

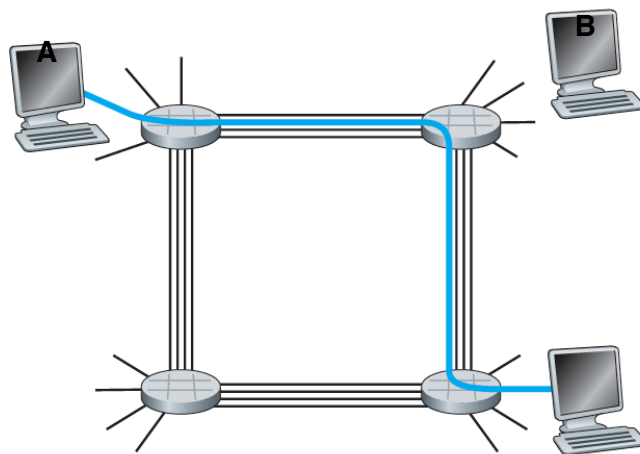
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TA1: Introduction to the Application Layer

Q1: It is advantageous to organize network protocols in layers because it allows for explicit structure and identification of relationships of complex systems. It also allows for easy maintenance and updating of parts of the system. A drawback is that one layer may duplicate lower-layer functionality. A second drawback is that functionality at one layer may need information that is located in another layer which violates the goal of separation.

Q2: Bandwidth is the raw capability of a system to move data through a communication channel at its most optimal whereas throughput is the actual measurement of data moving through a system based on its capabilities. Throughput can never be higher than bandwidth because essentially bandwidth is the highest possible throughput a system can send data. An example of this is a highway. Bandwidth would be the maximum amount of cars that can travel on this highway at any given time but throughput would be how many cars can actually use the highway given its capability. Things like lane closures, accidents, construction, etc. would affect throughput.

Q3: Yes, it is possible that circuit switching can end up hurting its users. Using the example graphic below, let's say computer B represents 4 devices. If Computer A wishes to communicate with those 4 devices at the same time, then 4 communication links would be used. Since there



are only 4 links per circuit switcher, no other connections by any other device or to more devices would be possible. Also, since each link has 4 circuits, dedicated transmission rate is limited to one fourth the total transmission rate.

Q4:

1. Propagation delay = $d/s = (10,000 \times 10^3)/(2.5 \times 10^8) = 0.04 \text{ s}$
2. Transmission delay for 8 packets = $L/R = (2 \text{ Mb} \times 8)/5 \text{ Mbps} = 16 \text{ Mb}/5 \text{ Mbps} = 3.2 \text{ s}$

Q5:

1. Propagation delay = $P(R1) + P(R2) = (5,000 \times 10^3)/(2.5 \times 10^8) + (5,000 \times 10^3)/(2.5 \times 10^8) = 0.04 \text{ s}$.
2. Transmission delay = $T(R1) + T(R2) = (2 \text{ Mb}/4 \text{ Mbps} + 2 \text{ Mb}/6 \text{ Mbps}) \times 8 = 6.666666667 \text{ s}$

Q6: I performed a **traceroute** for the Cambridge University web site cam.ac.uk.

```
mike — -bash — 81x31
Michael's-MacBook-Pro-2:~ mike$ traceroute cam.ac.uk
traceroute to cam.ac.uk (131.111.150.25), 64 hops max, 52 byte packets
 1 172.31.128.1 (172.31.128.1) 3.231 ms 1.829 ms 1.840 ms
 2 10.98.1.34 (10.98.1.34) 2.734 ms 3.308 ms 3.791 ms
 3 gw-conu-lb8.concordia.ca (132.205.236.2) 4.287 ms 3.761 ms 3.750 ms
 4 uconcordia-canarie.risq.net (206.167.253.81) 4.044 ms 3.619 ms 3.870 ms
 5 concordia-canet.dmtrl-uq.risq.net (132.202.30.201) 3.920 ms 3.877 ms 3.497
ms
 6 dmtrl-rq.risq.net (132.202.80.9) 3.983 ms 3.722 ms 4.513 ms
 7 canet-rq.risq.net (132.202.80.10) 3.755 ms 3.682 ms 3.675 ms
 8 mtrl2rtr1.canarie.ca (205.189.32.250) 5.276 ms 4.276 ms 4.139 ms
 9 canarie-bckp.mx1.par.fr.geant.net (62.40.124.225) 85.948 ms 85.673 ms 87.4
33 ms
10 ae1.mx1.lon2.uk.geant.net (62.40.98.76) 92.198 ms 93.396 ms 92.839 ms
11 ae3.mx1.lon.uk.geant.net.geant.net (62.40.98.78) 93.326 ms 93.781 ms 93.82
3 ms
12 janet-gw.mx1.lon.uk.geant.net (62.40.124.198) 93.088 ms 93.099 ms 93.712 m
s
13 ae29.londpg-sbr2.ja.net (146.97.33.2) 119.036 ms 94.756 ms 94.371 ms
14 ae30.londtw-sbr2.ja.net (146.97.33.6) 94.496 ms 95.271 ms 93.513 ms
15 146.97.38.18 (146.97.38.18) 103.774 ms 99.240 ms 102.367 ms
16 146.97.65.117 (146.97.65.117) 96.773 ms 97.391 ms 97.606 ms
17 university-of-cambridge.cambab-rbr1.eastern.ja.net (146.97.130.2) 99.539 ms
114.944 ms 148.758 ms
18 d-dw.s-dw.net.cam.ac.uk (193.60.88.2) 101.071 ms 99.190 ms 99.585 ms
19 d-dw.s-dw.net.cam.ac.uk (193.60.88.2) 98.141 ms 98.255 ms 99.525 ms
20 outside.fw-srv.net.cam.ac.uk (128.232.128.6) 97.685 ms 97.144 ms 97.485 ms
21 link-srv.uis.fw-srv.net.cam.ac.uk (128.232.129.2) 98.824 ms 97.652 ms 97.6
57 ms
22 primary.admin.cam.ac.uk (131.111.150.25) 97.581 ms 101.401 ms 101.027 ms
Michael's-MacBook-Pro-2:~ mike$
```

There does not appear to be anything unusual with the returned results although at line 15 and 16, those IP addresses did not resolve to domain names. In total there were 22 routes in between my machine and the target.

Q7: The Mirai malware targets Linux systems, specifically devices like remote cameras and home routers turning them into bots and thereby creating a botnet. The botnet is then used to attack internet hosts by means of DDoS attacks in which a host is bombarded with packets thereby denying service to actual users and rendering the host useless.

Q8: To retrieve a web page, the browser client first needs to find out the IP of the domain so it does a DNS request over UDP. Once the IP is known, a HTTP GET request is done over TCP to retrieve the relevant web page. So the Application Layer protocols are HTTP and DNS and the Transport Layer protocols are TCP and UDP.

Q9: Alice's email will remain in Alice's mail server message queue until Bob's mail server is back up and able to retrieve it. Alice's mail server will retry to send the email every 30 mins for a few days. If by then Bob's server is not back up, the server removes the message from the queue and notifies Alice.

Q10:

$$Des = \max \{ NF/us, F/dp \}$$

$$Dp2p = \max \{ F/us, F/dp, NF / (us + \sum_{i=1}^N (u_i)) \}$$

file size $F = 15 \text{ Gbits} * 1024 = 15360 \text{ Mbits}$

peers $N = 10, 100, 1000$

server upload rate $us = 30 \text{ Mbps}$

peer download rate $dp = 2 \text{ Mbps}$

peer upload rate $up = 300, 700, 2000 \text{ Kbps}$

Client Server

up / N	10	100	1000
300 Kbps	7680	51200	512000
700 Kbps	7680	51200	512000
2 Mbps	7680	51200	512000

Peer to Peer

up / N	10	100	1000
300 Kbps	7680	25600	46545.45455
700 Kbps	7680	15360	21041.09589
2 Mbps	7680	7680	7680