

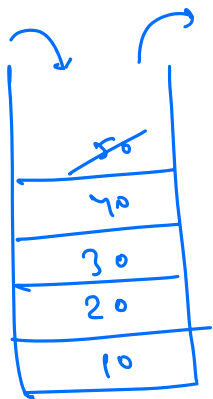
→ Quick Revision

→ Remove Equal Pair of consecutive elements

→ Largest Area in Histogram

→ Sum of max-min for all subarrays.

① Implement stack using linkedlist

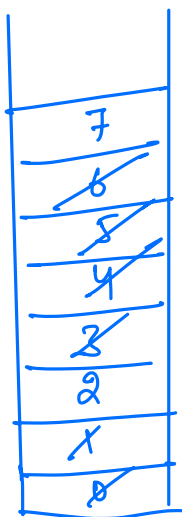


- ~~① Add at tail, remove from tail $O(n)$~~
- ② Add at head, remove from head.

② Nearest smaller on l.h.s

arr[] → [5 4 2 7 6 9 15 3]
 [0 1 2 3 4 5 6 7]

ans[] → [-1 -1 -1 2 2 4 5 2]



- pop all greater elements.
- update ans
- push current index in stack.

Code -

```
Stack<int> st;
```

```
ans[N];
```

```
for (i = 0; i < N; i++) {
```

```
    // pop all greater elements
```

```
    while (st.size() > 0 && arr[st.peek()] >= arr[i]) {
```

```
        st.pop();
```

```
    // update ans
```

```
    if (st.size() == 0) { ans[i] = -1;
```

```
    else { ans[i] = st.peek(); }
```

```
    // push current index in stack
```

```
    st.push(i);
```

```
}
```

```
return ans;
```

```
[ T.C  $\rightarrow O(N)$   
  S.C  $\rightarrow O(N)$  ]
```

③ Nearest smaller on r.h.s

⇓

iterate from r.h.s.

Q → Given a string, remove equal pair of consecutive elements till possible

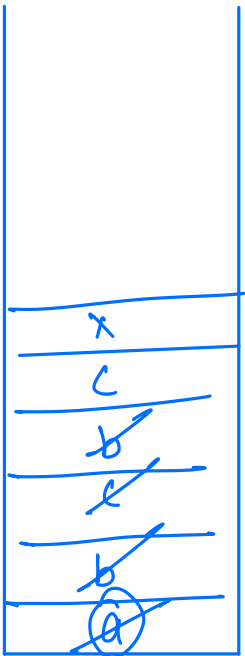
str → abcd d ans → abc

str → abcd dc → abff ans → ab

str → abbcbbaacx → ddcx → cx

str → abbb → ab

✓✓✓✓✓
abbcbbaacx



ans → cx

code →

```
Stack < character > st;
```

```
for ( i = 0; i < N; i++) {
```

```
    if ( st.size() == 0 || st.peek() != str[i] ) {
        |
        |         st.push( str[i] );
        |
        |
    }
    else {
        |         st.pop();
        |
        |
    }
}
```

```
StringBuilder sb;
```

```
while ( st.size() != 0 ) {
    |
    |         sb.append( st.pop() );
    |
    |
}
```

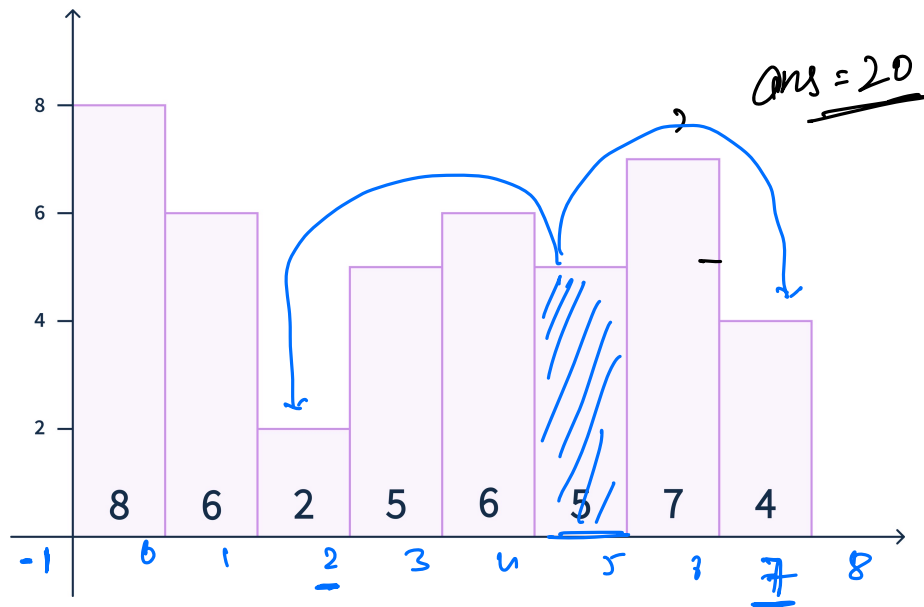
```
reverse ( sb )
```

```
return sb.toString()
```

$\left[\begin{array}{l} T.C \rightarrow O(N) \\ S.C \rightarrow O(N) \end{array} \right]$

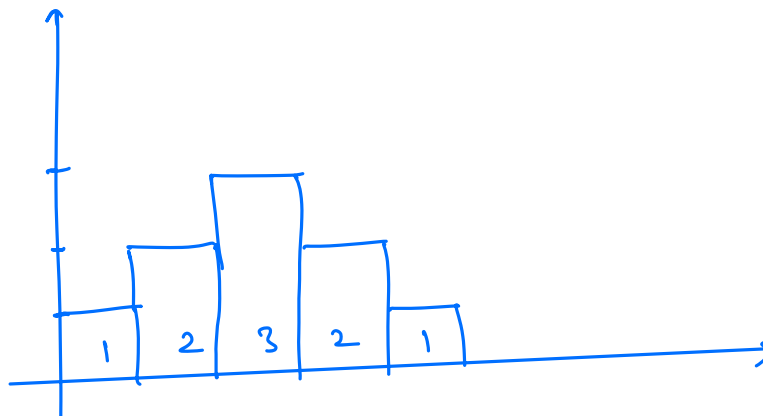


< **Question** > : Find the largest area of rectangle (formed by continuous bars) in histogram.



$$\text{arr}[i] \times (7 - 2 - 1)$$

$$\text{arr}[7] \rightarrow [1, 2, 3, 2, 1]$$



$$\underline{5, 6, 3, 6, 5}$$

max

$$\underline{\underline{\text{ans} = 6}}$$



B.f → for every bar,

→ find first bar with height less than or equal to current bar on l.h.s

→ find first bar with height less than or equal to current bar on r.h.s.

T.C → $O(N^2)$

S.C → $O(1)$

optimisation → find nearest smaller on l.h.s,
nearest smaller on r.h.s.

→
$$\frac{arr[i] \times (nsl[i] + nsr[i] - 1)}{\text{consider this factor for every bar.}}$$

```

public class Solution {
    public int largestRectangleArea(int[] arr) {
        int[] nsl = new int[arr.length];
        int[] nsr = new int[arr.length];

        Stack<Integer> st = new Stack<>();
        for(int i = 0; i < arr.length; i++){
            while(st.size() > 0 && arr[st.peek()] >= arr[i]){
                st.pop();
            }
            if(st.size() == 0){
                nsl[i] = -1;
            }else{
                nsl[i] = st.peek();
            }
            st.push(i);
        }

        st = new Stack<>();
        for(int i = arr.length - 1; i >= 0; i--){
            while(st.size() > 0 && arr[st.peek()] >= arr[i]){
                st.pop();
            }
            if(st.size() == 0){
                nsr[i] = arr.length;
            }else{
                nsr[i] = st.peek();
            }
            st.push(i);
        }

        int ans = 0;
        for(int i = 0; i < arr.length; i++){
            ans = Math.max(ans, arr[i] * (nsr[i] - nsl[i] - 1));
        }
        return ans;
    }
}

```

$$\left[\begin{array}{l} T.C \rightarrow O(N) \\ S.C \rightarrow O(N) \end{array} \right]$$



< **Question** > : Find sum of (max-min) for all subarrays.

$1 \leq N \leq 10^5$

arr : [1 2 3]
0 1 2

$$\sum \text{max-min} \\ \Rightarrow \sum \text{max} - \sum \text{min}$$

	max	min	max-min
[1]	1	1	0
[1,2]	2	1	1
[1,2,3]	3	1	2
[2]	2	2	0
[2,3]	3	2	1
[3]	3	3	0
	<hr/> 14	<hr/> 10	<hr/> 4



BF Idea

Consider all subarrays & for every subarray find min and max.

$$\left[\begin{array}{l} \text{T.C} \rightarrow O(N^3) \\ \text{S.C} \rightarrow O(1) \end{array} \right]$$

↓ (carry forward)
T.C $\rightarrow O(N^2)$



```

ans = 0;
for( i = 0 ; i < N ; i++ ) {
    min = arr[i], max = arr[i]
    for( j = i ; j < N ; j++ ) {
        min = Math.min( min, arr[j] )
        max = Math.max( max, arr[j] )
        ans += max - min;
    }
}
return ans;

```

$T.C \rightarrow O(N^2)$
 $S.C \rightarrow O(1)$

< Question > : In how many subarrays, ith element is the maximum element?

arr[] : [1 8 3 5 4 2 11 7 2]
0 1 2 3 4 5 6 7 8
↑
i

options for si

2

2

options for ei

3

4

5

[3, 5]

[3, 5, 4]

[3, 5, 4, 2]

[5]

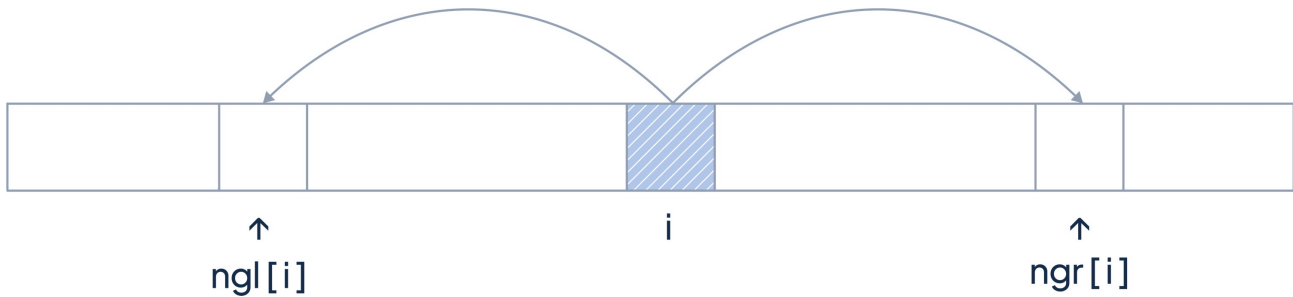
[5, 4]

[5, 4, 2]

$(i - \text{ngl}(i)) \times (\text{ngr}(i) - i)$

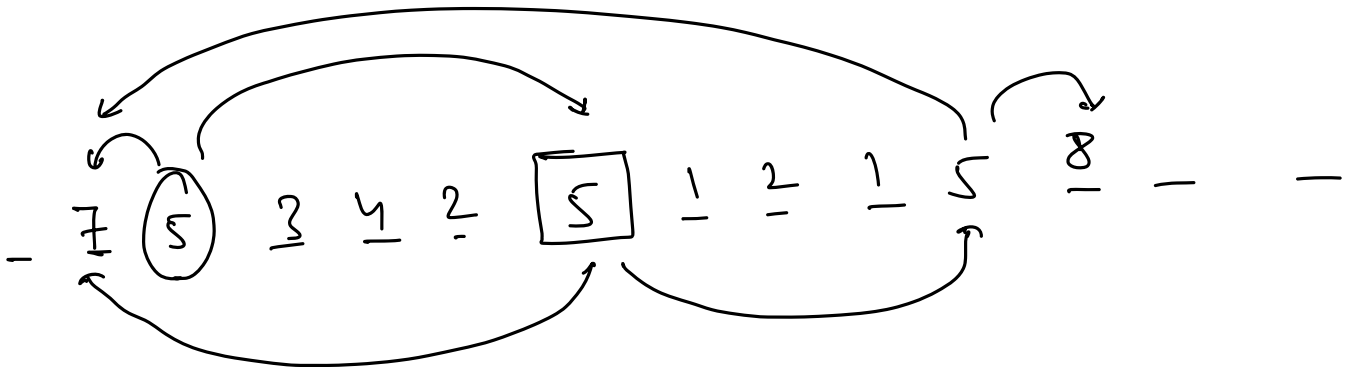


• Generalisation



total subarrays where $arr(i)$ is max $\rightarrow (i - ngl[i]) \times (ngr[i] - i)$

total subarrays where $arr(i)$ is min $\rightarrow (i - nsl[i]) \times (nsr[i] - i)$



nsl
 nsr

ngl
 ngr



< / > Code

```

st = new Stack<>();
for(int i = arr.length - 1; i >= 0; i--){
    while(st.size() > 0 && arr[st.peek()] <= arr[i]){
        st.pop();
    }
    if(st.size() == 0){
        ngr[i] = arr.length;
    }else{
        ngr[i] = st.peek();
    }
    st.push(i);
}

long ans = 0;
int mod = (int)1e9 + 7;
for(int i = 0; i < arr.length; i++){
    long maxContribution = arr[i] * (i - ngl[i]) * (ngr[i] - i);
    long minContribution = arr[i] * (i - nsl[i]) * (nsr[i] - i);
    // ans = (ans + (maxContribution - minContribution) + mod) % mod; // check
}
return (int)ans;
}

```

↑
chuck it out

$\left[\begin{array}{l} T.C \rightarrow O(N) \\ S.C \rightarrow O(N) \end{array} \right]$

Contest 2: Searching, Linked List, Stacks,
Queues & Trees,