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ABSTRACT

This paper gives general algorithm to extract n^{th} root of any real number. Algorithm uses multinomial expansion and intelligent guessing to compute digits of root. Paper also discusses various novel approaches to expand multinomial.

KEYWORD

 $\label{eq:local_problem} Algorithm, \ N^{th} \ Root, Multinomial \ Expansion \ , Intelligent \ Guessing \ , Integer \ Partition \\ \hline \textbf{PREREOUISITE}$

1) Binomial and multinomial theorem

GENERAL METHOD

Different method or approaches, based on multinomial expansion, can be used to extract n^{th} root of any real number. But first we restrict ourselves only with computing n^{th} root of perfect n^{th} power. Afterword general algorithm is given to extract n^{th} root of any number.

1) BACKWARD APPROACH

Example
$$(149721291)^{\frac{1}{3}} = ?$$
 (Given 149721291 perfect cube)

To calculate cube root, first divide given cube into '3 digit group' from right to left.

$$\therefore$$
 No. of digit = No. of groups = 3

Let integer a,b,c be unit, tens and hundred place of 3 digit cube root respectively.

1

$$\implies 0 \le a, b \le 9, 1 \le c \le 9$$

Now expanding cube of cba = a + 10b + 100c in power of 10's, we get

i.e.
$$(a+10b+100c)^3 = (a^3+3a^2b*10+(3ab^2+3a^2c)10^2+...$$

Let higher terms in RHS that are multiple of 10^3 be 1000Y, then

$$(a+10b+100c)^3 = a^3 + 3a^2b*10 + (3ab^2 + 3a^2c)10^2 + 1000Y$$

Now given cube = $150568768 \Rightarrow a^3 + 3a^2b * 10 + (3ab^2 + 3a^2c)10^2 + 1000Y = 150568768$

$$a^3 + 3a^2b * 10 + (3ab^2 + 3a^2c)10^2 + 1000Y = 149721291$$

$$\Rightarrow 3a^2b*10 + (3ab^2 + 3a^2c)10^2 + 1000Y = 149721291 - a^3 \qquad -----(1)$$

As LHS is multiple of 10, So RHS $149721291-a^3$ must be divisible by 10

 \Rightarrow 149721291 – a^3 must ends with 0 i.e. a^3 must ends with 1.

Now there is only one integer i.e. 1 between $\,0$ to $\,9\,$ whose cube end with $\,1$, $\,\dot{}$ a=1

Substituting a=1 in above equation (1)

$$3a^2b*10+(3ab^2+3a^2c)10^2+1000Y=149721291-1=149721290$$

$$3a^2b + (3ab^2 + 3a^2c)10 + 100Y = 14972129$$
 (Dividing both side by 10)

Similarly,

$$(3ab^2 + 3a^2c)10 + 100Y = 14972129 - 3a^2b$$

LHS is multiple of $10 \Rightarrow 3a^2b = 3b$ must ends with $9 \Rightarrow b = 3$

$$(3ab^2 + 3a^2c)10 + 100Y = 14972129 - 3*3 = 14972120$$

$$\therefore$$
 $(3ab^2 + 3a^2c) + 10Y = 1497212$

$$\Rightarrow 10Y = 1497212 - (3ab^2 + 3a^2c) = 1497212 - (3*1*3^2 + 3*1^2 *c)$$
$$10Y = 1497185 - 3c$$

LHS divisible by $10 \Rightarrow 3c$ must ends with 5 and $1 \le c \le 9 \Rightarrow c = 5$

∴ Cube root= cba= 531

In short,

1) a^3 ends with 1 (End digit of 149721291) \Rightarrow a = 1 $\therefore a^3 = 1$, Subtracting it from 149721291 and eliminating last digit i.e. 14972129 $\emptyset = 14972129$ —1

149721290

2) $3a^2b = 3b$ ends with 9 (End digit of 14972129) \Rightarrow b = 3 \therefore $3a^2b = 9$, Subtracting it from 14972129 and eliminating last digit i.e. 1497212 $\emptyset = 1497212$ —9

1497212

3) $3ab^2 + 3a^2c = 27 + 3c$ ends with 2 (End digit of 1497212) —42 \Rightarrow 3c ends with 5 \Rightarrow c = 5 \therefore $3ab^2 + 3a^2c = 42$.

:. Cube Root(150568768) = cba = 531

Applicability

Above method is convenient to extract nth root of perfect nth power that satisfies

$$GCD(10, n*U^{n-1}) = 1$$

 \Rightarrow GCD(10, n) = 1 and GCD(10, U) = 1 Where $U = Unit digit of n^{th} root$.

i.e. Unit digits of " n and n^{th} root " must be relative prime to 10 (base). (Why? Think).

2) FORWARD APPROACH

Example 1) $(20352513413376)^{\frac{1}{4}} = ?$ (Given number is perfect 4th power)

1) Arrange given perfect n^{th} power into *n digit group* from right to left. Here n = 4.

∴ No. of groups = $4 \implies$ No. of digit in root = 4Let root = 1000a + 100b + 10c + d 2) Except leftmost group, separate digits of all groups

i.e. 20 3 5 2 5 1 2 4 1 3 3 7 6

3) To extract 4 digit of perfect 4^{th} power, Only most significant '4+1' term are sufficient. Expanding $(1000a+100b+10c+d)^4$ according to descending power of ten using multinomial theorem, we get most significant 5 terms of expansion as

$$(1000a + 100b + 10c + d)^{4} = (a^{4}) * 10^{12} + (\frac{4!}{3!1!}a^{3}b) * 10^{11} + (\frac{4!}{2!2!}a^{2}b^{2} + \frac{4!}{3!1!}a^{3}c) * 10^{10}$$

$$+ (\frac{4!}{1!3!}a^{1}b^{3} + \frac{4!}{2!1!1!}a^{2}bc + \frac{4!}{3!1!}a^{3}d +) * 10^{9} + \dots$$

$$= (a^{4}) * 10^{12} + (4a^{3}b) * 10^{11} + (6a^{2}b^{2} + 4a^{3}c) * 10^{10}$$

$$+ (4a^{1}b^{3} + 12a^{2}bc + 4a^{3}d) * 10^{9} + \dots$$

4) Now highest a^4 less than or equal to 20 is $2^4 = 16$ $\therefore a = 2$.

Subtract a^4 from leftmost group and attach subtraction before next digit to get next gross dividend.

i.e. 2 2 5 1 2 4 1 3 3 7 6

Why This Works?

x is 4th root of N

$$f(x) = x^4 - 203525 1241 3376$$

(f(x)) has only one real positive root -According to Decarte's rule)

(Decartes Rule- Maximum no. of positive root of

f(x) = No. of sign (coefficient sign) changes in f(x))

$$2^4 \le 20 < 3^4 \implies 2^4 * 10^{12} \le 20 * 10^{12} < 3^4 * 10^{12}$$

 $2000^4 \le N < 3000^4$

$$\Rightarrow 2000^4 - N \le 0$$
 and $3000^4 - N > 0$

i.e.
$$f(2000) = -ve$$
 & $f(3000) = +ve$

Between two opposite signs of polynomial there always exist at least one real root.

.. Root (x) must lies between 2000 and 3000

Therefore most significant digit of root is 1

Similar reasoning can be given for subsequent steps of this method.

Aryabhata's square root method and many other method indirectly use this concept for intelligent guessing. 5) Now look up for highest integer b such that $(4a^3b) = 4*8*b = 32b$ less than or equal to 43

i.e.
$$32b \le 43 \Rightarrow b \le 1$$
 :: Highest b=1

subtracting $4a^3b = 32$ from 43, we get 11. Attach 11 before next digit 5 to get next gross dividend 115.

6) Look for highest integer c such that $(6a^2b^2 + 4a^3c)$ less than or equal to gross dividend 115.

i.e.
$$(6a^2b^2 + 4a^3c) \le 115 \implies 24 + 32c \le 115$$

 $32c \le 91 \implies c=2$

Subtracting $(6a^2b^2 + 4a^3c)$ from 115, we get 91-64 = 27, attaching it before next digit 2 to get next gross dividend.

7) Similarly,
$$4a^1b^3 + 12a^2bc + 4a^3d \le 272$$

 $\Rightarrow 8 + 96 + 32d \le 272$ $\Rightarrow 32d \le 168$
 $\Rightarrow d = 4$

Subtracting $(4a^{1}b^{3} + 12a^{2}bc + 4a^{3}d)$ from 272, we get 168-128 = 40, attaching it before next digit 5 to get next gross dividend.

$$\therefore (20352513413376)^{\frac{1}{4}} = 2124$$

Applicability

Above method is applicable to extract nth root of any number.

3) FORWARD -BACKWARD APPROACH

1) $(19477056866406726525123713030989808121)^{\frac{1}{7}} = ?$ (Given number is perfect 7th power)

Arranging number into 7 digit group from right to left

194 7705686 6406726 5251237 1303098 9808121

 \therefore No. of digit in 7th root = No. of groups = 6.

Let abcdef = f + 10e + 100d + 1000c + 10000b + 100000a be root

Here we can extract 4 least significant digit cdef by backward approach as below

$$(f+10e+10^{2}d+10^{3}c+10^{4}b+10^{5}a)^{7} = f^{7} + (7f^{6}e)*10 + (21f^{5}e^{2}+7f^{6}d)*10^{2} + (28f^{4}e^{3}+21f^{5}ed+7f^{6}d)*10^{3} + \dots + a^{7}*10^{35}$$

1) f^7 ends with 1 (End digit of 19477056866406726525123713030989808121)

$$\Rightarrow$$
 f = 1 : f^7 = 1, Subtracting it from ...08121, and eliminating last digit -1

.....0812Ø

2) $7f^6e = 7e$ ends with 2 (End digit of ...812)

$$\Rightarrow$$
 $e = 6$: $7f^6e = 42$, Subtracting it from ...8120812 and eliminating last digit0812

....077Ø

3) $21f^5e^2 + 7f^6d = 756 + 7d$ ends with 7 (End digit of ...77)

$$\Rightarrow 7d \text{ ends with } 7\text{-}6=1 \Rightarrow d = 3 \therefore 756+7d = 777, \qquad \dots 8077$$

Subtracting it from $\dots 8077$ and eliminating last digit . $\qquad \qquad -777$

.....370Ø

4) $28f^4e^3 + 21f^5ed + 7f^6c = 28*216 + 21*18 + 7d$ ends with 0 (End digit of ...370)

$$\Rightarrow$$
 7c ends with =4 \Rightarrow c = 2

$$\therefore$$
 cdef = 2361 -----(1)

Now we extract most significant 2 digit i.e. ab by forward approach as below

$$(10^5 a + 10^4 b + 10^3 c + 10^2 d + 10e + f)^7 = (a^7) * 10^{35} + (7a^6 b) * 10^{34} + (21a^5 b^2 + 7a^6 c) * 10^{33} + \dots + e^7$$

i)
$$a^7 \le 194 \implies a = 2$$

 $194 - a^7 = 194 - 128 = 66$

ii)
$$7a^6b \le 667 \implies b \le \left[\frac{667}{7*2^6}\right] \implies b = 1$$

 $\implies 667 - 7a^6b = 667 - 448 = 219$

iii)
$$21a^5b^2 + 7a^6c \le 2197$$

 $\Rightarrow 21*32*1 + 448c \le 2197 \Rightarrow 448c \le (2197 - 672)$
 $\Rightarrow c \le \left[\frac{1525}{448}\right] \Rightarrow c = 3$
 $2197 - 21a^5b^2 + 7a^6c = 1525 - 448*3$
 $= \text{Remainder}(\frac{1525}{448}) = 181$

$$0 < (c = 3) < 10 \implies \text{value of b is correct.}$$

:
$$ab = 21$$
 ----(2)

$$\therefore (19477056866406726525123713030989808121)^{\frac{1}{7}} = abcdef = 212361$$

(From equation (1) and (2))

GENERIC ALGORITHM

I. Initialization

1)
$$N = \text{Number to extract n}^{\text{th}} \text{ root}$$

= $D_{mn+n-1}D_{mn+n-2}D_{mn+n-3}.....D_2D_1D_0$

(Number in decimal notation)

$$= G_m G_{m-1} G_{m-2} ... G_1 G_0$$

(Arranging digits of number N into 'n digit group' from right to left)

2) Assume R =
$$(N)^{\frac{1}{n}} = a_m a_{m-1} a_{m-2} ... a_1 a_0 = (10^m a_m + 10^{m-1} a_{m-2} + + 10^2 a_2 + 10 a_1 + a_0)$$

- 3) MSB (R) = a_m = Highest nth power of integer that is less than or equal to G_m
- 4) $R_0 = G_m (a_m)^n$
- 5) Divisor = $n * (a_m)^{n-1}$

II. Iteration

```
For (k = 1 to m)  \{ \\ i) \ GD_k = R_{k-1} * 10 + D_{mn-k} \\ ii) \ ND_k = GD_k - C_k (a_m \ , a_{m-1} \ , ....., a_{m-k+1}) \\ where \ C_k (a_m \ , a_{m-1} \ , ....., a_{m-k+1}) \ denotes total contribution of <math>(a_m \ , a_{m-1} \ , ....., a_{m-k+1})  to coefficient of 10^{mn-k} in expansion (10^m \ a_m + 10^{m-1} \ a_{m-1} + 10^{m-2} \ a_{m-2} + .... + 10^2 \ a_2 + 10 \ a_1 + a_0)^n It is given as below
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$$\sum_{\substack{p_m+p_{m-1}+\ldots+p_{m-k+1}=n\\ mp_m+\ (m-1)p_{m-1}+\ (m-2)p_{m-2}+\ldots+\ (m-k+1)p_{m-k+1}=\min-k}}\frac{n!}{p_m!p_{m-1}!p_{m-2}!\ldots p_{m-k+1}!}(a_m)^{p_m}(a_{m-1})^{p_{m-1}}(a_{m-2})^{p_{m-2}}\ldots (a_{m-k+1})^{p_{m-k+1}}$$

iii) if
$$(ND_k < 0)$$
 // Validation
$$\{ a_{m-k+1} = (a_{m-k+1}) - 1$$
 // Decrease Last computed Quotient by 1
$$R_{k-1} = R_{k-1} + \text{Divisor}$$
 Continue same iteration // Go to step (i) to recompute ND_k
$$\}$$
 iv)
$$a_{m-k} = \text{Quotient } (\frac{ND_k}{Divisor}) = [\frac{ND_k}{Divisor}]$$

$$R_k = \text{Remainder } (\frac{ND_k}{Divisor}) = (ND_k) \% Divisor$$
 $\}$

FUTURE ENHANCEMENT

In above algorithm, we need to expand multinomial $(10^m a_m + 10^{m-1} a_{m-1} + 10^{m-2} a_{m-2} + + 10^2 a_2 + 10 a_1 + a_0)^n \ \text{to compute}$ $C_k(a_m, a_{m-1},, a_{m-k+1})$. Multinomial can be expanded using various approach that are listed as below

1) Recursive Approach

Use binomial theorem or Pascal triangle recursively to expand multinomial.

2) Algorithmic Approach

Develop or use existing integer processing based faster algorithm to break given integer power i.e. n into at most m positive integer i.e. $(p_m, p_{m-1},, p_0)$ such that

$$0*P_0+1P_1+2P_2+3P_3+...+(m-1)P_{m-1}+mP_m=L$$

In other word, develop efficient algorithm to determine all possible tuple

$$(p_0, p_1, p_2,, p_m)$$
 to solve
$$p_0 + p_1 + p_2 + ... + p_m = n$$

$$P_1 + 2P_2 + 3P_3 + + (m-1)P_{m-1} + mP_m = L \forall k \ 0 \le P_k \le \frac{L}{k}$$

After computing tuple, $C_L(a_m, a_{m-1},, a_{m-L+1})$ can be easily obtained by using multinomial theorem. Organise tuple and use relation between two tuple to avoid repetitive factorial computation.

3) Elimination and Retention Approach

Elimination and retention approach provides elegant, systematic way to expand multinomial. This method will be explored in upcoming book 'Modern Approach to Speed Math Secrets'. Note that all these algorithm and method for expansion of multinomial are also useful to compute any integer power of any number. Thus in future one can enhance this method by developing efficient algorithm to expand multinomial.

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