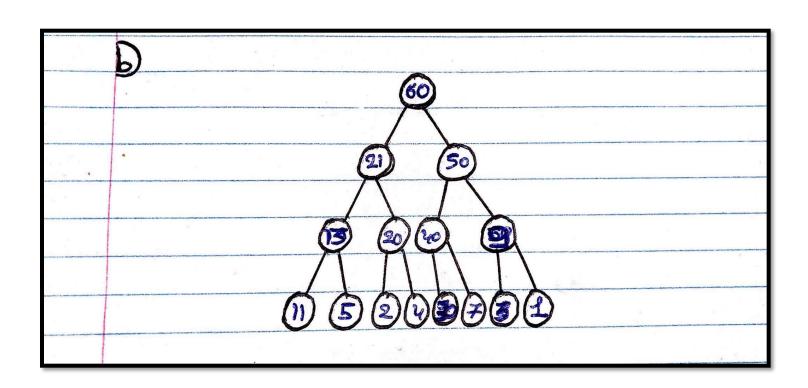
## CSI2110-D - Written Assignment #4

Alae Boufarrachene (300188539)

## Question 1:

a)

index	0	1	2	3	4	/ 5	6	7	8	9	10	11	12	13	14
Keys of A (before the loop)	1	2	3	5	4	<u>7</u> /	9	11	13	20	21	30	40	50	60
Keys of A	せ	2		Ŋ	£	7	60	H	13	20	ข	30	40	50	J
Keys of A	4	2	K	Ŋ	4	૫૦	60	H	13	20	21	30	7	5	LD
Keys of A	1	2	5	Ŋ	21	40	6ව	H	乜	20	Ч	30	干	50	صا
Keys of A	4	2	Ŋ	13	21	40	8	H	5	20	ų	B	7	50	រិ
Keys of A	Ą	27	60	13	ध	40	50	H	5	10	4	59	7	2	صا
Keys of A	¥	य	8	13	30	S	50	H	7	2	4	B	7	3	ف
Keys of A	60	21	50	13	20	40	9	H	5	2	4	30	子	3	ተ



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c)

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
index Keys of A	ᅱ	2	3	7	J	4	ø٦	5	13	20	21	30	<b>9</b>	5	60	业

d)

i	0	1	2	3	4	5	6	7 8	9	10	11	12	13	14
Keys of A	2	4	3	5	20	子	១	41 13	60	21	30	40	50	مىلالا

Question 2:
25
D'Einding Kth smallest Key element in a given away, where Kon:
- We use the array to constuct a minute (retich sorts the elements)
- We remove the root (must be repeated K-1 times)
 - The Kith smallest element should be the root now-
He simply return the root
My solution achines a complexity of O(n+Klog(n)) because the number of operations is directly protestional to the Kingut, inablition, my solution
only requires one loop (the one used to securous and swap the root K-1
times)

We use the away to implement a Max Heap of the Same size.

- We insert the first K elements of the away into the Max Heap.

- Once we reach the Kth element, we use it to compare with the remaining elements of the away, if it here are inferior to the root.

The element replaces the root (repeated until there are no more elements left).

- We will obtain a heap that only contains the K smallest elements.

- We return the root of the Max Heap.

Ity solution achieves a complaint of O (mlog (k)) because only one loop is required to create a Max Heap and to insert the first K element of the array into the newly—created Max Heap.