

MATH1061  
Discrete Mathematics I

Problem Set 3  
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Due: 5pm, 4<sup>th</sup> of October, 2024

**Question 1: (10 marks)**

Prove the following set identities

$$(1) (A \cup B) \times C = (A \times C) \cup (B \times C)$$

$$(2) (A' \cap B)' \cap D = (D \setminus A') \cup (D \setminus B)$$

**Solution:** (a)

$$\begin{aligned} (A \cup B) \times C &= \{(x, c) \mid x \in A \cup B, c \in C\} && \text{(Definition of Cartesian Product)} \\ &= \{(x, c) \mid (x \in A \vee x \in B), c \in C\} && \text{(Definition of Union)} \\ &= \{(x, c) \mid (x \in A, c \in C) \vee (x \in B, c \in C)\} && \text{(Distributivity of } \wedge \text{ over } \vee) \\ &= \{(a, c) \mid a \in A, c \in C\} \cup \{(b, c) \mid b \in B, c \in C\} && \text{(Definition of Union)} \\ &= (A \times C) \cup (B \times C) && \text{(Definition of Cartesian Product)} \end{aligned}$$

Which is what we wanted to show.

**Solution:** (b)

$$\begin{aligned} (A' \cap B)' \cap D &= (A \cup B') \cap D && \text{(De Morgan's)} \\ &= D \cap (A \cup B') && \text{(Commutativity)} \\ &= (D \cap A) \cup (D \cap B') && \text{(Distributivity)} \\ &= (D \setminus A') \cup (D \setminus B) && \text{(Set Difference Law)} \end{aligned}$$

Which is what we wanted to show.

**Question 2: (15 marks)**

Suppose  $f : A \rightarrow B$  and  $g : B \rightarrow A$  are functions, and  $\iota_A$  is the identity function on  $A$ ,  $\iota_B$  is the identity function on  $B$ . In particular,  $\iota_A(x) = x$ ,  $\forall x \in A$  and similarly  $\iota_B(x) = x$ ,  $\forall x \in B$ .

(1) Suppose  $f \circ g = \iota_B$  and  $g \circ f = \iota_A$ . Prove that  $f$  and  $g$  are bijections.

(2) Suppose  $g$  is surjective and  $f \circ g = \iota_B$ . Prove that  $g \circ f = \iota_A$ .

**Solution:** (a)

**Proposition.**  $f$  is a bijection

*Proof.* Show that  $f$  is injective.

Suppose  $x, y \in A$ ,  $f(x) = f(y)$ .

We can apply  $g$  to both sides,  $g(f(x)) = g(f(y))$ .

Which is the same as writing  $(g \circ f)(x) = (g \circ f)(y)$ .

Given that  $g \circ f = \iota_A$ , the previous expression can be rewritten  $\iota_A(x) = \iota_A(y)$ .

Evaluating the identity function,  $x = y$ .

Therefore,  $f$  is injective.

Show that  $f$  is surjective

Suppose  $b \in B$ . Take  $a = g(b) \in A$ , since  $g : B \rightarrow A$ .

Hence,  $f(a) = f(g(b)) = (f \circ g)(b) = \iota_B(b) = b$ .

Thus,  $\forall b \in B, \exists a \in A : b = f(a)$ .

Therefore,  $f$  is surjective.

$f$  is surjective and injective.

Therefore,  $f$  is bijective. □

**Proposition.**  $g$  is a bijection

*Proof.* Show that  $g$  is injective.

Suppose  $x, y \in B$ ,  $g(x) = g(y)$ .

We can apply  $f$  to both sides,  $f(g(x)) = f(g(y))$ .

Which is the same as writing  $(f \circ g)(x) = (f \circ g)(y)$ .

Given that  $f \circ g = \iota_B$ , the previous expression can be rewritten  $\iota_B(x) = \iota_B(y)$ .

Evaluating the identity function,  $x = y$ .

Therefore,  $g$  is injective.

Show that  $g$  is surjective

Suppose  $a \in A$ . Take  $b = f(a) \in B$ , since  $f : A \rightarrow B$ .

Hence,  $g(b) = g(f(a)) = (g \circ f)(a) = \iota_A(a) = a$ .

Thus,  $\forall a \in A, \exists b \in B : a = g(b)$ .

Therefore,  $g$  is surjective.

$g$  is surjective and injective.

Therefore,  $g$  is bijective. □

**Corollary.**  $f$  is a bijection and  $g$  is a bijection.

**Solution:** (b)

**Proposition.**  $g \circ f = \iota_A$

*Proof.* Suppose  $g$  is surjective and  $f \circ g = \iota_B$ .

$g$  is surjective, which means  $\forall a \in A, \exists b \in B : a = g(b)$ .

Take any  $\alpha \in A$ . There is a  $\beta \in B$  such that  $\alpha = g(\beta)$ .

Apply  $f$  to both sides,  $f(\alpha) = f(g(\beta))$ .

Apply  $g$  to both sides,  $g(f(\alpha)) = g(f(g(\beta)))$ .

Simplifying,  $(g \circ f)(\alpha) = g((f \circ g)(\beta))$ .

Simplify further,  $(g \circ f)(\alpha) = g(\iota_B(\beta))$ .

We see,  $(g \circ f)(\alpha) = g(\beta)$ .

Hence,  $(g \circ f)(\alpha) = \alpha$ , which is the definition of  $\iota_A$ .

Therefore  $g \circ f = \iota_A$ .

□

**Question 3: (15 marks)**

- (1) Show that  $\mathbb{Z}^+ \times \mathbb{Z}^+$  is in bijection with  $\mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+$ . Deduce that  $\mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+$  is countable.
- (2) Show that  $\mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+$  is in bijection with  $\mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+$ . Deduce that  $\mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+ \times \mathbb{Z}^+$  is countable.
- (3) Is the set  $\mathbb{Z}^+ \times \cdots \times \mathbb{Z}^+$  (Cartesian product  $n$  times) countable? A yes or no would suffice.

*Remark.* In this this question, we'll use  $\mathbb{Z}^+ =: \mathbb{N} = \{1, 2, 3, 4, \dots\}$ . We will denote the Cartesian product  $n$  times,  $\mathbb{N} \times \cdots \times \mathbb{N}$ , as  $\mathbb{N}^n$ . We'll also denote an injective function from set  $A$  to  $B$ ,  $f : A \rightarrow B$ , and a bijective function  $f : A \leftrightarrow B$ .

**Solution:** (a)

**Proposition.**  $\mathbb{N}^2$  is in bijection with  $\mathbb{N}^3$ .

*Proof.* We'll utilise the Schröder-Bernstein theorem, which states (with notation adapted for our specific problem)

$$f : \mathbb{N}^2 \rightarrow \mathbb{N}^3, g : \mathbb{N}^3 \rightarrow \mathbb{N}^2 \implies h : \mathbb{N}^2 \leftrightarrow \mathbb{N}^3$$

So, to prove that  $\mathbb{N}^2$  is in bijection with  $\mathbb{N}^3$ , we'll find two functions which map from one set to the other, and show that those functions are injective.

Show that there exists an injective function  $f : \mathbb{N}^2 \rightarrow \mathbb{N}^3$ .

Let's propose the function  $f : \mathbb{N}^2 \rightarrow \mathbb{N}^3$ , defined by  $f((a, b)) := (a, b, 0)$ .

Suppose  $(a_1, b_1), (a_2, b_2) \in \mathbb{N}^2$ , and  $f(a_1, b_1) = f(a_2, b_2)$ .

Then  $(a_1, b_1, 0) = (a_2, b_2, 0)$

Hence  $a_1 = a_2, b_1 = b_2, 0 = 0$ .

Therefore our proposed function  $f : \mathbb{N}^2 \rightarrow \mathbb{N}^3$  is injective.

Show that there exists an injective function  $g : \mathbb{N}^3 \rightarrow \mathbb{N}^2$ .

Let's propose the function  $g : \mathbb{N}^3 \rightarrow \mathbb{N}^2$ , defined by  $g((a, b, c)) := (2^a 3^b 5^c, 0)$ .

Suppose  $(a_1, b_1, c_1), (a_2, b_2, c_2) \in \mathbb{N}^3$  and  $g((a_1, b_1, c_1)) = g((a_2, b_2, c_2))$ .

Then  $(2^{a_1} 3^{b_1} 5^{c_1}, 0) = (2^{a_2} 3^{b_2} 5^{c_2}, 0)$

Hence,  $2^{a_1} 3^{b_1} 5^{c_1} = 2^{a_2} 3^{b_2} 5^{c_2}$ .

Which implies that  $2^{a_1} = 2^{a_2} \iff a_1 = a_2$ ,

$3^{b_1} = 3^{b_2} \iff b_1 = b_2$ ,

and  $5^{c_1} = 5^{c_2} \iff c_1 = c_2$ .

Therefore our proposed function  $g : \mathbb{N}^3 \rightarrow \mathbb{N}^2$  is injective.

There exists an injection  $f : \mathbb{N}^2 \rightarrow \mathbb{N}^3$ , namely  $f((a, b)) = (a, b, 0)$ , and an injection  $g : \mathbb{N}^3 \rightarrow \mathbb{N}^2$ , namely  $g((a, b, c)) = (2^a 3^b 5^c, 0)$ .

Therefore, by Schröder-Bernstein theorem, there exists a bijection,  $h : \mathbb{N}^2 \leftrightarrow \mathbb{N}^3$ .

Therefore  $\mathbb{N}^2$  is in bijection with  $\mathbb{N}^3$ . □

**Corollary.**  $\mathbb{N}^3$  is countable.

*Proof.*

$$\exists f : \mathbb{N}^2 \leftrightarrow \mathbb{N} \iff |\mathbb{N}^2| = |\mathbb{N}|.$$

$$\exists g : \mathbb{N}^2 \leftrightarrow \mathbb{N}^3 \iff |\mathbb{N}^2| = |\mathbb{N}^3|.$$

$\therefore |\mathbb{N}^3| = |\mathbb{N}|$ , by transitivity. This is the definition of countable.

Therefore  $\mathbb{N}^3$  is countable. □

**Solution:** (b)

**Proposition.**  $\mathbb{N}^3$  is in bijection with  $\mathbb{N}^4$ .

*Proof.* We'll utilise the Schröder-Bernstein theorem, which states (with notation adapted for our specific problem)

$$f : \mathbb{N}^3 \rightarrow \mathbb{N}^4, g : \mathbb{N}^4 \rightarrow \mathbb{N}^3 \implies h : \mathbb{N}^3 \leftrightarrow \mathbb{N}^4$$

So, to prove that  $\mathbb{N}^3$  is in bijection with  $\mathbb{N}^4$ , we'll find two functions which map from one set to the other, and show that those functions are injective.

Show that there exists an injective function  $f : \mathbb{N}^3 \rightarrow \mathbb{N}^4$ .

Let's propose the function  $f : \mathbb{N}^3 \rightarrow \mathbb{N}^4$ , defined by  $f((a, b, c)) := (a, b, c, 0)$ .

Suppose  $(a_1, b_1, c_1), (a_2, b_2, c_1) \in \mathbb{N}^3$ , and  $f(a_1, b_1, c_1) = f(a_2, b_2, c_2)$ .

Then  $(a_1, b_1, c_1, 0) = (a_2, b_2, c_2, 0)$

Hence  $a_1 = a_2, b_1 = b_2, c_1 = c_2, 0 = 0$ .

Therefore our proposed function  $f : \mathbb{N}^3 \rightarrow \mathbb{N}^4$  is injective.

Show that there exists an injective function  $g : \mathbb{N}^4 \rightarrow \mathbb{N}^3$ .

Let's propose the function  $g : \mathbb{N}^4 \rightarrow \mathbb{N}^3$ , defined by  $g((a, b, c, d)) := (2^a 3^b 5^c 7^d, 0, 0)$ .

Suppose  $(a_1, b_1, c_1, d_1), (a_2, b_2, c_2, d_2) \in \mathbb{N}^4$  and  $g((a_1, b_1, c_1, d_1)) = g((a_2, b_2, c_2, d_2))$ .

Then  $(2^{a_1} 3^{b_1} 5^{c_1} 7^{d_1}, 0, 0) = (2^{a_2} 3^{b_2} 5^{c_2} 7^{d_2}, 0, 0)$

Hence,  $2^{a_1} 3^{b_1} 5^{c_1} 7^{d_1} = 2^{a_2} 3^{b_2} 5^{c_2} 7^{d_2}$ .

Which implies that  $2^{a_1} = 2^{a_2} \iff a_1 = a_2$ ,

$3^{b_1} = 3^{b_2} \iff b_1 = b_2$ ,

$5^{c_1} = 5^{c_2} \iff c_1 = c_2$ ,

and  $7^{d_1} = 7^{d_2} \iff d_1 = d_2$ .

Therefore our proposed function  $g : \mathbb{N}^4 \rightarrow \mathbb{N}^3$  is injective.

There exists an injection  $f : \mathbb{N}^3 \rightarrow \mathbb{N}^4$ , namely  $f((a, b, c)) = (a, b, c, 0)$ , and an injection  $g : \mathbb{N}^4 \rightarrow \mathbb{N}^3$ , namely  $g((a, b, c, d)) = (2^a 3^b 5^c 7^d, 0, 0)$ .

Therefore, by Schröder-Bernstein theorem, there exists a bijection,  $h : \mathbb{N}^3 \leftrightarrow \mathbb{N}^4$ .

Therefore  $\mathbb{N}^3$  is in bijection with  $\mathbb{N}^4$ . □

**Corollary.**  $\mathbb{N}^4$  is countable.

*Proof.*

$$\exists f : \mathbb{N}^2 \leftrightarrow \mathbb{N} \iff |\mathbb{N}^2| = |\mathbb{N}|.$$

$$\exists g : \mathbb{N}^2 \leftrightarrow \mathbb{N}^3 \iff |\mathbb{N}^2| = |\mathbb{N}^3|.$$

$$\exists h : \mathbb{N}^3 \leftrightarrow \mathbb{N}^4 \iff |\mathbb{N}^3| = |\mathbb{N}^4|.$$

$\therefore |\mathbb{N}^4| = |\mathbb{N}|$ , by transitivity. This is the definition of countable.

Therefore  $\mathbb{N}^4$  is countable. □

**Solution:** (c) Yes.

I can always construct an injective function  $f : \mathbb{N}^{n-1} \rightarrow \mathbb{N}^n$ ,  $f((a_1, \dots, a_{n-1})) = (a_1, \dots, a_{n-1}, 0)$  and a second injective function  $g : \mathbb{N}^n \rightarrow \mathbb{N}^{n-1}$ ,  $g(a_1, \dots, a_n) = (\prod_{i=1}^n p_i^{a_i}, 0, \dots, 0)$ , which will inductively be in bijection with  $\mathbb{N}$

**Question 4: (10 marks)**

Let  $A$  be the set of all logical statements. Define a relation on  $A$ : for  $p, q \in A$ ,  $p$  is related to  $q$  if and only if  $p \wedge q$  and  $p \vee q$  have the same truth value.

Determine if the above relation is reflexive, symmetric, transitive. If your answer is yes for any of the three properties, please prove your answer; if your answer is no, please find a counterexample.

**Solution:**

Let  $\sigma : A \rightarrow A$  be defined By

$$\begin{aligned}
 \forall p, q \in A, p \sigma q &\iff p \wedge q \leftrightarrow p \vee q && \text{(Given)} \\
 &\iff ((p \wedge q) \rightarrow (p \vee q)) \wedge ((p \vee q) \rightarrow (p \wedge q)) && \text{(Definition of } \leftrightarrow \text{)} \\
 &\iff (\sim(p \wedge q) \vee (p \vee q)) \wedge (\sim(p \vee q) \vee (p \wedge q)) && \text{(Definition of } \rightarrow \text{)} \\
 &\iff ((\sim p \vee \sim q) \vee (p \vee q)) \wedge ((\sim p \wedge \sim q) \vee (p \wedge q)) && \text{(De Morgan's)} \\
 &\iff ((\sim p \vee p) \vee (\sim q \vee q)) \wedge ((\sim p \wedge \sim q) \vee (p \wedge q)) && \text{(Associativity)} \\
 &\iff (\top \vee \top) \wedge ((\sim p \wedge \sim q) \vee (p \wedge q)) && \text{(Negation Law)} \\
 &\iff \top \wedge ((\sim p \wedge \sim q) \vee (p \wedge q)) && \text{(Tautology)} \\
 &\iff (\sim p \wedge \sim q) \vee (p \wedge q) && \text{(Identity)} \\
 &\iff ((\sim p \vee p) \wedge (\sim q \vee p)) \wedge ((\sim p \vee q) \wedge (\sim q \vee q)) && \text{(Distributivity, twice)} \\
 &\iff (\top \wedge (\sim q \vee p)) \wedge ((\sim p \vee q) \wedge \top) && \text{(Negation Laws)} \\
 &\iff (\sim q \vee p) \wedge (\sim p \vee q) && \text{(Identity)} \\
 &\iff (q \rightarrow p) \wedge (p \rightarrow q) && \text{(Definition of } \rightarrow \text{)} \\
 &\iff p \leftrightarrow q && \text{(Definition of } \leftrightarrow \text{)}
 \end{aligned}$$

Is  $\sigma$  reflexive?

Suppose  $p \in A$ .

Then  $p \equiv \text{True}$  or  $p \equiv \text{False}$ .

In either case,  $p \leftrightarrow p$ .

So  $\forall p \in A, p \sigma p$

Therefore  $\sigma$  on  $A$  is reflexive.

Is  $\sigma$  symmetric?

Suppose  $p, q \in A$ .

Then  $p \equiv \text{True}$  or  $p \equiv \text{False}$ . And  $q \equiv \text{True}$  or  $q \equiv \text{False}$ .

Without any loss of generality, assume  $p \equiv q$ .

(if  $p \not\equiv q$ , then  $p \not\leftrightarrow q$  hence  $p \sigma q$  does not hold, and doesn't need to be considered.)

Suppose  $p \sigma q$ .

Then  $p \leftrightarrow q$ . Hence,  $q \leftrightarrow p$ , by commutativity of  $\leftrightarrow$ . Therefore,  $q \sigma p$ .

Therefore  $\sigma$  on  $A$  is symmetric.

Is  $\sigma$  transitive?

Suppose  $p, q, r \in A$ .

Then  $p \equiv \text{True}$  or  $p \equiv \text{False}$ .  $q \equiv \text{True}$  or  $q \equiv \text{False}$ . And  $r \equiv \text{True}$  or  $r \equiv \text{False}$ .

Without any loss of generality assume  $p \equiv q$  and  $q \equiv r$

(If  $p \not\equiv q$  then  $p \sigma q$  does not hold and need not be considered. If  $q \not\equiv r$  then  $q \sigma r$  does not hold and need not be considered.)

Suppose  $p \sigma q$  and  $q \sigma r$ .

Then  $p \leftrightarrow q$  and  $q \leftrightarrow r$ . Therefore  $p \leftrightarrow r$ , by transitivity of  $\leftrightarrow$ . Hence  $p \sigma r$ .

Therefore  $\sigma$  on  $A$  is transitive.

Therefore,  $\sigma$  is an equivalence relation on  $A$ .