

Hintikka’s World: Scalable Higher-order Knowledge

Tristan Charrier¹, Sébastien Gamblin², Alexandre Niveau² and François Schwarzentruher^{3*}

¹Univ Rennes, CNRS, IRISA, France

²Université de Caen Normandie, GREYC, Caen, France

³Univ Rennes, CNRS, IRISA, France

tristan.charrier@irisa.fr, sebastien.gamblin@unicaen.fr, alexandre.niveau@unicaen.fr,
francois.schwarzentruher@ens-rennes.fr

Abstract

Hintikka’s World is a graphical and pedagogical tool that shows how artificial agents can reason about higher-order knowledge (agent a knows that agent b knows that...). In this demonstration paper, we present the implementation of symbolic models in Hintikka’s World. They enable the tool to scale, by helping it to face the state explosion, which makes it possible to provide examples featuring real card games, such as Hanabi.

1 Introduction

Constructing intelligent programs that play games with imperfect information is challenging — two notorious examples that have recently stirred some interest being Hanabi [Bard *et al.*, 2019] and Starcraft 2 [Hu *et al.*, 2018]. As far as we know, an important ingredient that misses from most approaches is the ability to reason about higher-order knowledge (an agent knows that another agent knows that...). In these systems, epistemic logic and its dynamic extension, *dynamic epistemic logic* (DEL) [Baltag *et al.*, 1998], [van Ditmarsch *et al.*, 2008] may offer formal tools for providing explanations in such AI programs.

The only pedagogical tool explaining these models that we are aware of is *Hintikka’s world*, which was presented at ECAI-IJCAI 2018 [Schwarzentruher, 2018]. *Hintikka’s world* is a proof of concept of a graphical user interface that represents Kripke models by comic strips, as shown in Figure 1. It enables the user to explore the mental states of agents. The tool is available at the following address: <http://hintikkasworld.irisa.fr/> and the source code is available here: <https://gitlab.inria.fr/fschwarz/hintikkasworld>

Kripke models are graphs, represented explicitly in memory in the first version of the tool. Explicit models are useful to learn how dynamic epistemic logic works, by means of toy examples: “Muddy Children”, “Sally and Anne” [Wimmer and Perner, 1983], etc. However, the first version of Hintikka’s world does not *scale*. For instance, in real card games, such as Hanabi, there is an exponential number of possible configurations of cards. In its standard version, Hanabi has 50 cards total, each player’s hand contains 4 cards, and the

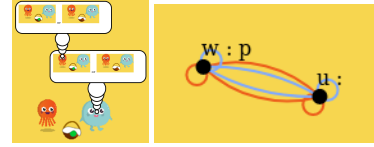


Figure 1: Graphical user interface of *Hintikka’s world*.

order of the cards is important. Therefore, with 4 players, the initial Kripke model features $50 \times 49 \times 48 \times \dots \times 35$ configurations, that is 1.03×10^{26} . Thus, it is impossible to explicitly represent the full graph in memory.

That is why, in this demonstration, we propose to represent Kripke models *symbolically* by using the approach of Charrier and Schwarzentruher [2017; 2018]. The implementation relies on Binary Decision Diagrams (BDDs) [Bryant, 1986]. There is another implementation of symbolic epistemic models, called DEMO [van Benthem *et al.*, 2015], but their implementation is difficult to use in a web application and has no graphical user interface.

2 Demonstration Outline

In the demonstration we run through a game of (a variant of) Hanabi. In Hanabi, each agent has cards with a color and a number, but cannot see his/her own hand. At each turn, in *Hintikka’s World*, the user can play the role of one of the agents: he/she can either give the information to some other agent about a number or a color, or play a card. The goal is to play the cards in increasing order for each color. During the process, the system keeps track on the knowledge of the agents. More precisely, the system shows the real world (the real distribution of the cards). When the user clicks on an agent, the system displays a *sampling* of some possible worlds for that agent (i.e., some distributions of cards he/she still considers as possible at this stage of the game). The agents also reason about knowledge of other agents, as shown in Figure 2 (two levels of knowledge are shown).

Note that in this demonstration, in order to explain models of DEL, the tool still presents examples that rely on explicit models, such as “Sally and Anne”, “Muddy Children”, “Consecutive Numbers”, etc.

*Contact Author

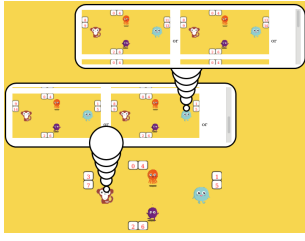


Figure 2: Screenshot of Hanabi in *Hintikka's World*.

3 Symbolic Models

In our tool, we definitely emphasize the use of model checking over theorem proving, as advocated by Halpern and Vardi [1991]. More precisely, we use the same ideas as in symbolic model checking, as defined for temporal logics [Burch *et al.*, 1990], adapted to DEL, as explained in papers by a subset of the contributors of this demo [Charrier and Schwarzenruber, 2017; Charrier and Schwarzenruber, 2018]. Our model checking procedure relies now on symbolic Kripke models, aimed at representing succinctly so-called Kripke models. A Kripke model is a graph whose nodes are *possible worlds*, edges are labeled by agents and an edge $w \rightarrow^a u$ means that agent a considers world u as possible in world w . Each world w is equipped with a valuation telling the true atomic propositions in w . The tool shows that graph on the right of the screen (in the example of Figure 1, the Kripke model has two possible worlds, w and u ; p is true in w but not in u ; \rightarrow_a is given in red and \rightarrow_b in blue).

A symbolic model gives a Boolean formula $\chi(\vec{x})$ that succinctly describes the set of possible worlds: a world is a valuation over Boolean variables \vec{x} satisfying $\chi(\vec{x})$. It also gives, for each agent a , a Boolean formula $\pi_a(\vec{x}, \vec{x}')$ that tells whether there is an edge labeled by agent a from a world described by a valuation over \vec{x} and a world described by a valuation over \vec{x}' . All these Boolean formula are then classically converted in BDDs. Typically, for Hanabi, $\chi(\vec{x})$ tells that \vec{x} describes an initial possible configuration. Formula $\pi_a(\vec{x}, \vec{x}')$ tells that agents other than a have the same cards in \vec{x} and \vec{x}' (this models the fact that agent a only sees his/her own cards).

Dynamic epistemic logic also provides so-called *event models* for describing actions (public announcements, public actions, private announcements/actions, etc.). The reader may refer to the textbook on DEL [van Ditmarsch *et al.*, 2008] and to Charrier and Schwarzenruber [2017] for symbolic event models, that we do not detail here.

4 System Description

Whereas the first version was written in JavaScript, in order to ease the development, the new version is written in TypeScript and relies on the Angular 7 framework.

4.1 Binary Decision Diagrams

As shown by Charrier and Schwarzenruber [2017], the symbolic model checking of DEL is PSPACE-complete, thus is critical. We manipulate sets of worlds as well as relations by means of Binary Decision Diagrams. To this aim, we wrote a

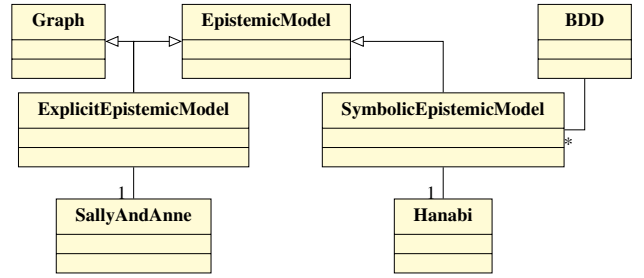


Figure 3: New architecture of *Hintikka's world*.

JavaScript wrapper of the C library CUDD (Colorado University Decision Diagram Package) [Somenzi, 2001]: we wrote a thin wrapper in C, then compiled into Web Assembly using Emscripten, in order to be usable from our JavaScript module.

In order to show possible worlds for a given agent a in some world w , we first construct the BDD of $\pi_a(descr(w), \vec{x}')$ where $descr(w)$ are the Boolean values of \vec{x} corresponding to world w . We then count the number of possible valuations \vec{x}' that make $\pi_a(descr(w), \vec{x}')$ true (BDDs are an efficient representation for counting valuations satisfying a Boolean formula). If the number of such valuations is small, we show all possible worlds, otherwise we randomly generate valuations for \vec{x}' that makes $\pi_a(descr(w), \vec{x}')$ true (we randomly select a branch that leads to the “true” leaf in the BDD of $\pi_a(descr(w), \vec{x}')$).

4.2 Class Architecture

Figure 3 shows the new architecture of *Hintikka's world*. *EpistemicModel* is an abstract class, used by the graphical user interface (GUI), that is independent from the concrete example (“Muddy Children”, “Sally and Anne”, “Hanabi”, etc.) but also, more interestingly, independent from the representation of the epistemic model itself. In particular, an epistemic model can be an *ExplicitEpistemicModel* (a graph) or a *SymbolicEpistemicModel* that relies on BDDs. To obtain a comic strips for a given example, it suffices to implement the method `draw` of a class that inherits from class *World*, that tells how a possible world is drawn.

4.3 Adding New Examples

Providing new examples is easy. Explicit epistemic models are directly described (sets of nodes and of edges). Symbolic epistemic models are described by a Boolean formula χ , and Boolean formulas for π_a . The system provides a way to easily describe how worlds are displayed in the comic strips.

5 Future Work

We plan to extend our implementation to Algebraic decision diagrams (ADD) [Bahar *et al.*, 1997] in order to add examples in robotics that need numerical values, not only Boolean values. We also aim to study counting and sampling specific techniques (see for instance Meel *et al.* [2016]) for generating possible worlds, in order to tackle even bigger examples.

Acknowledgments. We thank Hai Trung Pham for his work during his internship and Arthur Queffelec for his valuable comments on the architecture.

References

- [Bahar *et al.*, 1997] R. Iris Bahar, Erica A. Frohm, Charles M. Gaona, Gary D. Hachtel, Enrico Macii, Abelardo Pardo, and Fabio Somenzi. Algebraic decision diagrams and their applications. *Formal Methods in System Design*, 10(2/3):171–206, 1997.
- [Baltag *et al.*, 1998] Alexandru Baltag, Lawrence S Moss, and Slawomir Solecki. The logic of public announcements, common knowledge, and private suspicions. In *Proceedings of the 7th conference on Theoretical aspects of rationality and knowledge*, pages 43–56. Morgan Kaufmann Publishers Inc., 1998.
- [Bard *et al.*, 2019] Nolan Bard, Jakob N. Foerster, Sarath Chandar, Neil Burch, Marc Lanctot, Francis Song, Emilio Parisotto, Vincent Dumoulin, Subhodeep Moitra, Edward Hughes, Iain Dunning, Shibli Mourad, Hugo Larochelle, Marc G. Bellemare, and Michael Bowling. The Hanabi challenge: A new frontier for AI research. *CoRR*, abs/1902.00506, 2019.
- [Bryant, 1986] Randal E. Bryant. Graph-based algorithms for Boolean function manipulation. *IEEE Trans. Computers*, 35(8):677–691, 1986.
- [Burch *et al.*, 1990] Jerry R. Burch, Edmund M. Clarke, Kenneth L. McMillan, David L. Dill, and L. J. Hwang. Symbolic model checking: 10^{20} states and beyond. In *Proceedings of the Fifth Annual Symposium on Logic in Computer Science (LICS '90)*, Philadelphia, Pennsylvania, USA, June 4-7, 1990, pages 428–439, 1990.
- [Charrier and Schwarzenruber, 2017] Tristan Charrier and François Schwarzenruber. A succinct language for dynamic epistemic logic. In *Proceedings of the 16th Conference on Autonomous Agents and MultiAgent Systems, AAMAS 2017, São Paulo, Brazil, May 8-12, 2017*, pages 123–131, 2017.
- [Charrier and Schwarzenruber, 2018] Tristan Charrier and François Schwarzenruber. Complexity of dynamic epistemic logic with common knowledge. In *Advances in Modal Logic 12, proceedings of the 12th conference on "Advances in Modal Logic," held in Bern, Switzerland, August 27-31, 2018*, pages 103–122, 2018.
- [Halpern and Vardi, 1991] Joseph Y. Halpern and Moshe Y. Vardi. Model checking vs. theorem proving: A manifesto. In *Proceedings of the 2nd International Conference on Principles of Knowledge Representation and Reasoning (KR'91)*. Cambridge, MA, USA, April 22-25, 1991., 1991.
- [Hu *et al.*, 2018] Yue Hu, Juntao Li, Xi Li, Gang Pan, and Mingliang Xu. Knowledge-guided agent-tactic-aware learning for StarCraft micromanagement. In *Proceedings of the Twenty-Seventh International Joint Conference on Artificial Intelligence, IJCAI 2018, July 13-19, 2018, Stockholm, Sweden.*, pages 1471–1477, 2018.
- [Meel *et al.*, 2016] Kuldeep S. Meel, Moshe Y. Vardi, Supratik Chakraborty, Daniel J. Fremont, Sanjit A. Sheth, Dror Fried, Alexander Ivrii, and Sharad Malik. Constrained sampling and counting: Universal hashing meets SAT solving. In *Beyond NP, Papers from the 2016 AAAI Workshop, Phoenix, Arizona, USA, February 12, 2016.*, 2016.
- [Schwarzenruber, 2018] François Schwarzenruber. Hintikka’s world: Agents with higher-order knowledge. In *Proceedings of the Twenty-Seventh International Joint Conference on Artificial Intelligence, IJCAI 2018, July 13-19, 2018, Stockholm, Sweden.*, pages 5859–5861, 2018.
- [Somenzi, 2001] Fabio Somenzi. Efficient manipulation of decision diagrams. *STTT*, 3(2):171–181, 2001.
- [van Benthem *et al.*, 2015] Johan van Benthem, Jan van Eijck, Malvin Gattinger, and Kaile Su. Symbolic model checking for dynamic epistemic logic. In *Logic, Rationality, and Interaction - 5th International Workshop, LORI 2015 Taipei, Taiwan*, pages 366–378, 2015.
- [van Ditmarsch *et al.*, 2008] Hans van Ditmarsch, Wiebe van der Hoek, and Barteld Kooi. *Dynamic Epistemic Logic*. Springer, Dordrecht, 2008.
- [Wimmer and Perner, 1983] Heinz Wimmer and Josef Perner. Beliefs about beliefs: Representation and constraining function of wrong beliefs in young children’s understanding of deception. *Cognition*, 13(1):103–128, 1983.