

# Module (1)

## Course Introduction:

- **Instructor Greeting:**
  - "Hello, everybody. Welcome to this course on algorithm design and analysis, DA111."
  - Instructor hopes students had a great weekend.
- **First Week Overview:**
  - Course logistics will be covered.
  - Introduction to basics of asymptotic analysis.
- **Lecture Goals:**
  - By the end of the lecture/week, students will:
    - Have a complete overview of the syllabus.
    - Meet the instructors and TAs.
    - Know about required textbooks and references.
    - Understand the grading schemes.
    - Gain insights into the job market.
      - Importance highlighted for those aiming for data scientist and machine learning roles.
- **Motivation for Learning Algorithms:**
  - Emphasis on the importance of learning algorithms for future roles in data science and machine learning.

### Week 1 Lecture: Meet the Instructor and TAs

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#### Instructor Introduction:

- **Name:** Chiranjib
- **Background:**
  - Joined IIT Guwahati 10 months ago.
  - Formerly worked for Microsoft and GE Healthcare in the US.

#### Teaching Assistants (TAs):

- **Mehar Khatoon:**
  - M.Tech student in Data Science Department.
- **Avnee Gaur:**

- PhD student.
- Works with Rural Technology Department and some faculties in Mechanical Department.
- **Nikhil Jaiswal:**
  - PhD student in Data Science and AI Department.
- **Jyotishman Bora:**
  - PhD student in Data Science and AI Department.

## **Textbooks and References:**

- **Importance of Textbooks:**
  - Instructor can provide examples and lectures, but personal development of intuition is crucial.
  - Reading and extracting information from books is essential.
- **Primary Textbooks:**
  - **Algorithm Design by Jon Kleinberg and Éva Tardos:**
    - Suitable for newcomers.
    - Written in an accessible manner.
  - **Introduction to Algorithms by Cormen, Leiserson, Rivest, and Stein (CLRS):**
    - Considered the "Bible" of algorithms.
    - Advanced level, more beneficial after completing the course.
- **Additional References:**
  - **Algorithms by Sanjay Dasgupta, Christos Papadimitriou, and Umesh Vazirani:**
  - **Algorithms Illuminated by Tim Roughgarden:**
    - Professor at Columbia University, formerly at Stanford.
    - Videos and PPTs available online and on YouTube.
    - Content is advanced, suitable after completing this course.

## **Course Importance and Objectives:**

- **Core Course Requirement:**
  - Mandatory for all students, including M.Tech (equivalent to MS) and PhD.
- **Relevance to Job Market:**
  - Key for roles in data science and machine learning.
  - Algorithms course is critical for problem-solving and algorithm analysis.
  - Employers assess problem-solving and analysis skills based on this knowledge.

## **Course Materials Access:**

- Books available in the library, some e-copies may be available.
- Reading assignments will guide which books to focus on.

## **Final Note:**

- Take the course seriously for both academic and professional success.
- Algorithm analysis is fundamental for data science roles.

## **Week 1 Lecture: What is an Algorithm**

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### **Definition of an Algorithm:**

- An algorithm is a set of procedures to break a task and solve a well-specified problem.
- In professional settings, your role often involves breaking down complex problems into smaller, solvable tasks using computational techniques.

### **Understanding the Problem:**

- Continuously ask questions to your manager or client to gain a deep understanding of the problem.
- Real-world problems are often described vaguely, requiring clarification and analysis.

### **Communication and Vocabulary:**

- Develop a common vocabulary related to algorithm analysis to communicate effectively with peers and clients.
- Understanding and using this terminology is crucial for collaborative problem-solving.

### **Practical Examples and Strategies:**

#### **1. Family Cooking Analogy:**

- Solving a single-person problem vs. a family or restaurant scenario requires different levels of planning and synchronization.
- Reflects how algorithm complexity and constraints change with the scale of the problem.

#### **2. Traveling Salesman Problem (TSP):**

- Visit multiple cities (e.g., Delhi, Kanpur, Kashi, etc.) with the goal of creating the shortest possible tour visiting each city once.
- The nearest neighbor heuristic might not always yield the optimal solution.
- Real-life application: optimizing travel routes with minimal distance.

#### **3. Sorting Problems:**

- Simple sorting (e.g., red and blue balls) vs. complex sorting (e.g., diamonds of different sizes).
- Different problems require varying levels of expertise and different algorithmic approaches.

### **Why Algorithms Matter in Data Science:**

- Data scientists often need to analyze large datasets to extract meaningful insights and determine appropriate algorithms.
- Even if clients are not aware of technical requirements, data scientists must process and analyze the data to provide solutions.

### **Brute Force Search Method:**

- Used when simpler heuristics like the nearest neighbor fail.
- Involves examining all possible permutations to find the optimal solution.
- Extremely time-consuming and computationally expensive for large datasets.

### **Example: Permutations in TSP:**

- Calculating permutations can lead to a massive number of possible solutions.
- Example: 7 cities  $\rightarrow 7!$  (5040 permutations); 50 cities  $\rightarrow 50!$  (a huge number).

### **Importance of Efficient Algorithms:**

- Solutions must be feasible within a reasonable time frame.
- NP-Hard problems, like TSP, require innovative approaches to find practical solutions.
- Inefficient algorithms (e.g., long payment processing times) render the solution impractical for real-world use.

### **Conclusion:**

- Developing a strong understanding of algorithms and their analysis is crucial for problem-solving in data science and software engineering.
- This course aims to equip you with the necessary skills to tackle complex real-world problems effectively.

## **An Example - Walmart**

In this lecture, the focus is on understanding the complexities involved in solving real-world problems using algorithms, specifically illustrated through the traveling salesman problem (TSP) in the context of a grocery delivery app like Walmart or Flipkart.

### 1. Problem Context:

- The TSP involves finding the optimal route for a delivery person to visit multiple locations without revisiting any location.
- The example given is of a grocery delivery app within IIT Guwahati campus, covering various locations like boys' hostels, faculty quarters, institute buildings, etc.

### 2. Constraints Beyond Distance:

- **Time Slot Constraints:** Deliveries may need to consider operational hours of different buildings, e.g., faculty quarters in the morning, shops after they open, and hostels at specific times.
- **Waiting Time:** The average time a recipient takes to receive a delivery can affect route planning. Historical data and data science can help predict these times.

### 3. Real-life Example:

- In the morning, faculty quarters might be accessible, but institute buildings might not be operational, and shops might not be open. This requires scheduling based on availability rather than just proximity.

### 4. Geodesic vs. Physical Distance:

- **Euclidean Distance:** The shortest straight paths b/w two points.
- **Geodesic Distance:** The shortest path between two points on the Earth's surface, considering the curvature of the Earth.
- **Physical Distance:** The actual travel path, which might be longer due to road layouts or flight paths.
- An example is the route from Delhi to Kanpur, which might include detours through other towns, making the actual travel distance longer than the straight-line distance.

### 5. Advanced Considerations:

- Algorithms need to account for non-linear and complex real-world conditions.
- **Riemannian Manifold:** Used to consider deviations from linear distances, applicable in complex geographical scenarios.

## Conclusion:

- Understanding and solving real-world problems using algorithms requires considering multiple constraints and real-world complexities.
- Factors like time slots, waiting times, and non-linear distances must be integrated into the algorithm design.
- Data science plays a crucial role in collecting and analyzing data to optimize these algorithms.

By applying these principles, you can develop more effective and efficient algorithms for practical applications, ensuring better performance and customer satisfaction.

## **Week 1 Lecture: Why Should You Learn Algorithms?**

### **Key Points:**

#### **1. Importance of Algorithms:**

- Algorithms provide efficient solutions to problems.
- They offer mathematical analysis of time and space efficiency, crucial for resource management in services (e.g., server capacity for millions of customers).

#### **2. Efficiency in Practice:**

- User experience depends on efficiency (e.g., slow loading times on Flipkart can drive users to competitors like Amazon).
- Understanding mathematical principles helps in designing efficient solutions.

#### **3. Foundational Role in Computer Science:**

- Algorithms are fundamental to computer science.
- They underpin different branches of mathematics that form the basis of computer science.

#### **4. Provable Bounds and Complexity:**

- Algorithms have provable execution time bounds.
- Recognize exponential time analysis and understand that simple problems can become complex (e.g., Traveling Salesman Problem).

#### **5. Challenges in Algorithm Design:**

- Simple problems can be hard (like TSP)
- Simple ideas may not always be effective (e.g., nearest neighbor heuristics).
- Simple algorithms (e.g., brute-force methods) can be impractically slow.
- For NP-hard problems, even the best solutions can be slow, necessitating approximate solutions.

#### **6. Course Benefits:**

- Analytical thinking: The course fosters analytical thinking and understanding of algorithmic concepts.
- Communication: Learning algorithmic jargon is essential for effective communication in the software engineering community.
- Industry Relevance: There is a high demand for skills in algorithms and data structures in the booming data science and software engineering industries.

## **7. Industry Demand:**

- Data scientists and software engineers with strong algorithmic skills are in high demand.
- Skills in churning large amounts of data and deriving business decisions are valuable.

## **8. Essential Skills:**

- Mastering algorithms and data structures is crucial for penetrating the industry.
- Proficiency in these areas is necessary for career advancement and meeting industry expectations.

## **Summary:**

Learning algorithms is crucial due to their role in providing efficient solutions and their foundational importance in computer science. This course aims to develop analytical thinking, communication skills, and industry-relevant expertise. Mastering algorithms and data structures is essential for success in the data science and software engineering fields.

## **Week 1 Lecture: Should Data Scientists Need Algorithms?**

### **Key Points:**

#### **1. Role of Data Scientists:**

- Data scientists combine software engineering skills with data analysis.
- They utilize computer science fundamentals to solve complex problems efficiently.

#### **2. Importance of Algorithms:**

- Developing algorithms helps solve problems within limited time and resources.
- Computer science fundamentals are crucial for tasks involving artificial intelligence, machine learning, and natural language processing (NLP).

#### **3. Real-World Applications:**

- Efficient computing is essential for applications like deep learning and large language models (e.g., ChatGPT).
- Timely responses in applications are dependent on efficient algorithms.

#### **4. Course Relevance:**

- This course provides the necessary algorithmic knowledge to be successful in data science and software engineering roles.
- It prepares students for technical interviews, where coding and algorithm skills are assessed.

#### **5. Interview Preparation:**

- Interviews often include coding problems to evaluate algorithm and data structure proficiency.
- Continuous practice is necessary to develop the skills required to compete with peers.

#### **6. Practical Learning Approach:**

- The course combines theory with practical applications.
- Regular practice and application of concepts are essential to retain and effectively use the knowledge gained.

#### **7. Long-Term Benefits:**

- Mastering algorithms and data structures can solve fundamental problems in both professional and personal contexts.
- The skills learned in this course are foundational and applicable beyond the duration of the course.

## **Summary:**

Algorithms are essential for data scientists due to their role in efficient problem-solving, especially in AI, ML, and NLP applications. This course equips students with the necessary skills to succeed in technical interviews and in their careers. Continuous practice and application of these concepts are crucial for proficiency and long-term retention.

## **Week 1 Lecture: The Journey and Syllabus**

### **Key Points:**

#### **1. Significance of Algorithms:**

- Algorithms play a critical role in the efficiency of various systems.
- Examples include:
  - Operating Systems: Efficient algorithms ensure smooth performance (e.g., iOS vs. other operating systems).
  - Cryptography: Fast encryption and decryption are crucial for usability.
  - Networking: Efficient algorithms improve data packet delivery, leading to faster internet speeds.
  - Databases: Algorithms enhance data access speed and efficiency.
  - Machine Learning and Data Science: Fundamental algorithms are vital for model training and data processing.

#### **2. Applications Across Fields:**

- Computational Biology
- Natural Language Processing (NLP)
- Computer Vision



- **Speech Recognition:** Timely responses from voice assistants (e.g., Siri, Google Voice) rely on efficient algorithms.

### 3. Course Structure and Syllabus:

- **Asymptotic Analysis:** Understanding time and space complexity, which is foundational for algorithm analysis.
- **Randomized Algorithms:** Techniques that use random numbers to make decisions within an algorithm.
- **Sorting Algorithms:** Various methods to arrange data in a particular order.
- **Data Structures:** Review of essential data structures, focusing on their use in algorithms.
- **Greedy Algorithms:** Approach that makes locally optimal choices at each step with the hope of finding a global optimum.
- **Dynamic Programming:** Solving problems by breaking them down into simpler subproblems and storing the results of these subproblems to avoid redundant work.
- **Graph Algorithms:** Techniques for solving problems related to graphs, including:
  - Topological Sort
  - Minimum Spanning Trees
- **NP-Hard Problems:** Discussing problems that are notably difficult to solve in a reasonable time.
- **Approximate Algorithms:** Strategies for finding near-optimal solutions to NP-hard problems when exact solutions are impractical.

### 4. Integration of Concepts:

- Importance of integrating data structure knowledge with algorithm skills.
- Practical application of concepts in real-world problems and interview scenarios.
- Focus on foundational algorithms and their analysis to prepare for technical interviews and industry demands.

### 5. Key Focus Areas for Interviews:

- Divide and Conquer
- Dynamic Programming
- Greedy Algorithms
- Graph Algorithms
- Proficiency in these areas is essential for successful interview preparation.

## Summary:

This course emphasizes the fundamental role of algorithms in various applications and prepares students for both academic and professional success. The syllabus

covers essential topics like asymptotic analysis, sorting, data structures, greedy algorithms, dynamic programming, and graph algorithms, all crucial for tackling real-world problems and excelling in technical interviews.

## Week 1 Lecture: Real-Life Problem 1

### Key Points:

#### 1. Understanding the Real-Life Scenario:

- **Problem Statement:** Tracking interactions among students at IIT Guwahati.
- **Objective:** Determine who is interacting with whom, based on their locations.

#### 2. Problem Definition:

- **Location-Based Tracking:**
  - Utilize mobile phone connections to the server to track positions.
  - **Challenges:**
    - Distinguishing interactions in different contexts (e.g., sitting next to someone in a classroom vs. in a marketplace).
    - Lack of direct observation to confirm interactions.

#### 3. Formulating Solutions:

- **Approach:**
  - **Location Data:** Using geolocation coordinates to infer interactions.
  - **Algorithm Development:** Create algorithms to process location data and identify potential interactions.
- **Considerations:**
  - Vicinity-based interactions (proximity of individuals).
  - Contextual differences in various settings (classroom vs. marketplace).

#### 4. Camera-Based Tracking:

- **Scenario Change:** Introducing CCTV cameras for visual data.
- **Objective:**
  - **Recognition:** Identify individuals and observe interactions.
  - **Advantages:** Direct visual confirmation of interactions.
- **AI/ML Integration:**
  - Using computer vision to recognize individuals and track interactions.
  - More sophisticated solution requiring AI/ML skills.

#### 5. Comparing Solutions:

- **Location-Based:**
  - **Data:** Geolocation coordinates.
  - **Simplicity:** Less complex, does not require AI.

- **Type:** Simple database and algorithm problem.
- **Camera-Based:**
  - **Data:** Visual footage from cameras.
  - **Complexity:** Requires AI and ML for recognition and interaction tracking.
  - **Type:** Advanced engineering problem incorporating AI/ML.

#### 6. Engineering Considerations:

- **Time and Space Complexity:** Essential to ensure solutions are efficient and scalable.
- **Feasibility:** Solutions must be deliverable within acceptable time bounds to be practical.

#### 7. Thought Exercise:

- **Task:** Define the best algorithm for location-based tracking.
- **Focus:**
  - Effectiveness of the solution.
  - Time complexity of the algorithm.
- **Objective:** Prepare for detailed discussion and solution formulation.

### Summary:

This lecture introduces a real-life problem of tracking interactions among students at IIT Guwahati using location and camera-based solutions. The key takeaway is to understand the differences in problem formulation and solution approaches based on the type of data (geolocation vs. visual). Emphasis is placed on developing efficient algorithms, considering time and space complexity, and integrating AI/ML for more sophisticated solutions. Students are encouraged to think critically about defining and solving the location-based tracking problem to prepare for future discussions.

## Week 1 Lecture: Real-Life Problem 2

### Key Points:

#### 1. Problem Statement:

- **Objective:** Increase sales in a retail store.
- **Scenario:** Working in a large company like Walmart with multiple physical locations.

#### 2. Strategic Placement of Products:

- **Concept:**
  - **Complementary Products:** Place related products together to encourage combined purchases (e.g., butter and bread).

- **High-Demand Products:** Place essential items (e.g., milk) at the back of the store to encourage customers to walk through other aisles and potentially buy more items.

### 3. Examples of Product Placement:

- **Butter and Bread:** Likely to be bought together, so they should be placed close to each other.
- **Milk:** Placed at the back of the store to ensure customers pass by other products, increasing the likelihood of additional purchases.
- **Party Supplies:** Place items like soda and chips together to cater to customers preparing for parties, leading to higher sales of related products.

### 4. Challenges in Decision Making:

- **Scale:** Walmart may have up to 100,000 different products.
- **Data-Driven Decisions:** Determining optimal product placement requires analyzing large amounts of data on customer buying habits.

### 5. Role of a Data Scientist:

- **Data Analysis:** Utilize customer purchasing data to inform product placement decisions.
- **Big Data Technologies:** Employ tools like Spark and Hadoop to handle and analyze large datasets.
- **Modeling and Extraction:** Develop models to extract meaningful patterns and insights from data to drive strategic decisions.

### 6. Computer Science and Data Science Integration:

- **Traditional Problem:** Rooted in computer science, involving data processing and algorithm development.
- **Data Science Skills:** Essential for handling and interpreting massive datasets to derive actionable insights.

### 7. Discussion and Engagement:

- **Interactive Learning:** Encourage students to think about and propose solutions to the problem.
- **Discussion Forums:** Provide a platform for students to share and discuss their ideas and solutions.

### 8. Future Discussions:

- **Advanced Strategies:** Detailed discussion on the problem will take place when students are more familiar with relevant algorithms and strategies.

## Summary:

This lecture explores a real-life problem of increasing sales in a retail store like Walmart through strategic product placement based on customer buying habits. The focus is on integrating computer science and data science skills to analyze

large datasets and derive optimal placement strategies. The lecture encourages interactive learning and discussion, emphasizing the importance of big data technologies and data modeling in solving such complex problems.

## Week 1 Lecture: Real-Life Problem 3

### Key Points:

#### 1. Problem Statement:

- **Objective:** Designing a phone that caters to everyone's needs.
- **Challenge:** Balancing factors like price, speed, battery life, and features to create an optimal solution.

#### 2. Pareto Front Analysis:

- **Concept:**
  - **Price vs. Delay:** Graphical representation where the x-axis represents price and the y-axis represents delay (speed).
  - **Optimal Solutions:** Higher-priced phones tend to have lower delays, suitable for gaming, while lower-priced phones may have higher delays, unsuitable for gaming.

#### 3. Budget Constraint:

- **Limited Budget:** Individuals have varying budgets, but the mean budget tends to be constant across a large population.
- **Optimal Choices:** Individuals tend to choose phones that offer the best compromise between price and speed within their budget.

#### 4. Multi-Objective Optimization:

- **Complex Decision-Making:** Designing a phone involves optimizing multiple objectives such as gaming performance, battery life, camera quality, and price.
- **Trade-offs:** Balancing different components to ensure that the phone meets the needs of a wide range of users without compromising on essential features.

#### 5. Examples of Optimization:

- **Battery Life vs. Performance:** High gaming performance may drain the battery quickly, leading to frequent recharging, which is undesirable for most users.
- **Display Quality:** While a 4K display may enhance gaming and multimedia experience, it may not be necessary for basic tasks like browsing or calling.

#### 6. Role of Data Science and Algorithms:

- **Optimization Techniques:** Utilize data science and algorithms to find the most efficient solutions that cater to a diverse user base.

- **Economical Design:** Ensure that the components chosen for the phone are cost-effective and provide maximum utility to users.

#### 7. Discussion and Engagement:

- **Interactive Platform:** Encourage students to share their solutions and insights in the discussion section.
- **Collaborative Learning:** Foster a collaborative environment where students can exchange ideas and perspectives on solving real-life problems.

## Week 1 Lecture: Overall Structure, Grades, and Logistics

### Key Points:

#### 1. Course Structure:

- **Reading Materials:** Provided for additional learning and understanding.
- **Lecture Videos:** Covering essential concepts and explanations.
- **Live Videos:** Interactive sessions to engage with the instructor and peers.
- **Practice Quizzes:** Crucial for reinforcing learning and preparing for exams.

#### 2. Importance of Practice:

- **Practice Heavy Course:** Success in the course heavily depends on regular practice.
- **Programming Assignments:** Coding and problem-solving are integral to learning algorithms.
- **Visualization:** Visualize problems and understand complexities to internalize concepts effectively.

#### 3. Assessment Breakdown:

- **Quizzes and Assignments:** 10% of the grade.
- **Programming Assignments:** 30% of the grade.
- **Exams:** 60% of the grade.
- **Benchmarking:** Coding assignments evaluated on efficiency and optimization.

#### 4. Collaboration and Integrity:

- **Encouragement of Collaboration:** Collaboration is encouraged for learning and discussion.
- **Individual Work:** Quizzes, assignments, and exams should be done individually.
- **Academic Integrity:** Cheating will not be tolerated, and severe consequences will be imposed.

#### 5. Learning Ethics:

- **Gratitude for Opportunities:** Recognize the privilege of accessing such educational programs.

- **Avoid Blind Copying:** Utilize resources like the Internet for learning but refrain from blindly copying solutions.
- **Importance of Individual Effort:** Engineering requires the ability to break down and create solutions independently.

#### 6. Engagement and Support:

- **Active Participation:** Engage in live sessions, ask questions, and seek clarification.
- **Communication:** Reach out to the instructor for assistance or clarification on concepts.
- **Continuous Learning:** Understand that mastering algorithms is a gradual process that requires consistent effort and practice.

#### 7. Work Ethic:

- **Hard Work and Smart Work:** Emphasize the importance of both hard work and strategic learning.
- **Seeking Help:** Encouragement to seek assistance when facing challenges, optimizing time and learning efficiency.

### Summary:

The lecture outlines the structure of the course, emphasizing the importance of regular practice, individual effort, and academic integrity. It highlights the breakdown of assessments and encourages collaboration while emphasizing the necessity of independent problem-solving skills. Students are reminded to engage actively, seek assistance when needed, and approach learning algorithms with dedication and perseverance.

## Asymptotic Analysis

### Lecture Overview

#### Key Learning Points:

##### 1. Asymptotic Analysis:

- Understanding the behavior of algorithms as input size grows.

##### 2. Types of Analysis:

- Methods to determine the efficiency of algorithms.

##### 3. Practical Applications:

- Examples of how software engineers optimize time complexity in programs.

### Learning Outcomes:

By the end of this lecture, you will:

### 1. Understand Algorithm Notations:

- Learn and recognize Big O, Big  $\Theta$ , and Big  $\Omega$  notations.

### 2. Comprehend Time and Space Complexity:

- Grasp the concepts of how algorithms utilize time and memory.

### 3. Differentiate Solutions:

- Compare and contrast different algorithms to identify optimal solutions.

### 4. Evaluate Solutions:

- Develop skills to determine the best algorithm for specific problems.

### 5. Self-Evaluation of Complexity:

- Gain the ability to independently calculate time and space complexities.

## Practical Relevance:

- **Real-World Application:**

- Implementing learned techniques to choose and optimize algorithms for practical use.

- **Algorithm Implementation:**

- Translating given steps or pseudocode from resources into functional algorithms.

## Importance of Optimal Solutions:

- **Internet Resources:**

- Recognize that while solutions are readily available online, their optimality must be evaluated.

- **Understanding Jargon:**

- Learn to understand and apply algorithm analysis terminology effectively.

## Asymptotic Analysis

### Definition and Purpose

Set of parameters to compare different algorithms, mainly time and space. Asymptotic analysis is a method of describing the running time and space complexity of an algorithm as the input size grows. It provides a high-level understanding of how an algorithm performs and scales with large input sizes. This is crucial for comparing the efficiency of different algorithms and ensuring they are feasible for real-world applications.

### Real-world Analogy

Consider making food:



- **Single Person:** Making food for one person is straightforward, and the waiting time can be quite long.
- **Restaurant:** When serving a large number of customers, efficiency is critical. Pre-processing (e.g., pre-cut vegetables) helps reduce waiting times. Similarly, in algorithms, pre-processing can improve performance.

## Key Concepts

- **Input Size ( $n$ ):** Number of inputs or the size of the problem.
- **Time Complexity ( $T(n)$ ):** Function representing the time an algorithm takes based on the input size.
- **Space Complexity ( $S(n)$ ):** Function representing the space an algorithm uses (RAM) based on the input size.

## Importance

- **Scalability:** Efficient algorithms handle larger inputs without significant performance degradation.
- **Resource Management:** Ensuring minimal use of computational resources (CPU time, memory).

## Notations

1. **Big O ( $O$ ):** Upper bound of the running time. It describes the worst-case scenario.
2. **Big Omega ( $\Omega$ ):** Lower bound of the running time. It describes the best-case scenario.
3. **Big Theta ( $\Theta$ ):** Tight bound of the running time. It describes both the upper and lower bounds, representing the average case.

## Examples

1. **Big O Notation**
  - Describes the maximum time an algorithm will take.
  - Example: If a restaurant tells you that the food will be ready in at most 20 minutes, this is analogous to  $O(20)$ , meaning the upper bound is 20 minutes.
2. **Big Omega Notation**
  - Describes the minimum time an algorithm will take.
  - Example: If a restaurant tells you that you need to wait at least 20 minutes, this is analogous to  $\Omega(20)$ , meaning the lower bound is 20 minutes.
3. **Big Theta Notation**

- Describes the average time an algorithm will take.
- Example: If a restaurant tells you that you may need to wait around 20 minutes, it can be either more or less, this is analogous to  $\Theta(20)$ , representing the average case.

## Types of Time Complexities

1. **Constant Time ( $O(1)$ )**
  - The running time is constant and does not change with the input size.
  - Example: Accessing an element in a hash table.
2. **Logarithmic Time ( $O(\log n)$ )**
  - The running time increases logarithmically with the input size.
  - Example: Binary search.
3. **Linear Time ( $O(n)$ )**
  - The running time increases linearly with the input size.
  - Example: Iterating through an array.
4. **Linearithmic Time ( $O(n \log n)$ )**
  - The running time increases in proportion to  $n \log n$ .
  - Example: Merge sort.
5. **Quadratic Time ( $O(n^2)$ )**
  - The running time increases quadratically with the input size.
  - Example: Bubble sort.
6. **Cubic Time ( $O(n^3)$ )**
  - The running time increases cubically with the input size.
  - Example: Certain dynamic programming algorithms.
7. **Exponential Time ( $O(2^n)$ )**
  - The running time doubles with each additional input.
  - Example: Solving the traveling salesman problem using brute force.
  - np hard problems use this
8. **Factorial Time ( $O(n!)$ )**
  - The running time grows factorially with the input size.
  - Example: Generating all permutations of a set.
  - np hard problems use this too

## Practical Implications

- **Servers and Large-Scale Systems:** Efficient algorithms ensure that servers, like those of Google or Amazon, handle millions of users without crashing or significant slowdowns.
- **Embedded Systems:** Devices with limited resources, such as Raspberry Pi, require highly optimized algorithms to function correctly.

# Conclusion

Understanding asymptotic analysis and using it to evaluate algorithms ensures that they are efficient, scalable, and suitable for both small and large-scale applications. The ability to determine the time and space complexity is crucial for developing algorithms that meet performance requirements in various contexts.

## Polynomial vs. Exponential Functions in Algorithms

When discussing the performance of algorithms, the time complexity determines how the runtime increases with the size of the input (denoted as  $(n)$ ). Here's a comparison of different time complexities and how they affect the size of  $(n)$  that can be handled in a fixed amount of time.

### Time Complexity Classes

	1 second	1 minute	1 hour	1 day	1 month	1 year	1 century
$\lg n$	$2^{10^6}$	$2^{6 \cdot 10^7}$	$2^{36 \cdot 10^8}$	$2^{864 \cdot 10^8}$	$2^{25920 \cdot 10^8}$	$2^{315360 \cdot 10^8}$	$2^{31556736 \cdot 10^8}$
$\sqrt{n}$	$10^{12}$	$36 \cdot 10^{14}$	$1296 \cdot 10^{16}$	$746496 \cdot 10^{16}$	$6718464 \cdot 10^{18}$	$994519296 \cdot 10^{18}$	$995827586973696 \cdot 10^{16}$
$n$	$10^6$	$6 \cdot 10^7$	$36 \cdot 10^8$	$864 \cdot 10^8$	$2592 \cdot 10^9$	$31536 \cdot 10^9$	$31556736 \cdot 10^8$
$n \lg n$	62746	2801417	133378058	2755147513	71870856404	797633893349	68654697441062
$n^2$	1000	7745	60000	293938	1609968	5615692	56175382
$n^3$	100	391	1532	4420	13736	31593	146677
$2^n$	19	25	31	36	41	44	51
$n!$	9	11	12	13	15	16	17

### Impact on Input Size

As the time complexity increases, the maximum input size  $(n)$  that can be handled within a fixed amount of time decreases dramatically. For instance:

- **Logarithmic:** Handles exponentially large inputs.
- **Linear:** Handles large inputs.
- **Quadratic:** Handles moderate inputs.
- **Cubic:** Handles smaller inputs.
- **Exponential and Factorial:** Handles very small inputs.

### Order of Growth and Practical Implications

When analyzing algorithms, we focus on the order of growth:

# Common order-of-growth classifications

order of growth	name	typical code framework	description	example	$T(2N) / T(N)$
1	constant	<code>a = b + c;</code>	statement	add two numbers	1
$\log N$	logarithmic	<pre>while (N &gt; 1) {   N = N / 2;  ... }</pre>	divide in half	binary search	$\sim 1$
$N$	linear	<pre>for (int i = 0; i &lt; N; i++) {   ... }</pre>	loop	find the maximum	2
$N \log N$	linearithmic	[see mergesort lecture]	divide and conquer	mergesort	$\sim 2$
$N^2$	quadratic	<pre>for (int i = 0; i &lt; N; i++)   for (int j = 0; j &lt; N; j++)   {     ...   }</pre>	double loop	check all pairs	4
$N^3$	cubic	<pre>for (int i = 0; i &lt; N; i++)   for (int j = 0; j &lt; N; j++)     for (int k = 0; k &lt; N; k++)     {       ...     }</pre>	triple loop	check all triples	8
$2^N$	exponential	[see combinatorial search lecture]	exhaustive search	check all subsets	$T(N)$

## Practical Example

Consider a nested loop:

```
for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    // constant time operations
  }
}
```

This has a time complexity of (  $O(n^2)$  ). If there's another segment of code with (  $O(n)$  ), the overall complexity is still (  $O(n^2)$  ), since (  $O(n^2)$  ) dominates (  $O(n)$  ).

## Conclusion

Understanding the time complexity of algorithms helps in selecting the appropriate algorithm for a given problem and predicting its performance for different input sizes. In industry, while complex mathematical proofs of time complexity are rare, a practical understanding of these concepts is essential for writing efficient code.

## 3-SUM Problem Algorithm

## Problem Statement

Given an array of integers, find three numbers such that they add up to zero (or a specified value (  $c$  )). The three numbers must be distinct and should not be at the same index.

## Example Inputs and Outputs

- Input: ([ -1, 0, 1, 2, -1, -4 ])
  - Output: ([ [-1, -1, 2], [-1, 0, 1] ])
- Input: ([ 0, 0, 0 ])
  - Output: ([ [0, 0, 0] ])
- Input: ([ 1, 2, -2, -1 ])
  - Output: ([ ])

## Brute Force Solution

1. **Approach:** Use three nested loops to iterate over all possible triplets and check if their sum is zero.
2. **Time Complexity:** ( $O(n^3)$ )
3. **Pseudocode:**

```
for (int i = 0; i < n - 2; i++) {
    for (int j = i + 1; j < n - 1; j++) {
        for (int k = j + 1; k < n; k++) {
            if (arr[i] + arr[j] + arr[k] == 0) {
                print(arr[i], arr[j], arr[k]);
            }
        }
    }
}
```

## Optimized Solution (Two-pointer technique)

1. **Approach:** First sort the array. Then, use two pointers to find the other two numbers for each element.
2. **Time Complexity:** ( $O(n^2)$ )
3. **Steps:**
  - Sort the array.
  - Fix one element and find the other two using two pointers.
4. **Pseudocode:**

```

vector<vector<int>> threeSum(vector<int>& nums) {
    vector<vector<int>> result;
    sort(nums.begin(), nums.end());
    for (int i = 0; i < nums.size() - 2; i++) {
        if (i == 0 || (i > 0 && nums[i] != nums[i-1])) {
            int lo = i + 1, hi = nums.size() - 1, sum = 0 - nums[i];
            while (lo < hi) {
                if (nums[lo] + nums[hi] == sum) {
                    result.push_back({nums[i], nums[lo], nums[hi]});
                    while (lo < hi && nums[lo] == nums[lo+1]) lo++;
                    while (lo < hi && nums[hi] == nums[hi-1]) hi--;
                    lo++; hi--;
                } else if (nums[lo] + nums[hi] < sum) lo++;
                else hi--;
            }
        }
    }
    return result;
}

```

Certainly, Anubhav! Here's how you can approach the 3-SUM problem with a time complexity of  $O(n^2 \log n)$  using the algorithm you described. The key steps are to sort the array, then for each pair of elements, use binary search to find the third element that completes the triplet summing to zero.

Here's the algorithm in C++:

```

#include <iostream>
#include <vector>
#include <algorithm>

using namespace std;

// Binary search function to find the target in the sorted array
bool binarySearch(const vector<int>& nums, int left, int right, int target) {
    while (left <= right) {
        int mid = left + (right - left) / 2;
        if (nums[mid] == target) {
            return true;
        } else if (nums[mid] < target) {
            left = mid + 1;
        } else {
            right = mid - 1;
        }
    }
    return false;
}

```

```

vector<vector<int>> threeSum(vector<int>& nums) {
    vector<vector<int>> result;
    int n = nums.size();

    // Step 1: Sort the array
    sort(nums.begin(), nums.end());

    // Step 2: For each pair of numbers, use binary search to find the third number
    for (int i = 0; i < n - 2; ++i) {
        for (int j = i + 1; j < n - 1; ++j) {
            int target = -(nums[i] + nums[j]);
            if (binarySearch(nums, j + 1, n - 1, target)) {
                result.push_back({nums[i], nums[j], target});
            }
        }
    }

    // Optional: Remove duplicates from the result
    sort(result.begin(), result.end());
    result.erase(unique(result.begin(), result.end()), result.end());

    return result;
}

int main() {
    vector<int> nums = {-1, 0, 1, 2, -1, -4}; // Example input
    vector<vector<int>> triplets = threeSum(nums);

    // Print the result
    for (const auto& triplet : triplets) {
        cout << "[";
        for (int i = 0; i < triplet.size(); ++i) {
            cout << triplet[i];
            if (i < triplet.size() - 1) {
                cout << ", ";
            }
        }
        cout << "]\n";
    }

    return 0;
}

```

## Explanation:

1. **Sorting the array:** The array is sorted first. Sorting has a time complexity of  $O(n \log n)$ .

2. **Nested loops to find pairs:** Two nested loops iterate over all possible pairs of elements in the array. This part has a time complexity of  $O(n^2)$ .
3. **Binary search for the third element:** For each pair of elements ( $a[i]$ ) and ( $a[j]$ ), the target third element is calculated as  $-(a[i] + a[j])$ . Binary search is then used to find this target in the remaining part of the array. Binary search has a time complexity of  $O(\log n)$ .
4. **Removing duplicates:** Since the array was sorted initially, sorting the results and removing duplicates using `unique` ensures there are no duplicate triplets in the final result.

This approach reduces the time complexity to  $O(n^2 \log n)$  as required.

## Using a dictionary for 3 SUM

d

Certainly, Anubhav! Here's the algorithm for solving the 3-SUM problem using the approach with a dictionary (hash map) for  $O(n^2 \log n)$  complexity:

### Algorithm:

1. **Sort the array:** Sort the array of  $(N)$  distinct numbers.
2. **Dictionary for quick lookup:** For each pair of numbers ( $(a[i], a[j])$ ), use a dictionary to quickly find if the third number that completes the triplet exists.

Here's how you can implement this algorithm in C++:

```
#include <iostream>
#include <vector>
#include <unordered_map>
#include <algorithm>

using namespace std;

vector<vector<int>> threeSum(vector<int>& nums) {
    vector<vector<int>> result;
    int n = nums.size();

    // Step 1: Sort the array
    sort(nums.begin(), nums.end());

    // Step 2: Use a dictionary to find the third number
    for (int i = 0; i < n - 2; ++i) {
        for (int j = i + 1; j < n - 1; ++j) {
            unordered_map<int, int> numMap;
            int target = -(nums[i] + nums[j]);
```



```

        for (int k = j + 1; k < n; ++k) {
            if (numMap.find(target) != numMap.end()) {
                result.push_back({nums[i], nums[j], target});
                break; // Since all elements are distinct, no need to continue for
the same pair
            }
            numMap[nums[k]] = k; // Store the number in the map
        }
    }
}

// Optional: Remove duplicates from the result
sort(result.begin(), result.end());
result.erase(unique(result.begin(), result.end()), result.end());

return result;
}

int main() {
    vector<int> nums = {-1, 0, 1, 2, -1, -4}; // Example input
    vector<vector<int>> triplets = threeSum(nums);

    // Print the result
    for (const auto& triplet : triplets) {
        cout << "[";
        for (int i = 0; i < triplet.size(); ++i) {
            cout << triplet[i];
            if (i < triplet.size() - 1) {
                cout << ", ";
            }
        }
        cout << "]\n";
    }

    return 0;
}

```

## Explanation:

1. **Sorting the array:** The array is sorted first. Sorting has a time complexity of  $O(n \log n)$ .
2. **Nested loops to find pairs:** Two nested loops iterate over all possible pairs of elements in the array. This part has a time complexity of  $O(n^2)$ .
3. **Dictionary for lookup:** For each pair of elements  $((a[i], a[j]))$ , a dictionary (hash map) is used to check if the third number  $-(a[i] + a[j])$  exists in the remaining part of the array. The dictionary search has an average time

complexity of  $O(1)$ ), but since we're iterating over all pairs, this step collectively has a time complexity of  $O(n^2)$ .

4. **Removing duplicates:** Since the array was sorted initially, sorting the results and removing duplicates using `unique` ensures there are no duplicate triplets in the final result.

This algorithm efficiently finds the triplets with the desired sum of zero and has the required time complexity.

## Further Optimization Using Hash Maps

1. **Approach:** Use a hash map to store and look up values quickly.
2. **Time Complexity:** Still  $O(n^2)$ , but with faster lookups.
3. **Steps:**
  - For each pair, use a hash map to check if the complement exists.
4. **Pseudocode:**

```
vector<vector<int>> threeSum(vector<int>& nums) {
    vector<vector<int>> result;
    unordered_map<int, int> num_map;
    sort(nums.begin(), nums.end());
    for (int i = 0; i < nums.size() - 2; i++) {
        if (i > 0 && nums[i] == nums[i-1]) continue;
        for (int j = i + 1; j < nums.size() - 1; j++) {
            if (j > i + 1 && nums[j] == nums[j-1]) continue;
            int complement = -nums[i] - nums[j];
            if (num_map.find(complement) != num_map.end() && num_map[complement] >
j) {
                result.push_back({nums[i], nums[j], complement});
            }
            num_map[nums[j]] = j;
        }
        num_map.clear();
    }
    return result;
}
```

## Summary

- **Brute Force:**  $O(n^3)$  time complexity, simple but inefficient for large arrays.
- **Two-pointer technique:**  $O(n^2)$  time complexity, efficient for larger arrays.
- **Hash Map Optimization:**  $O(n^2)$  time complexity with faster lookups, but slightly more complex to implement.

For most practical scenarios, the two-pointer technique strikes a good balance between simplicity and efficiency.

## Optimizing Integer Multiplication with Karatsuba Algorithm

### Problem Statement

Given two large integers, multiply them efficiently. Traditionally, we use a naive ( $O(n^2)$ ) approach where each digit of one number is multiplied with each digit of the other number. However, there exists a more efficient approach known as the Karatsuba algorithm.

### Karatsuba Algorithm

The Karatsuba algorithm is a divide-and-conquer algorithm that reduces the multiplication of two ( $n$ )-digit numbers to at most ( $O(n^{\log_2 3})$ ) operations, which is approximately ( $O(n^{1.585})$ ). This is achieved by breaking down the multiplication into smaller multiplications and additions.

### Steps of Karatsuba Algorithm

1. **Split each number into two halves:**

- Let  $(x)$  be split into  $(a)$  and  $(b)$
- Let  $(y)$  be split into  $(c)$  and  $(d)$
- Such that  $(x = a \cdot 10^{m/2} + b)$  and  $(y = c \cdot 10^{m/2} + d)$ , where  $(m)$  is the number of digits of the largest number, and  $(m/2)$  is half of that number.

2. **Compute the following products recursively:**

- $(ac)$
- $(bd)$
- $(a + b) \cdot (c + d)$

3. **Use Gauss's trick** to compute the middle term  $(ad + bc)$ :

- Calculate  $(ad + bc = (a + b) \cdot (c + d) - ac - bd)$

4. **Combine the results:**

- The product  $(x \cdot y)$  is given by:  
[  
 $x \cdot y = ac \cdot 10^m + (ad + bc) \cdot 10^{m/2} + bd$   
]

### Implementation

Here is a C++ implementation of the Karatsuba algorithm:

```
def karatsuba(x, y):
    if x < 10 or y < 10:
        return x * y
    else:
        n = max(len(str(x)), len(str(y)))
        half = n // 2
        a = x // (10 ** (half)) # left part of x
        b = x % (10 ** (half)) # right part of x
        c = y // (10 ** (half)) # left part of y
        d = y % (10 ** (half)) # right part of y
        ac = karatsuba(a, c)
        bd = karatsuba(b, d)
        ad_plus_bc = karatsuba(a+b, c+d)-ac-bd
        return ac * (10 ** (2 * half)) + (ad_plus_bc * (10 ** half)) + bd
```

## Explanation

1. **Base Case:** If either string is a single digit, multiply directly.
2. **Recursive Split:** Split the input strings into two halves.
3. **Recursive Multiplications:** Recursively calculate the three products:  $(ac)$ ,  $(bd)$ , and  $((a + b)(c + d))$ .
4. **Combining Results:** Use string addition and subtraction to combine the intermediate results correctly.

This implementation leverages string manipulation to handle arbitrarily large integers, making it suitable for competitive programming and other applications involving large numbers.

## Asymptotic Analysis in Algorithm Comparison

### Importance of Time and Space Complexity

When comparing different algorithms, two critical parameters are time complexity and space complexity. Efficient algorithms ensure that software performs well under varying conditions, avoiding failures due to excessive memory usage or slow execution times.

### Practical Example: Comparing Two Implementations on Different Machines

Let's consider a scenario where two machines, A and B, run different implementations of the same algorithm on increasing input sizes. This example demonstrates the significance of choosing the right algorithm for different input sizes.

## Example Analysis

- **Small Input Sizes:** For small input sizes, Machine A is more efficient.
  - **Input Size 10:**
    - Machine A: 2 seconds
    - Machine B: 1 hour
- **Large Input Sizes:** As the input size increases, Machine B becomes more efficient.
  - **Input Size ( $10^6$ ):**
    - Machine A: 55 hours
    - Machine B: 5 hours

## Trade-Offs and Data Structures

- **Simple vs. Advanced Data Structures:**
  - Simple data structures are easy to use and maintain but may not perform well with large data sets.
  - Advanced data structures, while more complex and harder to maintain, offer better performance for large data sets due to optimized time and space complexity.

## Big Data Era Considerations

- **Scaling Solutions:**
  - In the current era of big data, algorithms must handle vast amounts of data efficiently.
  - Solutions must be scalable, balancing initial complexity with long-term efficiency.

## Summary

Understanding the trade-offs between different algorithms and data structures is crucial. While simpler solutions might be sufficient for small-scale problems, more complex data structures and algorithms are necessary for handling large-scale data efficiently. Always consider the potential growth of data and choose the right tools to ensure your solutions remain effective as they scale.

## Conclusion

In this week's content, we've explored the importance of asymptotic analysis and how it helps in choosing the right algorithms and data structures for different scenarios. As we move forward, we'll delve deeper into these concepts with

practical examples and assignments to solidify your understanding. See you in the live session and next week's content.