
100: 0x0000000c Miss 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
14 3 7 6 5 13 12 11 10 9 8 4 2 1 0 15
101: 0x00000010 Hit5 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
5 14 3 7 6 13 12 11 10 9 8 4 2 1 0 15
102: 0x00000014 Hit6 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
6 5 14 3 7 13 12 11 10 9 8 4 2 1 0 15
103: 0x00000018 Hit7 0x00000044 0x0000004E 0x0000004C 0x0000000C 0x00000050 0x00000010 0x00000014 0x00000018 0x00000054 0x0000005E 0x00000000
5c 0x00000060 0x00000064 0x00000068 0x0000003c 0x00000040
7 6 5 14 3 13 12 11 10 9 8 4 2 1 0 15
Hits = 76, Accesses = 103, Hit ratio = 0.737864

77: 0x00000034 Hit10 0x00000000 0x00000004 0x00000000 0x00000000 0x0000000c 0x0000001c 0x00000020 0x00000024 0x00000028 0x0000001c 0x00000030 0x00000000
034 0x00000038 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
10 9 8 7 6 11 5 4 3 2 1 0 15 14 13 12
78: 0x00000038 Hit11 0x00000000 0x00000004 0x00000000 0x00000000 0x0000001c 0x00000020 0x00000024 0x00000028 0x0000001c 0x00000030 0x00000000
034 0x00000038 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
11 10 9 8 7 6 5 4 3 2 1 0 15 14 13 12
79: 0x00000024 Hit6 0x00000000 0x00000004 0x00000000 0x00000000 0x0000001c 0x00000020 0x00000024 0x0000001c 0x00000030 0x00000000
34 0x00000038 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
6 11 10 9 8 7 5 4 3 2 1 0 15 14 13 12
80: 0x00000028 Hit7 0x00000000 0x00000004 0x00000000 0x00000000 0x0000001c 0x00000020 0x00000024 0x0000001c 0x00000030 0x00000000
34 0x00000038 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
7 6 11 10 9 8 5 4 3 2 1 0 15 14 13 12
81: 0x0000003c Miss 0x00000000 0x00000004 0x00000000 0x00000000 0x0000001c 0x00000020 0x00000024 0x0000001c 0x00000030 0x00000000
34 0x00000038 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
12 7 6 11 10 9 8 5 4 3 2 1 0 15 14 13
82: 0x00000040 Miss 0x00000000 0x00000004 0x00000000 0x00000000 0x0000001c 0x00000020 0x00000024 0x0000001c 0x00000030 0x00000000
34 0x00000038 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000 0x00000000
13 12 7 6 11 10 9 8 5 4 3 2 1 0 15 14
Hits = 66, Accesses = 82, Hit ratio = 0.829268

```

Set Associative Mapping – 16 bit**Code:**

```

#include <stdio.h>
#include <stdlib.h>
#include <math.h>
void main() {
int** tags = malloc(2 * sizeof(int*));
tags[0] = malloc(16 * sizeof(int));
tags[1] = malloc(16 * sizeof(int));
int mru[8] = {1,1,1,1,1,1,1,1};
int sets = 8;
int addr;

int hits, accesses;
FILE *fp;
fp = fopen("trace.txt", "r");
hits = 0;
accesses = 0;
while (fscanf(fp, "%x", &addr) > 0) {
accesses++;
int setNum = (addr >> 2) & ((int)(pow(2, log2(sets)) - 1));

```



```

int addrTag = addr >> (2 + (int)log2(sets));
printf("\n%3d: 0x%08x ", accesses, addr);
int i;
for(i = 0; i < 2; ++i) {
if(tags[i][setNum] == addrTag) {
printf("Hit %d", i);
hits++;
mru[setNum] = i;
break;
}
}
if(i == 2) {
printf("Miss");
tags[(mru[setNum] + 1) % 2][setNum] = addrTag;
mru[setNum] = (mru[setNum] + 1) % 2;
}
printf("\nTags:\n");
for(int i = 0; i < sets; ++i) {
printf("Set %2d: 0x%08x 0x%08x ", i, tags[0][i], tags[0][1]);
}
printf("\n");
}
printf("\nHits = %d, Accesses = %d, Hit ratio = %f\n", hits, accesses,
((float)hits)/accesses);
fclose(fp);
}

```

Output: Set Associative Mapping – 16 bit

```

Tags:
Set 0: 0x00000002 0x00000002 Set 1: 0x00000002 0x00000002 Set 2: 0x00000002 0x00000002 Set 3: 0xfc000000 0x00000002 Set 4: 0xfc000000
0x00000002 Set 5: 0xfc000000 0x00000002 Set 6: 0xfc000000 0x00000002 Set 7: 0x00000001 0x00000002

101: 0x00000010 Hit 0
Tags:
Set 0: 0x00000002 0x00000002 Set 1: 0x00000002 0x00000002 Set 2: 0x00000002 0x00000002 Set 3: 0xfc000000 0x00000002 Set 4: 0xfc000000
0x00000002 Set 5: 0xfc000000 0x00000002 Set 6: 0xfc000000 0x00000002 Set 7: 0x00000001 0x00000002

102: 0x00000014 Hit 0
Tags:
Set 0: 0x00000002 0x00000002 Set 1: 0x00000002 0x00000002 Set 2: 0x00000002 0x00000002 Set 3: 0xfc000000 0x00000002 Set 4: 0xfc000000
0x00000002 Set 5: 0xfc000000 0x00000002 Set 6: 0xfc000000 0x00000002 Set 7: 0x00000001 0x00000002

103: 0x00000018 Hit 0
Tags:
Set 0: 0x00000002 0x00000002 Set 1: 0x00000002 0x00000002 Set 2: 0x00000002 0x00000002 Set 3: 0xfc000000 0x00000002 Set 4: 0xfc000000
0x00000002 Set 5: 0xfc000000 0x00000002 Set 6: 0xfc000000 0x00000002 Set 7: 0x00000001 0x00000002

Hits = 76, Accesses = 103, Hit ratio = 0.737864

```



```
78: 0x00000038 Hit 0
Tags:
Set 0: 0x00000000 0x00000000 Set 1: 0x00000000 0x00000000 Set 2: 0x00000000 0x00000000 Set 3: 0x00000000 0x00000000 Set 4: 0x00000001
0x00000000 Set 5: 0x00000001 0x00000000 Set 6: 0x00000001 0x00000000 Set 7: 0x00000000 0x00000000

79: 0x00000034 Hit 1
Tags:
Set 0: 0x00000000 0x00000000 Set 1: 0x00000000 0x00000000 Set 2: 0x00000000 0x00000000 Set 3: 0x00000000 0x00000000 Set 4: 0x00000001
0x00000000 Set 5: 0x00000001 0x00000000 Set 6: 0x00000001 0x00000000 Set 7: 0x00000000 0x00000000

80: 0x00000038 Hit 1
Tags:
Set 0: 0x00000000 0x00000000 Set 1: 0x00000000 0x00000000 Set 2: 0x00000000 0x00000000 Set 3: 0x00000000 0x00000000 Set 4: 0x00000001
0x00000000 Set 5: 0x00000001 0x00000000 Set 6: 0x00000001 0x00000000 Set 7: 0x00000000 0x00000000

81: 0x0000003c Miss
Tags:
Set 0: 0x00000000 0x00000000 Set 1: 0x00000000 0x00000000 Set 2: 0x00000000 0x00000000 Set 3: 0x00000000 0x00000000 Set 4: 0x00000001
0x00000000 Set 5: 0x00000001 0x00000000 Set 6: 0x00000001 0x00000000 Set 7: 0x00000000 0x00000000

82: 0x00000040 Miss
Tags:
Set 0: 0x00000000 0x00000000 Set 1: 0x00000000 0x00000000 Set 2: 0x00000000 0x00000000 Set 3: 0x00000000 0x00000000 Set 4: 0x00000001
0x00000000 Set 5: 0x00000001 0x00000000 Set 6: 0x00000001 0x00000000 Set 7: 0x00000000 0x00000000

Hits = 68, Accesses = 82, Hit ratio = 0.829268
```

Conclusion

A cache is a smaller, quicker memory that stores copies of software instructions and data that are utilized regularly in the operation of programs and is positioned closer to the processor core. The program's overall speed is improved by having quick access to these commands. In this experiment, we learnt about fully associative mapping and set associative mapping and implemented them in a C program.



EXPERIMENT - 7

Aim: Study and Simulate (Stepwise) Experiments of MIPS Programming available at MIPS simulator. Modify the given programs to implement

- 1) Add 10 nos.
- 2) to print message "Hello MIPS"
- 3) To Reverse the input string (e.g."ABC" - "CBA").

Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	26/11/21	26/11/21	

Theory:

MIPS assembly language simply refers to the assembly language of the MIPS processor. The term MIPS is an acronym for Microprocessor without Interlocked Pipeline Stages. It is a reduced-instruction set architecture developed by an organization called MIPS Technologies. The MIPS assembly language is a very useful language to learn because many embedded systems run on the MIPS processor.

MIPS Architecture

Data Types

1. All the instructions in MIPS are 32 bits.
2. A byte in the MIPS architecture represents 8 bits; a halfword represents 2 bytes (16 bits) and a word represents 4 bytes (32 bits).
3. Each character used in the MIPS architecture requires 1 byte of storage. Each integer used requires 4 bytes of storage.

Literals

In the MIPS architecture, literals represent all numbers (e.g. 5), characters enclosed in single quotes (e.g. 'g') and strings enclosed in double quotes (e.g. "Deadpool").

Code 1: To add 10 numbers

Registers

MIPS architecture uses 32 general-purpose registers. Each register in this architecture is preceded by '\$' in the assembly language instruction. You can address these registers in one of two ways. Either use the register's number (that is, from \$0 to \$31), or the register's name (for example, \$t1).



MIPS Structure

General structure of a program created using the MIPS assembly language

A typical program created using the MIPS assembly language has two main parts. They are the data declaration section of the program and the code section of the program.

Data declaration section of a MIPS assembly language program

The data declaration section of the program is the part of the program identified with the assembler directive `.data`. This is the part of the program in which all the variables to be used in the program are created and defined. It is also the part of the program where storage is allocated in the main memory (RAM). The MIPS assembly language program declares variables as follows: `name: .storage_type value(s)`.

The "name" refers to the name of the variable being created. The "storage_type" refers to the type of data that the variable is meant to store. The "value(s)" refers to the information to be stored in the variable being created.

Code section of the MIPS assembly language program

The code section of the program is the part of the program in which the instructions to be executed by the program are written. It is placed in the section of the program identified with the assembler directive `.text`. The starting point for the code section of the program is marked with the label "main" and the ending point for the code section of the program is marked with an exit system call. This section of a MIPS assembly language program typically involves the manipulation of registers and the performance of arithmetic operations.

CODE 1: Addition of 10 numbers

```
.data
array: .word 5,6,12,4,19,8,4,21,9,1
length: .word 10
sum: .word 0
myMessage: .ascii "Sum of 10 nos. in the given array is: "
.text
main:
la $t0, array # Base address
li $t1, 0 # i = 0
lw $t2, length # $t2 = length
li $t3, 0 # sum = 0
sumLoop:
lw $t4, ($t0) # $t4 = array[i]
add $t3, $t3, $t4 # sum = sum + array[i]
```



```
add $t1, $t1, 1 # i = i + 1
add $t0, $t0, 4 # Updating the array address.
blt $t1, $t2, sumLoop
sw $t3, sum
li $v0, 4
la $a0, myMessage
syscall
li $v0, 1
move $a0, $t3
syscall
```

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File Edit Run Settings Tools Help

Run speed at max (no interaction)

Edit Execute

Text Segment

Bkpt	Address	Code	Basic	Source
	0x00400000	0x3e011001	lui \$1, 0x000001001	8: la \$t0, array # Base address
	0x00400004	0x34200000	ori \$9, \$1, 0x00000000	
	0x00400008	0x24090000	addiu \$9, \$0, 0x00000000	9: li \$t1, 0 # i = 0
	0x0040000c	0x3c011001	lui \$1, 0x000001001	10: lw \$t2, length # \$t2 = length
	0x00400010	0x8c2a0028	lw \$10, 0x00000028(\$1)	
	0x00400014	0x240b0000	addiu \$11, \$0, 0x00000000	11: li \$t3, 0 # sum = 0
	0x00400018	0x8d000000	lw \$12, 0x00000000(\$9)	13: lw \$t4, (\$t0) # \$t4 = array[i]
	0x0040001c	0x016c5920	add \$11, \$11, \$12	14: add \$t3, \$t3, \$t4 # sum = sum + array[i]
	0x00400020	0x21290001	addi \$9, \$9, 0x00000001	15: add \$t1, \$t1, 1 # i = i + 1
	0x00400024	0x21080004	addi \$9, \$9, 0x00000004	16: add \$t0, \$t0, 4 # Updating the array address.
	0x00400028	0x012a002a	blt \$1, \$9, \$10	17: blt \$t1, \$t2, sumLoop
	0x0040002c	0x1420ffff	hane \$1, \$0, 0xffffffff	

Data Segment

Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	0x00000005	0x00000006	0x0000000e	0x00000004	0x00000013	0x00000008	0x00000004	0x00000015
0x10010004	0x00000005	0x00000001	0x00000004	0x00000000	0x206a7553	0x3102466f	0xffff2030	0x6302e73
0x10010008	0x6874206a	0x69672065	0x206a6576	0x61727261	0x73692079	0x00000203a	0x00000000	0x00000000
0x1001000c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010010	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010014	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010018	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001001c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010020	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010024	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010028	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001002c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010030	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010034	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010038	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001003c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010040	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010044	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010048	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001004c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010050	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010054	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010058	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001005c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010060	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010064	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010068	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001006c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010070	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010074	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010078	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001007c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010080	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010084	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010088	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001008c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010090	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010094	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010098	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001009c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100ac	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100b0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100b4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100b8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100bc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100cc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100d0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100d4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100d8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100dc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100ec	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100f0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100f4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100f8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100fc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010100	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010104	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010108	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001010c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010110	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010114	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010118	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001011c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010120	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000

Registers

Name	Number	Value
\$zero	0	0x00000000
\$at	1	0x00000000
\$v0	2	0x00000000
\$v1	3	0x00000000
\$a0	4	0x00000000
\$a1	5	0x00000000
\$a2	6	0x00000000
\$a3	7	0x00000000
\$a0	8	0x00000000
\$t1	9	0x00000000
\$t2	10	0x00000000
\$t3	11	0x00000000
\$t4	12	0x00000000
\$t5	13	0x00000000
\$t6	14	0x00000000



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File Edit Run Settings Tools Help

Run speed at max (no interaction)

Edit Execute

Text Segment

Block	Address	Code	Basic	Source
	0x00400000	0x24020004	addiu \$2,\$0,0x00000000	5: li \$v0, 4
	0x00400004	0x3c011001	lui \$1,0x00001001	6: la \$a0, string
	0x00400008	0x34240000	ori \$4,\$1,0x00000000	
	0x0040000c	0x0000000c	syscall	7: syscall

Data Segment

Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	0x6c65480a	0x202c6f6c	0x6c726f57	0x000a2164	0x00000000	0x00000000	0x00000000	0x00000000
0x10010020	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010040	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010060	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010080	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010100	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010120	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000

Mars Messages Run I/O

Hello, World!

Clear

-- program is finished running (dropped off bottom) --

Registers Coproc 1 Coproc 0

Name	Number	Value
\$zero	0	0x00000000
\$at	1	0x10010000
\$v0	2	0x00000004
\$v1	3	0x00000000
\$a0	4	0x10010000
\$a1	5	0x00000000
\$a2	6	0x00000000
\$a3	7	0x00000000
\$t0	8	0x00000000
\$t1	9	0x00000000
\$t2	10	0x00000000
\$t3	11	0x00000000
\$t4	12	0x00000000
\$t5	13	0x00000000
\$t6	14	0x00000000
\$t7	15	0x00000000
\$a0	16	0x00000000
\$a1	17	0x00000000
\$a2	18	0x00000000
\$a3	19	0x00000000
\$a4	20	0x00000000
\$a5	21	0x00000000
\$a6	22	0x00000000
\$a7	23	0x00000000
\$s0	24	0x00000000
\$s1	25	0x00000000
\$s2	26	0x00000000
\$s3	27	0x00000000
\$s4	28	0x10000000
\$s5	29	0x77777777
\$s6	30	0x00000000
\$s7	31	0x00000000
\$pc		0x00400010
\$hi		0x00000000
\$lo		0x00000000

Mars Messages Run I/O

Hello, World!

Clear

-- program is finished running (dropped off bottom) --

CODE 3: To Reverse the input string (e.g."ABC" – "CBA").

```
.data
str: .ascii "India"
str_msg1: .ascii "Original string: "
str_msg2: .ascii "Reversed string: "
str_nl: .ascii "\n"
str_len: .word 0

.text
main:
    #print original string
    la $a0,str_msg1 #leading text
    li $v0,4
    syscall
    la $a0,str #original string
    li $v0,4
```



```
syscall
la $a0,str_nl #new Line
li $v0,4
syscall
#get length
add $t0,$zero,$zero #initialize registers when needed
add $a0,$zero,$zero
add $a1,$zero,$zero
la $a0,str #loads the address of the string
la $a1,str
getLen:
    lb $t0,0($a0) #load first byte
    beqz $t0,saveLen
    addi $a0,$a0,1
    j getLen #jump back to start of this loop
saveLen:
    subu $t0,$a0,$a1 #len = address of null terminator - str address
    sw $t0,str_len
#reverse the string
add $t0,$zero,$zero #address of the beginning of str
add $t1,$zero,$zero #address of the end of str
add $t2,$zero,$zero
add $t3,$zero,$zero
revString:
    #find the index of the last character before the end of the string
    la $t0,str #loads the address of the start of the string
    lw $t1,str_len #loads the length of the string
    addu $t1,$t0,$t1
    subi $t1,$t1,1
    loop:
        lb $t2,0($t0) #load the first character
        lb $t3,0($t1) #load the last character
        ble $t1,$t0,printRev
        sb $t3,0($t0)
        sb $t2,0($t1)
        addi $t0,$t0,1
        subi $t1,$t1,1 #and loop until we reach the middle
        j loop
#print the reversed version of the text
printRev:
    add $a0,$zero,$zero #initialize the a0 registry
    la $a0,str_msg2 #leading text
    li $v0,4
```



```
syscall  
la $a0,str #reversed string  
li $v0,4  
syscall  
li $v0,10 #exit program  
syscall
```

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Run speed at max (no interaction)

Text Segment

Bkpt	Address	Code	Basic	Source
	0x00400000	0x3e011001	lui \$1,0x00001001	10: la \$a0,str msg1 #leading text
	0x00400004	0x34240006	ori \$4,\$1,0x00000006	
	0x00400008	0x24020004	addiu \$2,\$0,0x00000004	11: li \$v0,4
	0x0040000c	0x0000000c	syscall	12: syscall
	0x00400010	0x3e011001	lui \$1,0x00001001	13: la \$a0,str \$original_string
	0x00400014	0x34240006	ori \$4,\$1,0x00000006	
	0x00400018	0x24020004	addiu \$2,\$0,0x00000004	14: li \$v0,4
	0x0040001c	0x0000000c	syscall	15: syscall
	0x00400020	0x3e011001	lui \$1,0x00001001	16: la \$a0,str nl \$new_line
	0x00400024	0x34240006	ori \$4,\$1,0x00000006	
	0x00400028	0x24020004	addiu \$2,\$0,0x00000004	17: li \$v0,4
	0x0040002c	0x0000000c	syscall	18: syscall

Data Segment

Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	0x6e646961	0x724f0049	0x6e696769	0x73206e61	0x6e697274	0x00203a67	0x69726652	0x64657372
0x10010004	0x27472320	0x3a676e69	0x00000020	0x00000005	0x00000000	0x00000000	0x00000000	0x00000000
0x10010008	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001000c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010010	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010014	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010018	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001001c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010020	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010024	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010028	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001002c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010030	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010034	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010038	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001003c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010040	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010044	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010048	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001004c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010050	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010054	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010058	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001005c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010060	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010064	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010068	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001006c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010070	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010074	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010078	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001007c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010080	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010084	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010088	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001008c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010090	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010094	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010098	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001009c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100ac	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100b0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100b4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100b8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100bc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100cc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100d0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100d4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100d8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100dc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100ec	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100f0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100f4	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100f8	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100fc	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010100	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010104	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010108	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001010c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010110	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010114	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010118	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x1001011c	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010120	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000

Registers

Name	Number	Value
\$zero	0	0x00000000
\$at	1	0x10010000
\$v0	2	0x0000000a
\$v1	3	0x00000000
\$a0	4	0x10010000
\$a1	5	0x10010000
\$a2	6	0x00000000
\$a3	7	0x00000000
\$a0	8	0x10010000
\$t1	9	0x10010002
\$t2	10	0x00000004
\$t3	11	0x00000004
\$t4	12	0x00000000
\$t5	13	0x00000000
\$t6	14	0x00000000
\$t7	15	0x00000000
\$a0	16	0x00000000
\$a1	17	0x00000000
\$a2	18	0x00000000
\$a3	19	0x00000000
\$a5	20	0x00000000
\$a5	21	0x00000000
\$a6	22	0x00000000
\$a7	23	0x00000000
\$t8	24	0x00000000
\$t9	25	0x00000000
\$t0	26	0x00000000
\$t1	27	0x00000000
\$gp	28</	



EXPERIMENT - 8

Aim : Run the simulator using the given link

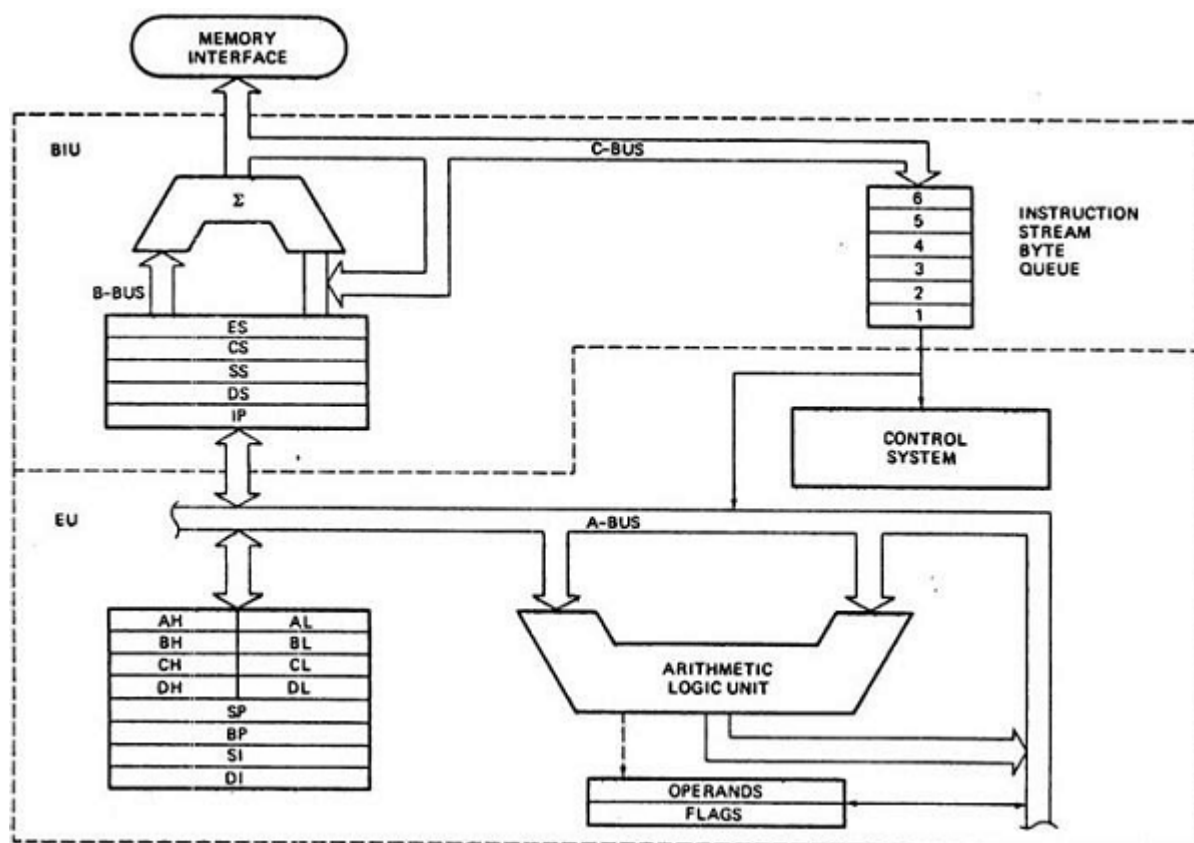
<https://ydoc2.github.io/8086-emulator-web/compile>. Study and execute First 6 examples from the repository <https://github.com/YJDoc2/8086-Emulator/tree/master/examples>

Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	17/12/21	17/12/21	

Architecture of 8086

The following diagram depicts the architecture of a 8086 Microprocessor –



Pins on 8086:



AD0-AD15 (Address Data Bus): Bidirectional address/data lines. These are low order address bus. They are multiplexed with data.

When these lines are used to transmit memory address, the symbol A is used instead of AD, for example, A0- A15.

A16 - A19 (Output): High order address lines. These are multiplexed with status signals.

A16/S3, A17/S4: A16 and A17 are multiplexed with segment identifier signals S3 and S4.

A18/S5: A18 is multiplexed with interrupt status S5.

A19/S6: A19 is multiplexed with status signal S6.

BHE/S7 (Output): Bus High Enable/Status. During T1, it is low. It enables the data onto the most significant half of the data bus, D8-D15. 8-bit devices connected to the upper half of the data bus use BHE signal. It is multiplexed with status signal S7. S7 signal is available during T3 and T4.

RD (Read): For read operation. It is an output signal. It is active when LOW.

Ready (Input): The addressed memory or I/O sends acknowledgment through this pin.

When HIGH, it denotes that the peripheral is ready to transfer data.

RESET (Input): System reset. The signal is active HIGH.

CLK (input): Clock 5, 8 or 10 MHz.

INTR: Interrupt Request.

NMI (Input): Non-maskable interrupt request.

TEST (Input): Wait for test control. When LOW the microprocessor continues execution otherwise waits.

VCC: Power supply +5V dc.

GND: Ground.

Operating Modes of 8086

There are two operating modes of operation for Intel 8086, namely the minimum mode and the maximum mode.

When only one 8086 CPU is to be used in a microprocessor system, the 8086 is used in the Minimum mode of operation.

In a multiprocessor system 8086 operates in the Maximum mode.

ADDRESSING:

The 8086 microprocessors have 8 addressing modes. Two addressing modes have been provided for instructions which operate on register or immediate data.

These two addressing modes are:

Register Addressing: In register addressing, the operand is placed in one of the 16-bit or 8-bit general purpose registers.

Example

- MOV AX, CX
- ADD AL, BL
- ADD CX, DX

Immediate Addressing: In immediate addressing, the operand is specified in the instruction itself.



Example

- MOV AL, 35H
- MOV BX, 0301H
- MOV [0401], 3598H
- ADD AX, 4836H

The remaining 6 addressing modes specify the location of an operand which is placed in a memory.

These 6 addressing modes are:

Direct Addressing: In direct addressing mode, the operand's offset is given in the instruction as an 8-bit or 16-bit displacement element.

Example

- ADD AL, [0301]

The instruction adds the content of the offset address 0301 to AL. the operand is placed at the given offset (0301) within the data segment DS.

Register Indirect Addressing: The operand's offset is placed in any one of the registers BX, BP, SI or DI as specified in the instruction.

Example

- MOV AX, [BX]

It moves the contents of memory locations addressed by the register BX to the register AX.

Based Addressing: The operand's offset is the sum of an 8-bit or 16-bit displacement and the contents of the base register BX or BP. BX is used as base register for data segment, and the BP is used as a base register for stack segment.

Effective address (Offset) = [BX + 8-bit or 16-bit displacement].

Example

- MOV AL, [BX+05]; an example of 8-bit displacement.
- MOV AL, [BX + 1346H]; example of 16-bit displacement.

Indexed Addressing: The offset of an operand is the sum of the content of an index register SI or DI and an 8-bit or 16-bit displacement.

Offset (Effective Address) = [SI or DI + 8-bit or 16-bit displacement]

Example

- MOV AX, [SI + 05]; 8-bit displacement.
- MOV AX, [SI + 1528H]; 16-bit displacement.

Based Indexed Addressing: The offset of operand is the sum of the content of a base register BX or BP and an index register SI or DI.

Effective Address (Offset) = [BX or BP] + [SI or DI]

Here, BX is used for a base register for data segment, and BP is used as a base register for stack segment.

Example

- ADD AX, [BX + SI]
- MOV CX, [BX + SI]



Based Indexed with Displacement: In this mode of addressing, the operand's offset is given by:

Effective Address (Offset) = [BX or BP] + [SI or DI] + 8-bit or 16-bit displacement

Example

- MOV AX, [BX + SI + 05]; 8-bit displacement
- MOV AX, [BX + SI + 1235H]; 16-bit displacement

1. Program to add two word length number

```
; Program to add two word length numbers
OPR1: DW 0x6459 ; declare first number
OPR2: DW 0x8420 ; declare second number
RESULT: DW 0 ; declare place to store result
; actual entry point of the program
start:
MOV AX, word OPR1 ; move first number to AX
MOV BX, word OPR2 ; move second number to BX
CLC ; clear the carry flag
ADD AX, BX ; add BX to AX
MOV DI, OFFSET RESULT ; move offset of result to DI
MOV word [DI], AX ; store result
print reg ; print result
```

Code Editor

```

1 ; Program to add two word length numbers
2 OPR1: DW 0x6459 ; declare first number
3 OPR2: DW 0x8420 ; declare second number
4 RESULT: DW 0 ; declare place to store result
5 ; actual entry point of the program
6 start:
7 MOV AX, word OPR1 ; move first number to AX
8 MOV BX, word OPR2 ; move second number to BX
9 CLC ; clear the carry flag
10 ADD AX, BX ; add BX to AX
11 MOV DI, OFFSET RESULT ; move offset of result to DI
12 MOV word [DI], AX ; store result
13 print reg ; print result

```

COMPILER RUN NEXT STOP

Reg	H	L
A	e8	79
B	84	20
C	00	00
D	00	00

Segments	
SS	0000
DS	0000
ES	0000

Pointers	
SP	0000
BP	0000
SI	0000
DI	0004

Flags:

OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	1	0	0	0	0

Memory

Start Address: 00000 SET

59	64	20	84	79	e8	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00

Input

Output



2. A Program to move data from one segment to another.

```
; A Program to move data from one segment to another
SET 0 ; set address for segment 1
src:DB 0x8 ; store data
DB 0x12
DB 0x13
SET 0x3 ; set address for segment 2
dest:DB [0,3] ; store data
; actual entry point of the program
start:
print mem 0:8 ; print initial state of segment 1
print mem 0x10:8 ; print initial state of segment 2
MOV AX, 0 ; move address of seg1
MOV DS,AX ; to ds
MOV AX, 0x1 ; move address of seg2
MOV ES,AX ; to es
MOV SI, OFFSET src ; move offset of source data
MOV DI, OFFSET dest ; move offset of destination data
MOV CX, 0x3 ; move number of data items
print reg ; print state of registers
_loop:
mov AH, byte DS[SI] ; move one byte from source to ah
mov byte ES[DI],AH ; move ah to destination
inc SI
inc DI
dec CX ; decrement count
jnz _loop ; if count is not zero jump back
print mem 0:8 ; print final state of segment 1
print mem 0x10:8 ; print final state of segment 2
```



8086 Compiler

Code Editor

COMPILE
RUN
NEXT
STOP

```

3  src:DB 0x8 ; store data
4  DB 0x12
5  DB 0x13
6  SET 0x3 ; set address for segment 2
7  dest:DB [0,3] ; store data
8  ; actual entry point of the program
9  start:
10 print mem 0:8 ; print initial state of segment 1
11 print mem 0x10:8 ; print initial state of segment 2
12 MOV AX, 0 ; move address of seg1
13 MOV DS,AX ; to ds
14 MOV AX, 0x1 ; move address of seg2
15 MOV ES,AX ; to es
16 MOV SI, OFFSET src ; move offset of source data
17 MOV DI, OFFSET dest ; move offset of destination data
18 MOV CX, 0x3 ; move number of data items
19 print reg ; print state of registers
20 _loop:
21 mov AH, byte DS[SI] ; move one byte from source to ah
22 mov byte ES[DI],AH ; move ah to destination
23 inc SI
24 inc DI
25 dec CX ; decrement count
26 jnz _loop ; if count is not zero jump back
27 print mem 0:8 ; print final state of segment 1
28 print mem 0x10:8 ; print final state of segment 2
          
```

Reg	H	L
A	13	01
B	00	00
C	00	00
D	00	00

Segments	
SS	0000
DS	0000
ES	0001

Pointers	
SP	0000
BP	0000
SI	0003
DI	0003

Flags:								
OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	0	1	0	1	0

Memory

Start Address: 00000 SET

08	12	13	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
08	12	13	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00

Input ✓

Output

3. Program to calculate factorial using looping.

```

; Program to calculate factorial using looping
NUM: DW 0x9 ; calculate factorial of 9
RESULT: DW 0 ; place to store the result
; actual entry point of the program
start:
MOV CX,word NUM ; move number into cx
MOV AX, 0x1 ; initialize accumulator with 1
NOTZEROLoop: ; label to jump back to
MUL CX ; multiple by the number
DEC CX ; decrement the number
JNZ NOTZEROLoop ; if not zero jump back
MOV word RESULT,AX ; store the result in memory
print reg ; print registers
          
```



8086 Compiler

8086 Compiler

Code Editor

1 ; Program to calculate factorial using looping
2 NUM: DW 0x9 ; calculate factorial of 9
3 RESULT: DW 0 ; place to store the result
4 ; actual entry point of the program

5 start:
6 MOV CX, word NUM ; move number into cx
7 MOV AX, 0x1 ; initialize accumulator with 1

8 NOTZEROLOOP: ; label to jump back to
9 MUL CX ; multiply by the number
10 DEC CX ; decrement the number

11 JNZ NOTZEROLOOP ; if not zero jump back
12 MOV word RESULT, AX ; store the result in memory
13 print reg ; print registers

14

Input ✓

Output

Reg	H	L
A	89	80
B	00	00
C	00	00
D	00	00

Segments	
SS	0000
DS	0000
ES	0000

Pointers	
SP	0000
BP	0000
SI	0000
DI	0000

Flags:

OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	0	1	0	1	0

Memory

Start Address: 00000 SET

09	00	80	89	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00

4. Program to show use of interrupts.

```

; Program to show use of interrupts
; Also, Hello World program !
hello: DB "Hello World" ; store string
; actual entry point of the program, must be present
start:
MOV AH, 0x13 ; move BIOS interrupt number in AH
MOV CX, 11 ; move length of string in cx
MOV BX, 0 ; mov 0 to bx, so we can move it to es
MOV ES, BX ; move segment start of string to es, 0
MOV BP, OFFSET hello ; move start offset of string in bp
MOV DL, 0 ; start writing from col 0
int 0x10 ; BIOS interrupt
  
```



8086 Compiler



Code Editor

COMPILE

RUN

NEXT

STOP

```

1 ; Program to show use of interrupts
2 ; Also, Hello world program I
3 hello: DB "Hello world"; store string
4 ; actual entry point of the program, must be present
5 start:
6 MOV AH, 0x13 ; move BIOS interrupt number in AH
7 MOV CX, 11 ; move length of string in cx
8 MOV BX, 0 ; mov 0 to bx, so we can move it to es
9 MOV ES, BX ; move segment start of string to es, 0
10 MOV BP, OFFSET hello ; move start offset of string in bp
11 MOV DL, 0 ; start writing from col 0
12 INT 0x10 ; BIOS interrupt
          
```

Input

Output

Hello World

Reg	H	L
A	13	00
B	00	00
C	00	0b
D	00	00

Segments	
SS	0000
DS	0000
ES	0000

Pointers	
SP	0000
BP	0000
SI	0000
DI	0000

Flags:

OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	0	0	0	0	0

Memory

Start Address: 00000

SET

48	65	6c	6c	6f	20	57	6f	72	6c	64	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00

5. Program to calculate LCM and GCD of two numbers.

```

; Program to calculate LCM and GCD of two numbers
no1: dw 0x6 ; number 1
no2: dw 0x5 ; number 2
gcd: dw 0 ; place to store gcd
lcm: dw 0 ; place to store lcm
; actual entry point of the program
start:
mov ax, word no1 ; move number 1 in accumulator
mov bx, word no2 ; move number 2 in register BX
loop0: mov dx, 0x0 ; place to loop back
; cannot use 'loop' as label, as loop is an opcode which will give error
when used with jumps
div bx ; divide accumulator by bx
mov ax, bx
mov bx, dx
cmp bx, 0x0 ; check if bx is 0
jnz loop0 ; if not loop back
mov word gcd, ax ; store gcd
mov cx, ax ; move ax in cx
mov ax, word no1 ; move number 1 in accumulator
mov bx, word no2 ; move number 2 in register BX
mul bx ; multiply accumulator by BX
div cx ; divide accumulator by CX
mov word lcm, ax ; store lcm
print mem :16 ; print memory
          
```



8086 Compiler

Code Editor

```

1 ; Program to calculate LCM and GCD of two numbers
2 n01: dw 0x6 ; number 1
3 n02: dw 0x5 ; number 2
4 gcd: dw 0 ; place to store gcd
5 lcm: dw 0 ; place to store lcm
6 ; actual entry point of the program
7 start:
8 mov ax, word n01 ; move number 1 in accumulator
9 mov bx, word n02 ; move number 2 in register BX
10 loop0: mov dx, 0x0 ; place to loop back
11 ; cannot use 'loop' as label, as loop is an opcode which will give error when used
12 div bx ; divide accumulator by bx
13 mov ax, bx
14 mov bx, dx
15 cmp bx, 0x0 ; check if bx is 0
16 jnz loop0 ; if not loop back
17 mov word gcd, ax ; store gcd
18 mov cx, ax ; move ax in cx
19 mov ax, word n01 ; move number 1 in accumulator
20 mov bx, word n02 ; move number 2 in register BX
21 mul bx ; multiply accumulator by BX
22 div cx ; divide accumulator by CX
23 mov word lcm, ax ; store lcm
24 print mem :16 ; print memory
          
```

Input

Output

Reg

Reg	H	L
A	00	1e
B	00	05
C	00	01
D	00	00

Segments

SS	0000
DS	0000
ES	0000

Pointers

SP	0000
BP	0000
SI	0000
DI	0000

Flags:

OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	0	1	0	1	0

Memory

Start Address: 00000

06	00	05	00	01	00	1e	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00

6. A program to sort the numbers in descending order, using bubble sort

```

; A program to sort the numbers in descending order, using bubble sort
vals: DB 0xF ; declaration of the numbers
DB 0x5A
DB 0x24
DB 0x2
DB 0x56
last: DB 0 ; declaring an element to get total number of elements later
; actual entry point of the program
start:
print mem :8 ; print initial state of memory
MOV CH, OFFSET last ; move number of elements to CH
outer: ; loop label for outer loop
MOV CL, OFFSET last ; move number of elements to CL
MOV SI, OFFSET vals ; move offset of values to si
inner: ; loop label for inner loop
MOV AX, word [SI] ; move two adjacent numbers to AX
CMP AL, AH ; compare both values
JNC skip ; jump to skip if first num is greater than second
XCHG AL, AH ; exchange both numbers
MOV word [SI], AX ; move exchanged numbers to memory
skip:
INC SI ; increment SI
DEC CL ; decrement inner loop counter
JNZ inner ; jump back to inner if counter is not zero
          
```



DEC CH ; decrement outer loop counter
JNZ outer ; jump back to outer, if counter is not zero
print mem :8 ; print final state of memory

8086 Compiler



Code Editor

COMPILE

RUN

NEXT

STOP

```
2  vals:DB 0xF ; declaration of the numbers
3  DB 0x5A
4  DB 0x24
5  DB 0x2
6  DB 0x56
7  last: DB 0 ; declaring an element to get total number of elements later
8  ; actual entry point of the program
9  start:
10 print mem :8 ; print initial state of memory
11 MOV CH, OFFSET last ; move number of elements to CH
12 outer: ; loop label for outer loop
13 MOV CL, OFFSET last ; move number of elements to CL
14 MOV SI, OFFSET vals ; move offset of values to SI
15 inner: ; loop label for inner loop
16 MOV AX, word [SI] ; move two adjacent numbers to AX
17 CMP AL, AH ; compare both values
18 JNC skip ; jump to skip if first num is greater than second
19 XCHG AL, AH ; exchange both numbers
20 MOV word [SI], AX ; move exchanged numbers to memory
21 skip:
22 INC SI ; increment SI
23 DEC CL ; decrement inner loop counter
24 JNZ inner ; jump back to inner if counter is not zero
25 DEC CH ; decrement outer loop counter
26 JNZ outer ; jump back to outer, if counter is not zero
27 print mem :8 ; print final state of memory
```

Input



Output

Reg	H	L
A	00	02
B	00	00
C	00	00
D	00	00

Segments	
SS	0000
DS	0000
ES	0000

Pointers	
SP	0000
BP	0000
SI	0005
DI	0000

Flags:

OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	0	1	0	1	0

Memory

Start Address

00000

SET

5a	56	24	0f	02	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00



EXPERIMENT - 9

AIM: Write Programs using Assembly Language:

1. To implement Macros for calculating Factorial of a number.
2. To study and Implement DOS interrupts. Calculate and display Sum of 2 user entered inputs using DOS interrupt

Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	17/12/21	17/12/21	

Theory :

Macros : Writing a macro is another way of ensuring modular programming in assembly language. A macro is a sequence of instructions, assigned by a name and could be used anywhere in the program. In NASM, macros are defined with %macro and %endmacro directives. The macro begins with the %macro directive and ends with the %endmacro directive. The Syntax for macro definition –

```
%macro macro_name number_of_params  
<macro body>  
%endmacro
```

Where, number_of_params specifies the number parameters, macro_name specifies the name of the macro. The macro is invoked by using the macro name along with the necessary parameters. When you need to use some sequence of instructions many times in a program, you can put those instructions in a macro and use it instead of writing the instructions all the time

Macros are just like procedures, but not really. Macros look like procedures, but they exist only until your code is compiled, after compilation all macros are replaced with real instructions. If you declared a macro and never used it in your code, compiler will simply ignore it. Unlike procedures, macros should be defined above the code that uses it.

The advantage of using Macro is that it avoids the overhead time involved in calling and returning (as in the procedures). Therefore, the execution of Macros is faster as compared to procedures. Another advantage is that there is no need for accessing stack or providing any separate memory to it for storing and returning the address locations while shifting the processor controls in the program.



But it should be noted that every time you call a macro, the assembler of the microprocessor places the entire set of Macro instructions in the mainline program from where the call to Macro is being made. This is known as Macro expansion. Due to this, the program code (which uses Macros) takes more memory space than the code which uses procedures for implementing the same task using the same set of instructions.

DOS Interrupts

Interrupt is the method of creating a temporary halt during program execution and allows peripheral devices to access the microprocessor. The microprocessor responds to that interrupt with an ISR (Interrupt Service Routine), which is a short program to instruct the microprocessor on how to handle the interrupt. The following image shows the types of interrupts we have in a 8086 microprocessor –

DOS Interrupt is a Software Interrupt. INT 21H is a DOS interrupt. It is one of the most commonly used interrupts while writing code in 8086 assembly language. To use the DOS interrupt 21H load AH with the desired sub-function. Load other required parameters in other registers, and make a call to INT 21H.

AH	Description	AH	Description
01	Read character from STDIN	02	Write character to STDOUT
05	Write character to printer	06	Console Input/Output
07	Direct char read (STDIN), no echo	08	Char read from STDIN, no echo
09	Write string to STDOUT	0A	Buffered input
0B	Get STDIN status	0C	Flush buffer for STDIN
0D	Disk reset	0E	Select default drive
19	Get current default drive	25	Set interrupt vector
2A	Get system date	2B	Set system date
2C	Get system time	2D	Set system time
2E	Set verify flag	30	Get DOS version
35	Get Interrupt vector		
36	Get free disk space	39	Create subdirectory



3A	Remove subdirectory	3B	Set working directory
3C	Create file	3D	Open file
3E	Close file	3F	Read file
40	Write file	41	Delete file
42	Seek file	43	Get/Set file attributes
47	Get current directory	4C	Exit program
4D	Get return code	54	Get verify flag
56	Rename file	57	Get/Set file date

CODE:

1. Factorial using Macros.

```
; Program to calculate factorial of a number using a macro
NUM: DW 0x7 ; calculate factorial of 7
RESULT: DW 0
; Macro for factorial
MACRO fact(no) -> MUL word no <-
; actual entry point of the program
start:
MOV AX, 0x0001 ; initialize accumulator with 1
NOTZEROLOOP:
fact(NUM) ;
DEC word NUM ; decrement number
JNZ NOTZEROLOOP ; jump back if number is not zero yet
MOV word RESULT,AX
print mem :16
```



8086 Compiler

8086 Compiler

Code Editor

```

1 ; Program to calculate factorial of a number using a macro
2 NUM: DW 0x7 ; calculate factorial of 7
3 RESULT: DW 0
4 ; Macro for factorial
5 MACRO fact(no) -> MUL word no <-
6 ; actual entry point of the program
7 start:
8 MOV AX, 0x0001 ; initialize accumulator with 1
9 NOTZEROLoop:
10 fact(NUM) ;
11 DEC word NUM ; decrement number
12 JNZ NOTZEROLoop ; jump back if number is not zero yet
13 MOV word RESULT,AX
14 print mem :10
          
```

Input

Output

Reg

Reg	H	L
A	13	b0
B	00	00
C	00	00
D	00	00

Segments

Segment	Address
SS	0000
DS	0000
ES	0000

Pointers

Pointer	Address
SP	0000
BP	0000
SI	0000
DI	0000

Flags:

OF	DF	IF	TF	SF	ZF	AF	PF	CF
0	0	0	0	0	1	0	1	0

Memory

Start Address: 00000

00	00	b0	13	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00

2. DOS Interrupt

```

DATA SEGMENT
NUM1 DB ?
NUM2 DB ?
RESULT DB ?
MSG1 DB 10,13,"ENTER FIRST NUMBER TO ADD : $"
MSG2 DB 10,13,"ENTER SECOND NUMBER TO ADD : $"
MSG3 DB 10,13,"RESULT OF ADDITION IS : $"
ENDS
CODE SEGMENT
ASSUME DS:DATA, CS:CODE
START:
MOV AX,DATA
MOV DS,AX
LEA DX,MSG1
MOV AH,9
INT 21H
MOV AH,1
INT 21H
SUB AL,30H
MOV NUM1,AL
LEA DX,MSG2
MOV AH,9
INT 21H
          
```



```
MOV AH,1
INT 21H
SUB AL,30H
MOV NUM2,AL
ADD AL,NUM1
MOV RESULT,AL
MOV AH,0
AAA
ADD AH,30H
ADD AL,30H
MOV BX,AX
LEA DX,MSG3
MOV AH,9
INT 21H
MOV AH,2
MOV DL,BH
INT 21H
MOV AH,2
MOV DL,BL
INT 21H
MOV AH,4CH
INT 21H
ENDS
END START
```



```
DOSBox 0.74-3, Cpu speed: 3000 cycles, Frameskip 0, Progra...
/z          Display source line with error message
/zi,/zd     Debug info: zi=full, zd=line numbers only
C:\>tasm Addition.asm
Turbo Assembler Version 2.51 Copyright (c) 1988, 1991 Borland International

Assembling file:   Addition.asm
Error messages:    None
Warning messages:  None
Passes:            1
Remaining memory:  491k

C:\>Addition
Illegal command: Addition.

C:\>tlink Addition
Turbo Link Version 4.0 Copyright (c) 1991 Borland International
Warning: No stack

C:\>Addition

ENTER FIRST NUMBER TO ADD : 1
ENTER SECOND NUMBER TO ADD : 0
RESULT OF ADDITION IS : 01
C:\>
```



EXPERIMENT - 10

Aim: Write a Program using ALP to Simulate Microcontroller interfacing with 7 segment display using
<http://vlabs.iitb.ac.in/vlabs-dev/labs/8051-Microcontroller-Lab/labs/exp1/simulation.php>. Display your SAP ID using this tool

Submission Sheet

SAP ID	Name of Student	Date of Experiment	Date of Submission	Remarks
60004190057	Junaid Girkar	3/10/21	3/10/21	

THEORY:

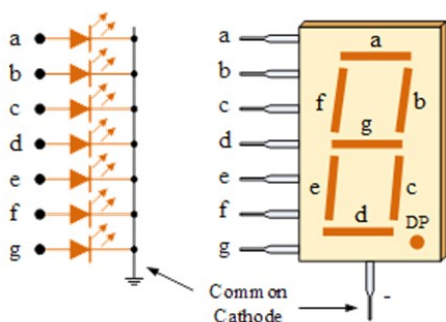
The 7-segment display consists of seven LEDs arranged in a rectangular fashion. Each of the seven LEDs is called a segment because when illuminated the segment forms part of a numerical digit (both Decimal and Hex) to be displayed. An additional 8th LED is sometimes used within the same package which is the indication of a decimal point(DP), when two or more 7-segment displays are connected together numbers greater than ten can be displayed.

So by forward biasing the appropriate pins of the LED segments in a particular order, some segments will be glowing and others will remain as it is, allowing the desired character pattern of the number to be generated on the display. This then allows us to display each of the ten decimal digits 0 to 9 on the same 7-segment display.

In the common cathode display, all the cathode connections of the LED segments are joined together to logic "0" or ground. The individual segments are illuminated by application of a "HIGH", or logic "1" signal via a current limiting resistor to forward bias the individual Anode terminals (a-g).



Common Cathode 7-segment Display



Decimal Digit	Common Cathode						
	Individual Segments Illuminated						
	a	b	c	d	e	f	g
0	1	1	1	1	1	1	
1		1	1				
2	1	1		1	1		1
3	1	1	1	1			1
4		1	1			1	1
5	1		1	1		1	1
6	1		1	1	1	1	1
7	1	1	1				
8	1	1	1	1	1	1	1
9	1	1	1	1		1	1

Common Cathode Decoding Table

CHAR	A	B	C	D	E	F	G	HEX
SAP ID = 60004190057								
6	1		1	1	1	1	1	5F
0	1	1	1	1	1	1		7E
4		1	1			1	1	33
1		1	1					30
9	1	1	1	1		1	1	7B
5	1		1	1		1	1	5B
7	1	1	1					70
HEXADECIMAL = DF8874769								
D	1	1	1	1	1	1		7E
F	1				1	1	1	47
8	1	1	1	1	1	1	1	7F



7	1	1	1					70
4		1	1			1	1	33
6	1		1	1	1	1	1	5F
9	1	1	1	1		1	1	7B
NAME = JUNAID								
J		1	1	1				38
U		1	1	1	1	1		3E
N	1	1	1		1	1		76
A	1	1	1		1	1	1	77
I		1	1					30
D		1	1	1	1		1	3D

CODE:

```

MOV P0,#5Fh    //to display 6
MOV P0,#7Eh    //to display 0
MOV P0,#7Eh    //to display 0
MOV P0,#7Eh    //to display 0
MOV P0,#33h    //to display 4
MOV P0,#30h    //to display 1
MOV P0,#7Bh    //to display 9
MOV P0,#7Eh    //to display 0
MOV P0,#7Eh    //to display 0
MOV P0,#5Bh    //to display 5
MOV P0,#70h    //to display 7

```

```

MOV P0,#3Dh    //to display d
MOV P0,#47h    //to display F
MOV P0,#7Fh    //to display 8
MOV P0,#7Fh    //to display 8
MOV P0,#70h    //to display 7
MOV P0,#33h    //to display 4
MOV P0,#70h    //to display 7
MOV P0,#5Fh    //to display 6
MOV P0,#7Bh    //to display 9

```

```

MOV P0,#38h    //to display J
MOV P0,#3Eh    //to display U

```



```
//to display n
//to display A
//to display I
//to display d
```

The image displays a 4x2 grid of screenshots from the Proteus 8051 IDE, illustrating the configuration of a microcontroller project. Each screenshot is divided into four main panels: Pin Diagram, Memory, Peripherals, and Editor.

- Pin Diagram:** Shows the 8051 microcontroller with its pins and their connections. The '8051' label is prominent in the center of each panel.
- Memory:** Displays the memory map, including locations and values for various components like P2, P0, SP, DP, and others.
- Peripherals:** Shows the configuration of various peripherals, including LEDs, switches, and a display. The '8051' label is also present in this panel.
- Editor:** Displays the assembly code for the 8051 microcontroller. The code is written in assembly language and includes comments.

The screenshots show the configuration of the 8051 microcontroller for a specific project, with the Editor view showing the assembly code. The top row shows the initial state, while the subsequent rows show the configuration of various pins and peripherals like LEDs, switches, and a display. The Editor view consistently shows assembly code for the 8051 microcontroller.



Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

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Pin Diagram **Memory** **Peripherals** **Editor**

Current Line: Next Line (PC) Error Line

1: MOV R0, #0 //to display 0
2: MOV R0, #1 //to display 1
3: MOV R0, #2 //to display 2
4: MOV R0, #3 //to display 3
5: MOV R0, #4 //to display 4
6: MOV R0, #5 //to display 5
7: MOV R0, #6 //to display 6
8: MOV R0, #7 //to display 7
9: MOV R0, #8 //to display 8
10: MOV R0, #9 //to display 9

Run Debug Reset

Pin Diagram **Memory** **Peripherals** **Editor**

Current Line: Next Line (PC) Error Line

1: MOV R0, #0 //to display 0
2: MOV R0, #1 //to display 1
3: MOV R0, #2 //to display 2
4: MOV R0, #3 //to display 3
5: MOV R0, #4 //to display 4
6: MOV R0, #5 //to display 5
7: MOV R0, #6 //to display 6
8: MOV R0, #7 //to display 7
9: MOV R0, #8 //to display 8
10: MOV R0, #9 //to display 9

Run Debug Reset

Pin Diagram **Memory** **Peripherals** **Editor**

Current Line: Next Line (PC) Error Line

1: MOV R0, #0 //to display 0
2: MOV R0, #1 //to display 1
3: MOV R0, #2 //to display 2
4: MOV R0, #3 //to display 3
5: MOV R0, #4 //to display 4
6: MOV R0, #5 //to display 5
7: MOV R0, #6 //to display 6
8: MOV R0, #7 //to display 7
9: MOV R0, #8 //to display 8
10: MOV R0, #9 //to display 9

Run Debug Reset

Pin Diagram **Memory** **Peripherals** **Editor**

Current Line: Next Line (PC) Error Line

1: MOV R0, #0 //to display 0
2: MOV R0, #1 //to display 1
3: MOV R0, #2 //to display 2
4: MOV R0, #3 //to display 3
5: MOV R0, #4 //to display 4
6: MOV R0, #5 //to display 5
7: MOV R0, #6 //to display 6
8: MOV R0, #7 //to display 7
9: MOV R0, #8 //to display 8
10: MOV R0, #9 //to display 9

Run Debug Reset

Pin Diagram **Memory** **Peripherals** **Editor**

Current Line: Next Line (PC) Error Line

1: MOV R0, #0 //to display 0
2: MOV R0, #1 //to display 1
3: MOV R0, #2 //to display 2
4: MOV R0, #3 //to display 3
5: MOV R0, #4 //to display 4
6: MOV R0, #5 //to display 5
7: MOV R0, #6 //to display 6
8: MOV R0, #7 //to display 7
9: MOV R0, #8 //to display 8
10: MOV R0, #9 //to display 9

Run Debug Reset

Pin Diagram **Memory** **Peripherals** **Editor**

Current Line: Next Line (PC) Error Line

1: MOV R0, #0 //to display 0
2: MOV R0, #1 //to display 1
3: MOV R0, #2 //to display 2
4: MOV R0, #3 //to display 3
5: MOV R0, #4 //to display 4
6: MOV R0, #5 //to display 5
7: MOV R0, #6 //to display 6
8: MOV R0, #7 //to display 7
9: MOV R0, #8 //to display 8
10: MOV R0, #9 //to display 9

Run Debug Reset



Editor

Current line Next line (PC) Error line

Pin Diagram Memory Peripherals

8051

1 RCV PB_#20h //to display d
2 RCV PB_#21h //to display e
3 RCV PB_#22h //to display f
4 RCV PB_#23h //to display 7
5 RCV PB_#24h //to display 8
6 RCV PB_#25h //to display 9
7 RCV PB_#26h //to display a
8

Run Debug Reset

Editor

Current line Next line (PC) Error line

Pin Diagram Memory Peripherals

8051

1 RCV PB_#20h //to display d
2 RCV PB_#21h //to display e
3 RCV PB_#22h //to display f
4 RCV PB_#23h //to display 7
5 RCV PB_#24h //to display 8
6 RCV PB_#25h //to display 9
7 RCV PB_#26h //to display a
8

Run Debug Reset

Editor

Current line Next line (PC) Error line

Pin Diagram Memory Peripherals

8051

1 RCV PB_#20h //to display d
2 RCV PB_#21h //to display e
3 RCV PB_#22h //to display f
4 RCV PB_#23h //to display 7
5 RCV PB_#24h //to display 8
6 RCV PB_#25h //to display 9
7 RCV PB_#26h //to display a
8

Run Debug Reset



Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

Pin Diagram

Memory

Peripherals

Editor

CONCLUSION: We learn about 7 segment display and simulations using the 8051 Microcontroller. We then simulated a few examples by writing the code.

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