

COMP 250

Lecture 9

mathematical induction

Sept. 26, 2016

$$1 + 2 + 3 + \dots + (n - 1) + n = \frac{n(n + 1)}{2}$$

How to prove such a statement ?

(By “proof”, we mean a formal logical argument that convincingly shows the statement is true.)

$$1 + 2 + \dots + (n - 1) + n$$

Write sum backwards:

$$n + (n - 1) + \dots + 2 + 1$$

Adding up n pairs gives $n * (n + 1)$.

Dividing by 2 gives the result.

You should be 100% convinced by this proof.

Mathematical Induction

Consider statements of the form:

“For all $n \geq n_0$, $P(n)$ ” where $P(n)$ is either true or false for each n , and n_0 is a constant.

Mathematical induction is a *general* technique for proving such statements.

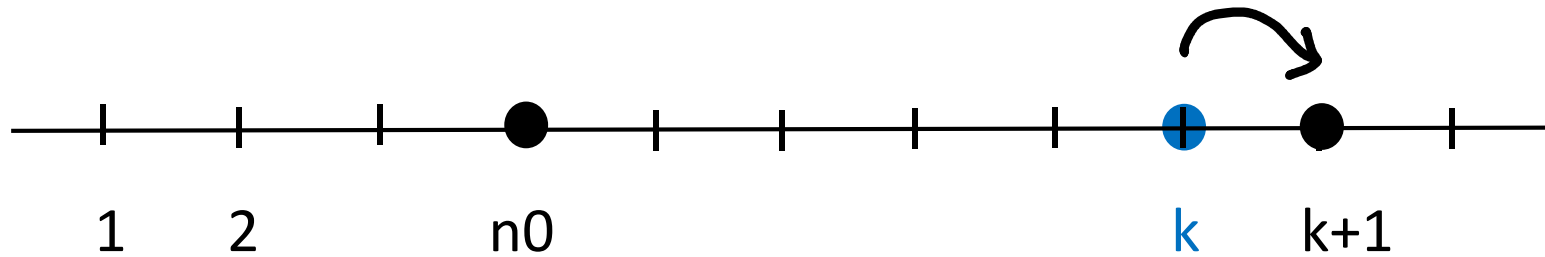
Mathematical induction requires proving two things:

Base case:

“ $P(n_0)$ is true.”

Induction step:

“For any $k \geq n_0$, if $P(k)$ is true, then $P(k+1)$ is also true.”



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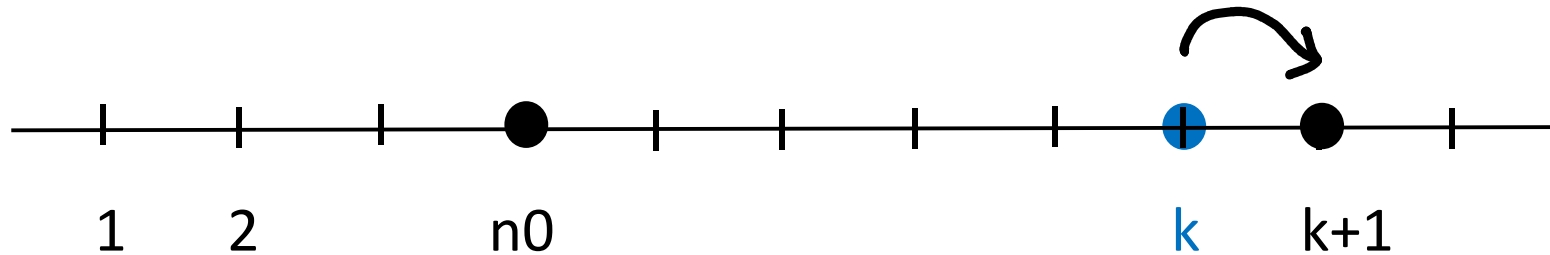
“ $P(k)$ is true” is called the “induction hypothesis”.

Base case:

$P(n_0)$ is true.

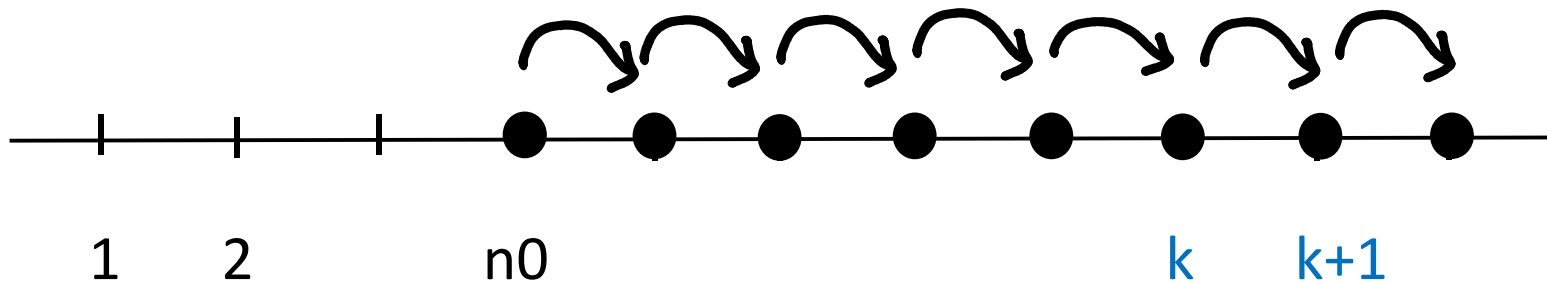
Induction step:

For any $k \geq n_0$, if $P(k)$ then $P(k+1)$.



Thus,

For any $n \geq n_0$, $P(n)$ is true.



Statement: For all $n \geq 1$,

$$1 + 2 + 3 + \dots + (n - 1) + n = \frac{n(n + 1)}{2}$$

Proof (base case, $n=1$):

$$1 = \frac{1(1+1)}{2} \quad (\text{true})$$

Proof of Induction Step:

$$(1 + 2 + 3 + \dots + k) + k + 1$$

$$= \frac{k(k+1)}{2} + k + 1$$

Why? Because of the induction hypothesis

$$(1 + 2 + 3 + \dots + k) = \frac{k(k+1)}{2}$$

Proof of Induction Step:

$$(1 + 2 + 3 + \dots + k) + k + 1$$

by induction hypothesis

$$= \frac{k(k+1)}{2} + k + 1$$

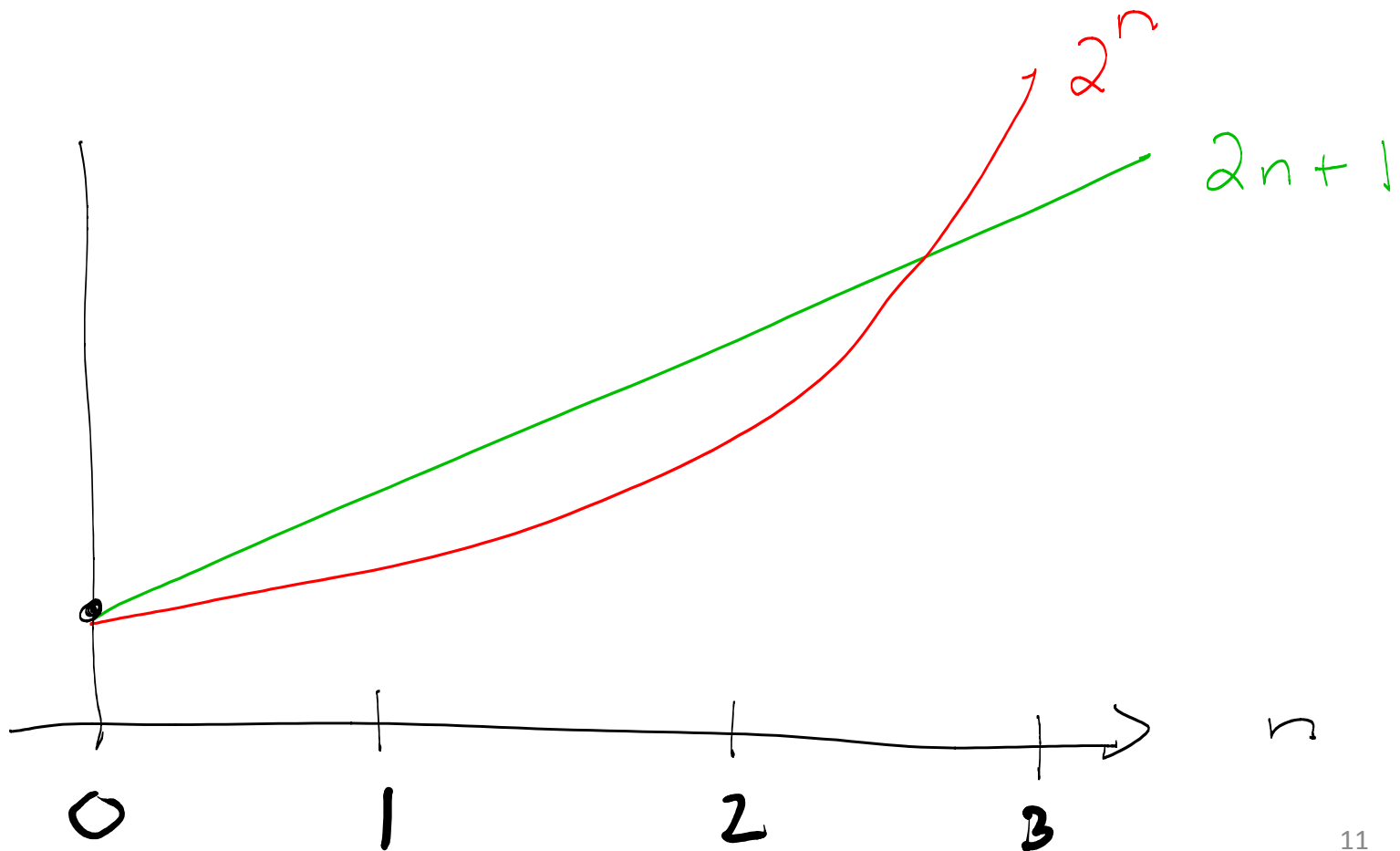
$$= \left(\frac{k}{2} + 1 \right) (k + 1)$$

$$= \frac{1}{2} (k + 2) (k + 1)$$

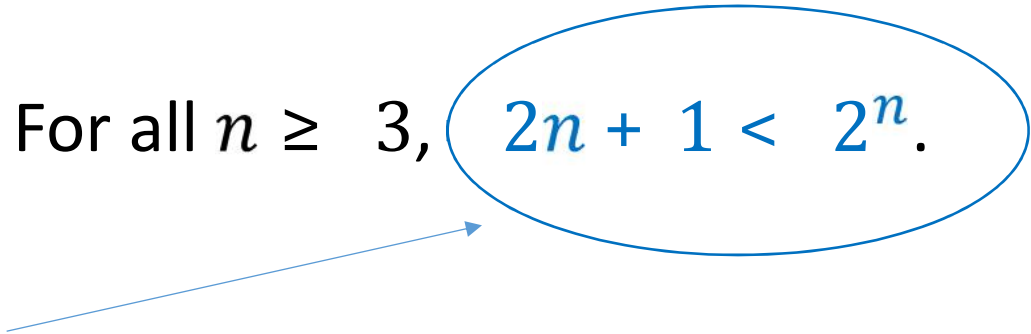
Thus, $P(k)$ is true implies $P(k + 1)$ is true.

Example 2

Statement: For all $n \geq 3$, $2n + 1 < 2^n$.



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Note: $P(n)$ is false for $n = 1, 2$.

Proof (base case, $n = 3$):

$$2*3 + 1 < 8 \quad (\text{true})$$

Proof of Induction Step:

We want to show that $P(k)$ implies $P(k+1)$.

$$2(k + 1) + 1 = 2k + 2 + 1$$

by induction hypothesis

$$< 2^k + 2$$

$$< 2^k + 2^k, \quad \text{for } k \geq 2$$

$$= 2^{k+1}$$



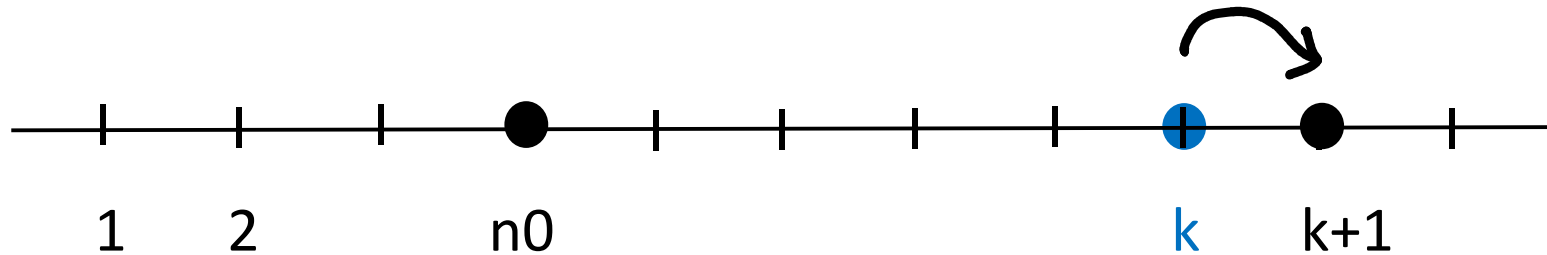
This condition is stronger than we need.

Base case:

$P(n_0)$ is true.

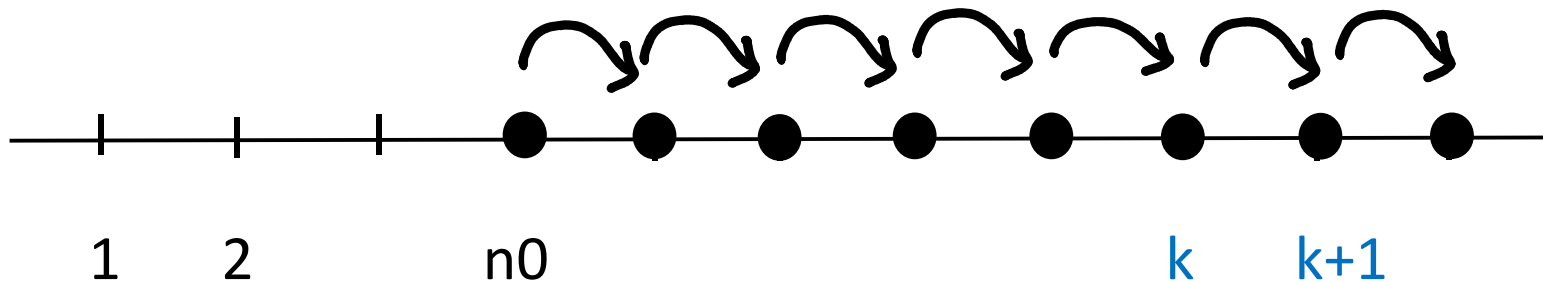
Induction step:

For any $k \geq n_0$, if $P(k)$ then $P(k+1)$.



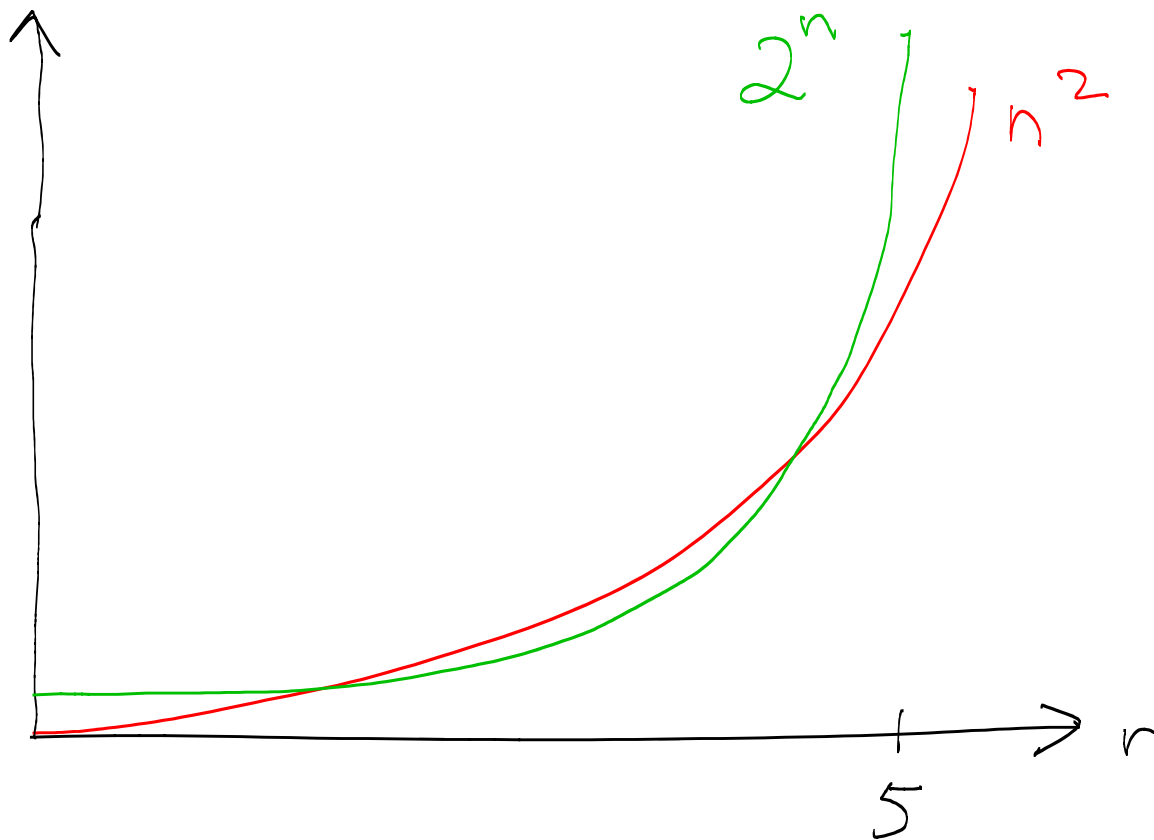
Thus,

For any $n \geq n_0$, $P(n)$ is true.



Example 3

Statement: For all $n \geq 5$, $n^2 < 2^n$.



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Base case ($n = 5$):

$$25 < 32$$

Induction step:

$$(k + 1)^2 = k^2 + 2k + 1$$

by induction hypothesis

$$< 2^k + 2k + 1$$

by Example 2

$$< 2^k + 2^k$$

$$= 2^{k+1}$$

Example 4 : Fibonacci Sequence

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55,

$$F(0) = 0$$

$$F(1) = 1$$

$$F(n + 2) = F(n + 1) + F(n) , \text{ for } n \geq 0.$$

Statement: For all $n \geq 0$, $F(n) < 2^n$

$$F(0) = 0$$

$$F(1) = 1$$

$$F(n + 2) = F(n + 1) + F(n) , \text{ for } n \geq 0.$$

Base case(s):

$$n = 0: \quad 0 < 2^0 \text{ is true.}$$

$$n = 1: \quad 1 < 2^1 \text{ is true.}$$

$$F(0) = 0$$

$$F(1) = 1$$

$$F(n + 2) = F(n + 1) + F(n) , \text{ for } n \geq 0.$$

Induction step:

$$F(k + 1) = F(k) + F(k - 1)$$

by induction hypothesis

$$< 2^k + 2^{k-1}$$

$$< 2^k + 2^k$$

$$= 2^{k+1}$$

How to use mathematical induction to prove an algorithm is correct ? e.g. insertion sort

```
insertion_sort( list ){  
  for k = 1 to n - 1 {  
    elementK = list[k]  
    i = k  
    while (i > 0) and (list[i - 1] > elementK ){  
      list[i] = list[i - 1]  
      i = i - 1  
    }  
    list[i] = elementK  
  }  
}
```

$P(k) =$

"At the start of pass k through the `for` loop, `list[0, ..., k-1]` contains the same elements as the original list, and these k elements are now sorted."

Q: What is the base case?

Q: What is the induction step?

Base case:

$P(1) =$

"At the start of pass 1 through the `for` loop, `list[0]` contains the same element as the original list, and these elements are now sorted."

Induction step

if

"At the start of pass k through the `for` loop,
`list[0, . . . k-1]` contains the same k elements as
the original list, and these elements are now sorted."

then

"At the start of pass $k+1$ through the `for` loop
(i.e. at the *end* of pass k through the `for` loop)
`list[0, . . . k]` contains the same elements as the
original list, and these elements are now sorted."

You may have already been convinced that this algorithm was correct.

I would argue, however, that your proof (or intuition) uses mathematical induction.

If not, then what *does* it use?