

Homework on Chap4

- 4.2-3,4.3-4
- Solve the recurrent $T(n) = T(\sqrt{n}) + \theta(n)$
- $A[1, 2, \dots, n]$ is an **increasing-ordered** array with all **distinct** integers (possibly be negative). Find an algorithm by Divide-and Conquer to find i that $A[i]=i$. The worst-case running time of your algorithm should be in $O(\lg n)$.

Homework of Chapter 5

- 5.2-4, 5.3-4
- Problem 5-1

4.2-3

How would you modify Strassen's algorithm to multiply $n \times n$ matrices in which n is not an exact power of 2? Show that the resulting algorithm runs in time $\theta(n^{\lg 7})$

将 $n \times n$ 的矩阵用0填充，使其维数成为大于n的最小的2的整数倍。假设原始矩阵算式为 $A * B = C$

经扩展得到 $A1, B1$, 可以应用Strassen算法，得到 $C1$, 此时我们可以从 $C1$ 中得到正确的答案 C

$$A1 = \begin{bmatrix} A & 0 \\ 0 & 0 \end{bmatrix} \quad B1 = \begin{bmatrix} B & 0 \\ 0 & 0 \end{bmatrix} \quad C1 = \begin{bmatrix} C & 0 \\ 0 & 0 \end{bmatrix}$$

时间复杂度: $\theta(m^{\lg 7})$, m 是最小的大于 n 的2的整数次幂, 故 $m < 2n$

$$\theta(m^{\lg 7}) < \theta((2n)^{\lg 7}) = \theta(2^{\lg 7} n^{\lg 7}) = \theta(n^{\lg 7})$$

4.3-4

Show that by making a different inductive hypothesis, we can overcome the difficulty with the boundary condition $T(1) = 1$ for recurrence (4.19) without adjusting the boundary conditions for the inductive proof.

$$T(n) = 2T(\lfloor n/2 \rfloor) + n$$

猜测: $T(n) \leq n \lg n + n$

代入:

$$T(n) \leq 2(c \lfloor n/2 \rfloor \lg \lfloor n/2 \rfloor + \lfloor n/2 \rfloor) + n \leq 2c(n/2) \lg(n/2) + 2(n/2) + n = cn \lg(n) + (2 - c)n$$

当 $c \geq 1$ 时, 满足 $T(n) \leq cn \lg n + n$

$$T(1) = 1 \leq cn \lg n + n = 0 + 1 = 1.$$

求解递归式 $T(n) = T(\sqrt{n}) + \theta(n)$

改变变量, 令 $m = \log(n)$ 得到 $T(2^m) = T(2^{\frac{m}{2}}) + \theta(2^m)$

重命名 $S(m) = T(2^m)$ 得到 $S(m) = S(m/2) + \theta(2^m)$

应用主定理, $a=1, b=2, f(m)=\theta(2^m)$

$$\log_b a = 0 \quad n^{\log_b a} = 1 \quad f(m) = \Omega(m^{\log_b a + \delta})$$

$$f\left(\frac{m}{2}\right) \leq cf(m), 2^{m/2} \leq c * 2^m, c = \frac{1}{2}, \text{显然, 当 } m \geq 2 \text{ 时成立}$$

由情况3, 递归式的解为 $\theta(2^m)$ 从S转换回T, $T(n) = T(2^m) = S(m) = \theta(2^m) = \theta(n)$

算法如下

与二分搜索原理类似。由于题目中明确了每个数字都是不重复的, 那么当我们发现某个数字大于他的索引, 这表明其右侧的所有数字都会大于其索引, 于是满足条件的数字一定在左侧。小于的情况同理。

$$T(n) = T(n/2) + c$$

最坏情况数字出现在边缘, $O(\lg n)$

```
1 public class HomeWork2_3 {
2     public static void main(String[] args) {
3         // index : 0 1 2 3 4 5 6 7 8 9
4         // value : -1 0 1 2 3 4 6 8 9 10
5         int[] array = {-1, 0, 1, 2, 3, 4, 6, 8, 9, 10};
6         Solution solution = new Solution();
7         solution.find(array, 0, array.length - 1);
8     }
9 }
10 class Solution{
11     public void find(int[] arr, int left, int right){
```

```
12     if(left <= right){
13         int mid = left + (right - left) / 2;
14         if(arr[mid] == mid){
15             System.out.println(mid);
16         }
17         else if(arr[mid] > mid){
18             right = mid - 1;
19             find(arr, left, right);
20         }
21         else{
22             left = mid + 1;
23             find(arr, left, right);
24         }
25     }
26 }
27 }
```

第五章作业老师上课说可以下周一起交。
