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## 1 Basic considerations

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1.1 Notation
\prec \bullet \succ indicates that the programmer must provide an appropriate replacement.
For example \langle expr \rangle can be replaced by 10.0*sin(pi/6)
[•] indicates that the programmer can either provide a replacement or opt to omit it.
For example \langle T \rangle / \langle aT \rangle = \langle expr \rangle can be replaced by either double d or double d=3.0
\bullet \star indicates an option for the programmer.
For example const \langle T \rangle \langle aT \rangle = \langle expr \rangle (\langle expr \rangle) can be replaced by either const bool b=true or const bool b(x==3)
Case-sensitivity and semicolon: C++ is a case-sensitive language, there are (usually) no continuation lines
and the end of a statement (except a block statement) is marked with a semicolon (;)
    Comments and main segment
// comment terminating with a new-line character
/* comment terminating with the symbol */
int main()
                                    // argument-free version main segment
\{[\prec statements \succ]\}
                                   // corresponding statement block
int main(int argc, char **argv) // main segment with arguments, argc is the number of arguments passed to program and, if not zero,
                                   // char[0]..char[argc-1] are pointers to character sequences containing command-line arguments
\{[\prec statements \succ]\}
                                   // corresponding statement block
Relevant utilities:
Premature termination can be forced by using the function exit() from the header #include<cstdlib>.
A system command can be executed by using the function system("≺command≻") available from the header #include<cstdlib>.
Assertions can be established by using the macro assert(\precboolean_expr\succ) from the header #include<cassert>.
1.3
    File inclusion and namespace use
#include <≺system_header≻>
                                      // inclusion of a system header
#include "≺file≻"
                                      // inclusion of a user header (the extension .hpp is common for ≺file≻)
using namespace ≺namespace_name≻; // provides access to the scope of a given namespace (in this case ≺namespace_name≻)
For example:
using namespace std;
                                      // provides access to the Standard namespace
                                      // ensures that the corresponding header is included once (non-standard but usually accepted)
#pragma once
Source code files and header files (combined files sent to linker):
≺file> with ≺cpp> extension
                                      // Source code file: contains stand alone source code to be compiled (other extensions possible!)
\precfile\succ with \prechpp\succ extension
                                      // Header file: contains source code intended for inclusion (other extensions possible!)
extern "C" {≺C functions≻}
                                      // functions compiled by C typically enclosed by the curly braces
For example:
extern "C" {int mult(int* i,int* j);} // example of a C function prototype
extern "C" { #include "header.h" }
                                           // example of a C header
   Types and declarations
     Fundamental types
≺C≻ is either a fundamental type or a type introduced by the user 1.
Common fundamental types:
void, char, int, bool, double, float
auto: in the context of variable initialization, leaves to the compiler the decision to determine the appropriate type
For sizes of containers, the header #include<cstdlib> includes the type size_t.
Type use: The keyword decltype can be used to determine the type of an expression:
decltype(<expr>) <avariable>;
decltype(\prec expr\succ) \prec areference \succ = \prec expr\succ;
For example<sup>2</sup>:
typename vector<int>::iterator aiterator=v.begin() // can be replaced by auto aiterator=v.begin()
decltype(adouble) anotherdouble;
                                                       // make use of an existing type
decltype((adouble)) adoublereference=adouble;
                                                       // another use of an existing type
Examples of declarations:
                                // integer variable declaration
              \precaint\succ;
size t
              \precaunsigned\succ;
                                // non-negative integer size_t is a type synonym (requires cstdlib) for a unsigned integer
                                // character variable declaration ('a', 'b', ...)
char
              \precachar\succ;
float
              \precafloat\succ;
                                // real (32 bits in many implementations) variable declaration
                                // real (64 bits in many implementations) variable declaration
double
              \precadouble\succ;
                                // boolean variable declaration (values are either true or false)
bool
              \precabool\succ;
              \precaauto\succ = \precexpr\succ;// the type is inferred by the compiler from the type of \precexpr\succ;
auto
```

<sup>&</sup>lt;sup>1</sup>In the present context, these can be struct and enum <sup>2</sup>auto can effectively replace most typename specifiers

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2.2 Typical declaration on the stack
[static] \prec T \succ \prec aT \succ [(\prec LE \succ) | =\prec C \succ() | =\prec expr \succ]; // declares a type where \prec C \succ() invokes the default constructor
Examples:
                                                            // global int variable initialized with the value 10
static int aint(10);
float afloat=30.0;
                                                            // float variable initialized with the value 30.0
double adouble=double();
                                                            // local double variable initialized with the default value, 0.0
complexclass acomplex(0.0,1.0);
                                                            // hypothetical complex number initialized as a unit imaginary
\prec T \succ is given in general as
[const] \prec C \succ [\clubsuit [const]]
The above keywords are:
                                                            // maintains the variable value after leaving the scope of declaration
static
                                                            // it automatically initializes to the default value
                                                            // qualifier that, as prefix, forces the variable to hold a constant value
const
                                                            // (attempts to change the value by de-referencing are flagged as errors at compile time)
                                                            // qualifier that, as suffix, declares the address stored in the pointer to be constant
const
           Temporary variables are forced to be constant (const) by the compiler
Warning:
The symbol & is one of the following:3
        // pointer, when an address (and type) of a variable is stored, only use for heap allocation
        // with 'new' and when a 0 can possibly be returned
&
        // reference, when an alias of a given variable is stored, requires initialization
&&
        // rvalue reference, similar to & but allowing the assignment of temporaries
*&
        // reference of a pointer, when an address stored at a pointer is allowed to change,
        // useful to pass pointers to functions with new and delete operations
The list of sent expressions ≺LE≻ is an ordered comma-separated set of N expressions⁴.
If N==0 parenthesis are omitted around \precLE\succ.
For N expressions, with N \neq 0,
\precLE\succ \equiv \precexpr_0\succ, \precexpr_1\succ, ..., \precexpr_N-1\succ
Examples:
const double adouble(3.0);
                                          // declares a constant double with the value 3
                                          // declares a pointer to a constant integer
const int* apint=&aint;
double* const apdouble=new double(); // declares a constant pointer to a double
bool& abool2=abool1;
                                        // declares and initializes a boolean reference
                                          // declares and initializes a constant boolean reference
const bool& abool3=false;
2.3 Array in the stack
\prec T \succ \prec array \succ [\prec n \succ]; // declares an array containing, sequentially, \prec n \succ elements of type \prec T \succ. For large arrays, heap allocation
                          // is recommended. Multidimensional arrays are also possible, but seldom needed.
2.4 Heap space management
Space in the heap will have to be accessed by variables which are declared in a given local scope.
\prec apT \succ and \prec arrayT \succ are pointers of type \prec T \succ and their declarations are:
\prec T \succ * \prec apT \succ = 0;
                               // declares a pointer to \prec T \succ and sets it to 0
\precT\succ* \precarrayT\succ=0;
                              // also declares a pointer to \prec T \succ but additional space will be reserved, the heap space is managed as:
[\prec apT \succ] = \text{new } \prec T \succ [(\prec LE \succ)]; // \text{ reserves space for a variable of type } \prec T \succ
                              // uses ≺apT≻
delete \prec apT \succ;
                              // frees the space which was previously reserved
                               // do not free twice, set to zero after deletion
\precapT\succ=0
\prec \operatorname{arrayT} \succ = \operatorname{new} \prec \operatorname{T} \succ [\prec \operatorname{n} \succ]; // reserves space for an array of type \prec \operatorname{T} \succ with \prec \operatorname{n} \succ elements
                              // uses ≺arrayT≻
delete[] \prec arrayT \succ;
                               // frees the space which was previously reserved
                               // do not free twice, set to zero after deletion
\precarrayT\succ=0
Smart pointers (the header #include<memory> is required, see also RAII for a generalization):
scoped_ptr\langle T \rangle \rightarrow \langle apT \rangle = new \langle T \rangle [(\langle LE \rangle)];
                                                          // declares a single-owned smart pointer with transference on copy
shared_ptr<T>> <apT>=new <T>[(<LE>)]; // declares a reference-counted smart pointer, container-friendly
Destroy utility function for standard pointers:
template<typename \prec T \succ inline void destroy(\prec T \succ * \& \prec apT \succ){if(\prec apT \succ)delete \prec apT \succ; \prec apT \succ = 0;}
2.5 Type synonyms or 'typedefs'
\texttt{typedef} \  \, {\prec} \texttt{T}{\succ} \  \, {\prec} \texttt{Tsynonym}{\succ}
                                   // this specifier declares the identifier \langle Tsynonym \rangle as a synonym of \langle T \rangle
                                   // as it is often easier to write ≺Tsynonym≻. Can also be used in structs. An example of use of synonyms:
typedef int* pint;
                                   // pointer to int also identified as pint
pint apint=new int;
                                   // typical use of pint
     Enumerations
\mathbf{2.6}
Enumerations are objects which can have a finite number of values, for which synonyms are defined.
Definition of a enumeration:
enum \prec E \succ
{
\precname_0\succ[=\precinteger_expr_0\succ],
\precname_1\succ[=\precinteger_expr_1\succ],
};
Declaration and use of a enum object:
\langle E \rangle \langle aE \rangle;
\prec aE \succ = \prec integer\_expr\_I \succ;
   3Double, triple, etc pointers, **, ***, also exist but are not necessary with C++ since Standard Library containers can be nested and are preferable.
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 $<sup>^4</sup>$ The ( $\prec$ LE $\succ$ ) option may fail in certain odd cases, in that case an extra set of parenthesis must be added.

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2.7 Unions
Unions are amalgamations of data. Memory sharing is allowed for certain data types (C++11).
Definition of a union:
union ≺U≻
\langle C_0 \rangle \langle aC_0 \rangle;
\langle C_1 \rangle \langle aC_1 \rangle;
};
Declaration and use of a union object:
\prec U \succ \prec aunion \succ;
\precaunion\succ=\preccompatible_expr\succ;
      Type identification during run time - RTTI
The header <typeinfo> provides the functions typeid(\langle T \rangle) and typeid(\langle expr \rangle). This should only be used for comparisons.
Examples of application:
if(typeid(complex)!=typeid(notacomplex))cout<< "Not a complex";</pre>
cout<<typeid(complex).name()<<endl; // compiler-dependent name</pre>
     Operators and expressions
      Nomenclature of Rexpr and Lexpr
                        // an expression that, const apart, \underline{\operatorname{can}} be used either in the left or right-hand-side of an attribution
\precLexpr\succ
                        // an expression which \underline{\text{cannot}} be used, even if non-constant, in the left-hand-side of an attribution
\prec \texttt{Rexpr} \succ
                        // a general expression, either \prec Rexpr \succ or \prec Lexpr \succ
\precexpr\succ
                        // an expression which has a boolean value or one that can be implicitly converted to a bool
\precboolean_expr\succ
\prec \mathtt{integer\_expr} \succ
                        // an expression which has an integer value
      Typical operations (\prec a\succ, \prec b\succ \text{ and } \prec c\succ \text{ are expressions}, \prec i\succ \text{ is an integer expression and } \prec l1\succ \text{ and } \prec l2\succ \text{ are boolean expressions})
                                                                 // assignment, the value stored in \langle a \rangle will be lost and
\prec a \succ = \prec b \succ ;
                                                                // \prec a \succ will now contain the value of \prec b \succ
++ \prec i \succ, -- \prec i \succ, \prec i \succ ++, \prec i \succ --;
                                                                // pre-increment, pre-decrement, post-increment and post-decrement
! \prec 11 \succ, \prec 11 \succ | | \prec 12 \succ, \prec 11 \succ \&\& \prec 12 \succ;
                                                                // negation, "or" and "and". The result is a boolean
\langle a \rangle = = \langle b \rangle, \langle a \rangle ! = \langle b \rangle;
                                                                // equal, not equal
\prec a \succ > = \prec b \succ, \prec a \succ < = \prec b \succ, \prec a \succ < \prec b \succ, \prec a \succ > \prec b \succ; // greater than or equal, less than or equal, less than, greater than
\forall a \succ + \forall b \succ, \forall a \succ - \forall b \succ, \neg \forall a \succ, + \forall a \succ;
                                                                // addition, subtraction, unary minus, unary plus
\forall a \succ * \forall b \succ, \forall a \succ / \forall b \succ, \forall a \succ \% \forall b \succ;
                                                                // multiplication, division, remainder
                                                                // composition of attribution and any binary operator @ this is \langle a \rangle = \langle a \rangle \langle b \rangle;
\prec a \succ 0 = \prec b \succ ;
\prec a \succ (\prec b \succ);
                                                                // use of () in a function invocation
\prec a \succ * (\prec b \succ + \prec c \succ);
                                                                // use of () in expression grouping
\prec a \succ . \prec b \succ ;
                                                                // access to a member variable
\prec a \succ . \prec b \succ (\prec c \succ);
                                                                // access to a member function
\langle a \rangle - \rangle \langle b \rangle (\langle c \rangle);
                                                                // access to a member function when \prec a \succ is a pointer
*≺a≻;
                                                                // value of \langle a \rangle: contents stored at address \langle a \rangle
                                                                // address of \prec a \succ
&≺a≻;
\prec a \succ [\prec i \succ] or *(\prec a \succ + \prec i \succ) or \prec i \succ [\prec a \succ];
                                                                // accesses the contents stored at address \langle a \rangle + \langle i \rangle
3.3 Block statement
Syntactically, when one statement is expected and the programmer needs to insert more than one,
curly braces are used {} to create a block statement. For N statements, the block statement has the form:
{\prec statement_0 \succ; \prec statement_1 \succ; \dots; \prec statement_N-1 \succ;}
Example:
if(a!=b)a=b;
                                     // one statement expected from the if condition
if(a!=b)
{ a=b;
c=b; }
                                     // when two or more statements are needed, a block statement must be used
      Old-fashioned type cast operator
For conciseness reasons, I prefer the old-fashioned C-style type cast operator.
The most common application of type cast occurs in pointers of structures (or classes) when inheritance is involved:
struct \langle C2 \rangle : \langle C1 \rangle and access member functions in \langle C2 \rangle not present in \langle C1 \rangle :
                                                 // base allocated as base
\prec C1 \succ * \prec aP1C1 \succ = new \prec C1 \succ;
\prec C1 \succ * \prec aP1C2 \succ = new \prec C2 \succ;
                                                 // base allocated as derived
\langle C2 \rangle * \langle aP2C2 \rangle = (\langle C2 \rangle *) \langle aP1C2 \rangle;
                                                // cast to a \precC2\succ* pointer
The cast operator can invoke the overloaded cast operator if it is defined for a given struct or class.
Type conversion operations are performed at compile-time.
3.5 Sizeof operator
                      // returns the number of bytes of type \langle C \rangle
sizeof(\langle C \rangle)
sizeof(\prec expr \succ) // returns the number of bytes resulting from the evaluation of \prec expr \succ
Examples:
sizeof(double) // returns 8 in most cases
float f=3.0;
sizeof(f-33.0) // returns 4 in most cases
      Conditional (ternary) operator
\precboolean_expr\succ?\precexpr_1\succ:\precexpr_2\succ;
                                                    // returns ≺expr_1≻ if ≺boolean_expr≻ evaluates to true and ≺expr_2≻ if not
                                                     // note that ≺expr_1≻ and ≺expr_2≻ evaluations may be, in general, of different types
```

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For example:
apint == 0?int(): *apint // returns the default value of int if apint is 0 or returns the value stored at apint's address.
    Functions
     Prototype and definition
\prec T \succ \prec name \succ (\prec LP \succ);
                                          // function prototype or declaration
[inline] \langle T \rangle \prec name \rangle (\langle LR \rangle)
                                          // definition (note the difference between \prec LR \succ and \prec LP \succ and the lack of semicolon)
\{//\longleftarrow \text{ beginning of the function block}\}
\precstatements\succ
[return ≺expr≻;]
                                           // return of \precexpr\succ only if \precT\succ is not void
}// \leftarrow end of the function block
\precLR\succ is the list of N received variables
If N==0 parenthesis are not omitted.
For N arguments, \prec LR \succ is given as:
\precLP\succ is the list of N received types
If N==0 parenthesis are not omitted.
For N arguments, \prec LP \succ is given as:
\precT_0\succ, \precT_1\succ, ..., \precT_N-1\succ
In \langle LR \rangle, the expressions \langle expr_i \rangle provide the default value of the corresponding argument, and must satisfy:
If ≺expr_i≻ is present then ≺expr_i+1≻ must also be present. Strict reference arguments (&) cannot have default values.
Examples<sup>5</sup>:
                                                             // declaration of a function named test
void test(const int&,char*,const double&);
void test(const int& i,char* c,const double& d=100.0);// definition of test
double* newvector(size_t dim=10)
\{\ldots\}
return new double[dim] ;
                                                              // with return value
}
4.2 Function overloading
We can declare and define functions with the same name but with different lists of received types.
For example:
void print(const int& val)
                                 // will be invoked when an int is used in \precLE\succ
void print(const double& val) // will be invoked when a double is used in \prec LE \succ
4.3 Argument passing
Each \prec T\_I \succ in \prec LP \succ and \prec LR \succ should have one of the following forms:
\prec C \succ
              // passing by value, can receive any ≺expr≻ and will
              // not modify the argument (avoid except with reference-counting smart pointers)
const \prec C \succ \& // passing by constant reference, can receive any \prec expr \succ and the expression will not be modified (inspector)
\precC\succ&
              // passing by reference, can only receive <Lexpr> and the expression can be modified (mutator)
              // passing by rvalue reference, can receive any <expr>, and the expression can be modified if it is a <Lexpr>
\prec C \succ \&\&
<C≻*&
              // a reference to pointer, for new and delete operations inside the function, also <Lexpr>
\precC\succ*
              // a pointer, also <Lexpr> but allowing a default value (useful for returning 0)
≺C≻* const // a pointer const, inhibiting reseat (including new and delete).
4.4 Return value
The return value (\prec T \succ) can be:
                // a copy of the calculated object is returned to the caller
[const] ≺C≻& // a reference to an argument or to a static or heap object is returned
[const] \prec C \succ * // a pointer to an argument or to a heap variable defined inside the function is returned
≺C≻* const // a pointer const, inhibiting new and delete of the result
4.5 Invocation
[\prec Lexpr \succ =] \prec name \succ (\prec LE \succ); // invocation of a function with a list of expressions
The list of sent N expressions, \langle LE \rangle, can be written as:
\precexpr_0\succ, \precexpr_1\succ,..., \precexpr_N-1\succ
When N==0 then parenthesis cannot be omitted in the invocation.
Each of the expressions must be \prec \text{Lexpr}\succ if the corresponding type in \prec \text{LP}\succ is either \prec \text{C}\succ \& or \prec \text{C}\succ \& (see previous sub-section).
When default values are defined, then the corresponding argument can be omitted in the invocation.
4.6 Function pointers and references
Using pointers:
\prec T \succ (* \prec name \succ) (\prec LP \succ);
                                                    // declares \precname\succ as a pointer to a function with \precT\succ as return type
                                                    // and \prec LP \succ as list of received arguments
                                                    // assigns an existing function (the & is not required)
~name>=&<name_of_function>;
[\prec Lexpr \succ =] (* \prec name \succ) (\prec LE \succ);
                                                    // invokes the function
Using references:
\prec T \succ (\& \prec name \succ) (\prec LP \succ) = \prec name\_of\_function \succ; // function reference
[\prec Lexpr \succ =] (\prec name \succ) (\prec LE \succ);
                                                   // invokes the function
  ^{5}One can declare a function inline to improve performance, but this is better left to the compiler
```

```
typedef \langle T \rangle (*\langle name \rangle)(\langle LP \rangle);
                                                     // this allows the direct use of name as a function name:
                                                     // declares \prec var \succ as a pointer to function
\precname\succ \precvar\succ;
4.7 Template function
Templatization provides type parametrization (i.e. in practice types will also be arguments)
Declaration and definition formats:
template <typename Typ1[=\precT1\succ], typename Typ2[=\precT2\succ],...> // note that Typ1 and Typ2 are type parameters
\prec T \succ \prec name \succ (\prec LR \succ)
\{...// Typ1 and Typ2 can now be used as types }
invocation: \langle name \rangle \langle \langle T1 \rangle, \langle T2 \rangle, ... \rangle (\langle LE \rangle); where \langle T1 \rangle, \langle T2 \rangle, ... are types.
The compiler is often able to determine the types, so that invocation may also be:
\precname\succ(\precLE\succ);
    Statements
    Condition (if)
if(≺boolean_expr≻)≺statement≻;// the shortest version of the "if" condition
if(\prec boolean\_expr\succ)
                                    // with a else branch
\precstatement_0\succ;
                                     // executes an statement if the ≺boolean_expr≻ is true
                                     // executes another statement if not
else
\precstatement_1\succ;
if (\prec boolean_expr_0 \succ)
                                     // if..else if..else version
\precstatement_0\succ;
else if(\precboolean_expr_1\succ)
\precstatement_1\succ;
else if(\precboolean_expr_2\succ)
\precstatement_2\succ;
else
\precstatement_N\succ;
5.2 Selection (switch)
switch(≺integer_expression≻)
{
case ≺const_integer_0≻:
\precstatement_0\succ;
break;
case \precconst_integer_1\succ:
\precstatement_1\succ;
break:
default: // this branch will be executed when none of the above are satisfied
≺default_statement>;
5.3 Loop (while)
while(\prec boolean\_expression\succ)
\precstatement\succ;
     Loop (do while)
do
\precstatement\succ;
while(≺boolean_expression≻)
5.5 Loop (for)
for([≺initialization≻];[≺continuation_condition≻];[≺incrementation≻])
\precstatement\succ
Note that ≺initialization≻ and ≺incrementation≻ can have multiple statements separated by commas.
Examples:
for(int i=0;i!=n;++i)
                                 // classical loop (note the prefix increment)
{...}
int i=1;
                                 // with separated initialization
for(;(i<n)&&(j>=0);++i,--j)
{...}
for(;;)
                                 // infinite loop with break condition
{...
if(i<n)break;}</pre>
5.6 Loop exit and skipping
                  // breaks out of the loop
break;
continue;
                  // skips the remaining part of the loop
goto <label≻; // goes to a line marked with a label
≺label≻:
                  // this marks a line with a label
```

With type synonyms:

#### General considerations

```
Struct (an extension of C-language structs) are user-defined types (user-defined \langle C \rangle).
Each struct definition should be valid for both const and non-const objects.
Struct may contain type synonyms, enumerations, variables and member functions.
Struct and class are similar: however, they have distinct default access for members and inheritance (public in struct and private in class).
Often, it is necessary to use pointers or references to a struct before it is defined. A forward declaration can then be used:
struct <name_of_used_struct>; // forward declaration of a struct
struct ≺name_struct≻
\{\dots // \text{ use of pointers or references}^6 \text{ of objects of type } \prec \text{name\_of\_used\_struct} > \}
6.2 Definition of a structure
struct \prec name\_struct \succ [:[virtual] \prec name\_ancestor\_struct\_1 \succ],...,[virtual] \prec name\_ancestor\_struct\_m \succ]
// struct name and ancestor inheritance (virtual keyword indicates merging of the ancestors)
\{//\longleftarrow beginning of the struct member declaration and definition
/* one ancestor type */
typedef <name_ancestor_struct_1> super;
/* friend structs and functions will have full access to the contents of the structure (they are not inherited) */
friend struct ≺name_of_another_struct≻;
                                                                                   // friend struct declaration
                                                                                   // friend function declaration
friend \langle T \rangle \prec name\_of\_a\_function \rangle (\langle LP \rangle);
/* typedef is a new type synonym defined by the class and can be accessed as ≺name_struct≻::≺type_name≻ */
typedef \prec T \succ \prec type\_name \succ;
/* enumerations can also be defined in the class, such as \prec E \succ, and can be accessed as \prec name_struct \succ :: \prec E \succ */
enum \prec E \succ
{...};
/* now the member variable and functions subsequent to access specifiers */
[public: | protected: | private:]
                                                                                   // access specifier
[mutable] \prec T \succ \prec var \succ;
                                                                                   // object variable
[virtual] \prec T \succ \prec name\_of\_function \succ (\prec LR \succ) [const] [=0][{...}|;]
                                                                                   // member function
[virtual] \prec T \succ operator @ (\prec LR \succ) [const] [=0] [{...}];]
                                                                                   // member operator
using ≺name_ancestor_struct≻::≺name_of_function≻;
                                                                                   // access to ancestor member
/* classical constructor and destructor */
[explicit] \precname_struct\succ(\precLR\succ)[:\precLI\succ]{...}
                                                                                   // constructor, \prec LI \succ is the list of direct initialization
[virtual] ~\langle name_struct \rangle [=0]{\ldots\}
                                                                                   // destructor - should be virtual with polymorphic structs
/* static variables and functions */
static \prec T \succ \prec var\_static \succ;
                                                                                   // static variable (prefer a
                                                                                   // static function returning a reference)
                                                                                   // static function
static \langle T \rangle \prec name\_static \rangle (\langle LR \rangle) \{...\}
\};// \leftarrow end of the struct member declaration and definition
\precT\succ \precname_struct::var_static\succ=\precexpr\succ;
                                                                                   // compulsory static variable initialization
Note that the Lazyc++ tool simplifies the process of struct creation and allows the initiation of static variables in the struct definition.
In the struct, we have:
≺LR≻ as the list of received variables
\prec LP \succ as the list of received types and
≺LI≻ as the direct list of initialization with the general form:
\precI1\succ,\precI2\succ,\precI3\succ,...
Where each item \prec IJ \succ in the list of initialization is either
\precvarJ\succ(\precLE_J\succ) or \precname_ancestor_struct\succ(\precLE\succ) where the constructor of the ascending structure is used.
Note that the compiler interprets the initialization list from the last to the first item.
In \forall \text{varJ} \succ (\forall \text{LE}_{\text{J}} \succ), \forall \text{varJ} \succ must be an object variable and \forall \text{LE}_{\text{J}} \succ is the list of arguments of a corresponding constructor.
Access specifiers:
                  Access is granted to member functions and variables under this scope
public:
                  Access is only granted to functions and variables from structs publicly derived <sup>7</sup>from the present one
protected:
private:
                  Access is only granted to functions and variables of objects of the same struct
In addition:
The virtual specifier indicates that the function is dynamically determined in a struct hierarchy.
In addition, when the =0 suffix is used (it requires the virtual keyword), the function must be redefined by the descendants when needed.
The virtual specifier only needs to be used in the base struct.
The post const keyword:
Post const means that the function cannot alter non-mutable variables. Objects declared with the const keyword will only be able to invoke the
const version of member functions. The mutable keyword indicates that the member variable is allowed to be changed by a const function.
Mutator functions can alter the values of a given object and inspector functions cannot. The latter should have the const keyword.
```

### Explicit keyword in the constructor:

The explicit keyword removes automatic type conversion effected by constructors that accept a single-argument, therefore eliminating possible cast ambiguities

### Not inherited:

Constructors, destructors and the copy operator overload are not inherited by derived structs. Friends are also not inherited.

### Automatically invoked by inherited objects:

Default constructor and destructor

## Virtual inheritance:

When the virtual keyword is used before the ancestor name, the following structs inheriting from more than one of the 'virtual-inheritance' structs will have a common base struct. The common base struct is called base node, the two or more 'virtual-inheritance' structs are called derived and the most derived class is the join struct.

<sup>&</sup>lt;sup>6</sup>Variables are not allowed

#### 6.3 Access to functions $\prec$ name\_ancestor\_struct $\succ$ :: $\prec$ name\_of\_function $\succ$ ( $\prec$ LE $\succ$ ); // access to ancestor function Access to virtual member functions: ≺name\_struct> ≺aname\_struct>; // declaration of a object of a struct // (does not allow use of polymorphism) $\prec$ name\_ancestor\_struct $\succ$ \* $\prec$ pname\_struct $\succ$ =& $\prec$ aname\_struct $\succ$ ; $\forall$ pname\_struct $\succ$ -> $\prec$ name\_of\_function $\succ$ ( $\prec$ LE $\succ$ ); // invokes name\_struct virtual function or ≺name\_ancestor\_struct≻ \* ≺pname\_struct≻=new ≺name\_struct≻; // the same effect with heap allocation 6.4 "this" const pointer this-> $\prec$ name\_of\_a\_function $\succ$ ( $\prec$ LE $\succ$ ); // invokes a function on the given object // accesses a variable of a given object this->≺var≻; // returns the object return \*this; 6.5 Use of a structure $\prec$ name\_struct $\succ$ $\prec$ aname\_struct $\succ$ ( $\prec$ LE $\succ$ ); // object declaration ≺name\_struct\_base>\* </ // use of a static variable $\prec$ name\_struct $\succ$ :: $\prec$ var\_static $\succ$ =...; $\prec$ name\_struct $\succ$ :: $\prec$ name\_static $\succ$ ( $\prec$ LE $\succ$ ); // invocation of a static function $\prec$ name\_struct $\succ$ :: $\prec$ type\_name $\succ$ $\prec$ aTinname\_struct $\succ$ ; // use of a struct type synonym in a declaration of a variable; // use of a struct enumeration in a declaration of a variable; $\prec$ name\_struct $\succ$ :: $\prec$ E $\succ$ $\prec$ aEinname\_struct $\succ$ ; $\prec$ aname\_struct $\succ$ . $\prec$ name\_of\_function $\succ$ ( $\prec$ LE $\succ$ ); // invocation of a function by an object $\prec$ aname\_struct $\succ$ . $\prec$ name\_static $\succ$ ( $\prec$ LE $\succ$ ); // an alternative invocation of a static function $\prec$ aname\_struct $\succ$ .operator@( $\prec$ LE $\succ$ ); // a form of invocation of a operator by an object $\prec$ pname\_struct $\succ$ -> $\prec$ name\_of\_function $\succ$ ( $\prec$ LE $\succ$ ); // invocation of a function by a pointer, same as // (\* $\prec$ pname\_struct $\succ$ ). $\prec$ name\_of\_function $\succ$ ( $\prec$ LE $\succ$ ); // use of a structure typedef (with disambiguation) typename ≺name\_struct≻::≺type\_name≻ ≺atypeinname\_struct≻; Other typical functions and operators - friends defined in struct scope #include≺iostream≻ // includes the input/output library using namespace std; // access the standard namespace struct... (see above) { // insertion operator (defined only in the base friend ostream& operator<<(ostream& $\prec$ out $\succ$ ,const $\prec$ name\_struct $\succ$ \*|& $\prec$ rhs $\succ$ ); // should invoke a virtual function // for each derived struct, similar for binary) friend istream& operator>>(istream& $\prec$ in $\succ$ , $\prec$ name\_struct $\succ$ \*|& $\prec$ rhs $\succ$ ); // extraction operator (defined only in the base // should invoke a virtual function // for each derived struct, similar for binary) bool operator==(const $\prec$ name\_struct $\succ$ & $\prec$ other $\succ$ ) const {...}; // equality comparison bool operator!=(const $\prec$ name\_struct $\succ$ & $\prec$ other $\succ$ ) const {...}; // inequality comparison bool operator<(const $\prec$ name\_struct>& $\prec$ other>) const {...}; // less-than comparison [virtual] $\prec T \succ \& operator[](size_t index){...}$ // access operator, non-constant [virtual] const ≺T≻& operator[](size\_t index) const {...} // access operator, constant [virtual] $\prec T \succ \& operator()( \prec LR \succ) \{...\}$ // invocation operator // invocation operator (constant version) [virtual] const $\prec T \succ \& operator()( \prec LR \succ) const {...}$ [virtual] ≺name\_struct≻\* to≺name\_struct≻{return(this);} // avoid dynamic\_cast, requires the same // function in the base class $\prec$ name\_struct $\succ$ (){...} // default constructor explicit $\prec$ name\_struct $\succ$ ( $\prec$ LR $\succ$ ): $\prec$ LI $\succ$ {...} // constructor with initialization list operator $\prec T \succ ()$ const{} // type cast (conversion) operator // will explicitly be invoked when a type cast is used $\prec$ name\_struct $\succ$ (const $\prec$ name\_struct $\succ$ & $\prec$ other $\succ$ ){...} // copy constructor $\prec$ name\_struct $\succ$ ( $\prec$ name\_struct $\succ$ && $\prec$ other $\succ$ ){...} // move constructor, will subtract the resources of other [virtual] void swap( $\prec$ name\_struct $\succ$ &){...} // swap function to be used in copy constructor [virtual] $\prec$ name\_struct $\succ$ \* clone() const[=0]{...} // virtual clone (uses copy constructor) [virtual] <name\_struct>\* create() const[=0]{...} // virtual create (uses default constructor) ≺name\_struct≻& operator=(const ≺name\_struct≻& ≺other≻) // assignment operator (don't forget self-assignment $\{if(this!=\& \prec other \succ) \{... \prec name\_ancestor\_struct \succ :: operator = ( \prec other \succ);... \}$ // and ancestor assignment) // the assignment operator is not inherited return \*this;} $\prec$ name\_struct $\succ$ & operator++(){...} // prefix increment (usually faster) // suffix increment (a dummy int argument is used) $\prec$ name\_struct $\succ$ operator++(int trash){...} ≺name\_object≻& operator\*(){...} // dereference (\*) operator const <\name\_object>& operator\*() const {...} // dereference (\*) operator, constant ≺name\_object≻\* operator->(){...} // dereference (->) operator const <name\_object>\* operator->() const {...} // dereference (->) operator, constant friend const <name\_struct≻ operator+ // addition operator as friend function // other binary operators can also be friend functions $(const \prec name\_struct \succ \& first, const \prec name\_struct \succ \& second) \{...\};$ }; Synthesized member functions - present if not defined by the programmer • Default constructor (empty) • Copy constructor (copies all member variables)

- Destructor (by default non-virtual if the base class does not possess a virtual destructor)
- Assignment operator (copies all member variables)
- Dereference operators

# 6.8 Template structs and specialization

Basic declaration and usage:

```
// Typ1 and Typ2 can now be used as types}
struct ≺name_struct≻{
Specialization. When a specific instantiation of a template for type \prec T \succ is needed:
template <typename Typ>
struct \prec name\_struct \succ {...};
                                                 // general template
template <>
struct \prec name\_struct \succ \prec T \succ {...};
                                               // specialized template
template <typename Typ>
struct \precname_struct\succ <Typ*> \{...\};
                                           // specialized template for pointer (partial specialization)
     Pointer to struct members
We can use a variable to select which member function to invoke. The syntax is:
typedef \langle T \rangle  \langle name\_struct \rangle :: * \langle synonym\_pointer \rangle;
                                                                             // defines ≺synonym_pointer≻ as a pointer to a member variable
 \mathsf{typedef} \prec \mathsf{T} \succ (\prec \mathsf{name\_struct} \succ : : * \prec \mathsf{synonym\_pointer} \succ) (\prec \mathsf{LP} \succ) [\mathsf{const}]; \ \ // \ \ \mathsf{defines} \ \ \prec \mathsf{synonym\_pointer} \succ \ \mathsf{as} \ \mathsf{a} \ \mathsf{pointer} \ \mathsf{to} \ \mathsf{a} \ \mathsf{member} \ \mathsf{variable} \ \mathsf{or} \ \mathsf{function} 

<synonym_pointer

<apointer

;</pre>
                                                                             // declares ≺apointer≻ as a object pointer to a member function
≺apointer>=& <name_struct>:: <name_member_function>;
                                                                             // assigns a specific member functions to ≺apointer≻
[\prec Lexpr \succ =] (\prec name \succ . * \prec apointer \succ) (\prec LE \succ);
                                                                             // with ≺name≻ being an
                                                                             // object of type ≺name_struct≻, invokes a member function
[\prec Lexpr \succ =] (\prec name \succ -> * \prec apointer \succ) (\prec LE \succ);
                                                                             // with ≺name≻ being a pointer to an
                                                                             // object of type ≺name_struct≻, invokes a member function
Access macro (to avoid the previous syntax): #define ACCESS_MEMBER(name, apointer) ((name).*(apointer))
Usage: [ \prec Lexpr \succ = ] ACCESS_MEMBER( \prec name \succ, \prec apointer \succ) ( \prec LE \succ)
    User-defined namespaces
To define a scope spanning one file or more, the user can create a named namespace:
namespace \precnamespace_name\succ{ \preccontents\succ } // \preccontents\succ is composed of typedefs, variables, structs, etc
The scope operator :: is used to access the contents of a namespace. For example, if a struct point was declared in namespace custom,
a variable can be declared as:
custom::point var;
In alternative, we can use the keyword using:
using namespace custom;
point var;
Specific ≺contents≻ use can also be prescribed:
using custom::point;
point var;
Global variables and functions can be used with the global scope operator ::
::var; // in alternative, they can be inserted in a unnamed namespace. Unnamed namespaces force translation unit scope.
Namespace aliases can be defined as follows:
namespace \prec namespace\_alias \succ = \prec existing\_namespace \succ;
8 Exceptions
{ // code that may throw exceptions or invocation of such code
throw ≺expr≻ }
\operatorname{catch}(\prec T \vdash \prec \operatorname{var} \vdash) // catches the type returned by a previous \prec \operatorname{expr} \vdash
{ // code that will invoked according to the value of \langle var \rangle}
catch(...) // catches any throw
{ // code that will invoked to handle any exception }
    Uniform initialization and initializer lists
The C++11 norm allows the use of a homogeneous container for inserting a braced-delimited list of values.
The following header inclusion is required: #include <initializer_list>
To use a initializer_list in a declaration (≺name_struct≻ must declare a initializer_list constructor):
\precname_struct\succ \precaname_struct\succ {\precLE\succ};
Common use of initializer_list is:

    Constructors

   • Functions with an arbitrary number of arguments
   • Conversion between containers
Example:
vector<string> vecstr {"First name", "Second name", "Third name", "Fourth name"};
10
     Tools
Tool to generate header and source from a .lzz file (Lazyc++): http://www.lazycplusplus.com/
Tool to improve the source-code readability: http://astyle.sourceforge.net/
Tool to create makefiles: http://www.cmake.org/
     More information, acronyms and patterns
C++FAQ: http://www.parashift.com/c++-faq/index.html
Curiously Recurring Template Pattern (CRTP)
Resource Acquisition Is Initialization (RAII)
```

template <typename Typ1,typename Typ2,...> // note that Typ1 and Typ2 are type parameters