**COMP 250** 

Lecture 13

mergesort, quicksort

Oct. 6, 2017

# Time complexity

$O(log_2 n)$	O(n)	$O(n^2)$
• convert to binary	<ul> <li>List operations: findMax, remove</li> </ul>	• insertion/selection/ bubble sort
• binary search	<ul> <li>grade school addition or subtraction</li> </ul>	• grade school multiplication
•	•	•

#### Computers perform $\sim 10^9$ operations per second.

$$2^{10} \approx 10^3$$

$$2^{20} \approx 10^6$$

$$2^{30} \approx 10^9$$

### Computers perform $\sim 10^9$ operations per second.

$log_2 n$	n	$n^2$
10	$2^{10} \approx 10^3$	$10^6$
20	$2^{20} \approx 10^6$	10 <sup>12</sup> minuteshours
30	$2^{30} \approx 10^9$	$10^{18}$

second

centuries

## Better sorting algorithms?

$$O(n) < ? < O(n^2)$$

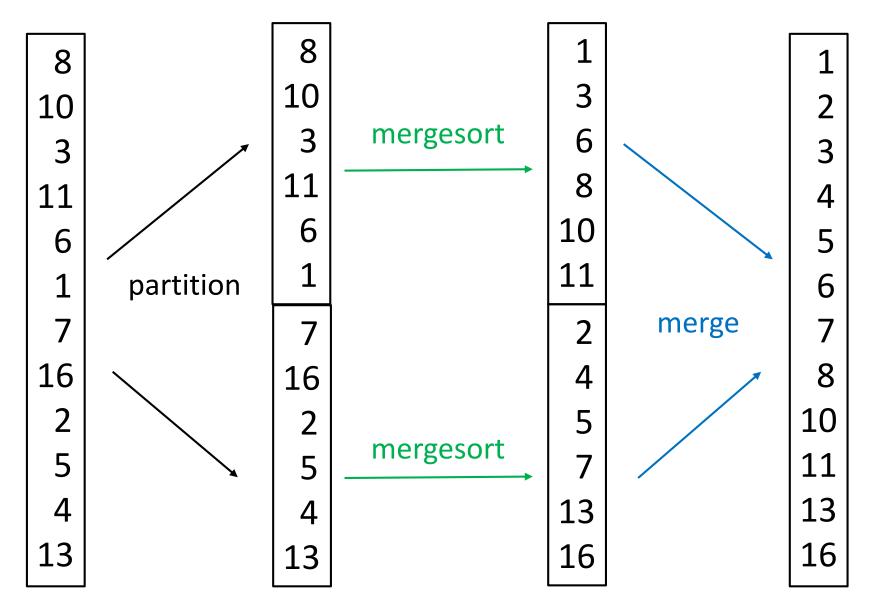
### Mergesort

Given a list, partition it into two halves (1st & 2nd).

Sort each half (recursively).

Merge the two halves.

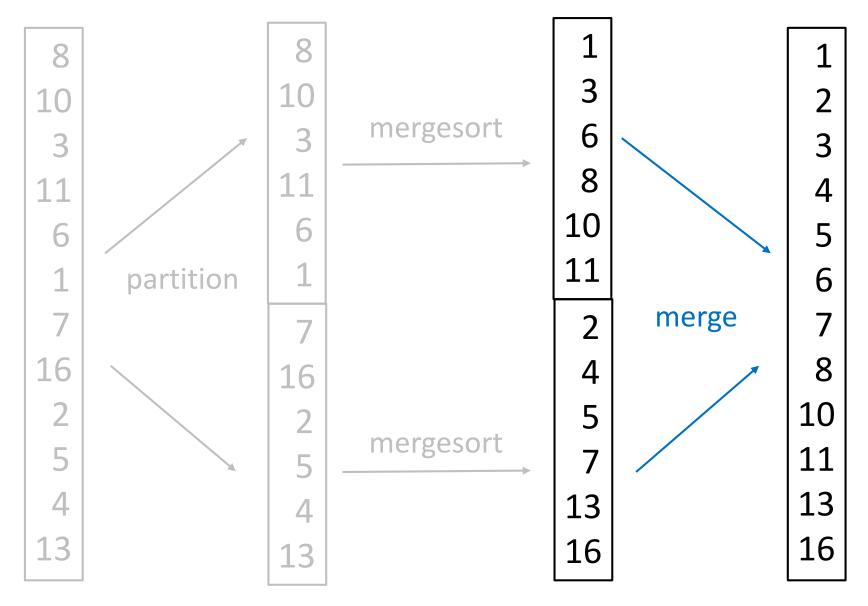
This turns out to be much faster than the other list sorting algorithms we have seen.

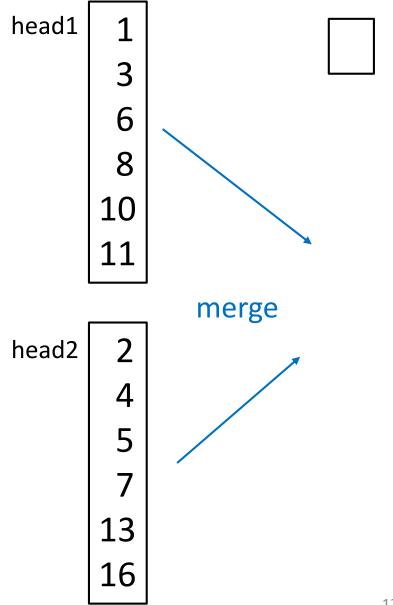


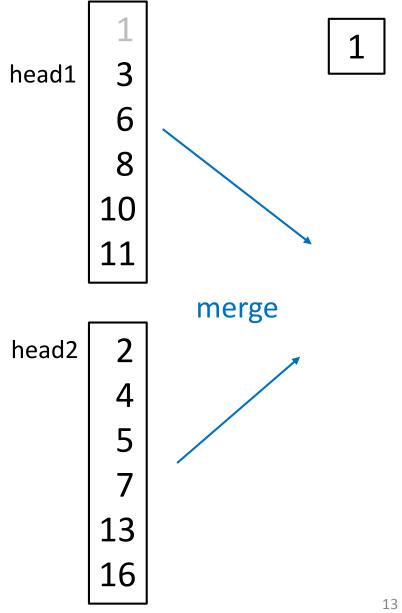
```
mergesort(list){
  if list.length == 1
     return list
  else{
     mid = (list.size - 1) / 2
     list1 = list.getElements(0,mid)
     list2 = list.getElements(mid+1, list.size-1)
     list1 = mergesort(list1)
     list2 = mergesort(list2)
     return merge(list1, list2)
```

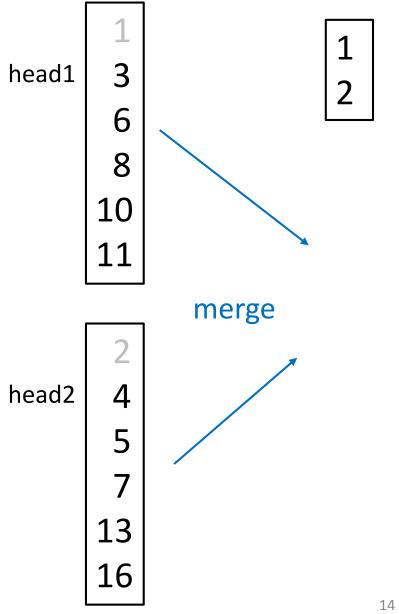
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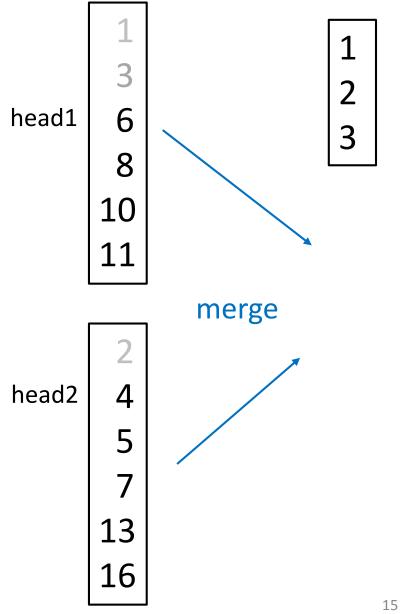
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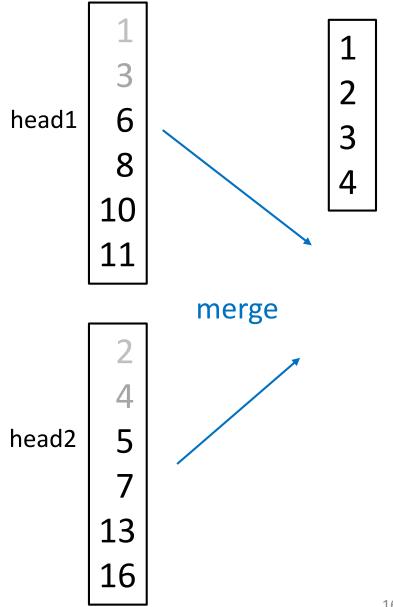


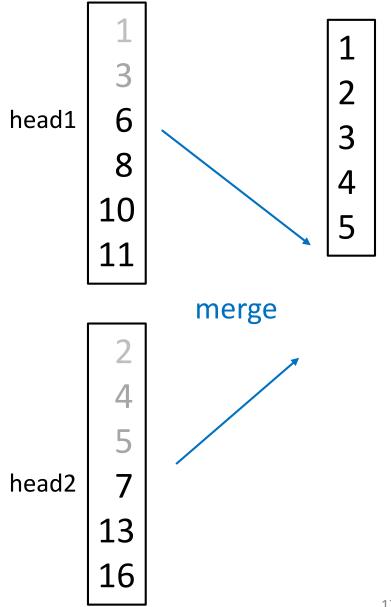


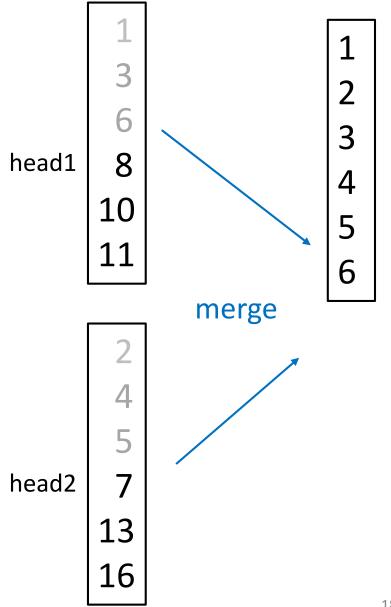


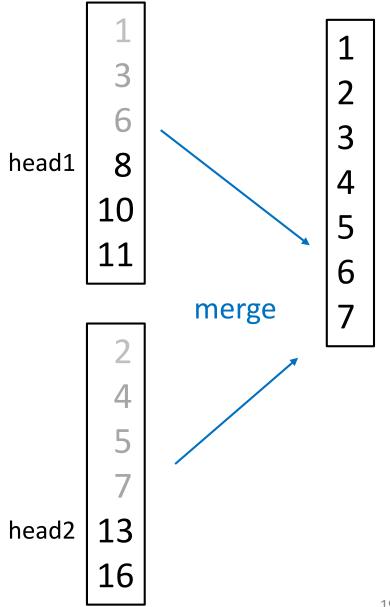




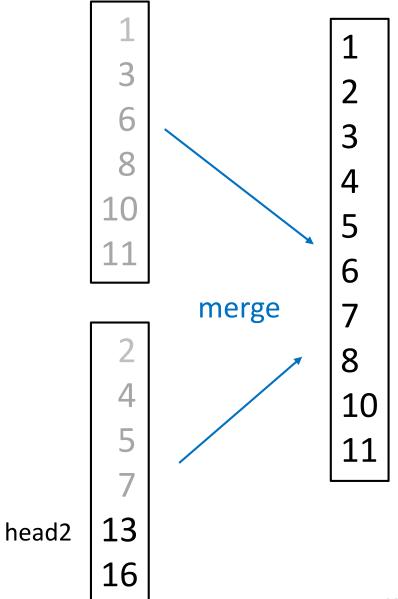




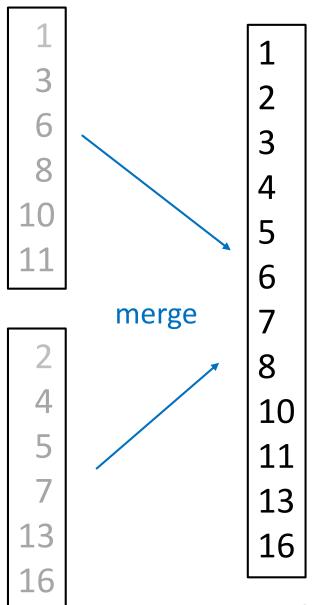




....and so on until one list is empty.



Then, copy the remaining elements.

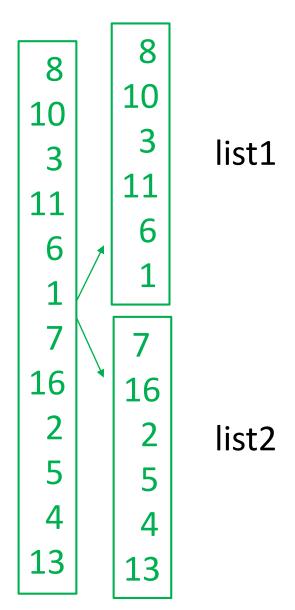


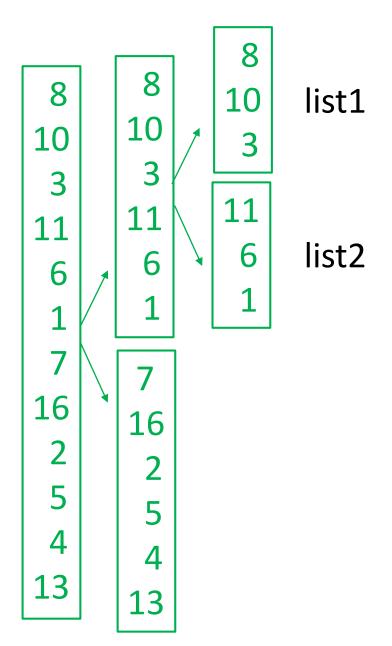
```
merge(list1, list2){
 initialize list to be empty
 while (list1 is not empty) & (list2 is not empty){
    if (list1.first < list2.first)</pre>
        list.addlast( list1.removeFirst(list1) )
    else
        list.addlast( list2.removeFirst(list2) )
 while list1 is not empty
    list.addlast(list1.removeFirst(list1))
 while list2 is not empty
    list.addlast(list2.removeFirst(list2))
 return list
```

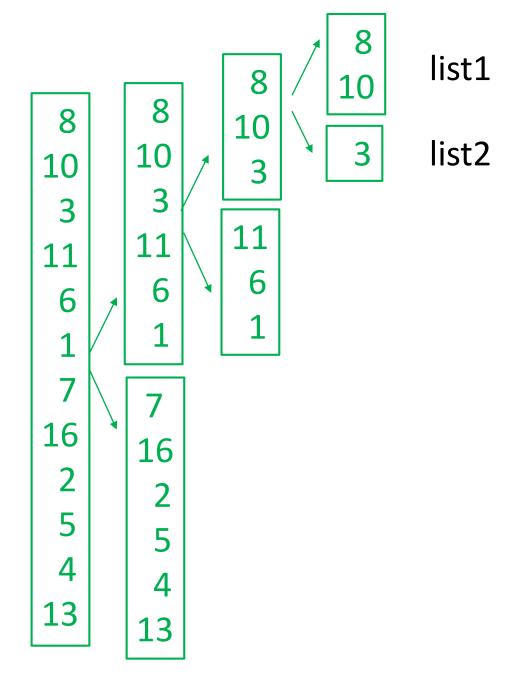
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    else
        list.addlast( list2.removeFirst(list2) )
 while list1 is not empty
    list.addlast( list1.removeFirst(list1) )
 while list2 is not empty
    list.addlast( list2.removeFirst(list2) )
 return list
```

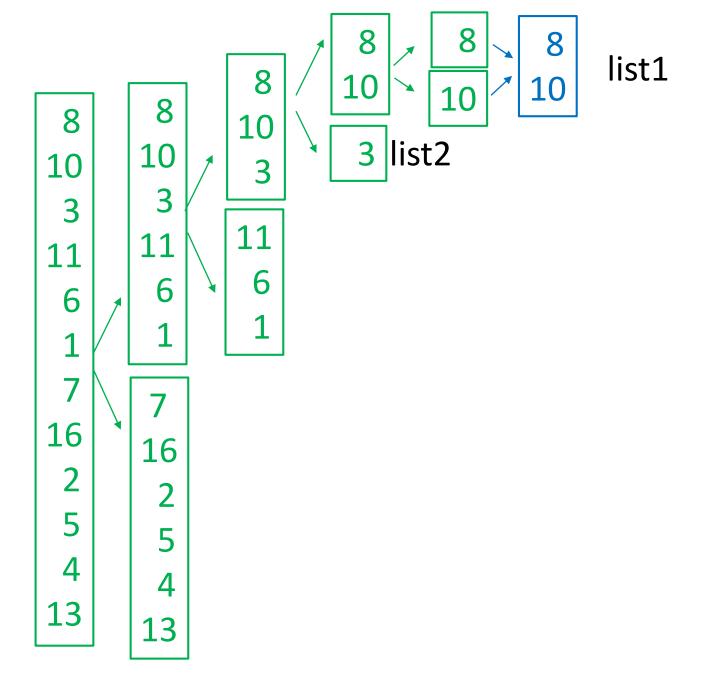
```
10
 6
16
```

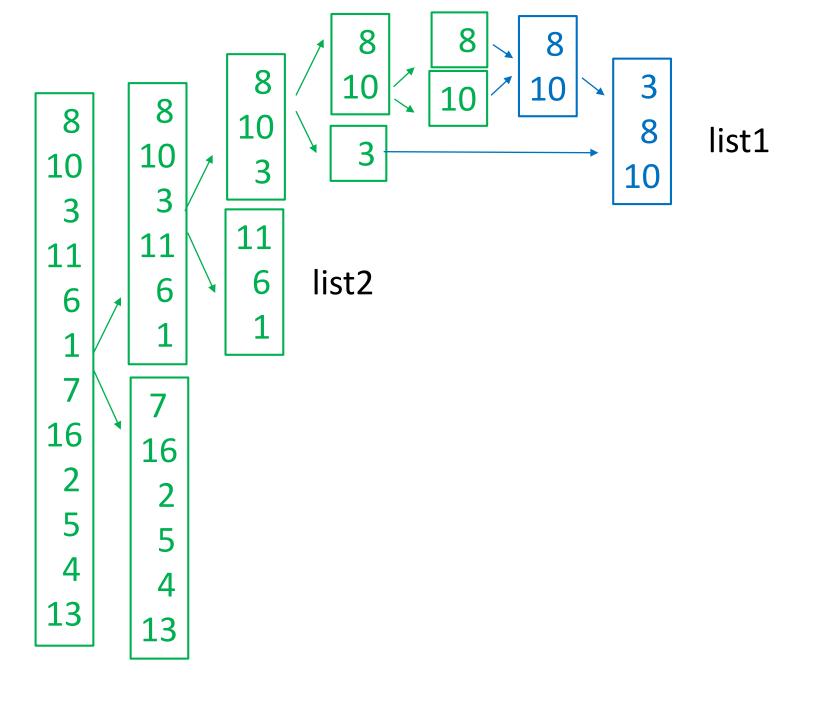
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mergesort(list){
  if list.length == 1
     return list
  else{
     mid = (list.size - 1) / 2
     list1 = list.getElements(0,mid)
     list2 = list.getElements(mid+1, list.size-1)
     list1 = mergesort(list1)
     list2 = mergesort(list2)
     return merge(list1, list2)
```

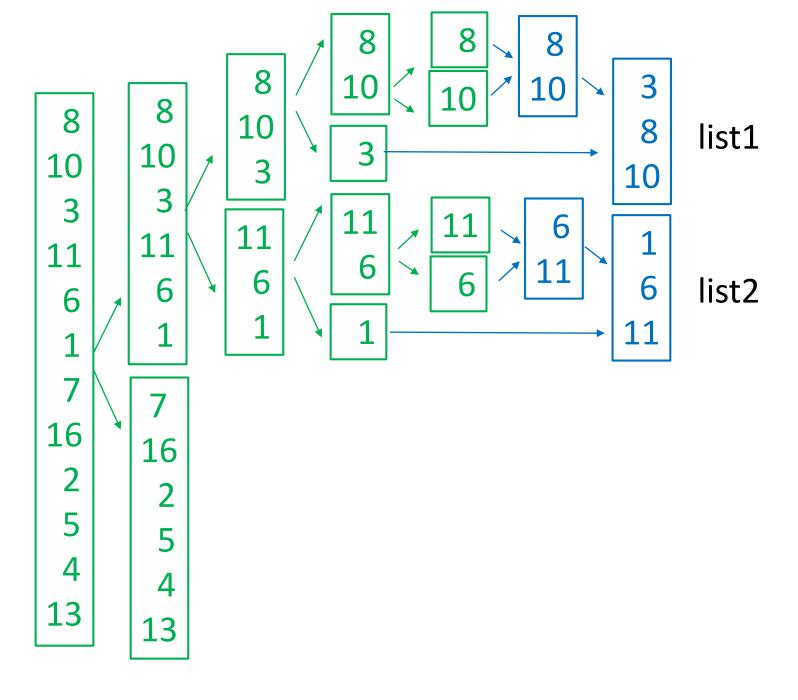


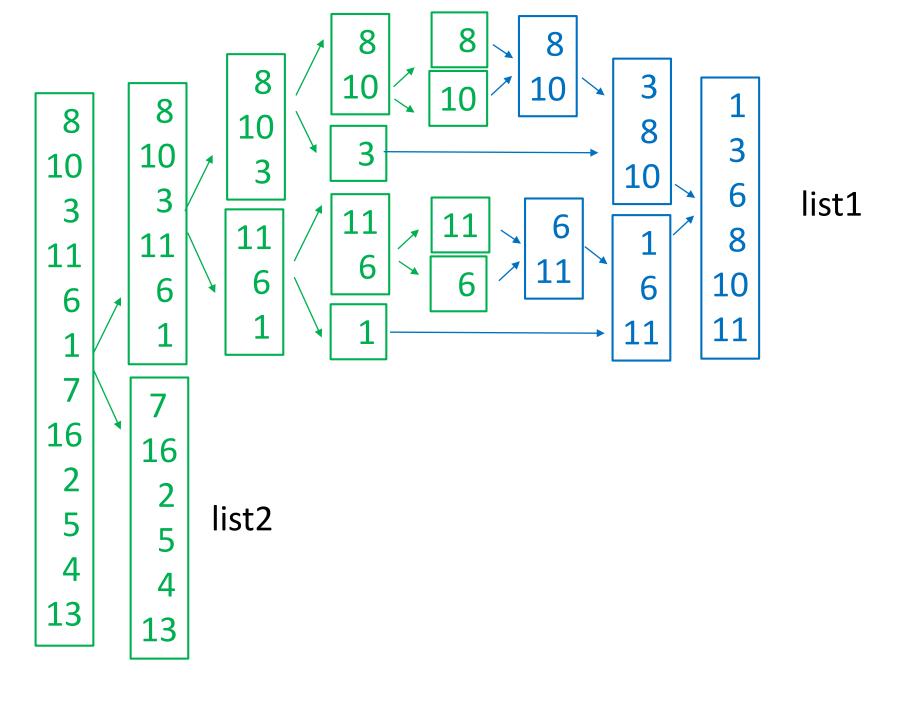


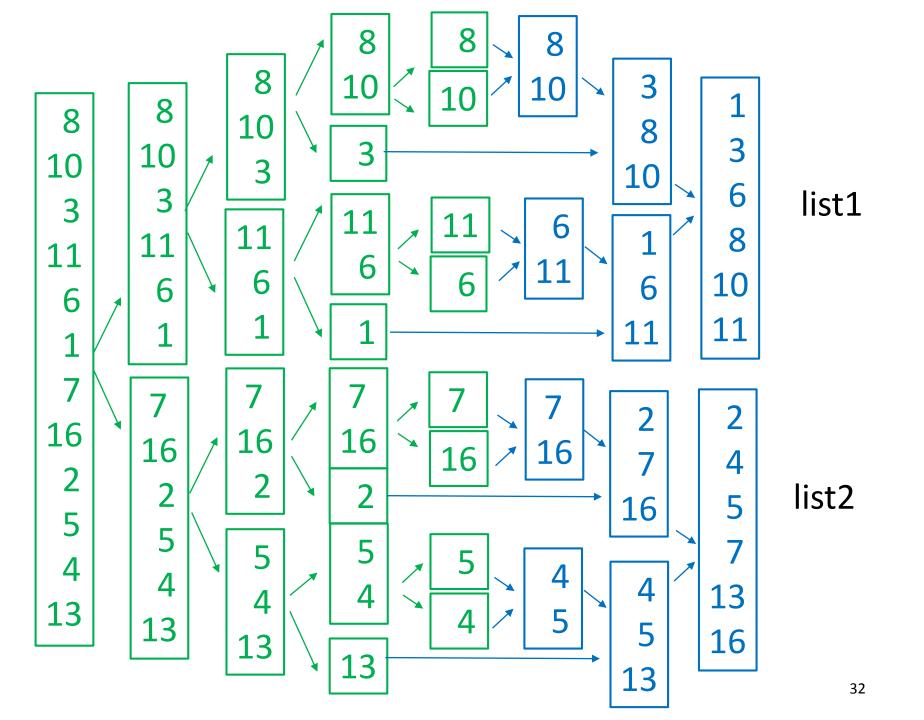


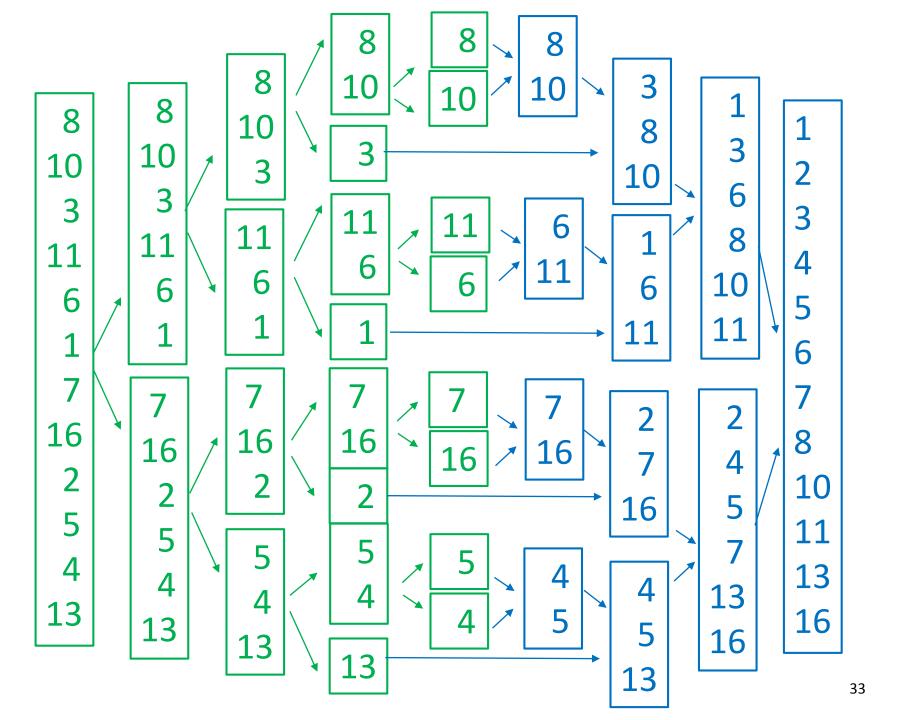








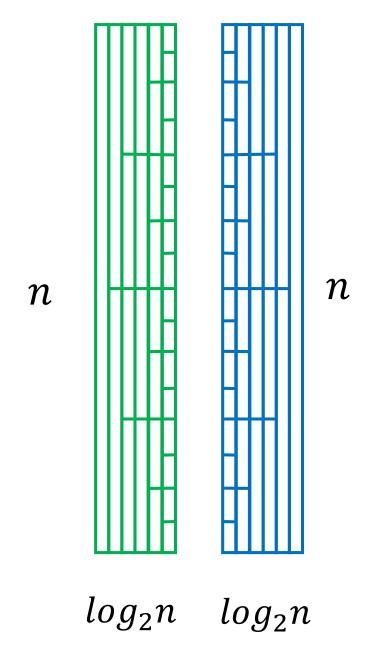




Q: How many operations are required to mergesort a list of size n?

A:  $O(n \log_2 n)$ 

This will become more clear a few lectures from now when we discuss recurrences.



 $n \log_2 n$  is much closer to n than to  $n^2$ 

$log_2 n$	n	$n log_2 n$	$n^2$
10	$2^{10} \approx 10^3$	$\mathbf{10^4}$	$10^6$
20	$2^{20} \approx 10^6$	$\sim 10^7$	$10^{12}$
30	$2^{30} \approx 10^9$	$\sim \! 10^{10}$	$10^{18}$

### Computers perform $\sim 10^9$ operations per second.

$log_2 n$	n	$n \log_2 n$	$n^2$
10	$2^{10} \approx 10^3$	<b>10</b> <sup>4</sup>	$10^6$
20	$2^{20} \approx 10^6$	~10 <sup>7</sup>	10 <sup>12</sup>
30	$2^{30} \approx 10^9$	~10 <sup>10</sup>	10 <sup>18</sup>

milliseconds minutes

hours centuries

$$O(n) < O(n \log_2 n) \ll O(n^2)$$

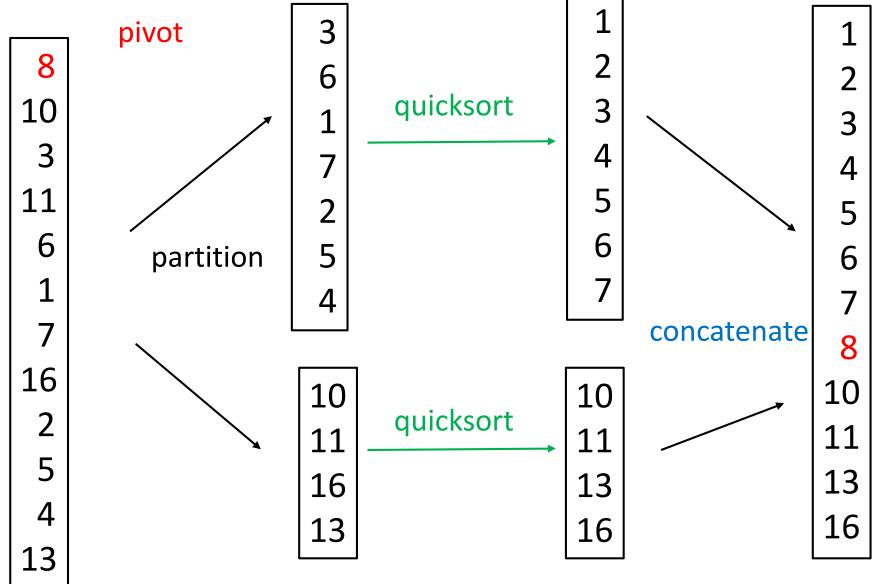
mergesort quicksort bubble sort selection sort insertion sort

### Quicksort

```
quicksort(list){
  if list.length <= 1
    return list
  else{
    pivot = list.removeFirst() // or some other element
    list1 = list.getElementsLessThan(pivot)
    list2 = list.getElementsGreaterOrEqual(pivot)
    list1 = quicksort(list1)
    list2 = quicksort(list2)
  return concatenate(list1, e, list2)
```

### Quicksort

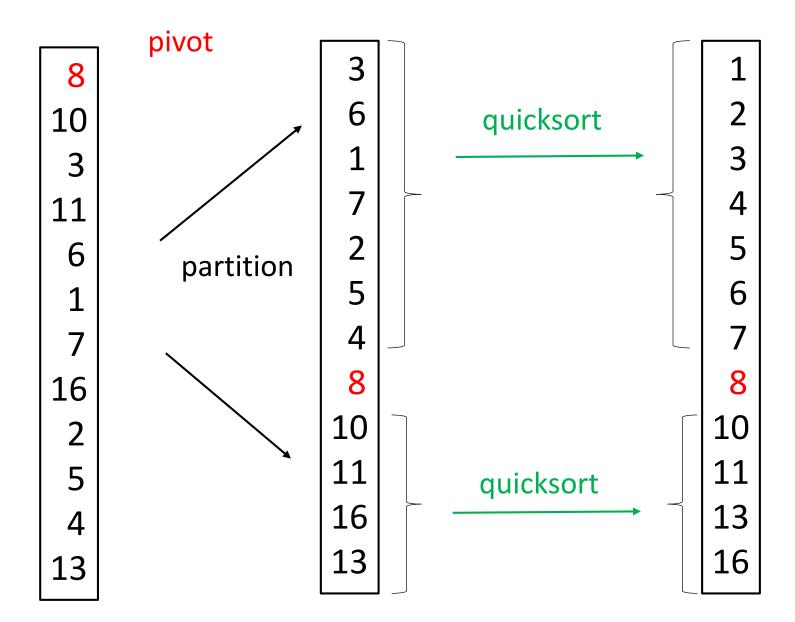
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    list1 = quicksort(list1)
    list2 = quicksort(list2)
  return concatenate(list1, e, list2)
```



# Quicksort can be done "in place"

```
quicksort( low, high ){
  if low > high
    return
  else{
    pivot = ; // select index in {low, ..., high}
    partitionIndex = makePartition (low, high, pivot)
    quicksort(low, partitionIndex - 1)
    quicksort(partionIndex + 1, high)
```

Quicksort partitioning can be done 'in place' using a clever swapping and scanning technique. (See web for details, if interested.)



### Mergesort vs. Quicksort

 Mergesort typically uses an extra list. More space can hurt performance for big lists.

• We will discuss worst case performance of quicksort later in the course.

 See stackoverflow if you want opinions on which is better.