COMP 250

Lecture 9

queue ADT

Sept. 27/28, 2017

ADT (abstract data type)

List add(i,e), remove(i), get(i), set(i),

Stackpush, pop(), ...

Queue enqueue(e), dequeue()

Queue

dequeue (remove from front)

enqueue (add at back)

front

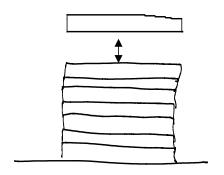
front

Clients

Examples

- keyboard buffer
- printer jobs
- CPU processes (applications do not run in parallel)
- web server

•

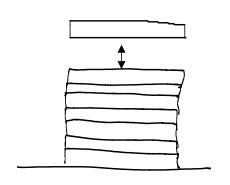


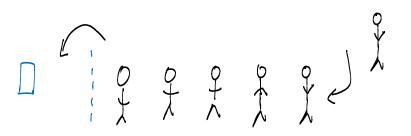
<u>Stack</u>

push(e)

pop()

LIFO (last in, first out)





<u>Stack</u>

push(e)

pop()

LIFO (last in, first out)

Queue

enqueue(e)

dequeue()

FIFO

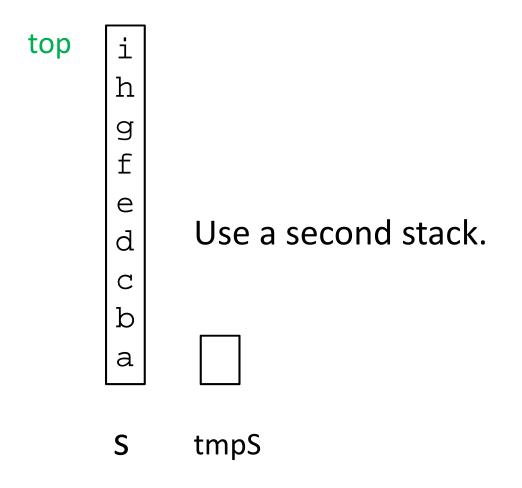
(first in, first out)

"first come, first serve"

Exercise: Use stack(s) to implement a queue.

Write pseudocode for these two methods that uses a stack, namely use the operations push(e), pop(), isEmpty().

Hint for Exercise



Hint for Exercise

top a h b g C while (!s.isEmpty()){ d tmpS.push(s.pop()) e £ d g h b i a

S tmpS s tmpS

Queue Example

```
enqueue(a) a ab dequeue() b
```

Queue Example

```
enqueue(a)
                          a
enqueue(b)
                          ab
dequeue()
enqueue(c)
                           bc
enqueue( d )
                           bcd
enqueue(e)
                          bcde
dequeue()
                          cde
enqueue(f)
                          cdef
enqueue(g)
                          cdefg
```

How to implement a queue?

	enqueue(e)	dequeue()
singly linked list		
doubly linked list		
array list		

How to implement a queue?

enqueue(e) dequeue()

singly linked list

doubly linked list

array list

addLast(e) removeFirst()

(unnecessary)

How to implement a queue?

enqueue(e) dequeue()

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(unnecessary)

addLast(e) removeFirst()

SLOW

Implementing a queue with an array list. (BAD)

Implementing a queue with an array list. (BAD)

Implementing a queue with an array list. (BAD)

0123 indices length = 4enqueue(a) enqueue(b) dequeue() enqueue (c) enqueue (d) bcdenqueue (e) bcde dequeue() cdeenqueue(f) cdef cdefg--requires expansion enqueue (g)

Implementing a queue with an **expanding array.** (also BAD)

Use head and tail indices (tail = head + size - 1)

```
enqueue( a ) | a--- | (0,0) | enqueue( b ) | ab-- | (0,1) | dequeue( ) | -b-- | (1,1) | enqueue( c ) | -bc- | (1,2) | enqueue( d ) | -bcd | (1,3) | enqueue( e ) | ?
```

Implementing a queue with an **expanding array**. (also BAD)

Use head and tail indices (tail = head + size - 1)

```
(0,0)
enqueue(a)
enqueue (b)
                              (0,1)
dequeue()
                              (1,1)
                  -b--
enqueue(c)
                  -bc-
                              (1,2)
                  -bcd
                              (1,3)
enqueue (d
                                        Make bigger
                  -bcde---
enqueue (e)
                              (1,4)
                                        array and
                                        copy to it.
                  --cde--
                              (2,4)
dequeue (
enqueue (f)
                  --cdef--
                              (2,5)
enqueue (g)
                  --cdefg-
                              (2,6)
                                      19
```

An expanding array is an inefficient usage of space.

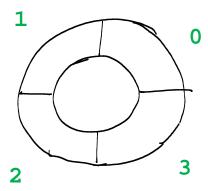
A better idea is....

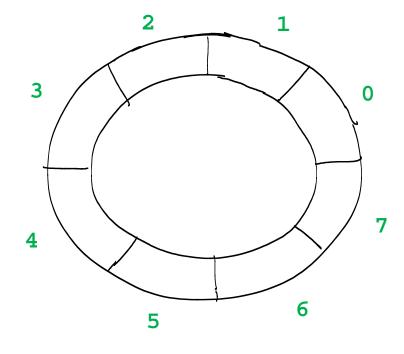
length = 4

0123

length = 8

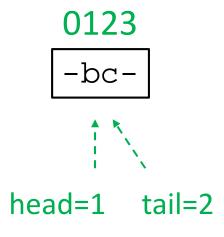
01234567

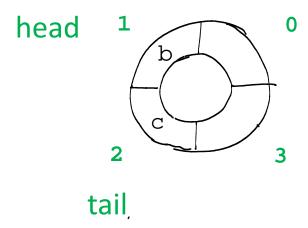




tail = (head + size - 1) % length

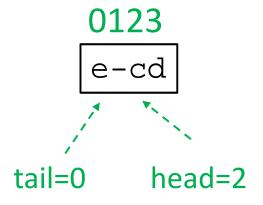
```
enqueue( a
enqueue( b )
dequeue()
enqueue( c )
```

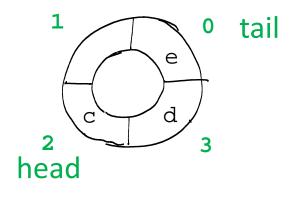




tail = (head + size - 1) % length

```
enqueue(a)
enqueue(b)
dequeue()
enqueue(c)
enqueue(d)
enqueue(e)
```





```
tail = (head + size - 1) \% length
```

```
0123

e-cd

tail=0 head=2
```

```
enqueue( element ){
    if (size < length)
       queue[ (tail + 1) % length] = element
    else .... // coming up
    size = size + 1
dequeue(){ // check if empty omitted
   element = queue[head]
    head = (head + 1) \% length
    size = size-1
    return element
```

Implementing a queue with a circular array (GOOD)

tail = (head + size - 1) % length

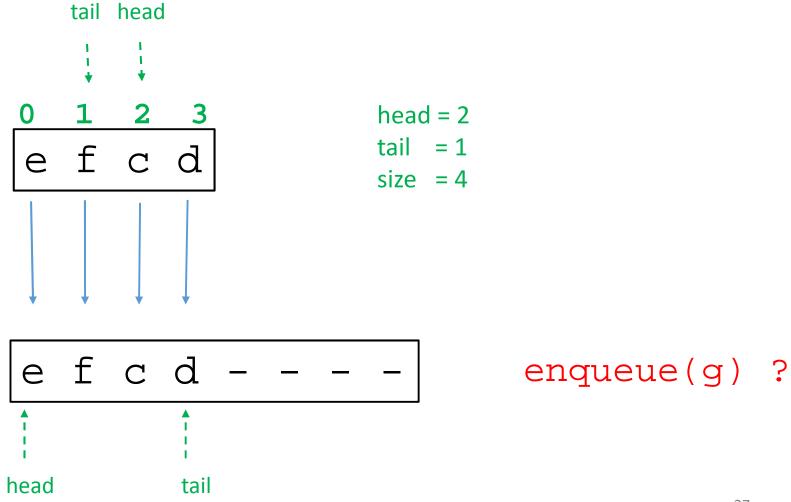
	array	(head, tail,	size)
enqueue(a)	a	(0,0,	1)
enqueue(b)	ab	(0, 1,	2)
dequeue()	-b	(1, 1,	1)
enqueue(c)	-bc-	(1, 2,	2)
enqueue(d)	-bcd	(1, 3,	3)
enqueue(e)	ebcd	(1, 0,	4)
dequeue()	e-cd	(2, 0,	3)
enqueue(f)	efcd	(2, 1,	4)

Implementing a queue with a circular array

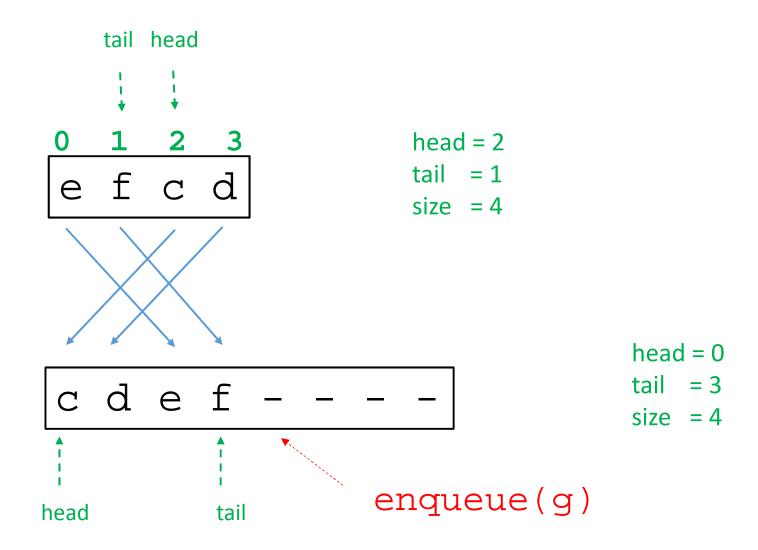
$$tail = (head + size - 1) \% length$$

```
(head, tail, size)
                   array
enqueue (a)
                            (0, 0, 1)
                   a---
enqueue (b)
                            (0, 1, 2)
                   ab--
                            (1, 1, 1)
dequeue()
                   -b--
enqueue(c)
                            (1, 2, 2)
                   -bc-
enqueue (d)
                            (1, 3, 3)
                   -bcd
                            (1, 0, 4)
enqueue (e)
                   ebcd
dequeue()
                   e-cd
                           (2, 0, 3)
                           (2, 1, 4)
                   efcd
enqueue(f)
                      ?
enqueue(g)
```

Increase length of array and copy? **BAD**



Increase length of array. Copy so that head moves to front. (GOOD)



```
enqueue( element ){
   if ( queue.size == queue.length) {
      // increase length of array
      create a bigger array tmp[] // e.g. 2*length
      for i = 0 to queue.length - 1
          tmp[i] = queue[ (head + i) % queue.length ]
      head = 0
      queue = tmp
   }
   queue[size] = element
   queue.size = queue.size + 1
```

What happens when size == 0?

$$tail = (head + size - 1) \% length$$

```
Initial state ---- (0, 3, 0)
enqueue(a) a--- (0, 0, 1)
enqueue(b) ab-- (0, 1, 2)
dequeue() -b-- (1, 1, 1)
dequeue() (2, 1, 0)
tail head
```

ADT's, API's & Java

The following are related, but quite different:

- ADT (abstract data type)
- Java API (application program interface)
- Java keyword interface

To be discussed much more at end of the course.

ADT (abstract data type)

Defines a data type by the values and operations from the user's perspective only. It ignores the details of the implementation.

Examples:

- list
- stack
- queue
- •

Java API

API = application program *interface*

Gives class methods and some fields, and comments on what the methods do. e.g.

https://docs.oracle.com/javase/7/docs/api/java/util/LinkedList.html

Java interface

reserved word (nothing to do with "I" in "API")

 like a class, but only the method signatures are defined

Example: List interface

```
interface List<T> {
   void add(T)
   void add(int, T)
   T remove(int)
   boolean isEmpty()
   T get(int)
   int size()
   :
}
```

class ArrayList<T> implements List<T> {

```
void add(T) { .... }
void add(int, T) { .... }
T remove(int) { .... }
boolean isEmpty() { .... }
T get(int) { .... }
int size() { .... }
:
```

Each of the List methods are implemented. (In addition, other methods may be defined and implemented.)

class LinkedList<T> implements List<T> {

```
void add(T) { .... }
void add(int, T) { .... }
T remove(int) { .... }
boolean isEmpty() { .... }
T get(int) { .... }
int size() { .... }
:
}
```

Each of the List methods are implemented. (In addition, other methods may be defined and implemented.)

More examples

• interface List add(i,e), remove(i), get(i), set(i),

• class Stack push, pop(), ..

• interface Queue offer(e), poll(),