

Q3:

.globl map

.text

main:

jal ra, create_default_list

add s0, a0, x0 # a0 = s0 is head of node list

#print the list

add a0, s0, x0

jal ra, print_list

print a newline

jal ra, print_newline

load your args

add a0, s0, x0 # load the address of the first node into a0

load the address of the function in question into a1 (check out la on the green sheet)

YOUR CODE HERE

la a1, square

issue the call to map

jal ra, map

print the list

add a0, s0, x0

jal ra, print_list

print another newline

jal ra, print_newline

addi a0, x0, 10

ecall #Terminate the program

map:

Prologue: Make space on the stack and back-up registers

YOUR CODE HERE

addi sp, sp, -12

sw s0, 8(sp)

sw s1, 4(sp)

sw ra, 0(sp)

beq a0, x0, done # If we were given a null pointer (address 0), we're done.

```
mv s0, a0 # Save address of this node in s0
mv s1, a1  # Save address of function in s1
```

Remember that each node is 8 bytes long: 4 for the value followed by 4 for the pointer to next.

What does this tell you about how you access the value and how you access the pointer to next?

```
# load the value of the current node into a0
# THINK: why a0?
```

```
### YOUR CODE HERE ###
```

```
lw a0, 0(s0)
```

Call the function in question on that value. DO NOT use a label (be prepared to answer why).

```
# What function? Recall the parameters of "map"
```

```
### YOUR CODE HERE ###
```

```
jalr s1
```

```
# store the returned value back into the node
```

```
# Where can you assume the returned value is?
```

```
### YOUR CODE HERE ###
```

```
sw a0, 0(s0)
```

```
# Load the address of the next node into a0
```

```
# The Address of the next node is an attribute of the current node.
```

```
# Think about how structs are organized in memory.
```

```
### YOUR CODE HERE ###
```

```
lw a0, 4(s0)
```

```
# Put the address of the function back into a1 to prepare for the recursion
```

```
# THINK: why a1? What about a0?
```

```
### YOUR CODE HERE ###
```

```
mv a1, s1
```

```
# recurse
```

```
### YOUR CODE HERE ###
```

```
jal ra, map
```

done:

```
# Epilogue: Restore register values and free space from the stack
```

```
### YOUR CODE HERE ###
```

```
lw s0, 8(sp)
```

```
lw s1, 4(sp)
```

```
lw ra, 0(sp)
```

```
addi sp sp 12
```

```
jr ra # Return to caller
```

square:

```

mul a0 ,a0, a0
jr ra

```

create_default_list:

```

addi sp, sp, -12
sw  ra, 0(sp)
sw  s0, 4(sp)
sw  s1, 8(sp)
li  s0, 0      # pointer to the last node we handled
li  s1, 0      # number of nodes handled

```

loop: #do...

```

li  a0, 8
jal ra, malloc      # get memory for the next node
sw  s1, 0(a0)      # node->value = i
sw  s0, 4(a0)      # node->next = last
add s0, a0, x0     # last = node
addi s1, s1, 1     # i++
addi t0, x0, 10
bne s1, t0, loop    # ... while i!= 10
lw  ra, 0(sp)
lw  s0, 4(sp)
lw  s1, 8(sp)
addi sp, sp, 12
jr ra

```

print_list:

```

bne a0, x0, printMeAndRecurse
jr ra      # nothing to print

```

printMeAndRecurse:

```

add t0, a0, x0  # t0 gets current node address
lw  a1, 0(t0)   # a1 gets value in current node
addi a0, x0, 1   # prepare for print integer ecall
ecall
addi a1, x0, ' '  # a0 gets address of string containing space
addi a0, x0, 11   # prepare for print string syscall
ecall
lw  a0, 4(t0)    # a0 gets address of next node
jal x0, print_list # recurse. We don't have to use jal because we already have where we
want to return to in ra

```

print_newline:

```

addi a1, x0, '\n' # Load in ascii code for newline
addi a0, x0, 11
ecall

```

jr ra

malloc:

addi a1, a0, 0

addi a0, x0 9

ecall

jr ra

The screenshot shows the Venus RISC-V simulator interface. At the top, there are tabs for 'Venus', 'Editor', 'Simulator', and 'Chocopy'. Below these are buttons for 'Run', 'Step', 'Prev', 'Reset', 'Dump', and 'Trace'. The main area displays assembly code with columns for PC, Machine Code, Basic Code, and Original Code. The registers panel on the right shows values for zero, ra (x1), sp (x2), gp (x3), tp (x4), t0 (x5), and t1 (x6). The memory panel at the bottom shows a hex dump of memory.

PC	Machine Code	Basic Code	Original Code
0x0	0x088000EF	jal x1 136	jal ra, create_default_list
0x4	0x00050433	add x8 x10 x0	add s0, a0, x0 # a0 = s0 is head of node list
0x8	0x00040533	add x10 x8 x0	add a0, s0, x0
0xc	0x0C8000EF	jal x1 200	jal ra, print_list
0x10	0x0F0000EF	jal x1 240	jal ra, print_newline
0x14	0x00040533	add x10 x8 x0	add a0, s0, x0 # load the address of the first node into a0
0x18	0x00000597	auipc x11 0	la a1, square
0x1c	0x06858593	addi x11 x11 104	la a1, square

Registers (x1-x6):

- zero: 0x00000000
- ra (x1): 0x00000030
- sp (x2): 0x7FFFFFF0
- gp (x3): 0x10000000
- tp (x4): 0x00000000
- t0 (x5): 0x10008000
- t1 (x6): 0x00000000

Memory (Hex):

9 8 7 6 5 4 3 2 1 0
81 64 49 36 25 16 9 4 1 0

Q4:

.globl map

.data

arrays: .word 5, 6, 7, 8, 9

.word 1, 2, 3, 4, 7

.word 5, 2, 7, 4, 3

.word 1, 6, 3, 8, 4

.word 5, 2, 7, 8, 1

start_msg: .asciiz "Lists before: \n"

end_msg: .asciiz "Lists after: \n"

.text

main:

jal create_default_list

mv s0, a0 # v0 = s0 is head of node list

#print "lists before: "

```
la a1, start_msg
li a0, 4
ecall
```

```
#print the list
add a0, s0, x0
jal print_list
```

```
# print a newline
jal print_newline
```

```
# issue the map call
add a0, s0, x0      # load the address of the first node into a0
la  a1, mystery     # load the address of the function into a1
```

```
jal ra, map
```

```
# print "lists after: "
la a1, end_msg
li a0, 4
ecall
```

```
# print the list
add a0, s0, x0
jal print_list
```

```
li a0, 10
ecall
```

map:

```
addi sp, sp, -12
sw ra, 0(sp)
sw s1, 4(sp)
sw s0, 8(sp)
```

```
beq a0, x0, done      # if we were given a null pointer, we're done.
```

```
add s0, a0, x0        # save address of this node in s0
add s1, a1, x0        # save address of function in s1
add t0, x0, x0        # t0 is a counter
```

```
# remember that each node is 12 bytes long:
# - 4 for the array pointer
# - 4 for the size of the array
```

- 4 more for the pointer to the next node

also keep in mind that we should not make ANY assumption on which registers
are modified by the callees, even when we know the content inside the functions
we call. this is to enforce the abstraction barrier of calling convention.

mapLoop:

lw t1, 0(s0) # load the address of the array of current node into t1
lw t2, 4(s0) # load the size of the node's array into t2

slli t3, t0, 2
add t1, t1, t3 # offset the array address by the count
lw a0, 0(t1) # load the value at that address into a0

mv t4, t1
jalr s1 # call the function on that value.
mv t1, t4

sw a0, 0(t1) # store the returned value back into the array
addi t0, t0, 1 # increment the count
bne t0, t2, mapLoop # repeat if we haven't reached the array size yet

lw a0, 8(s0) # load the address of the next node into a0
mv a1, s1 # put the address of the function back into a1 to prepare for the

recursion

jal ra, map # recurse

done:

lw s0, 8(sp)
lw s1, 4(sp)
lw ra, 0(sp)
addi sp, sp, 12
jr ra

mystery:

mul t1, a0, a0
add a0, t1, a0
jr ra

create_default_list:

addi sp, sp, -4
sw ra, 0(sp)
li s0, 0 # pointer to the last node we handled
li s1, 0 # number of nodes handled
li s2, 5 # size

```

    la s3, arrays
loop: #do...
    li a0, 12
    jal malloc      # get memory for the next node
    mv s4, a0
    li a0, 20
    jal  malloc      # get memory for this array

    sw a0, 0(s4)     # node->arr = malloc
    lw a0, 0(s4)
    mv a1, s3
    jal fillArray    # copy ints over to node->arr

    sw s2, 4(s4)     # node->size = size (4)
    sw s0, 8(s4)     # node-> next = previously created node

    add s0, x0, s4   # last = node
    addi s1, s1, 1   # i++
    addi s3, s3, 20  # s3 points at next set of ints
    li t6 5
    bne s1, t6, loop # ... while i!= 5
    mv a0, s4
    lw ra, 0(sp)
    addi sp, sp, 4
    jr ra

```

```

fillArray: lw t0, 0(a1) #t0 gets array element
           sw t0, 0(a0) #node->arr gets array element
           lw t0, 4(a1)
           sw t0, 4(a0)
           lw t0, 8(a1)
           sw t0, 8(a0)
           lw t0, 12(a1)
           sw t0, 12(a0)
           lw t0, 16(a1)
           sw t0, 16(a0)
           jr ra

```

```

print_list:
    bne a0, x0, printMeAndRecurse
    jr ra    # nothing to print
printMeAndRecurse:
    mv t0, a0 # t0 gets address of current node
    lw t3, 0(a0) # t3 gets array of current node

```


Q5:

```
.globl f
```

```
.data
```

```
neg3: .ascii "f(-3) should be 6, and it is: "
```

```
neg2: .ascii "f(-2) should be 61, and it is: "
```

```
neg1: .ascii "f(-1) should be 17, and it is: "
```

```
zero: .ascii "f(0) should be -38, and it is: "
```

```
pos1: .ascii "f(1) should be 19, and it is: "
```

```
pos2: .ascii "f(2) should be 42, and it is: "
```

```
pos3: .ascii "f(3) should be 5, and it is: "
```

```
output: .word 6, 61, 17, -38, 19, 42, 5
```

```
.text
```

```
main:
```

```
    la a0, neg3
```

```
    jal print_str
```

```
    li a0, -3
```

```
    la a1, output
```

```
    jal f                # evaluate f(-3); should be 6
```

```
    jal print_int
```

```
    jal print_newline
```

```
    la a0, neg2
```

```
    jal print_str
```

```
    li a0, -2
```

```
    la a1, output
```

```
    jal f                # evaluate f(-2); should be 61
```

```
    jal print_int
```

```
    jal print_newline
```

```
    la a0, neg1
```

```
    jal print_str
```

```
    li a0, -1
```

```
    la a1, output
```

```
    jal f                # evaluate f(-1); should be 17
```

```
    jal print_int
```

```
    jal print_newline
```

```
    la a0, zero
```

```
    jal print_str
```

```
    li a0, 0
```

```

la a1, output
jal f          # evaluate f(0); should be -38
jal print_int
jal print_newline

```

```

la a0, pos1
jal print_str
li a0, 1
la a1, output
jal f          # evaluate f(1); should be 19
jal print_int
jal print_newline

```

```

la a0, pos2
jal print_str
li a0, 2
la a1, output
jal f          # evaluate f(2); should be 42
jal print_int
jal print_newline

```

```

la a0, pos3
jal print_str
li a0, 3
la a1, output
jal f          # evaluate f(3); should be 5
jal print_int
jal print_newline

```

```

li a0, 10
ecall

```

f takes in two arguments:
 # a0 is the value we want to evaluate f at
 # a1 is the address of the "output" array (defined above).
 # Think: why might having a1 be useful?
 f:

```

# YOUR CODE GOES HERE!
addi t0, a0, 3      # index
slli t0, t0, 2
add t0 a1, t0
lw a0, 0(t0)
jr ra               # Always remember to jr ra after your function!

```

```

print_int:
    mv a1, a0
    li a0, 1
    ecall
    jr    ra

```

```

print_str:
    mv a1, a0
    li a0, 4
    ecall
    jr    ra

```

```

print_newline:
    li a1, '\n'
    li a0, 11
    ecall
    jr    ra

```

PC	Machine Code	Basic Code	Original Code
0x0	0x10000517	auipc x10 65536	la a0, neg3
0x4	0x00050513	addi x10 x10 0	la a0, neg3
0x8	0x120000EF	jal x1 288	jal print_str
0xc	0xFFD00513	addi x10 x0 -3	li a0, -3
0x10	0x10000597	auipc x11 65536	la a1, output
0x14	0x0CB58593	addi x11 x11 203	la a1, output
0x18	0x0EC000EF	jal x1 236	jal f # evaluate f(-3); should be 6
0x1c	0x0Fc000EF	jal x1 252	jal print_int
0x20	0x118000EF	jal x1 280	jal print_newline

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```

f(-3) should be 6, and it is: 6
f(-2) should be 61, and it is: 61
f(-1) should be 17, and it is: 17
f(0) should be -38, and it is: -38
f(1) should be 19, and it is: 19
f(2) should be 42, and it is: 42
f(3) should be 5, and it is: 5

```

Integer (R)
Floating (F)

zero

0x00000000

ra

(x1)

0x000000FC

sp

(x2)

0x7FFFFFF0

gp

(x3)

0x10000000

tp

(x4)

0x00000000

t0

(x5)

0x100000F3

t1

...

0x00000000

Display Settings
Hex