Lecture 16 — Recursive Functions

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A function that includes a call to itself is said to be recursive.

At first glance, this might seem like a really strange idea. Why would a function call itself?

Recall that we said earlier that a common engineering strategy is to break a big problem down into some smaller problems.

Recursion is useful if we are breaking down a problem so that the subproblems are smaller versions of the same problem.

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Recursion Example: Factorial

You may have learned in math class about the factorial.

Commonly written n! in mathematical notation for a number n.

The factorial of a non-negative integer n is defined as the product of all positive integers less than or equal to n.

Examples:

```
1! = 1

2! = 2 \times 1 = 2

3! = 3 \times 2 \times 1 = 6

4! = 4 \times 3 \times 2 \times 1 = 24

5! = 5 \times 4 \times 3 \times 2 \times 1 = 120
```

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Recursion Example: Factorial

Observe that 5! could be rewritten as $5 \times 4!$.

The problem of 5! can be broken down into two subproblems:

- 1 4!
- $\mathbf{2}$ 5× the result of (1).

The subproblem 4! is a smaller version of the problem we're working on (5!), so this problem is a good candidate for recursion.

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Recursion Example: Factorial

Suppose now you are going to implement the factorial function.

```
int factorial ( int n )
{
    return n * factorial( n - 1 );
}
```

There is a problem with this implementation.

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The Factorial Problem

What happens when the expression n - 1 becomes zero or negative?

This loop will continue forever...

Except in practice this will be stopped by an error.

Running this program produces an error called a Stack Overflow.

In a later lecture we will discuss what a stack is and how it works. For now, a simplified view of the problem.

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Stack Overflow

When main() calls factorial(), the computer needs to keep track of where it was in main() at the time that it went to the other function.

It puts that information in a designated memory area called the stack.

Each time factorial() calls itself, more information is added to the stack to keep track of where it was in factorial().

If we do this too many times, the stack gets "full" (exceeds available memory for it) and this results in a stack overflow error.

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```
int factorial ( int n )
{
    return n * factorial( n - 1 );
}
```

The above implementation lacks a stopping condition.

The function will keep calling itself until a stack overflow occurs.

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Factorial Problem

Even if program execution did not terminate abnormally as a result of a stack overflow, there's another problem, mathematically.

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

When we reach 1, we no longer multiply by the next smallest integer.

Our stopping case is therefore when n is 1.

The way the code is written, when n is 1 the evaluation goes on and the expression to the right of return is 1 * factorial(\emptyset).

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```
int factorial ( int n )
{
    if ( n == 1 )
    {
       return 1;
    }
    return n * factorial( n - 1 );
}
```

The revised implementation has the stopping condition of n equals 1.

The function will keep calling itself until n is 1. When that happens, 1 is returned.

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To make proper use of recursion, we need:

- 1 One or more cases in which the function calls itself; and
- 2 One or more cases in which the function does not call itself.

Point (2) is called the stopping case or base case.

As we saw, without a properly defined stopping case, recursion will result in a stack overflow

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Recursion versus Iteration

Recursion can be a difficult or confusing topic. Is it strictly necessary?

Any task that can be accomplished with recursion can be done without using recursion, such as using a loop (iteratively).

As a small note on efficiency, a recursively written function may run slower than an iteratively written one. Recursion has two advantages:

- 1 Write less code; and
- 2 The computer tracks state on the stack; iteratively, we must keep track of the state ourselves.

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Here's the factorial function implemented iteratively:

```
int factorial ( int n )
{
    int product = 1;
    for ( int i = n; i > 1; i-- )
    {
        product *= i;
    }
    return product;
}
```

The iterative implementation uses a for loop, but it could have been written using a while loop.

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Thinking Recursively

When designing a recursive function, there are three important criteria to consider:

- There is no infinite recursion;(A chain of recursive calls eventually reaches a stopping case)
- 2 Each stopping case returns the correct value for that case; and
- The final returned value is correct if the recursive call(s) returns the correct value(s).

This is an example of *mathematical induction*. If you have not yet learned about it, ignore this for now; you will see it next term in ECE 103 (Discrete Mathematics).

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Applying Recursive Thinking

Consider another mathematical problem, exponentiation: x^y .

If we wrote a function signature for exponentiation: int pow (int \mathbf{x} , int \mathbf{y})

Is this problem a good candidate for recursive solution?

Mathematically, $a^b = a \times a^{b-1}$.

So, yes, this is the kind of problem that lends itself well to recursion.

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Applying Recursive Thinking

Now, let's write an implementation for pow().

```
// Precondition: y >= 0
// Postcondition: returns x to the power of y
int pow ( int x, int y )
{
    if ( y == 0 )
    {
        return 1;
    }
    return pow( x, y - 1 ) * x;
}
```

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Applying the Test

Let's examine the design criteria and see if this is going to work.

1. There is no infinite recursion.

The second argument to pow(x, y) is decreased by 1 in each call, so eventually we must get to pow(x, 0), a stopping case. (As long as the precondition is not violated.)

Thus, there is no infinite recursion; criterion 1 is satisfied.

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Applying the Test

2. Each stopping case returns the correct value for that case. Yes. $x^0 = 1$ is mathematically correct.

3. The final returned value is correct if the recursive call returns the correct value.

pow(x, y - 1) * x follows the rule that
$$a^b = a^{b-1} \times a$$
.

Criteria 2 and 3 are satisfied.

Having checked those things, we can now be satisfied that the implementation of pow() is correct.

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Here's the exponential function implemented iteratively:

```
int pow ( int x, int y )
{
    int result = 1;
    for ( int i = 0; i < y; i++ )
    {
        result *= x;
    }
    return result;
}</pre>
```

Like the factorial function, the iterative implementation uses a for loop, but it could have been written using a while loop.

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