# Steganography Assisted Tor

Michael Freyberger ELE 454 Spring 2015 Prof. Pruchal

#### What is Steganography?

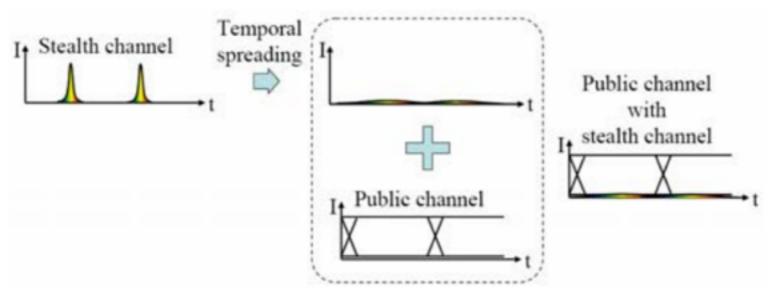


Figure 1: Schematic illustration of optical steganography using group velocity dispersion.

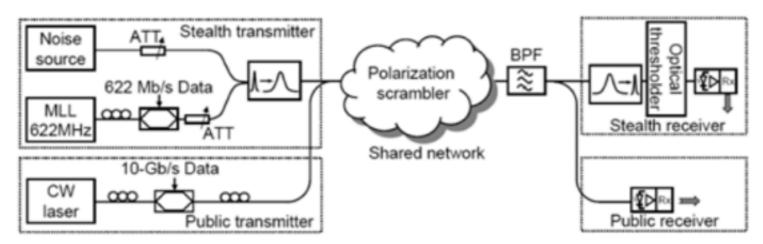
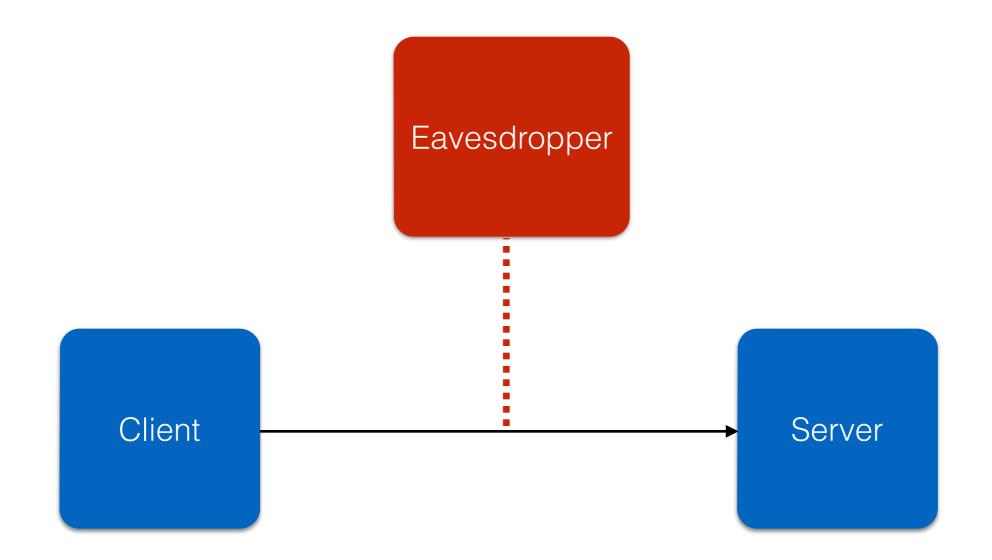
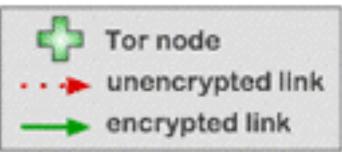


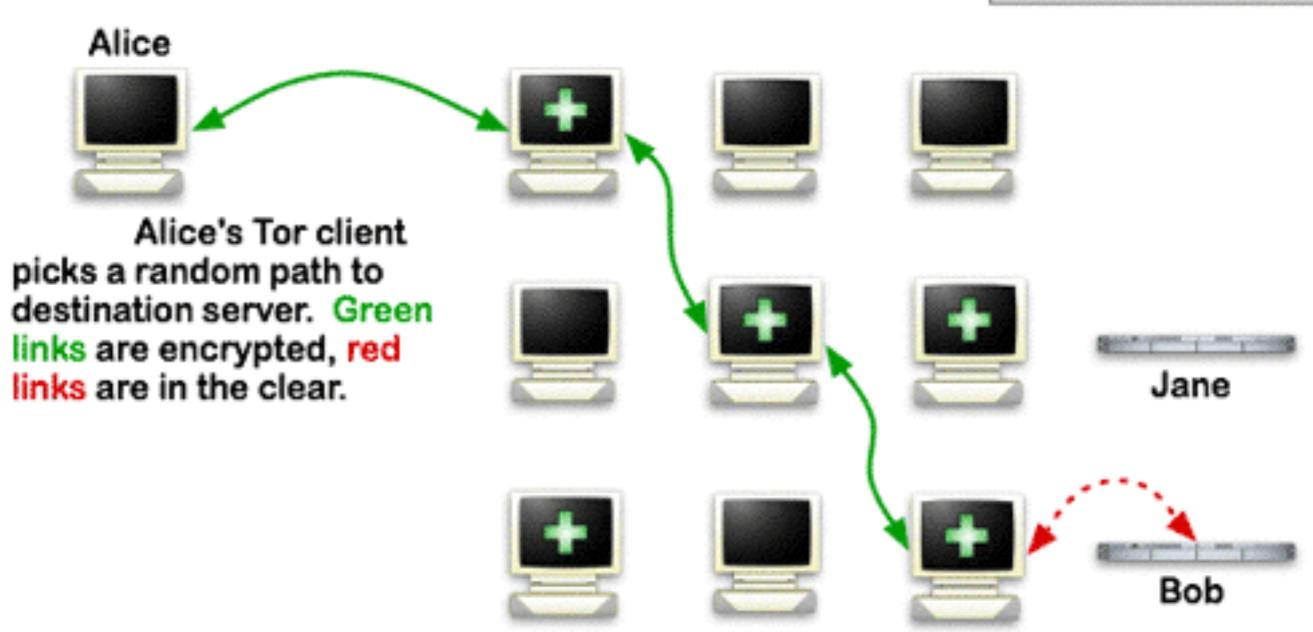
Figure 2: Schematic illustration of optical steganography. MLL: picosecond pulsed laser; ATT: variable optical attenuator; BPF: 3-nm optical band pass filter.

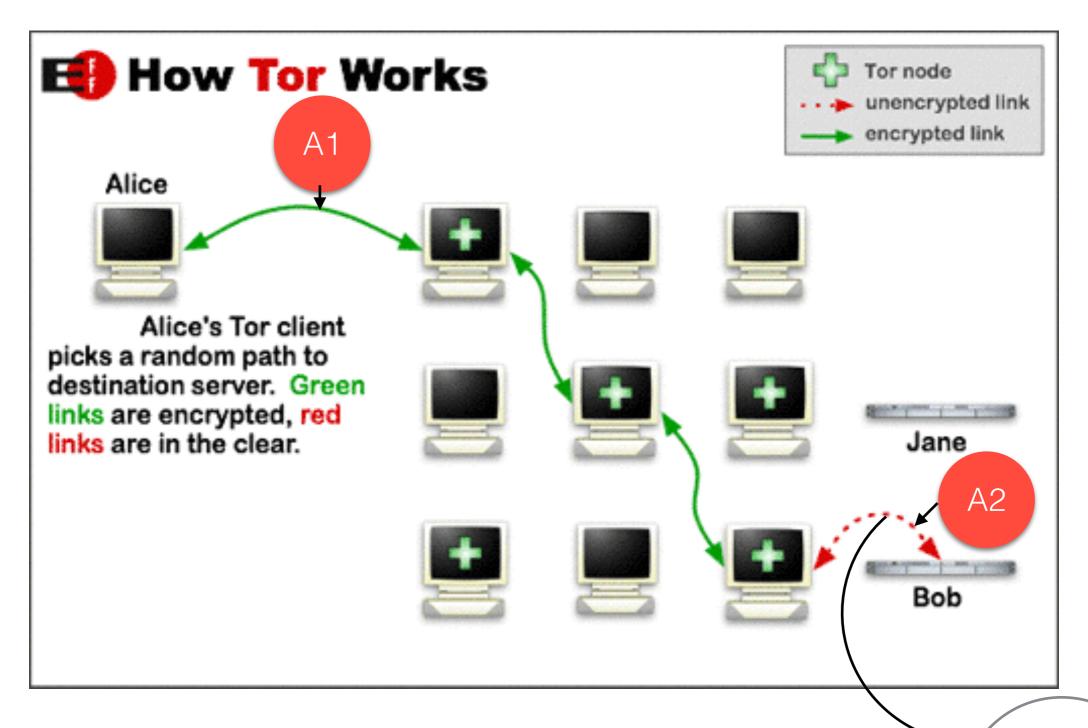


## The Problem Tor Solves

#### How Tor Works







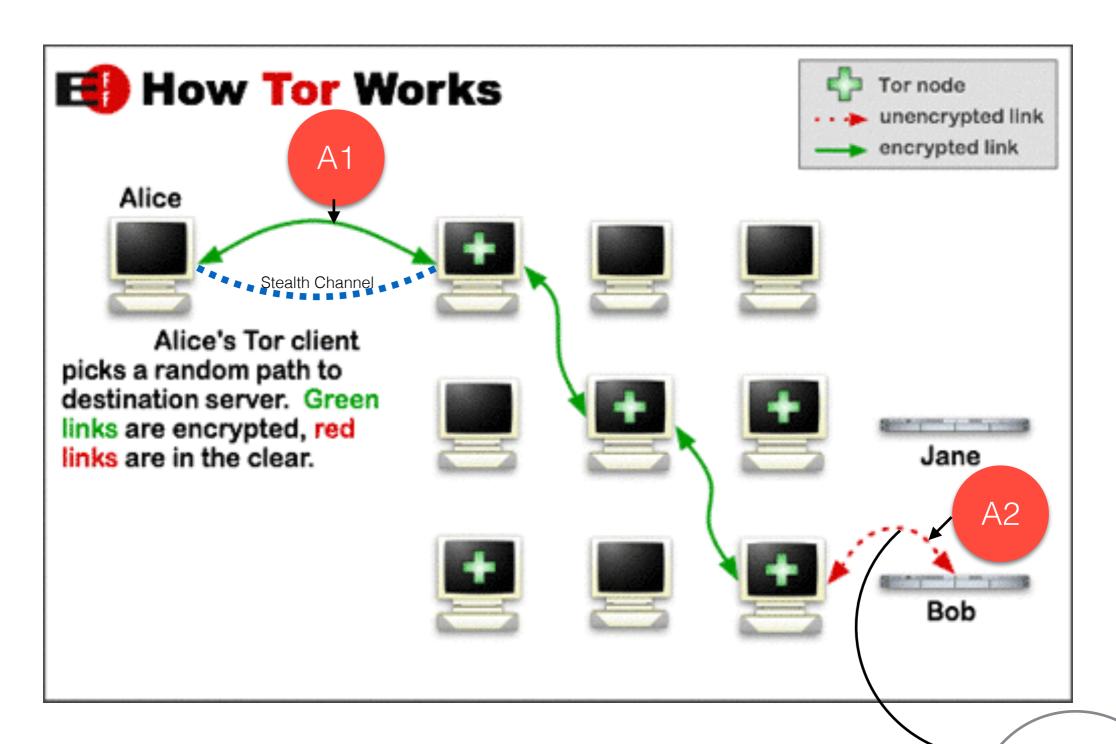
## The Attack

Assume
this is encrypted
with something like
https

## Possible Timing Analysis Attack

- 1. Each attacker measures the amount of packets that occur during a time window.
- 2. The attackers share the information and determine the cross correlation between the two packet streams.
- 3. If the correlation is above a certain threshold, the attackers conclude that they are on the same path

	Number of Packets	
Time Window: (Each window is 1 Second Long)	Client Link	Server Link
0	15	15
1	40	42
2	189	187
3	18	19
4	74	76
5	50	50
6	52	49
7	74	73
8	94	95
9	112	113
10	99	97



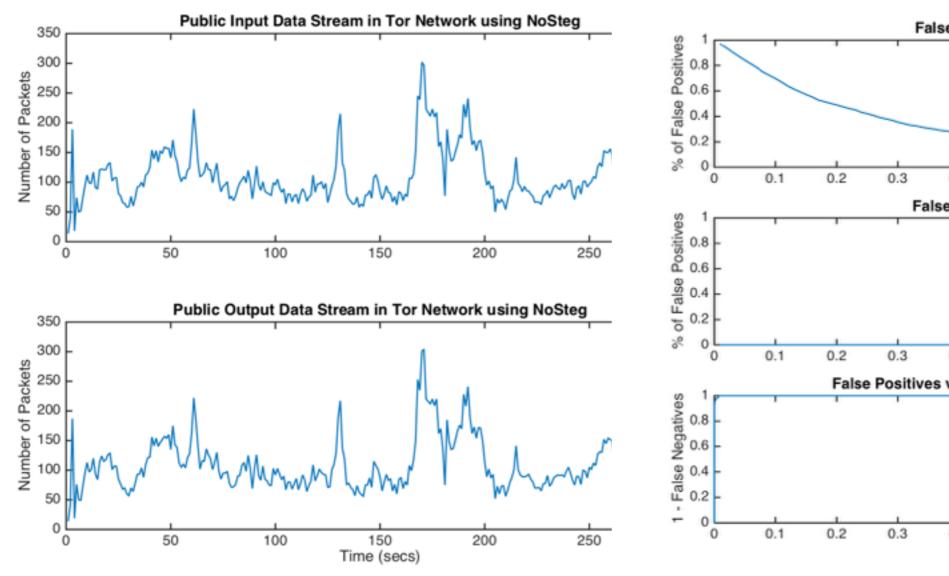
The Solution

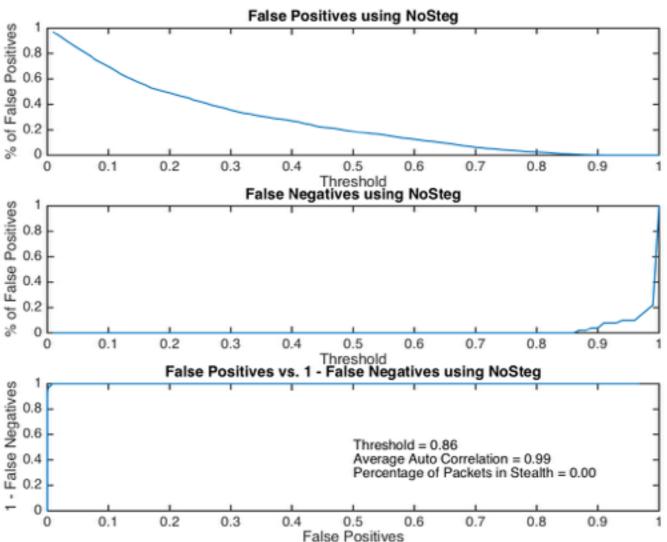
Assume
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#### Distinguishing Characteristics

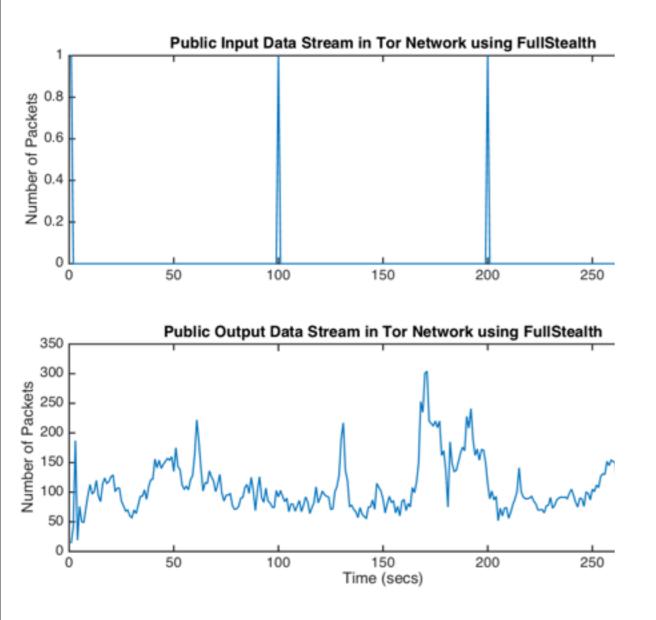
- Percentage of False Positives
- Percentage of False Negatives
- Percentage of Packets using Steganography

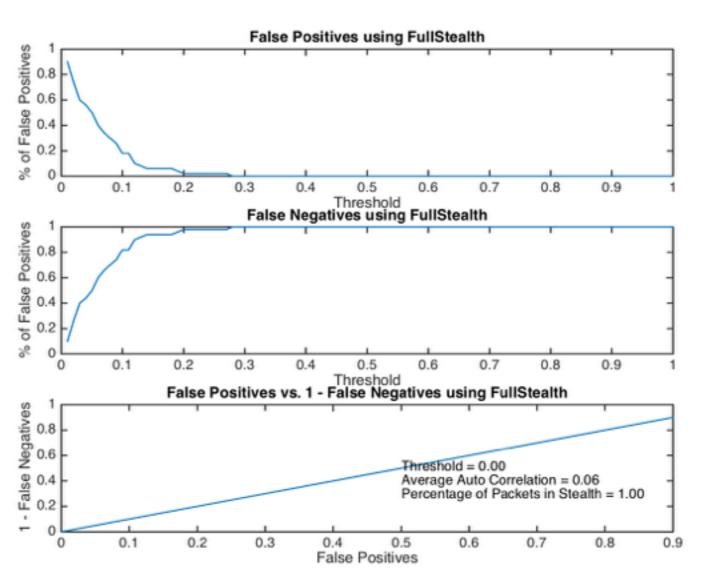
#### No Steganography



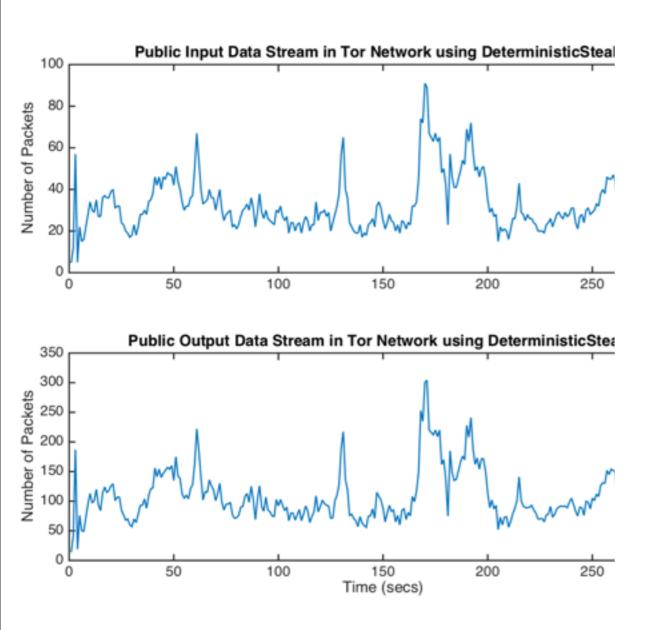


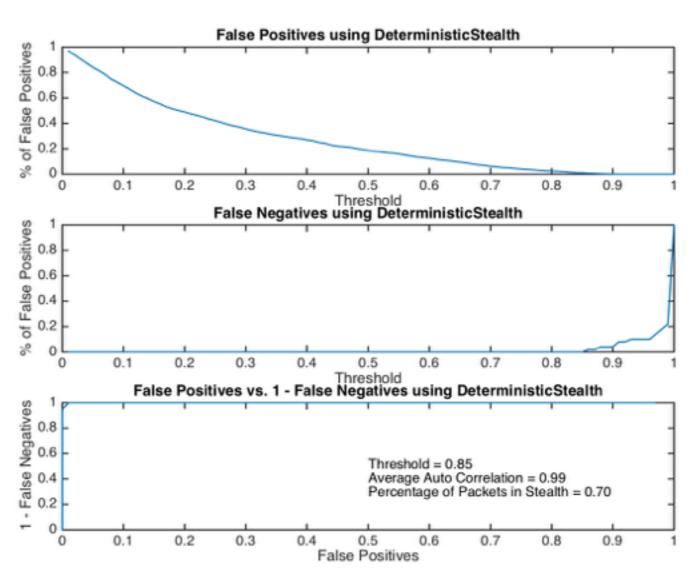
#### Full Steganography



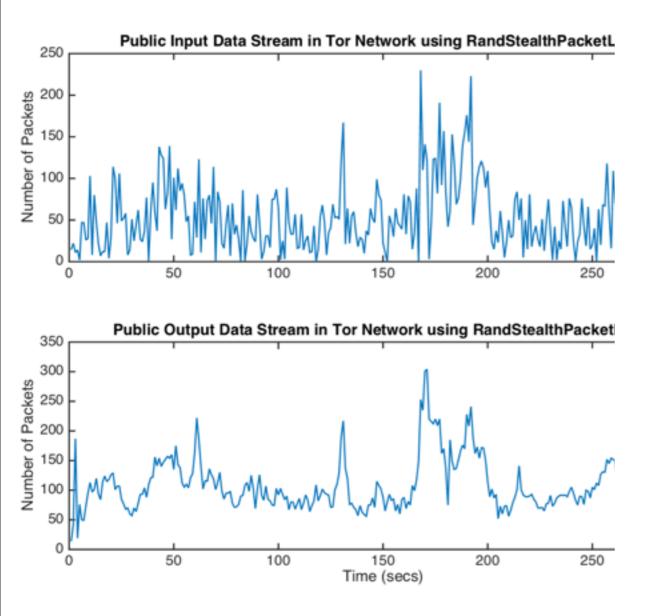


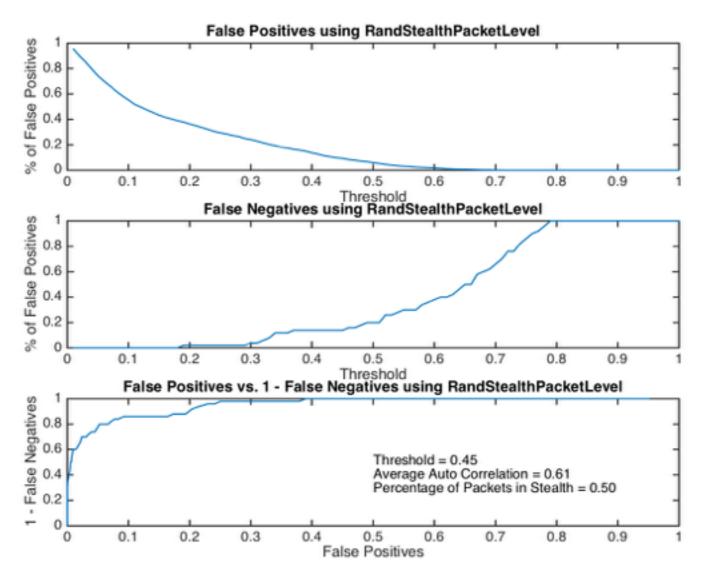
#### Deterministic 70% Steganography



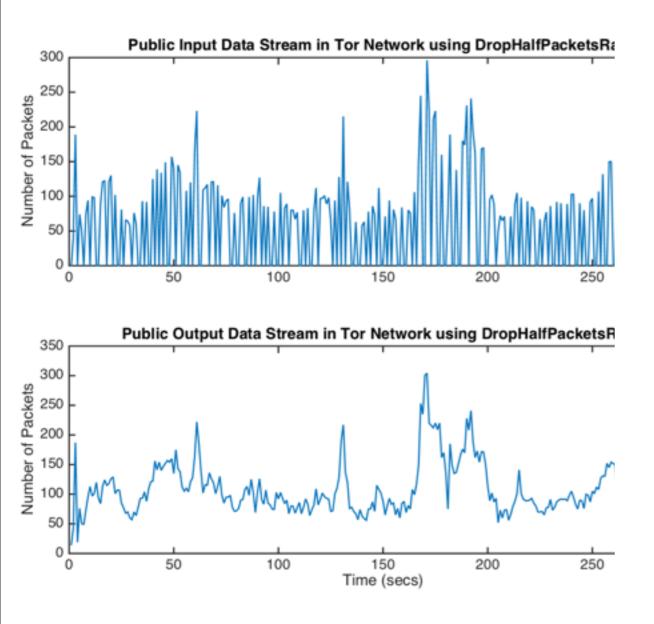


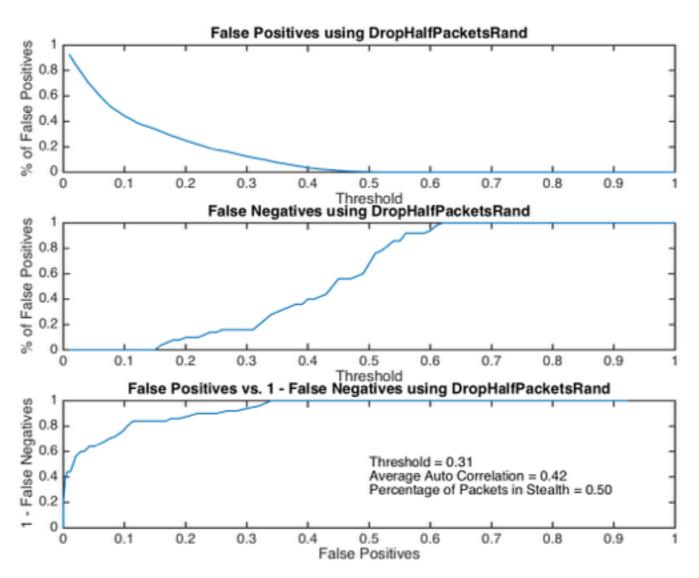
#### Random 50% Stealth: Limit Burst Sizes



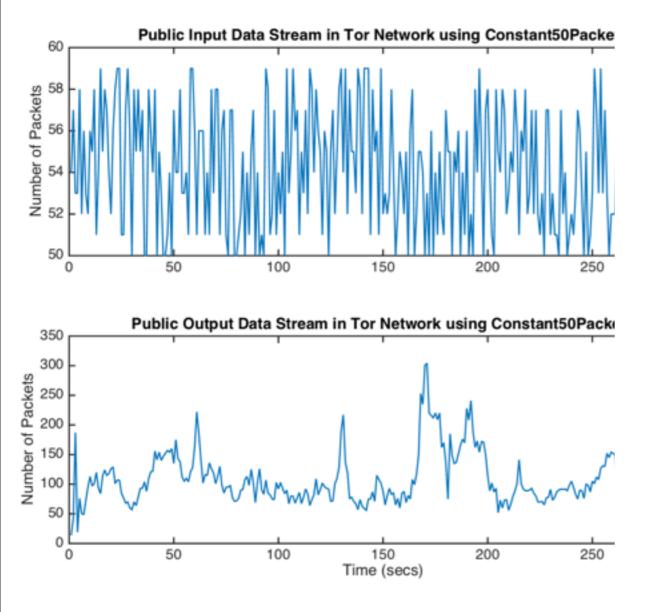


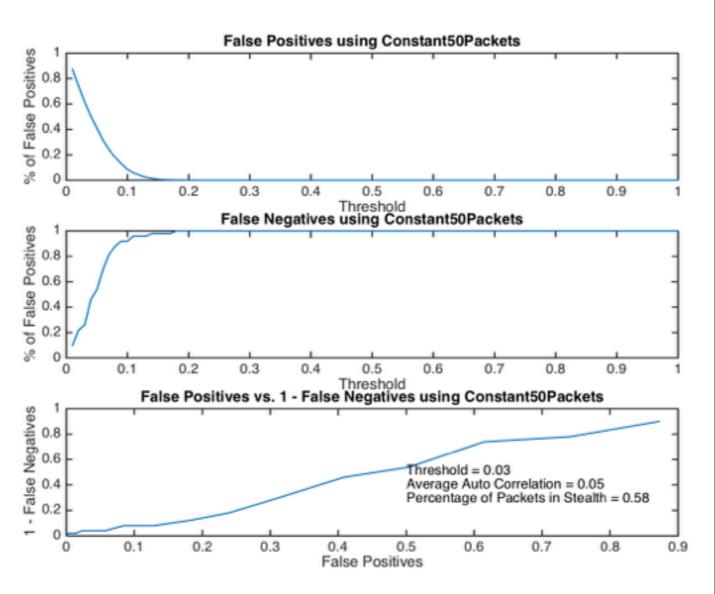
#### Random 50% Stealth: Drop Full Bursts



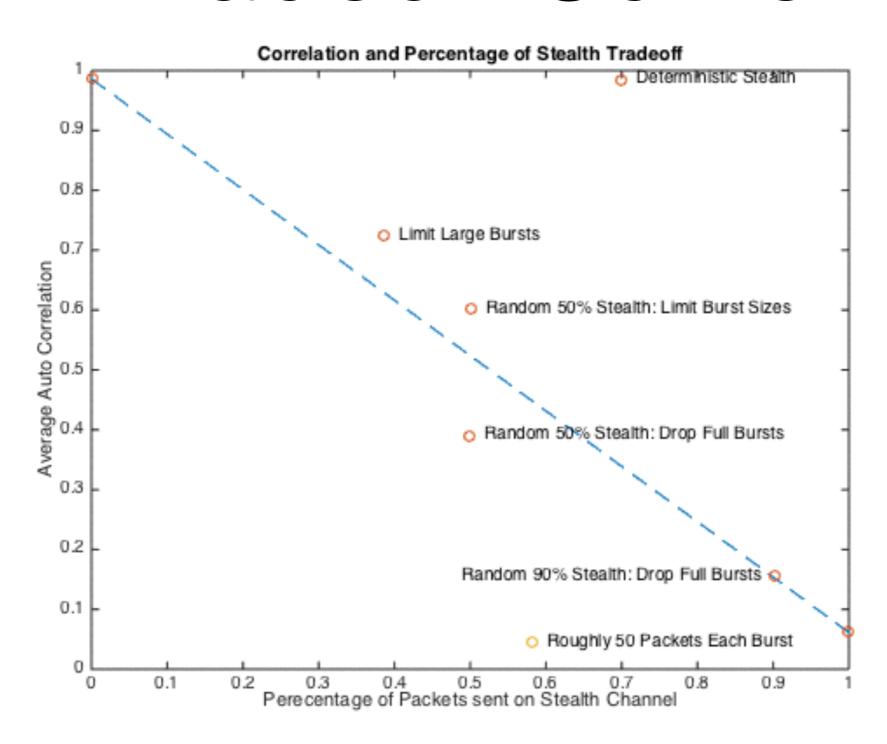


#### Roughly 50 Packets Each Burst





## Tradeoff Curve



### Conclusion

- Timing analysis attacks can be avoided with steganography, if applied correctly
- Bandwidth impact must be determined before selecting the scheme
- Implementation details TBD

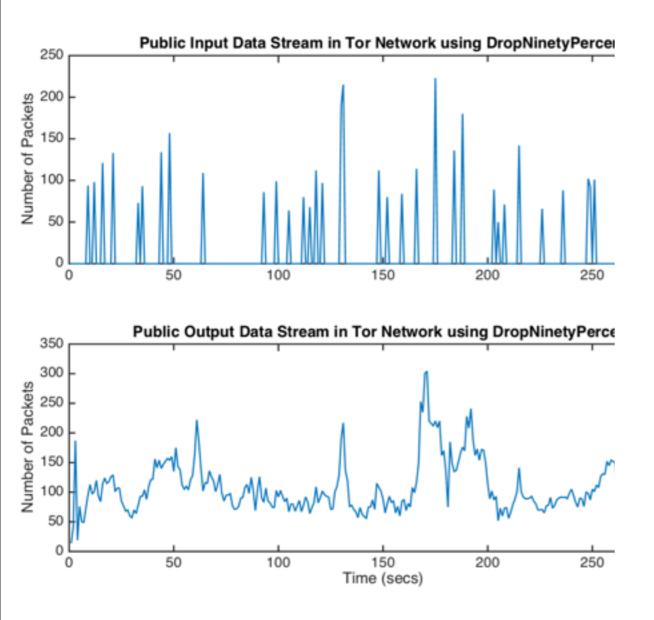
## References

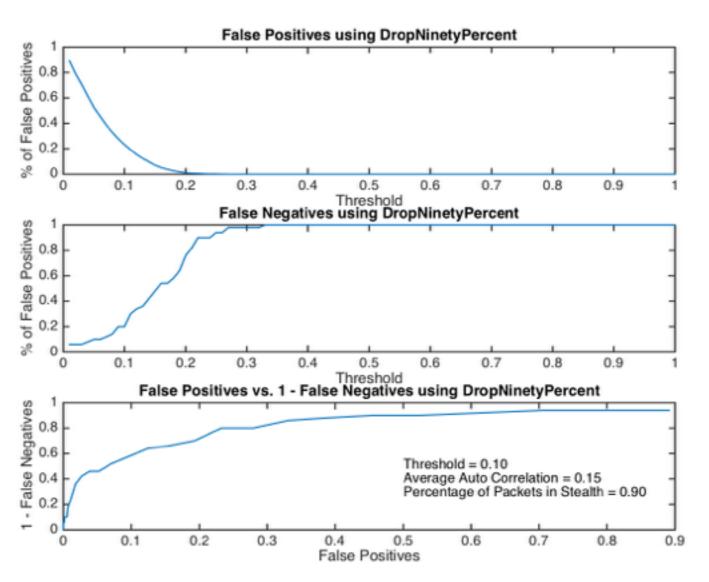
- Physical Layer Security Based on Optical Steganography and Optical Encryption. Ben Wu
- Optical steganography based on amplified spontaneous emission noise. Ben Wu,\* Zhenxing Wang, Yu, Tian, Mable P. Fok, Bhavin J. Shastri, Daniel R. Kanoff, and Paul R. Prucnal.
- Studying Timing Analysis on the Internet with SubRosa. Hatim Daginawala and Matthew Wright
- Timing analysis in low-latency mix networks: attacks and defenses. Vitaly Shmatikov and Ming-Hsi. Wang
- Timing Attacks in Low-Latency Mix Systems (Extended Abstract). Brian N. Levine, Michael K. Reiter. Chenxi Wang, and Matthew Wright

# Questions? Other possible schemes?

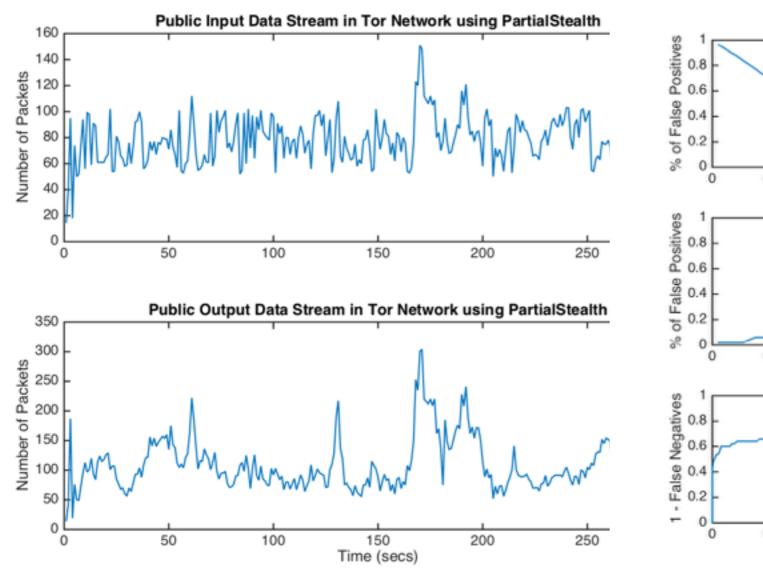
(We can demo your scheme now, if you'd like)

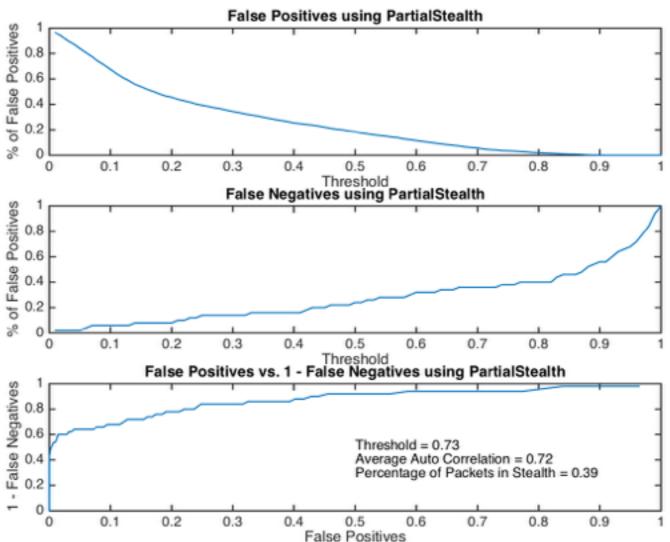
#### Random 90% Stealth: Drop Full Bursts





#### Limit Large Bursts





#### How does Steganography work?

- Hide the data within the noise of the signal
- Mask the information in the time domain
- Mask the information in the spectral domain

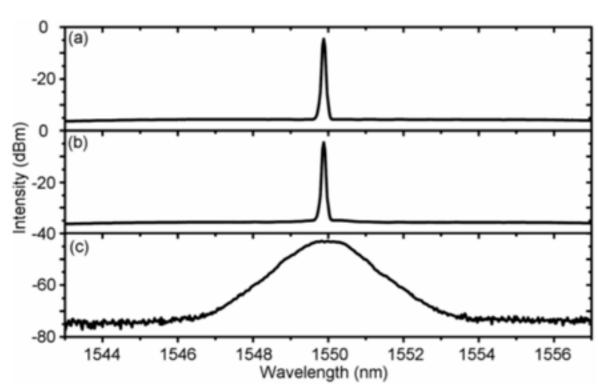


Figure 3: Optical spectra (a) public channel (b) public channel with the stealth channel (c) stealth signal alone.

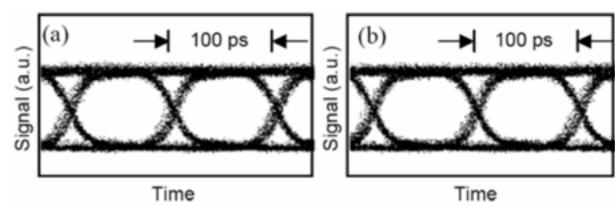
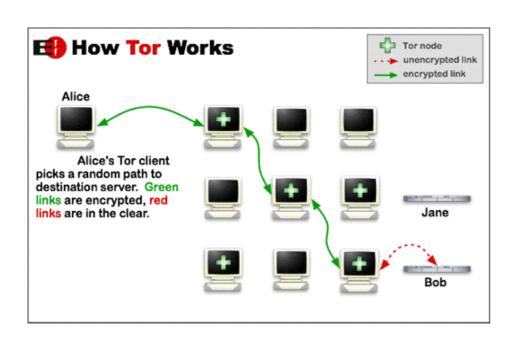


Figure 4: Eye diagrams (a) without stealth transmission (b) with stealth signal in the network.

## The Onion Router



- 1. Determine a path between 3 Tor nodes
- 2. Use asymmetric keys to determine 3 shared keys, one key for each node in the path
- 3. Encrypt the message with all three shared keys
- 4. Each node will decrypt the message with their shared key. This essentially peels off a layer of the onion.
- 5. Deliver the unencrypted packet to the server.
- 6. Use the same path for the return, but rather than peeling off a layer, add a layer. Each node encrypts the message with their shared secret.
- Decrypt using all 3 shared secret keys in the correct order. Completely peel open the onion.

### Correlation Formula

$$r(d) = \frac{\sum_{i} ((x_{i} - \mu) (x'_{i+d} - \mu'))}{\sqrt{\sum_{i} (x_{i} - \mu)^{2}} \sqrt{\sum_{i} (x'_{i+d} - \mu')^{2}}}$$

Xi: The amount of input packets in the ith window.

Xi': The amount of outgoing packets in the ith window.

μ: The mean packet size on the input stream

μ': The mean packet size on the output stream

d: The delay used to determine the cross correlation.

This value is typically set to 0.