

Q4

ANS

let L be a regular language and $M = (Q, \Sigma, \delta, q_0, F)$ be a finite automata with n -states. language L is accepted by M . let $w \in L$ and $|w| \geq n$, then w can be written as xyz where

(i) $|y| > 0$

(ii) $|ny| \leq n$

(iii) $xy^iz \in L$ for all $i \geq 0$ here y^i denotes that y is repeated or pumped i times

Interpretation of pumping lemma :

- Pumping lemma gives a necessary condition for an input string to belong to a regular set
- Pumping lemma does not give sufficient condition for a language to be regular.
- Pumping lemma should not be used to establish that a given language is regular.
- Pumping lemma should be used to establish that a given language is not regular.
- It uses the pigeonhole principle.

$$L = \{w = w^R \mid w \in \{a, b\}^*\}$$

let $L = \{w \mid w^R \mid w \in \{a, b\}^*\}$ is a regular language.

let n be a pumping lemma constant

let $z = w$

Represent Z in P.L.C

$$\text{let } Z = a^n b a^n = w^R$$

Z can also be written as

$$Z = uvw \text{ such that } |uv| \leq n \text{ and } |v| \leq 1$$

select values of u, v, w satisfying above two conditions

$$\text{let } u = a^{n-1}$$

$$v = a$$

$$w = ba^n$$

then

$$uv^i w = a^{n-1} (a)^i b a^n$$

For $i=0$, we get

$$uv^0 w = a^{n-1} b a^n \notin L$$

\therefore Resultant $w (a^{n-1} b a^n)$ cannot be reverse of $w(w^R)$.

Hence it contradicts our assumption.

$\therefore L$ is not a regular language.

Q5

~~Q5~~

ANS

Convert given grammar to CNF

$$S \rightarrow ABC \mid BAB$$

$$A \rightarrow aA \mid BAC \mid aaa$$

$$B \rightarrow bBb \mid a \mid D$$

$$C \rightarrow CA \mid AC$$

$$D \rightarrow \epsilon$$

(i) Removing ϵ transition

$$D \rightarrow B$$

$$B \rightarrow D$$

$$\therefore B \rightarrow \epsilon$$

$$S \rightarrow ABC$$

$$B \rightarrow \epsilon$$

$$S \rightarrow AC$$

$$S \rightarrow BAB$$



1: $B \rightarrow \epsilon$

2: $B \rightarrow \epsilon$ (last)

3: both $B \rightarrow \epsilon$

$$S \rightarrow AB$$

$$S \rightarrow BA$$

$$S \rightarrow A$$

$$A \rightarrow BAC$$

$$B \rightarrow \epsilon$$

$$A \rightarrow ac$$

$$B \rightarrow bBb$$

$$\therefore B \rightarrow \epsilon$$

$$B \rightarrow bb$$

$$S \Rightarrow ABC | BAB | AC | AB | BA | A$$

$$A \rightarrow aA | BAC | aqa | aC$$

$$B \rightarrow bBb | a | bb$$

$$C \rightarrow CA | AC$$

(ii) Removing unit production

$$S \rightarrow A$$

$$S \rightarrow aA | BAC | aqa | aC$$

$$\therefore S \rightarrow ABC | BAB | AC | AB | BA | aA | BAC | aqa | aC$$

$$A \rightarrow aA | BAC | aqa | aC$$

$$B \rightarrow bBb | a | bb$$

$$C \rightarrow CA | AC$$

(iii) Removing useless productions

There are no useless productions

For CNF,

$$\text{let } c_a = a$$

$$c_b = b$$

$$\therefore S \rightarrow ABC | BAB | AC | AB | BA | c_a A | B c_a C | c_a c_a c_a | c_a C$$

$$A \rightarrow c_a A | B c_a C | c_a c_a c_a | c_a C$$

$$B \rightarrow c_b B c_b | a | c_b c_b$$

$$C \rightarrow CA | AC$$

$$\text{let } F_1 \rightarrow AB$$

$$F_2 \rightarrow B c_a$$

$$F_3 \Rightarrow a c_a$$

$$F_4 \rightarrow c_b B$$

∴ CNF form :

$$S \rightarrow F_1 C \mid B F_1 (A C \mid A B \mid B A \mid C a) \mid F_2 C \mid F_3 (a \mid C a) C$$

$$A \rightarrow C a A \mid F_2 C \mid F_3 C a \mid C a C$$

$$B \rightarrow F_4 C b \mid a \mid C b C b$$

$$C \rightarrow C a \mid A C$$

$$C a \rightarrow a$$

$$C b \rightarrow b$$

$$F_1 \rightarrow A B$$

$$F_2 \rightarrow B C a$$

$$F_3 \rightarrow C a C a$$

$$F_4 \rightarrow C b R$$

Q6

ANS

A	B	S	C
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

LOGIC: N_c q_{nc} no carry (state with)
 c q_c carry (state with)

$\Sigma \backslash Q$	00	01	10	11
q_{nc}	0	1	1	0
q_c	1	0	0	1

$\Sigma \backslash Q$	00	01	10	11
q_{nc}	q_{nc}	q_{nc}	q_{nc}	q_{nc}
q_c	q_{nc}	q_c	q_c	q_c

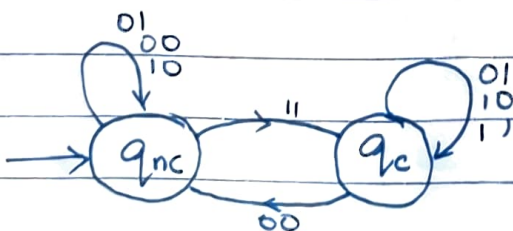
mealy
machine

λ

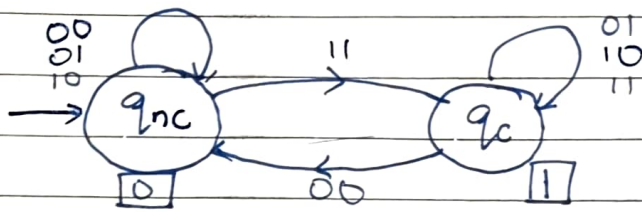
ORIGINAL TRANSITION

Moore machine transition diagram

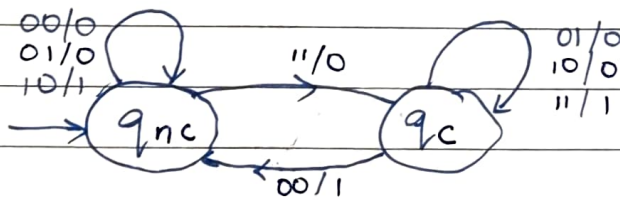
Q	λ
q_{nc}	0
q_c	1



ORIGINAL TRANSITION DIAGRAM



Moore transition diagram.



Mealy transition Diagram