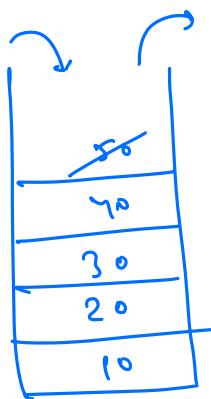


- Quick Revision
- Remove Equal Pair of consecutive Elements
- Largest Area in Histogram
- Sum of max-min for all subarrays.

## ① Implement stack using linkedlist



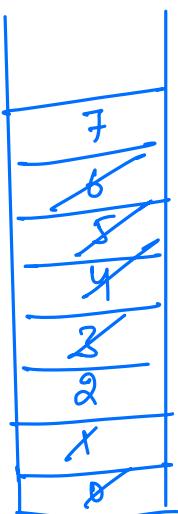
✗ Add at tail, remove from tail  $O(n)$

✓ Add at head, remove from head.

## ② nearest smaller on l.h.s

$$\text{arr}[9] = [5, 4, 2, 7, 6, 9, 15, 3]$$

$$\text{nsl}[9] = [-1, -1, -1, 2, 2, 4, 5, 2]$$



→ pop all greater elements.

→ update ans

→ push current index in stack.

Code →

```
Stack<int> st;  
nsd[N];  
  
for( i=0; i<N; i++) {  
    // pop all greater elements  
    while( st.size() > 0 && arr[st.peek()] ≥ arr[i] )  
        st.pop();  
  
    // update ans  
    if( st.size() == 0 ) { nsd[i] = -1 }  
    else { nsd[i] = st.peek() }  
  
    // push current index in stack  
    st.push(i);  
  
}  
return nsd[T];
```

[ T.C → O(N) ]  
[ S.C → O(N) ]

③ Nearest smaller on r.h.s

↓

iterate from r.h.s.

Q → Given a string, remove equal pair of consecutive elements till possible

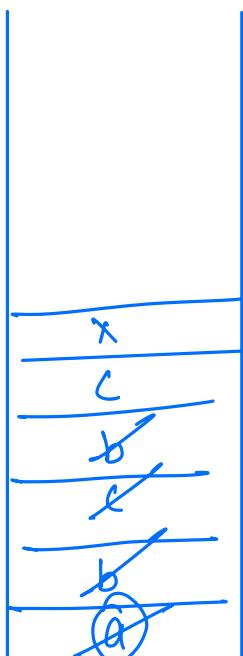
Str → abc dd      ans → abc

Str → ab c dd c → abf4      ans → ab

Str → abbbcbbcacx → dfdgdx → cx

Str - abffb → ab

↑↑↑↑↑  
abbcbbcacx



ans → cx

code ->

Stack < character > st;

```
for( i=0; i<N; i++) {  
    if( st.size() == 0 || st.peek() != str[i] ) {  
        st.push(str[i]);  
    }  
    else {  
        st.pop();  
    }  
}
```

StringBuilder sb;

```
while( st.size() != 0 ) {  
    sb.append( st.pop() );  
}
```

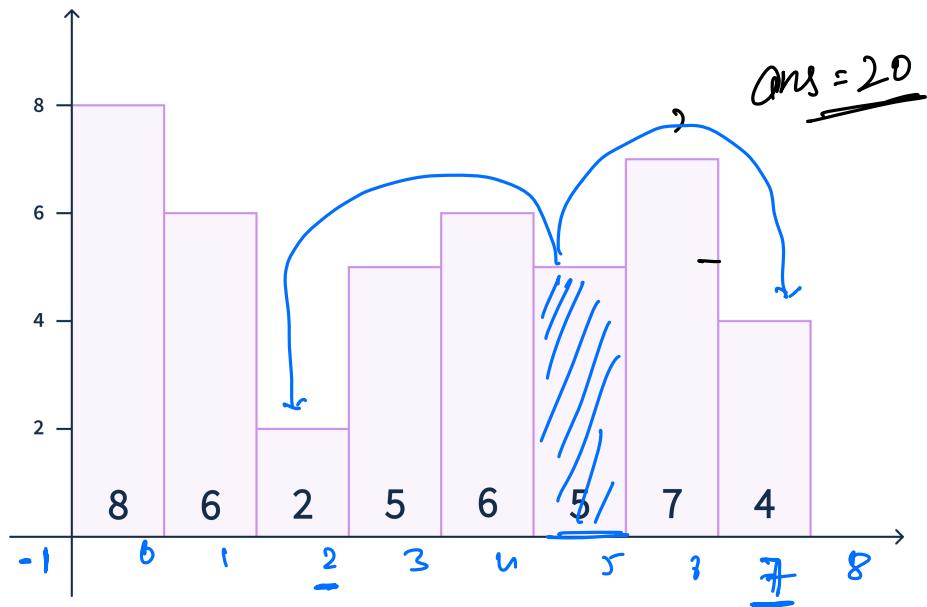
reverse (sb)

return sb.toString()

T.L  $\rightarrow \Theta(N)$   
S.C  $\rightarrow \Theta(N)$

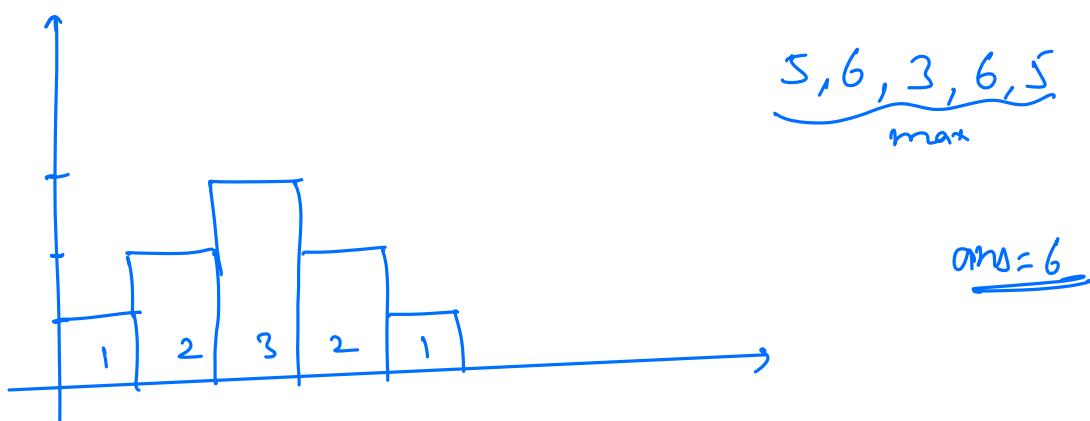


< Question > : Find the largest area of rectangle ( formed by continuous bars ) in histogram.



$$\text{arr}[i] \times (7-2-1)$$

$\text{arr}[7] \rightarrow [1, 2, 3, 2, 1]$





B.F - for every bar,

- find first bar with height less than or equal to current bar on l.h.s
- find first bar with height less than or equal to current bar on r.h.s.

$$\begin{aligned} T.C &\rightarrow O(N^2) \\ S.C &\rightarrow O(1) \end{aligned}$$

optimisation → find nearest smaller on l.h.s,  
nearest smaller on r.h.s.

$$\rightarrow \underbrace{\text{arr}[i] \times (\text{nsr}[i] + \text{nsl}[i] - 1)}_{\text{consider this factor for every bar.}}$$

```

public class Solution {
    public int largestRectangleArea(int[] arr) {
        int[] nsl = new int[arr.length];
        int[] nsr = new int[arr.length];

        Stack<Integer> st = new Stack<>();
        for(int i = 0; i < arr.length; i++){
            while(st.size() > 0 && arr[st.peek()] >= arr[i]){
                st.pop();
            }
            if(st.size() == 0){
                nsl[i] = -1;
            }else{
                nsl[i] = st.peek();
            }
            st.push(i);
        }

        st = new Stack<>();
        for(int i = arr.length - 1; i >= 0; i--){
            while(st.size() > 0 && arr[st.peek()] >= arr[i]){
                st.pop();
            }
            if(st.size() == 0){
                nsr[i] = arr.length;
            }else{
                nsr[i] = st.peek();
            }
            st.push(i);
        }

        int ans = 0;
        for(int i = 0; i < arr.length; i++){
            ans = Math.max(ans, arr[i] * (nsr[i] - nsl[i] - 1));
        }
        return ans;
    }
}

```

$$\begin{cases} T.C \rightarrow O(N) \\ S.L \rightarrow O(N) \end{cases}$$



< Question > : Find sum of ( max-min ) for all subarrays.

$$1 \leq N \leq 10^5$$

arr : [ 1 2 3 ]  
0 1 2

$$\begin{aligned} & \sum \text{max-min} \\ & \Rightarrow \sum \text{max} - \sum \text{min} \end{aligned}$$

	max	min	max-min
[ 1 ]	1	1	0
[ 1, 2 ]	2	1	1
[ 1, 2, 3 ]	3	1	2
[ 2 ]	2	2	0
[ 2, 3 ]	3	2	1
[ 3 ]	3	3	0
	<hr/> 14	<hr/> 10	<hr/> 4



## BF Idea

Consider all subarrays & for every subarray find min and max.

$$\left[ \begin{array}{l} T.C \rightarrow O(N^3) \\ S.C \rightarrow O(1) \end{array} \right]$$

↓ carry forward

$$T.C \rightarrow O(N^2)$$



```

ans = 0;

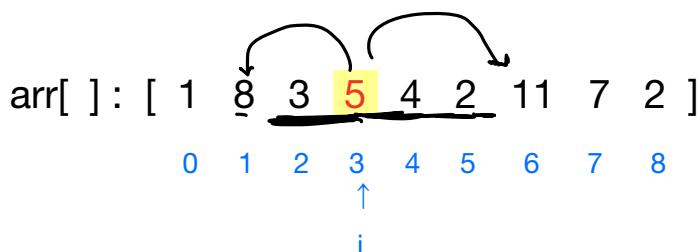
for( i = 0; i < N; i++) {
    min = arr[i], max = arr[i]

    for( j = i; j < n; j++) {
        min = Math.min( min, arr[j])
        max = Math.max( max, arr[j])
        ans += max - min;
    }
}
return ans;

```

$$\begin{cases} T.C \rightarrow O(N^2) \\ S.C \rightarrow O(1) \end{cases}$$

**< Question > :** In how many subarrays, ith element is the maximum element?



options for si

2
2

options for ei

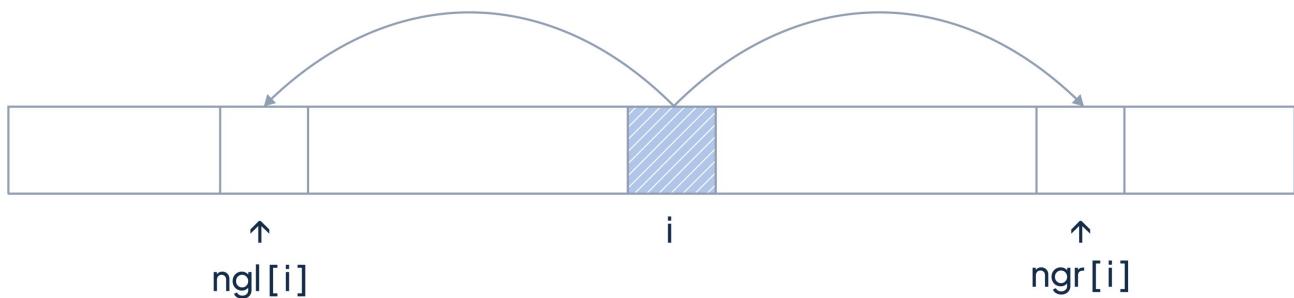
3
4
5

$$(i - \text{ngl}(i)) \times (\text{ngr}(i) - i)$$

- (3, 5)
- (3, 5, 4)
- (3, 5, 4, 2)
- (5)
- (5, 4)
- (5, 4, 2)

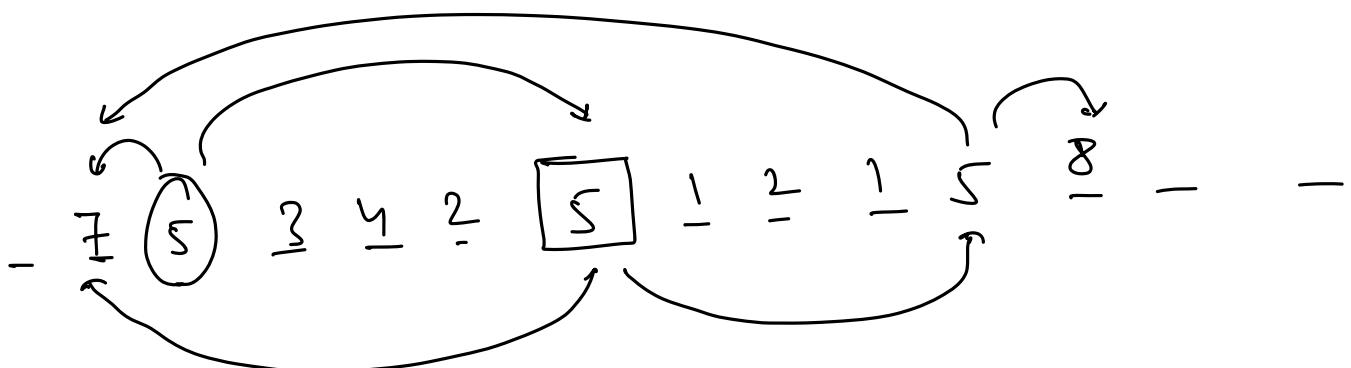


- **Generalisation**



total subarrays where  $\text{arr}(i)$  is max  $\rightarrow (i - ngl(i)) \times (ngr(i) - i)$

total subarrays where  $\text{arr}(i)$  is min  $\rightarrow (i - nsr(i)) \times (nsr(i) - i)$



nsr

nsr

ngl

ngl

**< / > Code**

```
st = new Stack<>();
for(int i = arr.length - 1; i >= 0; i--){
    while(st.size() > 0 && arr[st.peek()] <= arr[i]){
        st.pop();
    }
    if(st.size() == 0){
        ngr[i] = arr.length;
    }else{
        ngr[i] = st.peek();
    }
    st.push(i);
}

long ans = 0;
int mod = (int)1e9 + 7;
for(int i = 0; i < arr.length; i++){
    long maxContribution = arr[i] * (i - ngl[i]) * (ngr[i] - i);
    long minContribution = arr[i] * (i - nsr[i]) * (nsr[i] - i);

    ans = (ans + (maxContribution - minContribution) + mod) % mod; // check
}
return (int)ans;
}
```

↑

Chuck it out

T.C  $\rightarrow O(N)$   
S.C  $\rightarrow O(N)$

Contest 2: Searching, Linked List, Stacks,  
Queues & Trees,