

Describe a programming language

Based on the slides of Maurizio Gabbrielli

Description of a language (artificial)

- The description of a language is on three dimensions (Morris, 1938):

- Syntax

Grammar rules
When a sentence is correct
Relationship between symbols

- Semantics


Attribution of Meaning
What does a correct sentence mean
Relationship between signs and meaning

- Pragmatic

How correct and sensible sentences are used
Relationships between signs, meaning and user

A fourth level

- For an executable language
 - Implementation



Execute a correct sentence
respecting its semantics

Tools to describe languages

- Syntax
 - Formal: Generative Grammars (Chomsky), BNF(Backus Naur Form)
 - Quasi-formal: typing, contextual constraints
- Semantics
 - Informal: Natural language, Manual
 - Formal: denotational, operational (axiomatic, algebraic)
- Pragmatic
 - Informal: examples
 - Semi-formal: programming methodologies
- Implementation
 - Derivation from semantics
 - Abstract machine

Syntax

- It tells us how to build the correct language phrases:

var A: integer;

var A: ~~intgr~~;

Error *Lexical*

if a > 0 then C else C

if a > 0 ~~else~~ C then C;

Error *Grammar*

- Same constructs expressed differently in different languages
 - Ex: vector of 10 elements
 - In C: **int V [10]**
 - In Pascal: **V: Array [0.. 9] of Integer**
 - Presumably these two objects have the same meaning (the same semantics).

Syntax

- How to **define and recognize** the correct language constructs?
This is an essential problem when compiling a program
 - Enumerating all well-formed programs? ... a little long
 - Various tools:
 - **Generative grammars**: rules for building language sentences
 - Regular grammars
 - Context-free grammars
 - **Automata**: formalisms to recognize the constructs of a language
 - **BNF**: Backus Naur Form first used for Algol in 1958 by Naur (Chair of ALGOL Committee) and Backus (Secretary)
 - Same thing as context-free grammars

BNF

- The form of Backus and Naur (BNF) is a notation for presenting free grammars, used for the first time in the syntax definition ALGOL 60.
 - $\langle W \rangle$ notation is used to indicate a symbol
 - We use the notation
 $\langle A \rangle ::= \langle B \rangle \mid \langle C \rangle$
To indicate two rules $\langle A \rangle ::= \langle B \rangle$ And $\langle A \rangle ::= \langle C \rangle$
 - In the extended version the symbol $\langle A \rangle^*$ (Kleene star) has the meaning that A is repeated 0 or more times

Examples of derivations

$\langle \text{exp} \rangle ::= \langle \text{num} \rangle \mid - \langle \text{exp} \rangle$

$(\langle \text{exp} \rangle) \mid \langle \text{exp} \rangle + \langle \text{exp} \rangle \mid \langle \text{exp} \rangle * \langle \text{exp} \rangle$

$\langle \text{num} \rangle ::= 0 \mid 1 \mid 2 \mid \dots \mid 9$

$\langle \text{exp} \rangle \Rightarrow \langle \text{exp} \rangle * \langle \text{exp} \rangle$

$\rightarrow \langle \text{exp} \rangle * \langle \text{num} \rangle$

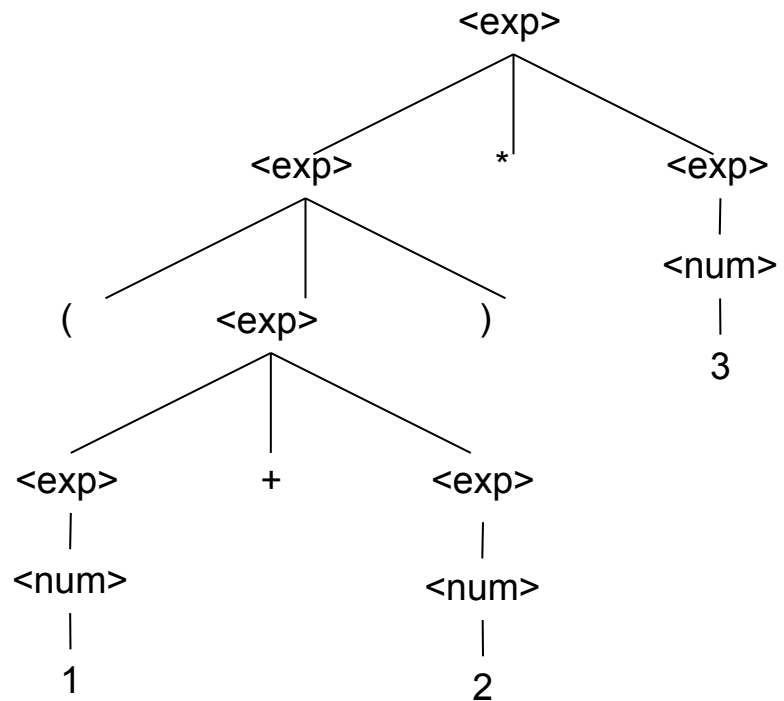
$\rightarrow \langle \text{exp} \rangle * 3$

$\rightarrow (\langle \text{exp} \rangle) * 3$

$\rightarrow (\langle \text{exp} \rangle + \langle \text{exp} \rangle) * 3 \rightarrow \dots$

$\rightarrow (1 + \langle \text{exp} \rangle) * 3 \rightarrow \dots$

$\rightarrow (1 + 2) * 3$



A simple imperative language example

$Num ::= 1 \mid 2 \mid 3 \mid \dots$

$Var ::= X_1 \mid X_2 \mid X_3 \mid \dots$

$AExp ::= Num \mid Var \mid AExp + AExp \mid AExp - AExp$

$BExp ::= \mathbf{tt} \mid \mathbf{ff} \mid AExp == AExp \mid \neg BExp \mid BExp \wedge BExp$

$Com ::= \mathbf{skip} \mid Var := AExp \mid Com; Com \mid$
 $\mathbf{if} BExp \mathbf{then} Com \mathbf{else} Com \mid \mathbf{while} BExp \mathbf{do} Com$

The *Dangling Else*

Grammar are ambiguous sometimes

Example in Java:

```
if (door-open)
  if (none-in-view)
    return "nobody?";
    else ring-bell();
```

which "if" the "else" is paired?

We look at the official definition of Java...

So...

The correct reading is:

```
if (door-open)
    if (None-in-view)
        return "Nobody?";
    else Play-camp ();
```

- "The Java programming language, like C and C++ and many programming languages before them, arbitrarily decided that an else clause belongs to the innermost if to which it might possibly belong"

Semantics

- Need to be able to define exactly "What does" a program
- The formal (semi) definition of semantics is part of a language standard definition of languages since the 80'
- Various ways to define the semantics of a programming language
 - Algebraic
 - Axiomatic
 - Operational
 - Denotational
- There is no universally accepted standard methodology

Operational semantics

- Obtained by defining an Interpreter of the language on a guest machine which components are mathematically described
 - State (describes altogether the system)
 - Transition (specify a computation step that changes the state and/or the program)
 - Computation (sequence of transitions)
- Especially useful because it directly provides an implementation model
- You can define it in a formal and syntax-driven way (SOS)

Example

$$\langle \mathbf{skip}, \sigma \rangle \rightarrow \sigma$$

$$\langle \mathbf{if\ tt\ then\ } c_1 \mathbf{\ else\ } c_2, \sigma \rangle \rightarrow \langle c_1, \sigma \rangle$$

$$\frac{\langle c_1, \sigma \rangle \rightarrow \langle c'_1, \sigma' \rangle}{\langle c_1; c_2, \sigma \rangle \rightarrow \langle c'_1; c_2, \sigma' \rangle}$$