The Perception Algorithme An online algorithme for learning threshold functions. f: R" -> {+1,-1} IER" $f(\vec{z}) = \begin{cases} +1 & \text{if } \vec{\omega} \cdot \vec{z} \geq 6 \\ \vec{\omega}, \vec{b} \end{cases}$ l -1 if w. ≠ < b This defines a hyperplane in R" with was the normal to the hyperplane. The goal is to learn which hyperplane separates the positive labelled points from the negative told labelled points. Without loss of generality we can set b=0. We can redefine the problem in \mathbb{R}^{n+1} $\overline{f}: \mathbb{R}^{n+1} \longrightarrow \{+1,-1\}$ and change $\vec{z} \in \mathbb{R}^n$ to $\vec{z}' = (\vec{z}) \in \mathbb{R}^{nH}$ Now $\vec{\pi}$ · $\vec{\omega}$ ' $\vec{\phi}$ $\vec{\phi}$ We want to bearn a from examples. The examples are presented to us in online fashion. So we have to make a prediction and there the true label is revealed. We seek bounds as our mistakes. BASIC PERCEPTRON ALGORITHM: Loop Init: $\vec{\omega}_i = 0$ t = 1... T (Obtain $\vec{\chi}_t$ (PREDICT) { if w. \$\frac{1}{2} \go Hun \geq = +1 else \geq = -1 Obtain y6 (UPDATE) L if ye + ye then We+1 + WE+ ye xe also We+1 = WE When it makes a mistake on a positive execuple it shifts the weight vector towards that point if it makes a mistake on a negative example it shifts the weight vector away.

Crucial concept: How close is a point to the hyperplane Let f be a threshold function determined by a hyperplane passing through the origin with unit normal is.

Set The margin of f at a labelled point (\vec{x}, \vec{y}) is defined to be \(\vec{x} \cdot \vec{u} \).

Remark If the margin is positive them of classifies (x, y) correctly and if it is negative there of classifies the point incorrectly. The numerical value of the margin gives the distance to the decision boundary. The margin of a set of points is the minimum margin of any point in the set. The margin measures vobustness of the classification.

Thun Suppose (\$\overline{\infty}_1, y_1), (\$\overline{\infty}_2, y_2) \cdots (\$\overline{\infty}_7, y_7) are such that I Till & D for some D>0 and all test.... T3. Suppose I I a unit vector & 7>0 such that Vt ye (x. u) > r. Then the perception algorithm makes at most (D/r) mistakes. Proof Let me be the number of mistakes just before round t. Thus m, = 0. We proceed by breaking the proof into 2 lemmas. hemma 1 YtoEI,..., T} We. U > MEY Proof By induction ont.

Base case t=01, m, =0 so clearly to, ti=0=m, r Induction step: suppose there is a mistake in round t. Then we have

Wer, · U = (we + ye xe) · U = We U + ye xe · U $\geq m_{t} \gamma + \gamma = m_{t+1} \gamma$ Find Margin Hyp assemption

So lemma I shaws that as mis-bles are made the projection of the or ti gets longer.

lemma 2 | | \vec{w}_e|^2 \le m_t D^2

Proof Induction on t; base case its trivials as is

the case colour there is no mistake. Suppose there
is a mistake in round t. Then we have

||\vec{w}_{t+1}||^2 = ||\vec{w}_t + y_t \vec{x}_t||^2

= ||\vec{w}_t||^2 + ||\vec{x}_t||^2 + 2y_t (\vec{w}_t \cdot \vec{x}_t)

N'expetive

\le ||\vec{w}_t||^2 + ||\vec{x}_t||^2 + 2y_t (\vec{w}_t \cdot \vec{x}_t)

N'expetive

\le ||\vec{w}_t||^2 + ||\vec{x}_t||^2

\le m_t D^2 + D^2 = m_{t+1} D^2

Proof of Thue ||\vec{w}_t||^2 \le m_t D^2 (herrora 2)

or D√m_t ≥ ||\vec{w}_t|| ||\vec{u}_t||

\le ||\vec{w}_t \vec{u}_t| ||\vec{w}_t||

\le m_t \vec{v}_t \vec{u}_t| ||\vec{w}_t||

\le m_t \vec{v}_t \vec{u}_t| ||\vec{w}_t||

\le m_t \vec{v}_t \vec{u}_t| ||\vec{w}_t||

\le m_t \vec{v}_t| ||\vec{v}_t||

\le m_t \vec{v}_t||

\le m_t

APPLICATION Suppose there are n financial analysts who predict every day whether the market will go up or down. We represent each prediction as a vector in {+1,-13°. We would like to use perception to figure out whom to follow. Suppose there are k "experts" within the group s.t. a majority vote among these k always gives the right answer.

Define $\vec{u} = \vec{t_k} (0, 0, 1, 1, 0, ... 0, 1...)$ where the entry is 1 if it corresponds to one of the k experts.

Then $y_k(\vec{x_k}, \vec{t_k}) \ge 1/k$ since this subset is always right.

Thus $||x_k|| \le \sqrt{n} = D \cdot x = 1/k$. So the mistake bound is $n \cdot k$.