



# Master's thesis

Mikkel Storgaard Knudsen

## FShark

Futhark programming in FSharp

Advisor: Cosmin eller Troels

Handed in: July 26, 2018



## **Abstract**

Here is a nice abstract  
prut prut

# Contents

<b>1</b>	<b>Introduction</b>	<b>3</b>
1.1	What FShark sets out to do . . . . .	4
<b>2</b>	<b>Background</b>	<b>8</b>
2.1	CUDA . . . . .	8
2.1.1	A simple CUDA program . . . . .	8
2.2	Futhark . . . . .	10
2.3	F# . . . . .	11
<b>3</b>	<b>The FShark subset</b>	<b>14</b>
3.1	FShark syntax . . . . .	15
3.2	F# operators available in FShark . . . . .	17
3.3	F# standard library functions available in FShark . . . . .	17
3.3.1	On selection the F# subset to include in FShark . . . . .	17
3.3.2	Missing arithmetic operators in FShark . . . . .	18
<b>4</b>	<b>FSharkPrelude</b>	<b>22</b>
4.1	Arrays in F# versus in Futhark . . . . .	27
4.2	Converting jagged arrays to Futhark's flat arrays, and back again . . . .	28
4.2.1	Analysis of FlattenArray . . . . .	28
4.2.2	Analysis of UnflattenArray . . . . .	29
<b>5</b>	<b>The FShark Compiler and Wrapper</b>	<b>32</b>
5.1	Design choices in writing the FShark Compiler . . . . .	37
<b>6</b>	<b>The Futhark C# backend</b>	<b>40</b>
6.0.1	Entry functions . . . . .	45
<b>7</b>	<b>Benchmarks and evaluation</b>	<b>49</b>
<b>8</b>	<b>Current limitations</b>	<b>52</b>
<b>9</b>	<b>Related work</b>	<b>53</b>
<b>10</b>	<b>Conclusion and future work</b>	<b>54</b>
<b>11</b>	<b>FSharks interoperability between F# and Futhark (C#)</b>	<b>55</b>
11.0.1	Pros and cons of the current design . . . . .	56
11.1	The future of FShark interoperability . . . . .	57



# Chapter 1

## Introduction

Developers worldwide are, and have always been, on the lookout for increased computing performance. Until recently, the increased performance could easily be achieved through advances within raw computing power, as CPU's had steadily been doubling their number of on-chip transistors, in rough accordance to Moore's Law (citér her).

As performance increases in single-CPU design has stalled due to the power wall[?] (among other things), developers are turning to multi-core processors instead. As the number of cores increases, so does the number of active threads available for parallel data processing.

Modern mainstream GPUs can run tens of thousands of threads in parallel. Modern mainstream CPUs, like the current Ryzen series by AMD, usually support between 10 and 20 simultaneous threads. This makes GPUs the optimal target for data-parallel programming.

GPU programming is complicated: GPU-targeting developers must not only write the computational kernels for the GPUs, but also often manually handle the memory allocations and -transfers between the main program and the GPU device. Such difficulties in GPU development, compared to normal (sequential) CPU development, severely hinders the adaption of GPU programming in general.

Even though most programming languages support GPU programming through various libraries, there are very solutions that offers GPU programming through high level programming - the users still have to write their own kernels in some form, and likewise declare their own buffers.

Two mainstream languages which lack high level GPU programming solutions are C# and F#.

It is safe to say that there exists plenty of C# and F# projects in the real world, which could greatly benefit from parallelizing parts of their algorithms, but current solutions would then demand that those parts in particular should be rewritten at least partly as GPU code, depending on the libraries used. Depending on someone to have non-mainstream GPU coding skills on a conventional developer team is not feasible, so the benefits from parallelizing are often left alone in favor of maintaining a more accessible code base.

Currently, there does exist plenty of high level solutions to this problem. In particular, numerous domain specific languages exists that allows programmers to solve their domain specific problems in a high level language, and compile it to standalone GPU accelerated libraries or programs.

Of such DSLs we have for instance:

- Forma
- Ebb
- one more

However, DSLs such as these are always either embedded in some host language (such as Accelerate), or compiled to standalone executable programs. Their compilers are designed to optimize performance, but rarely to support interoperability to any significant degree.

Projects like APLtail[?] have shown new ways to obtain GPU-accelerated executables and libraries from source code written in high level language. APLtail parses and compiles APL<sup>1</sup> code into redistributable C or Python libraries.

In summary, the hardware for massively parallel programming is widely available. Furthermore, solutions exists for writing efficient GPU programs in high level languages, but these have weak interoperability support with mainstream languages.

## 1.1 What FShark sets out to do

This thesis takes inspiration from APLtail[?], and creates a solution that lets users compile efficient GPU programs from a high level programming language, whilst at the same time supporting a high level of interoperability with a mainstream language. Whereas APLtail allows integration of GPU Futhark-written computation kernels in C- and Python programs (by means of code generation), we would like to use code generation to make Futhark-written kernels available for use in C# and F# programs.

To show that this is feasible, we first design a C# code generator for the Futhark compiler. This code generator must be able to generate C# source files, that can be compiled and used either as standalone executables, or as importable libraries in any other C# or F# program. MERE HER

The code generator alone should not convince anyone that we are creating GPU kernels from a high level language, which is why we also design and implement a compiler which takes source code written in a mainstream language, compiles it as efficient GPU kernels, and (together with the code generator) makes the resulting GPU program immediately operable from the mainstream language itself.

Empirical evaluation demonstrates that this approach is feasible. We both show that that unit tests written in a high level language can be compiled and executed correctly as computational kernels on the GPU, just as we also take complex benchmark programs written in a mainstream language, compile them into computational kernels for the GPU, and use them directly in the mainstream language afterwards.

---

<sup>1</sup>more accurately a subset of the APL language

## Motivation

FShark is intended to be a way of writing and utilizing Futhark, without actually having to write or interact with the Futhark language and compiler itself. Besides some tooling and an F# SOAC library, it primarily consists of the FShark compiler that compiles from F# source code to Futhark source code, and the Futhark C# generator, which compiles Futhark programs as either standalone C# programs or -libraries.

As much as most developers are happy to increase performance on big computations, it is not always an option to incorporate an extra language into an already existing programming language. At this moment, using Futhark in either a C#- or F# project is a contrived process that usually requires spawning a subprocess with a `futhark-opencl` C program from inside one of the .NET projects.

In order to use Futhark natively in .NET languages, it is therefore necessary to write a backend for Futhark in a .NET language. For FShark, I have chosen to implement this backend in C#, as the Futhark intermediate code `ImpCode`<sup>2</sup> is trivial to translate into imperative C# statements and expressions. Also, there are C# libraries available which supply OpenCL bindings, which are needed to implement the necessary OpenCL constructs from `ImpCode`.

It is my belief that exporting Futhark programs as .NET executables and -libraries will lower the barrier to Futhark usage in .NET projects significantly, hopefully increasing the all-round number of Futhark users, and in the long term, increasing utilization of GPU programming and making it more widely available.

However, one could do even more than just exporting Futhark to .NET, to increase accessibility:

As tens of thousands of programmers worldwide (CHECK NUMBER JEEEEZ) are already writing F# programs, and that most of F#'s functional language features can be directly translated into equivalent Futhark features, it became worthwhile to investigate whether it was possible to design a way for users to both write and utilize Futhark in F# projects, without ever actually touching the Futhark language or compiler themselves. Instead, users can write their data parallel F# modules in FShark, and compile these modules into Futhark libraries automatically.

In this case, it would be possible to get Futhark speeds in F# programs, without doing much more than installing the Futhark compiler locally, and adding the required FShark libraries to the F# project.

It is my belief that being able to achieve Futhark performance in regular F# programs almost automatically, will make it significantly easier for people to adapt to Futhark programming.

(SOME MORE MORE SOME MORE)

---

<sup>2</sup>which stands for Imperative Code



## The contributions of this thesis

The contributions of this thesis are as follows:

1. The `FSharkPrelude`:

The `FSharkPrelude` is a subset of the Futhark SOACs, ported to `F#`. To write an `FShark` program, the user is directed to limit himself to the SOACs in the `FSharkPrelude`. This means exchanging `Array.map` for `FSharkPrelude.Map`, `Array.foldBack` for `FSharkPrelude.foldr`, and so on. However, the `FSharkPrelude` carries the guarantee, that all `FSharkPrelude` functions works equivalently to their Futhark SOAC namesakes. This prelude, together with the `F#` subset chosen for `FShark`, makes it possible to write `F#` programs which, when translated to Futhark code, are equivalent to their Futhark counterparts.

2. An `F#` subset translatable with `FShark`:

As `F#` is not only a multi-paradigm language, but also has access to the entire standard .NET library, it was required to make `FShark` support only a subset of `F#`. This has been implemented by whitelisting only the `F#`-to-Futhark translatable types, constructs and expressions in the `FShark` compiler. Furthermore, no other module imports than `FSharkPrelude` are allowed. This subset is of course documented for users.

3. The `FShark` Compiler and Wrapper:

The `FShark` Compiler and Wrapper takes a module written in `FShark` as input, converts the `FShark` module into a compiled Futhark `C#` module, and makes it available to the `F#` program, all at runtime. The pipeline is described in sec ??

4. A `C#` backend for Futhark:

To actually use Futhark in `C#` projects (and transitively `F#` projects), it was necessary to develop and add a `C#` backend to the Futhark compiler. This backend is equivalent in functionality to the `C-` and the `Python` backends that are already available.

## Vocabulary

Unless otherwise specified, these are the terms used in the thesis:

### For FShark

- The FShark *subset* is the subset of the F# language that is supported by the FShark compiler.
- The FShark Prelude is the library of F#-ported Futhark array functions and SOACs, and is included with FShark.
- FShark code is F# code which exclusively uses the FShark subset and FSharkPrelude.
- FShark modules are F# modules written entirely in FShark code.
- FShark projects are F# projects which uses FShark and FShark modules.

### For Futhark

- Futhark code is code written in Futhark.
- Futhark C-, Python- or C# code refers to Futhark code that has been compiled into C-, Python or C# source code.

## Roadmap

The main part of this thesis is split in four parts. blaaaah

## Chapter 2

# Background

In this chapter we will first show two languages for GPU programming, namely CUDA and Futhark. Then we will show C# and its interoperability with Futhark. Finally, we will take a look at F#, and how we can expect Futhark/F# interoperability.

### 2.1 CUDA

GPU programming is in principle easily available for everyone. As long as the user has access to a GPU and a reasonable PC for developing software, it just takes a bit of effort and reading to get started with CUDA, OpenCL or similar programming. Realistically however, it takes much more than just a little effort to start writing one's own GPU programs.

#### 2.1.1 A simple CUDA program

Take for instance the function  $f(x, y) = ax + y$ . In figure 2.1 we see the function implemented as a CUDA program. In this program, we are defining the kernel `saxpy` itself, and also manually copying data back and forth between the GPU.

##### The CUDA kernel

Line 3's `__global__` signifies that the following function is a CUDA kernel. Line 4 to 10 contains the computational kernel for the GPU. Line 4 contains the kernel's name and arguments. The kernel takes as arguments  $a$  is a scalar constant, the pointers  $x$  and  $y$  are references to floating point arrays in the GPU device memory, and  $n$  denotes the length of the arrays  $x$  and  $y$ .

Line 6 computes the current computational thread's global id  $i$ . If we compare a parallel computational kernel with a sequential for-loop, this global id is the iterator variable. Line 8 performs the actual  $f(x, y)$  calculation, and stores the result in the  $y$  array, but only if the if-clause in line 7 is true. As we have tens of thousands of threads running simultaneously on the CUDA device, we only want to perform any array operations if

we know that our current global id is within the length of the array.

### **The CUDA main function**

We also need a main function to run the kernel:

Line 14 sets  $N$  to  $1 \ll 20$  (or  $2^{20}$ ).

Line 16 and 17 allocates memory for two arrays  $x$  and  $y$  in system memory.

Line 19 and 20 allocates memory for two arrays  $x$  and  $y$  on the CUDA device (the GPU.)

Line 22 to 25 initializes the arrays  $x$  and  $y$  with scalar values 1.0 and 2.0.

Line 27 and 28 copies our arrays from system memory to the corresponding arrays on the CUDA device.

Line 31 executes the `saxpy` kernel.

Line 33 copies the result from CUDA device back to the  $y$  array in the system memory.

The remaining lines free the allocated memory, first from the CUDA device and then from the system memory.

---

```

1  #include <stdio.h>
2
3  __global__
4  void saxpy(int n, float a, float *x, float *y)
5  {
6      int i = blockIdx.x*blockDim.x + threadIdx.x;
7      if (i < n){
8          y[i] = a*x[i] + y[i];
9      }
10 }
11
12 int main(void)
13 {
14     int N = 1<<20;
15     float *x, *y, *d_x, *d_y;
16     x = (float*)malloc(N*sizeof(float));
17     y = (float*)malloc(N*sizeof(float));
18
19     cudaMalloc(&d_x, N*sizeof(float));
20     cudaMalloc(&d_y, N*sizeof(float));
21
22     for (int i = 0; i < N; i++) {
23         x[i] = 1.0f;
24         y[i] = 2.0f;
25     }
26
27     cudaMemcpy(d_x, x, N*sizeof(float), cudaMemcpyHostToDevice);
28     cudaMemcpy(d_y, y, N*sizeof(float), cudaMemcpyHostToDevice);
29
30     // Perform SAXPY on 1M elements
31     saxpy<<<(N+255)/256, 256>>>(N, 2.0f, d_x, d_y);
32
33     cudaMemcpy(y, d_y, N*sizeof(float), cudaMemcpyDeviceToHost);
34
35     cudaFree(d_x);
36     cudaFree(d_y);
37     free(x);
38     free(y);
39 }

```

---

Figure 2.1:  $ax + y$  in CUDA

## 2.2 Futhark

Whereas the CUDA program and kernel contained large amounts of memory handling and bounds checking, a similar program written in Futhark spares us for a lot of the manual labor above. Figure 2.2 contains a Futhark program that is semantically equivalent to the CUDA program, in regards to the computational result.

---

```

1 let saxpy (a : f32) (x : f32) (y : f32) : f32 =
2   a*x+y
3
4 let main =
5   let N = 1<<20
6   let a = 2f32
7   let xs = replicate N 1f32
8   let ys = replicate N 2f32
9   let ys' = map2 (saxpy a) xs ys
10  in ys'
11

```

---

Figure 2.2:  $ax + y$  in Futhark

Line 1 to 2 defines a function that takes three floats ( $a$ ,  $x$  and  $y$ ) and returns  $ax + y$ .

Line 4 to 10 defines our main function. Line 5 sets  $N$  to  $1 \ll 20$  (or  $2^{20}$ ).

Line 6 sets  $a$  to 2.0.

Line 7 uses the built-in function `replicate`<sup>1</sup> to generate an  $N$  element array of 1.0

Line 8 uses the built-in function `replicate` to generate an  $N$  element array of 2.0

Line 9 uses the built-in function `map2` to apply the curried function `(saxpy a)` to the arrays `xs` and `ys`.

`map f xs` has the type  $((a \rightarrow b) \rightarrow [a] \rightarrow [b])$ , and returns the array of  $f$  applied to each element of `xs`.

`map2 f xs ys` is very similar, but has the type  $((a \rightarrow b \rightarrow c) \rightarrow [a] \rightarrow [b] \rightarrow [c])$ , and applies  $f$  to the elements of `xs` and `ys` pairwise.

In this case, we are calling `map2` with the function `(saxpy a)`, which is just `saxpy` with the first argument  $a$  already defined.

When we compare the program in figure 2.1 to the same program written in Futhark (figure 2.2), we quickly see how Futhark's high level declarative approach is simpler and less verbose than CUDA's. The Futhark compiler does the heavy lifting, by parsing Futhark source code and generating OpenCL code and wrapping them in standalone C- or Python programs.

## 2.3 F#

F# is a high level multi-paradigm programming language in the .NET family. F#'s syntax follows a classical functional programming style. For instance, this means we can take (some programs) written in Futhark, and port them to F# in a way that very closely resembles the original Futhark code.

Figure 2.3 shows the Futhark program from ?? written in F#.

---

<sup>1</sup>`replicate` has the type `int -> a -> [a]`

---

```

1  let saxpy (a : single) (x : single) (y : single) : single =
2      a*x+y
3
4      let main =
5          let N = 1<<<20
6          let a = 2.0f
7          let xs = Array.replicate N 1.0f
8          let ys = Array.replicate N 2.0f
9          let ys' = Array.map2 (saxpy a) xs ys
10         in ys'

```

---

Figure 2.3:  $ax + y$  in Futhark

F# also supports object oriented programming, and has seamless interoperability with the rest of the .NET language family. We can therefore readily use C# libraries and classes in F#, and vice versa.

## C#

C# C# C# C#

## Chapter 3

# The FShark subset

In the following chapter, we will describe the FShark subset. Although the subset is just a part of F#, we will describe it as if it was a language itself.



### 3.1 FShark syntax

Figures 3.1 to 3.5 shows the complete FShark syntax.

$prog$	$::=$	$module\ prog$
	$ $	$prog'\ prog$
	$ $	$\epsilon$
$prog'$	$::=$	$typealias$
	$ $	$fun$
$progs'$	$::=$	$prog'\ progs'$
	$ $	$\epsilon$
$typealias$	$::=$	$type\ v = t$
$module$	$::=$	$module\ v = prog'\ progs'$
$fun$	$::=$	$[<FSharkEntry>]\ let\ id\ (v_1 : t_1) \dots (v_n : t_n) : t = e$
	$ $	$\let\ v\ (v_1 : t_1) \dots (v_n : t_n) : t' = e,$
		$(for\ any\ i \in 1..n, t_i\ is\ not\ a\ tuple)$

Figure 3.1: FShark statements

$e$	$::=$	$(e)$	(Expression in parenthesis)
	$ $	$k$	(Constant)
	$ $	$v$	(Variable)
	$ $	$(e_0, \dots, e_n)$	(Tuple expression)
	$ $	$\{id_0 = e_0; \dots; id_n = e_n\}$	(Record expression)
	$ $	$[e_0; \dots; e_n]$	(Array expression)
	$ $	$v.[e_0] \dots [e_n]$	(Array indexing)
	$ $	$v.id$	(Record indexing)
	$ $	$v.id$	(Module indexing)
	$ $	$e_1 \odot e_2$	(Binary operator)
	$ $	$-e$	(Prefix minus)
	$ $	$not\ e$	(Logical negation)
	$ $	$if\ e_1\ then\ e_2\ else\ e_3$	(Branching)
	$ $	$let\ p = e_1\ in\ e_2$	(Pattern binding)
	$ $	$fun\ p_0 \dots p_n \rightarrow e$	(Anonymous function)
	$ $	$e_0\ e_1$	(Application)

Figure 3.2: FShark expressions

$$\begin{array}{ll}
p ::= id & \text{(Name pattern)} \\
| (p_0, \dots, p_n) & \text{(Tuple pattern)}
\end{array}$$

Figure 3.3: FShark patterns

$$\begin{array}{ll}
t ::= \text{int8} \mid \text{int16} \mid \text{int} \mid \text{int64} & \text{(Integers)} \\
| \text{uint8} \mid \text{uint16} \mid \text{uint} \mid \text{uint64} & \text{(Unsigned integers)} \\
| \text{single} \mid \text{double} & \text{(Floats)} \\
| \text{bool} & \text{(Booleans)} \\
| (t_0 * \dots * t_n) & \text{(Tuples)} \\
| \{id_0 : t_0; \dots; id_n : t_n\} & \text{(Records)} \\
| t \text{ array} & \text{(Arrays)}
\end{array}$$

Figure 3.4: FShark types

$$\begin{array}{ll}
k ::= n\mathbf{Y} \mid n\mathbf{S} \mid n \mid n\mathbf{L} & \text{(8-, 16-, 32- and 64 bit signed integers)} \\
| n\mathbf{UY} \mid n\mathbf{US} \mid n \mid n\mathbf{UL} & \text{(8-, 16-, 32- and 64 bit unsigned integers)} \\
| d\mathbf{f} \mid d & \text{(Single and double precision floats)} \\
| \text{true} \mid \text{false} & \text{(Boolean)} \\
| (k_0, \dots, k_n) & \text{(Tuple)} \\
| \{id_0 = k_0; \dots; id_n = k_n\} & \text{(Record)} \\
| [k_0; \dots; k_n] & \text{(Array)}
\end{array}$$

Figure 3.5: FShark literals

## 3.2 F# operators available in FShark

The F# subset chosen for FShark is described in 3.6. Note that all of these operators are overloaded and defined for all integer and floating point types in F#, except for modulus which in Futhark is only defined for integers.

### Arithmetic operators

The set of supported arithmetic operators is addition (+), binary subtraction and unary negation (-), multiplication (\*), division (/) and modulus (%).

### Boolean operators

FShark currently supports logical AND (&&), logical OR (||), less- and greater-than (<, >), less- and greater-or-equal (<=, >=), equality (=), inequality (<>) and logical negation (not).

### Special operators

FShark also supports some of F#'s syntactic sugar. These operators might not have direct Futhark counterparts, but their applications can be rewritten in Futhark for equivalent functionality. The supported operators are back- and forward pipes (<| and |>), and the range operator ( $e_0 \dots e_1$ ), which generates the sequence of numbers in the interval  $[e_0, e_1]$ .

Note that the range operator must be used inside an array (as so  $[|e_0..e_1|]$ ), so the expression generates an array instead of a list.

Figure 3.6: FShark operators

## 3.3 F# standard library functions available in FShark

FShark supports a subset of the F# standard library. These are readily available in all F# programs, without having to open other modules. The standard library subset is shown in figure 3.7.

### 3.3.1 On selection the F# subset to include in FShark

For selecting the F# subset to support in FShark, I chose to look at what functions that were included in F#'s prelude. That is, the functions that are available in an F# program without having to open their containing module first. Fortunately, F# opens several modules by default of which I only needed to look in two different ones, to be able to support a reasonable amount of F# built-ins in FShark.

The primary module used in my supported F# subset is the module `FSharp.Core.Operators`. This module contained not only the standard arithmetic described in figure 3.6, but also most<sup>1</sup> of the functions shown in the figure 3.7. Except for unit type functions like `failwith`, `exit` and `async`, most of the functions and operators `FSharp.Core.Operators` have direct counterparts in Futhark's prelude, with equivalent functionality: All except for four of operators and functions chosen for FShark are in fact implemented in Futhark's `math.fut` library. It was therefore an obvious decision to support these functions and operators in FShark.

<sup>1</sup>except for some conversion functions, found in `FSharp.Core.ExtraTopLevelOperators`

**id**

The identity function.

**Common math function**

The square root function (`sqrt`), the absolute value (`abs`), the natural exponential function (`exp`), the natural- and the decimal logarithm (`log` and `log10`).

**Common trigonometric functions**

Sine, cosine and tangent functions (both standard and hyperbolic): `sin`, `cos`, `tan`), `sinh`, `cosh` and `tanh`. Also one- and two-argument arctangent: `atan` and `atan2`.

**Rounding functions**

FShark supports all of F#'s rounding functions: `floor`, `ceil`, `round` and `truncate`.

**Number conversion functions**

FShark supports all of F#'s number conversion functions. For all the following functions  $t, te = e', e : t_0, e' : t$ , barring exceptions like trying to convert a too large 64-bit integer into a 32-bit integer.

The conversion functions available are `int8`, `int16`, `int`, `int64`, `uint8`, `uint16`, `uint`, `uin64`, `single`, `double`, `bool`.

**Various common number functions**

`min`, `max`, `sign` and `compare`.

Figure 3.7: FShark operators

However, for the remaining four functions, that didn't have equivalents in Futhark's `math.fut`, their function calls are replaced with their identities instead. In example, the FShark code in figure 3.8 can be expressed in Futhark directly (figure 3.9). However, the hyperbolic cosine function is available in F# (figure 3.10), but not in Futhark.

Therefore, the compiled FShark Futhark code is just the hyperbolic function inlined (figure 3.11).

These rewritings are not pretty to look at from a programmer's perspective, but FSharks Futhark code is not meant to be read by humans anyhow.

(MAYBE INVESTIGATE WHETHER INLINING THESE HAS PERFORMANCE PENALTIES)

### 3.3.2 Missing arithmetic operators in FShark

Currently, bitwise operators like bitwise-AND and bitwise-OR are missing, but they should be relatively simple to add to the FShark subset, by adding them to the set of supported operators in the FShark compiler.

---

```
exp x
```

---

Figure 3.8:  $e^x$  in F#

---

```
exp x
```

---

Figure 3.9:  $e^x$  in Futhark

---

```
cosh x
```

---

Figure 3.10:  $\cosh x$  in F#

---

```
((exp x) + (exp (-x))) / 2.0
```

---

Figure 3.11: FSharks  $\cosh x$  compiled to Futhark

## The correctness of the FShark subset.

When transpiling code from one language to another, it is absolutely vital that the programmer can trust that the target language result is semantically equivalent to the source language code.

In FSharks case, it means that any program written using the FShark subset must have the same result no matter whether it is run natively as F# code, or run as FShark compiled Futhark code.

Take for example the logarithm function `log`. In both F# and Futhark, this function denotes the natural logarithm. If we were translating from a careless language *l*, where `log` instead denoted the binary logarithm, the translation would still be succesfull and not trigger any type errors. However, the native result and the Futhark result would be wildly different.

To ensure that every operator and function in the FShark subset has equivalent results, no matter whether the FShark code is run as native F# code, or compiled into Futhark, I have written a test suite with unit tests for each element in the F# subset.

An FShark test is an FShark module, but with two extra values added, namely an input and an output value for the test. For example, the test written for the division operator is shown in figure 3.12.

---

```
module Div
open FSharkPrelude.FSharkPrelude
open FShark.TestTypes.TestTypes
open System

[<FSharkEntry>]
let div (fiveByte : int8) (fiveShort : int16) (five : int)
      (fiveLong : int64) (fiveSingle : single) (fiveDouble : double)
      : (int8 * int16 * int * int64 * single * double) =
    (fiveByte / 2y, fiveShort / 2s, five / 2,
     fiveLong / 2L, fiveSingle / 2.0f, fiveDouble / 2.0)

[<FSharkInput>]
let value = [|5y; 5s; 5; 5L; 5.0f; 5.0|] : obj array

[<FSharkOutput>]
let sameValue =
    (2y, 2s, 2, 2L, 2.5f, 2.5) : (int8 * int16 * int * int64 * single *
    ↪ double)
```

---

Figure 3.12: The unit test for the FShark division operator

## Testing arithmetic operators

For all arithmetic operators available in FShark, I have written an accompanying test, suitably located in the directory `FSharkTests/UnitTests/Operators` in the FShark project.

The test suite for the operators are quite rudimentary, and are primarily designed to confirm that the types of the operands are preserved when passed to Futhark C# and

back again.

The test suite does not test edges (such as integer overflows or dividing decimals by zero) behaves equally in natively executed `FShark` code and translated Futhark C# code. These edge cases has been left out for SOME reasons.

(THESE ARE NOT ACTUALLY DONE YET) (Are unit tests enough?)

all conversion functions pass through i64. this might be a mistake, as real supports f32 to f64

thoughts on correctness of translations testing correctness of these translations

**Testing F# standard library functions**

**On comparing floating point values**

## Chapter 4

# FSharkPrelude

Besides defining an `F#` subset suitable for Futhark translation, it was also imperative to create a library of SOACs and array functions for `FShark`, to make it possible to write programs with parallel higher-order array functions.

Similarly to how the subset of math functions chosen from `F#` to include in the `FShark` was chosen, the SOACs and array function included in the `FSharkPrelude` has been picked directly from the Futhark libraries `futlib/array.fut` and `futlib/soacs.fut`. The `FSharkPrelude` doesn't discriminate between array functions and SOACs, as maintaining and importing two different prelude files in `FShark` was needlessly complicated.

The `FSharkPrelude` consists of functions which are directly named after their Futhark counterparts, and have equivalent functionality. This prelude, together with the `FShark` subset, is what makes up the `FShark` language. When `FShark` developers are writing modules in `FShark`, they are guaranteed that their `FShark` programs has the same results, no matter whether their programs are executed like native `F#` code, or compiled and executed as Futhark.

The `FSharkPrelude` versions of Futhark functions are defined in three different ways.

1. Functions like `map` and array functions like `length` have direct `F#` equivalents. The `FSharkPrelude` versions therefore simply pass their arguments on to the existing functions. In example, `FSharkPrelude.Map` is shown in figure 4.1. Other functions, like the `map` functions which takes multiple arrays as arguments, require a bit of assembly first. For those `map` functions, we zip the arguments before using `Array.map` as usual:
2. Some Futhark SOACs have `F#` counterparts that are very close to their original definition. I.e., Futhark's `reduce` takes a neutral element<sup>1</sup> as one of the arguments in their function calls, whilst their `F#` counterparts (`Array.reduce`) does only take an operator and an array as arguments. In such cases, the `FSharkPrelude` version changes the input slightly before passing it on to the existing function. See figure ??.

---

<sup>1</sup>For parallelization purposes



3. Lastly, some functions does not have F# counterparts at all, such as `scatter`. In these cases, we manually implement an equivalent function in F#. Note that we are not limited to the `FShark` subset in the `FSharkPrelude`, as the prelude functions are not translated by the `FShark` compiler, but detected caught and exchanged for `Futhark` functions during the `FShark` compilation (see ??). `FSharkPrelude.Scatter` is shown in figure 4.3.

The complete list of available SOACs and array functions is available in appendix ??.  
MAYBE

Note that calls to `FSharkPrelude` functions

---

```
let Map (f : 'a -> 'b) (xs : 'a array) : 'b array =  
    Array.map f xs
```

---

Figure 4.1: FSharkPrelude.Map

---

```
let Reduce (op: 'a -> 'a -> 'a) (neutral : 'a) (xs : 'a array) =  
    let xs' = Array.append [|neutral|] xs  
    in Array.reduce op xs'
```

---

Figure 4.2: FSharkPrelude.Reduce

---

```
let Scatter (dest : 'a array) (is : int array) (vs : 'a array) : 'a array =  
    for (i,v) in Zip is vs do  
        dest.[i] <- v  
    dest
```

---

Figure 4.3: FSharkPrelude.Scatter

## Why is FSharkPrelude part of the FShark language?

Although plenty of functions in the Futhark library already has F# counterparts, we have chosen not to allow these F# counterparts to be used directly in FShark programs. Besides basic differences like different naming in F# and Futhark for equivalent functions<sup>2</sup> like, there are multiple other reasons.

1) From a user experience point of view, it is awkward to maintain a whitelist of accepted functions from certain classes. In example, `Array.map` is exchangeable with Futhark's `map`, but there are no immediate Futhark version of F#'s `Array.sortInPlace`. Therefore, the FShark compiler would successfully exchange a call to `Array.map` with a call to `map`, but it would have to halt with an error message, if the user tried to use `Array.sortInPlace`.

2) Some `Array` functions have subtle differences compared to their Futhark counterparts. As shown in figure 4.2, `Array.reduce` is slightly different in F#.

We are slightly hypocritical, as we DO let users use a subset of F#'s standard library functions. However, there is a whitelist available for this subset in the FShark language specification, and the standard library functions are not visibly called as a method from another module.

## How FShark SOACs differ from Futhark's ditto

On a surface level, FShark and Futhark SOACs are the same. After all, they have equivalent functionality. However, Futhark's SOACs gets special treatment in the Futhark compiler, and are fused together where applicable. Take for instance the short code example in figure 4.4.

---

```
entry main : [] f32 =  
  let xs = iota 100  
  let ys = map (f32.i32) xs  
  let zs = map (+ 4.5 f32) ys  
  in zs
```

---

Figure 4.4: A short Futhark program consisting of just SOACs

For non-OpenCL programs, Futhark's compiler fuses all three expressions into one for-loop, as described in simplified Futhark C# code in figure 4.5. Similarly, in an OpenCL program, the short code example is translated into a single kernel, as shown in figure 4.6.

In both of the compiled examples, we must first allocate a target array for our result, but note that although we obtain three different arrays in the original Futhark code, both of the compiled versions transform the `iota` expression into a for-loop instead, and inserts the operators from the two subsequent `maps` into the loop.

This is a concrete implementation Futhark fusion rules as defined in [?]; which states that  $(\text{map } f) \circ (\text{map } g) \equiv \text{map}(f \circ g)$

---

<sup>2</sup>like `fold` and `foldBack` vs. `foldl` and `foldr`

---

```

float[] mem = new float[100];
for (int i = 0 ; i < 100 ; i++)
{
    float res = int_to_float(i);
    res = res + 4.5f;
    mem[i] = res;
}

```

---

Figure 4.5: Figure 4.4 compiled as (simplified) non-OpenCL C# code.

---

```

__kernel void map_kernel(__global unsigned byte *mem)
{
    int global_thread_id = get_global_id();
    bool thread_active = global_thread_id < 100;

    float res;

    if (thread_active) {
        res = int_to_float(global_thread_id);
        res = res + 4.5f;
    }
    if (thread_active) {
        *(__global float *) &mem[global_thread_id * 4] = res;
    }
};

```

---

Figure 4.6: Figure 4.4, but compiled as a simplified OpenCL kernel.

However, executing FShark code as native F# code will execute the expressions as written, which means that we are allocating and writing to an array three times, once for each line in the program.

### Futhark and nested maps

Futhark’s compiler also specializes in parallelizing nested SOAC calls[?], which in example transforms nested map expressions into one single map expression. For Futhark programs like the one in figure 4.7, the resulting OpenCL program contains a single map kernel with  $i * j$  active threads.

---

```

let xss = map (λrow ->
    map (fun col ->
        row * col
    ) <| iota j
) <| iota i

```

---

Figure 4.7: A nested FShark program

The F# compiler doesn’t make any such transformations for FShark programs.

## Arguing for Futhark-equivalent functionality

using a test suite with both positive and negative testing

### 4.1 Arrays in F# versus in Futhark

As Futhark is an array language, designing the array handling for FSharp was a non-inconsequential part of the design process. Whereas multidimensional arrays in Futhark are written as i.e. `[[ ] i32` for a two dimensional integer array, their actual representation in the compiled code is a flat array of bytes, and an array of integers denoting the lengths of the dimensions. Accessing the array at runtime can be done in  $O(1)$ , whether it's either at some constant or a variable index (i.e. `let second_x = xs[2]` or `let n = xs[i, j]`). The indexes are resolved during the Futhark compilation, either as scalars, or as a variable calculated from other variables.

RESEARCH WHETHER “RANDOM READS” ARE COALESCED IN A NICE WAY IN FUTHARK.

Functional languages like Haskell and F# mainly works with lists. In F#, lists are implemented as singly linked lists. Nodes in the list are dynamically allocated on the heap, and lookups take  $O(n)$  time. We cannot make multidimensional lists, but we can make lists of lists: If we were to emulate a two dimensional list of integers in F#, we could use the type `int list list`. At runtime, the type would then be realized as a singly linked list of references to singly linked list of integers. For an `int list list` of  $i \times j$  integers, we therefore have lookups in  $O(i + j)$  time.

F# does also have arrays. The `System.Array` class itself is reference type. If we initialize an integer array in FSharp like so: `let arr = Array.create 10 0`, the type of `arr` is a reference to where it's corresponding array is located in memory. As the integers contained in the array are value types, the layout of the array referenced by `arr` is some initial array metadata, and then the ten integers stored in sequence.

We can access the array elements on  $O(1)$  time, as indexing into the array is just done by accessing the array reference plus an index offset. If we want to emulate multidimensional arrays with these elements, we can create arrays of arrays (in .NET terms, these are called “jagged arrays”). In figure 4.8 we initialize a jagged array of integers.

`xss` is an array of arrays, so `xss` is a reference to an array in memory, which itself contains references to other arrays. To retrieve the variable `some_two`, we first follow the reference to the array `xss` in memory. There we get the second element, which is a reference to another array in memory. In this array, we read the third element, which in this case is the 2 that we wanted.

The lookup takes  $O(d)$  time<sup>3</sup>, as we access arrays in  $O(1)$  time, and have to follow  $d$  references to get to our element. If we just wanted a reference to the second array in `xss`, we would be chasing the first reference to `arr`, and then return one of the references stored within.

---

<sup>3</sup>where  $d$  is the number of references we we are chasing to get our element.

---

```

let i = 8
let j = 5
let xss = Array.init i <| (Array.create j)

(* xss = [
    [|0;0;0;0;0;0|];
    [|1;1;1;1;1;1|];
    [|2;2;2;2;2;2|];
    [|3;3;3;3;3;3|];
    [|4;4;4;4;4;4|];
    [|5;5;5;5;5;5|];
    [|6;6;6;6;6;6|];
    [|7;7;7;7;7;7|];
  ]
*)

let some_two = xss.[2].[3]

```

---

Figure 4.8: Initializing a jagged array of integers in FSharp

FSharp also offers actual multidimensional arrays. Instead of by `initabb` DO MORE HERE

comparison between arrays: multidims are represented as single objects in memory, less cache misses. and so forth

WE HAVE NO WAY OF KNOWING WHERE THINGS ARE ALLOCATED

SHOULD WE LOOK AT PERFORMANCE OR LANGUAGE DESIGN FIRST?

## 4.2 Converting jagged arrays to Futhark's flat arrays, and back again

As mentioned in section 6.0.1, we cannot just pass jagged arrays as arguments to the Futhark C# entry functions. Instead, we must convert our jagged array into a flat array and an array of integers, and pass these two objects as arguments instead.

### 4.2.1 Analysis of FlattenArray

The simple algorithm for this flattening is described in pseudocode in figure 4.9. The implemented algorithm is slightly more complex, as it has to perform various type castings, and also checks for invalid arrays such as irregular arrays. The implemented algorithm is available in the appendices.

When `FlattenArray` first is called with a jagged array as input, we don't know how many dimensions this array has. Therefore, we recursively call `FlattenArray` on the subarrays of the arrays, until these recursive calls reach a base case. The base case is the array that does not contain array references, but primitive values.

- L2** : For a one dimensional jagged array, this branch is taken once. For a jagged array of  $d$  dimensions, it's taken  $\prod_{n=1}^{d-1}(\text{subarrays at } d_n)$  times.
- L3** is the base case, which takes  $O(1)$  time. This is because we are just returning a tuple with the original array, and singleton array that holds the length of the array (creating the singleton array is also  $O(1)$ .)
- L4** : For a jagged array of  $d$  dimensions, this branch is taken  $\prod_{n=1}^{d-1}(\text{subarrays at } d_n)$  times.
- L5** is the start of the recursive case. This line is called  $O(d)$  times,  $d$  being the number of dimensions in the jagged array. The result of `map FlattenArray array` is an array of a array references and integer array references.
- L6** MORE HERE
- L7** simply retrieves a reference to the first array in the array of subarray lengths. This is  $O(1)$ .
- L8** is by far the most costly line in the function. `F#s Array.concat` function takes a sequence of arrays, allocates a new array, and copies each element of the old arrays into the new array. Each of the  $n$  elements in the jagged array is copied to a new array a maximum of  $d$  times, which means we are performing  $O(n * d)$  reads and writes.
- L9** retrieves the length of an array, and is  $O(1)$ .
- L10** appends a singleton array to the accumulated array of subarray dimensions, by first creating a singleton array, and then copy both the single element and the contents of the accumulated array to a third array of their collected length.

All in all, the upper bound on the `FlattenArray` algorithm is  $O(n * d)$ . This is a far cry from the performance of flattening in Futhark. Flattening is done in  $O(1)$ , as flattening merely calculates the product of the dimensions of the array, and returns the result as the new single dimension of the array.

---

```

1 FlattenArray (array : Array of a) : (Array of b * Array of int) =
2   if a is not (Array of a):
3     return (array, [len(array)])
4   else:
5     subarrays_and_lengths = map FlattenArray array
6     (subarrays, subarrays_lengths) = unzip subarrays_and_lengths
7     subarray_lengths = head(subarrays_lengths)
8     concatenated_subarrays = concat subarrays
9     this_length = len(array)
10    lengths = [this_length] @ subarray_lengths
11    return (concatenated_arrays, lengths)

```

---

Figure 4.9: Flattening jagged arrays, pseudocode

## 4.2.2 Analysis of UnflattenArray

The algorithm `UnflattenArray` in figure 4.10 restores the flat array from the Futhark C# program, to a jagged array in F#. Like in `FlattenArray`, the most expensive

line in the function is the array-manipulating one. In `UnflattenArray`, it is line 7: For each dimension in the `lengths` array, we chunk our data array into multiple smaller arrays. Each of the  $n$  elements in the initial array is moved to a new and smaller array  $d$  times, which makes the complexity of this algorithm  $O(n * d)$ .

### Why `UnflattenArray` hinders a specific tuple type

When an `FShark` function is invoked, it's arguments are prepared by an argument converter first. For scalar arguments, the argument is simply returned. But for array arguments, we must flatten the jagged array into a tuple that follows Futhark's array representation.

When the Futhark function returns, we then have to unflatten the Futhark arrays back into jagged arrays. To do this, we naively look at all the values returned by the Futhark function, and whenever we encounter a tuple of type `('a [] * int64 [])`, we assume that this is a flat array that needs to be unflattened. This procedure works fine, but has one side effect: `FShark` doesn't support entry functions that has `(('a [] * int64 []))` tuples in their return types, because this type is reserved.

To circumvent this, the user is instead encouraged to return the tuple as two separate values.

---

```

1 UnflattenArray (lengths : Array of int) (data : Array of a) =
2   if len(lengths) = 1:
3     return data
4   else:
5     length = head(lengths)
6     lengths' = tail(lengths)
7     data' = chunk_array length data
8     data'' = map (UnflattenArray lengths') data'
9     return data''

```

---

Figure 4.10: Recreating a jagged array from flat array with dimensions

### An alternative solution (`FSharkArrays`)

Instead of using jagged arrays (or even multidimensional arrays), we initially considered using an `FShark` specific array type, which could be directly translated to Futhark's flat array structure.

How they work

How they would alleviate the problem

Why they weren't chosen anyhow (hint; needing to pepper `FSharkArray` all over code, would stand in way of idiomatic `FSharp` style)



## Conclusion on arrays

Ultimately, choosing between jagged arrays, multidimensional arrays and `FSharkArrays` became a question of simplicity vs. performance. For `FShark`, I had the liberty to focus solely on simplicity, as `FShark` code is neither intended or even efficient when executed as native FSharp code. Therefore I could choose to let `FShark` use jagged arrays, instead of any of the other options.

The syntax for declaring a jagged array type closely resembles Futhark's multidimensional array syntax (take for instance FSharp's `int [] []` versus Futhark's `[] [] i32` for declaring two-dimensional integer arrays). The close similarities between Futhark and `FShark` code means that `FShark` generated Futhark code is easier to read for debugging purposes, and likewise makes Futhark code easier to port to `FShark`.

## Chapter 5

# The FShark Compiler and Wrapper

Although `FShark` code can be executed directly in `F#` as normal `F#` code, our benchmarks in chapter ?? shows that compiling our `FShark` code to Futhark OpenCL modules gives us performance increases by several orders of magnitudes (from  $\times 100$  to  $\times 1000$ ).

### Introduction

Parsing and building a regular `F#` program is trivial when using official build tools like `msbuild` or `fsharpc`. But in the case of `FShark`, we are not interested in the final result from the `F#` compiler, but merely its half-finished product.

As the `F#` Software Foundation offers the official `F# Compiler` as a freely available NuGet package for `F#` projects, we can use this package `FSharp.Compiler.Services` to parse the entire input `FShark` program and give us a Typed Abstract Syntax Tree of the `FSharp` expressions therein.

The `F#` Software Foundation actively encourages developers to create projects using the `F#` compiler library, they have published the collected `F#` compiler as a NuGet package, alongside a tutorial??on the usage of the various compiler parts.

For `FShark`, the `Compiler Services` package is used to compile a Typed Abstract Syntax Tree from a wellformed `FShark` source code file, which we then convert into- and print as a valid Futhark program. The Typed Abstract Syntax Tree is merely an AST that already has tagged all the contained expressions with their respective types.

We'll start with a detailed explanation of the `FShark` Compiler Pipeline.

## The FShark Compiler Pipeline in practice

To examine the compiler pipeline in action, we'll go through the motions with the small example program displayed in figure 5.1.

---

```
1 module FSharkExample
2 open FShark.Main
3
4 [<EntryPoint>]
5 let main argv =
6     let wrapper =
7         new FSharkWrapper(
8             libName="ExampleModule",
9             tmpRoot="/home/mikkel/FShark",
10            preludePath=
11                ↪ "/home/mikkel/Documents/fshark/FSharkPrelude/bin/Debug/FSharkPrelude.dll",
12            openCL=true,
13            unsafe=true,
14            debug=false
15        )
16    wrapper.AddSourceFile "../srcs/ExampleModule.fs"
17    wrapper.CompileAndLoad
18    let xs = [|1;2;3;4|]
19    let input = [|xs|] : obj array
20    let xs' = wrapper.InvokeFunction "MapPlusTwo" input :?> int
21    ↪ array
22    printfn "Mapping (+2) over %A gives us %A" xs xs'
23    0
```

---

Figure 5.1: An F# program using FShark

We begin by constructing an instance of the `FSharkWrapper`. It has the following mandatory arguments:

### **libName**

This is the library name for the FShark program. In the final Futhark .cs and .dll files, the main class will have the same name as the `libName`. This doesn't really matter if FShark is just used as a JIT compiler, but it's good to have a proper name if the user only wants to use the compiler parts of FShark.

### **tmpRoot**

The FShark compiler works in its own temporary directory. This argument must point to a directory where F# can write files and execute subprocesses (Futhark- and C# compilers) which also has to write files.

### **preludePath**

The FShark compiler needs the FShark prelude available to compile FShark programs.

### **openCL**

Although Futhark (and therefore FShark) is most effective on OpenCL-enabled computers, the benchmarks in ?? still show a significant speed increase for non-OpenCL Futhark over native F# code. Therefore, FShark is also available for

non-OpenCL users. Use this flag to tell FShark whether Futhark should compile C# with or without OpenCL.

#### unsafe

For some Futhark programs, the Futhark compiler itself is unable to tell whether certain array operations or SOAC usages are safe, and will stop the compilation, even though the code should (and does) indeed work. To enable these unsafe operations, pass a `true` flag to the compiler.

#### debug

Passing the debug flag to the FShark compiler enables various runtime debugging features, for instance benchmarking the time it takes to run various parts of the compiler.

Now, we can pass a source file to the FShark wrapper, compile<sup>1</sup> it and load it into the FShark wrapper object.

To use the compiled FShark function, we must first wrap our designated input in an `obj` array. In this case, our chosen FShark function takes one argument, an `int` array. We define this array, and construct an argument array containing this single element. If the FShark function takes two arguments, we define an input `obj` array with two elements, and so forth. It is important to declare the input array as an `obj` array. Otherwise, F#'s own type checker might very well faultily infer the input array as something else. In this particular case, `input` would've been inferred as being an `int array array`, until we declared its type specifically.

We then invoke the desired function through the wrapper. As all reflection-invoked functions return a value of type `obj`, we need to downcast this object manually. In this example, we use F#'s downcast operator (`:?>`) to declare the return value as an `int` array. The actual return type is always the same as the return type declared in the source FShark file.

## When FShark Wrapper Compiles

The general way to compile and load an FShark program into the FShark Wrapper, is by adding FShark source files to the wrapper object by calling the `AddSourceFile` method, and followingly calling the `CompileAndLoad` method. Although the FShark wrapper also offers other methods of loading and compilation, this is the primary one, as it initiates the entire FShark compilation pipeline.

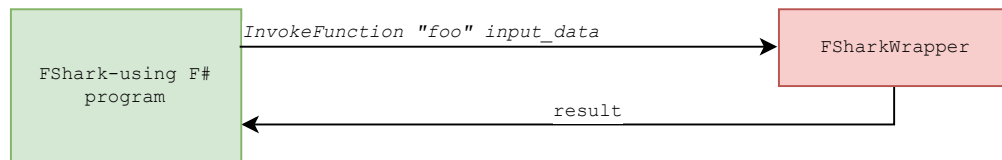


Figure 5.2: The FShark compilation pipeline

When calling `CompileAndLoad`, the supplied FShark source files are concatenated into one long source file, and written to a temporary location. An `FSharpChecker` is

<sup>1</sup> See subsection ??

then initialized, so we can parse and type check the concatenated source code. The `FSharpChecker` is a class exported by the FSharp Compiler Services, and is a class that lets developers use part of the F# compilation pipeline at runtime.

We supply the `FSharpChecker` with the path to our precompiled `FSharkPrelude` assembly, and then call its `ParseAndCheckProject` method on to receive an assembly value, which contains the complete Typed Abstract Syntax Tree of our FShark program, in the form of an `FSharpImplementationFileDeclaration`.

If the FShark developer followed the guidelines to write a well-formed FShark module, the main declaration of the program, the `FSharpImplementationFileDeclaration`, should contain a single `FSharpEntity`, which in turn contains all the remaining declarations in the program.

### The declaration types within F#'s Typed AST

The `FSharpImplementationFileDeclaration` type has three union cases.

#### **InitAction of FSharpExpr**

`InitActions` are `FSharpExprs` that are executed at the initialization of the containing entity. These are not supported in FShark.

#### **Entity of FSharpEntity \* FSharpImplementationFileDeclaration list**

An `Entity` is the declaration of a type or a module. In the case of FShark, three different kinds of entities are supported:

**FSharpRecords** are standard record types, and can be translated to Futhark records with ease. This entity has an empty `FSharpImplementationFileDeclaration list`.

**FSharpAbbreviations** are type abbreviations, and are easily translated into Futhark type aliases. This entity has an empty `FSharpImplementationFileDeclaration list`.

**FSharpModules** are named modules which contains subdeclarations. In this case, we retrieve the subdeclarations from the `FSharpImplementationFileDeclaration list`. The FShark compiler supports building FShark modules, but current limitations demands that modules are flattened when compiled to Futhark. This also means that function name prefixes in function calls are stripped when compiled to Futhark.

#### **MemberOrFunctionOrValue of FSharpMemberOrFunctionOrValue \* FSharpMemberOrFunctionOrValue**

F# doesn't differ between functions and values, which means that a function is merely a value with arguments. A pattern matched `MemberOrFunctionOrValue` value has the form `MemberOrFunctionOrValue (v, args, exp)`, where `v` contains the name and the type of the variable. If the `args` list is empty, `v` is simply a variable. If not, `v` is a function. `exp` contains the `FSharpExpr` that `v` is bound to. An `FSharpExpr` can be anything from a numeric constant to a very long function body.

In figure 5.3 we see a small but valid FShark program. It reads like a regular F# program, but contains the three vital parts that makes it usable as an FShark program.

---

```
1  module ExampleModule
2  open FSharkPrelude
3
4  module SomeValues =
5      let Four : int = 4
6
7      let SomePlus (x : int) (y : int) : int = x + y
8
9      [<FSharkEntry>]
10     let TimesTwo (x : int) : int =
11         SomeValues.SomePlus x x
12
13     [<FSharkEntry>]
14     let MapPlusTwo (xs : int array) : int array =
15         Map ((+)2) xs
16
17     let PlusSeven (x : int) : int =
18         SomeValues.SomePlus x 7
```

---

Figure 5.3: A valid FShark program

- The module declaration on the first line declares that the following code is inside a module. In this case, we are declaring the module `ExampleModule`, although we could use any valid F# module name. As shown in figure 5.4, the top module declaration falls away during compilation, so only the top module contents are left.
- This `open` statement ensures that the F# Compiler Services has access to the `FSharkPrelude` during the compilation. It is possible to write an FShark program which doesn't use the `FSharkPrelude`, but this removes access to the SOACs that we use to write our data parallel programs.
- The `[<FSharkEntry>]` attributed function `TimesTwo` ensures that the resulting Futhark library from the FShark compiler has at least one entry point function. Without any entry point functions, we won't have any functions in the final compiled FShark program.

In figure 5.4 we see the resulting Futhark program. For now, we will ignore the transformations that have happened, except for two things: The `Map` function (called from `FSharkPrelude`) has been rewritten as the plain Futhark SOAC `map` in lowercase, and the module `SomeValues` has been flattened (see sec ?? for future plans.)

This Futhark program is then stored in a temporary location in the user's file system, and compiled into as a library, using Futhark's C# compiler, either with or without OpenCL support. Finally after this compilation, we can invoke the resulting .dll file from within the FShark-using F# program.

---

```

let Four : i32 = 4i32
let SomePlus (x : i32) (y : i32) : i32 =
  ((x i32.+ y))
entry TimesTwo (x : i32) : i32 =
  unsafe SomePlus(x) (x)
entry MapPlusTwo (xs : []i32) : []i32 =
  unsafe map (let x = 2i32 in
    (\(y : i32) -> ((x i32.+ y)))) (xs)
let PlusSeven (x : i32) : i32 =
  SomePlus(x) (7i32)

```

---

Figure 5.4: A valid FShark program, compiled to Futhark

## Building FShark from the Typed AST

Only the supported FSharpExpr's has been listed here, but the full range of FSharpExpr's are available on [?].

### FSharp-to-FSharkIL rules

INTRODUCTION HERE

For these translations, we will disregard that all FSharpExprs are union cases of the F# data type BasicPatterns.

## 5.1 Design choices in writing the FShark Compiler

$$\begin{aligned}
& \llbracket \text{Entity}(\text{IsFSharpRecord}, [(field_0 : \tau_0), \dots, (field_n : \tau_n)]) \rrbracket \\
& = \text{FSharkRecord}([ (field_0 : \llbracket \tau_0 \rrbracket), \dots, (field_n : \llbracket \tau_n \rrbracket) ]) \\
\\
& \llbracket \text{Entity}(\text{IsFSharpTypeAbbreviation}, alias, \tau) \rrbracket \\
& = \text{FSharkTypeAlias}(alias, \llbracket \tau \rrbracket) \\
\\
& \llbracket \text{Entity}(\text{IsFSharpModule}, [decl_0, \dots, decl_n]) \rrbracket \\
& = [\llbracket decl_0 \rrbracket, \dots, \llbracket decl_n \rrbracket] \\
\\
& \llbracket \text{MemberOrFunctionOrValue}((name, \tau^*, \text{IsEntryFunction}), [(arg_0 : \tau_0), \dots, (arg_n : \tau_n)], e) \rrbracket \\
& = \text{FSharkVal}(\text{IsEntryFunction}, \text{FSharkFunction}([\llbracket \tau_0 \rrbracket, \dots, \llbracket \tau_n \rrbracket], \llbracket \tau^* \rrbracket), name, [arg_0, \dots, arg_n], \llbracket e \rrbracket)
\end{aligned}$$

Figure 5.5: Rules for translating FSharp declarations to FShark declarations

$$\begin{aligned}
\llbracket \text{System.Int8} \rrbracket & = \text{FInt8} \\
\llbracket \text{System.Int16} \rrbracket & = \text{FInt16} \\
\llbracket \text{System.Int32} \rrbracket & = \text{FInt32} \\
\llbracket \text{System.Int64} \rrbracket & = \text{FInt64} \\
\llbracket \text{System.UInt8} \rrbracket & = \text{FUInt8} \\
\llbracket \text{System.UInt16} \rrbracket & = \text{FUInt16} \\
\llbracket \text{System.UInt32} \rrbracket & = \text{FUInt32} \\
\llbracket \text{System.UInt64} \rrbracket & = \text{FUInt64} \\
\llbracket \text{System.Single} \rrbracket & = \text{FSingle} \\
\llbracket \text{System.Double} \rrbracket & = \text{FDouble} \\
\llbracket \text{System.Boolean} \rrbracket & = \text{Bool} \\
\llbracket \text{System.Array } \tau \rrbracket & = \text{FSharkArray } \llbracket \tau \rrbracket \\
\llbracket \text{System.Tuple } (\tau_0 \times \dots \times \tau_n) \rrbracket & = \text{FSharkTuple } (\llbracket \tau_0 \rrbracket \times \dots \times \llbracket \tau_n \rrbracket)
\end{aligned}$$

INSERT NOTE ON RULE FOR TUPLE ('a [] \* long [])

Figure 5.6: Rules for translating .NET types to FSharkIL types

$$\begin{aligned}
\llbracket \text{FInt8} \rrbracket & = \text{i8} \\
\llbracket \text{FInt16} \rrbracket & = \text{i16} \\
\llbracket \text{FInt32} \rrbracket & = \text{i32} \\
\llbracket \text{FInt64} \rrbracket & = \text{i64} \\
\llbracket \text{FUInt8} \rrbracket & = \text{u8} \\
\llbracket \text{FUInt16} \rrbracket & = \text{u16} \\
\llbracket \text{FUInt32} \rrbracket & = \text{u32} \\
\llbracket \text{FUInt64} \rrbracket & = \text{u64} \\
\llbracket \text{FSingle} \rrbracket & = \text{f32} \\
\llbracket \text{FDouble} \rrbracket & = \text{f64} \\
\llbracket \text{Bool} \rrbracket & = \text{bool} \\
\llbracket \text{FSharkArray } \tau \rrbracket & = [] \llbracket \tau \rrbracket \\
\llbracket \text{FSharkTuple } (\tau_0 \times \dots \times \tau_n) \rrbracket & = (\llbracket \tau_0 \rrbracket, \dots, \llbracket \tau_n \rrbracket)
\end{aligned}$$

Figure 5.7: FSharkIL types to Futhark types



$\llbracket Const(obj, \tau) \rrbracket$	$= Const(obj, \llbracket \tau \rrbracket)$
$\llbracket Value(v) \rrbracket$	$= Var(v)$
$\llbracket AddressOf(v) \rrbracket$	$= \llbracket v \rrbracket$
$\llbracket NewTuple(., [e_0, \dots, e_n]) \rrbracket$	$= Tuple([\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket])$
$\llbracket NewRecord((v_0 : \tau_0 * \dots * v_n : \tau_n), [e_0, \dots, e_n]) \rrbracket$	$= Record([\llbracket v_0, [e_0] \rrbracket, \dots, \llbracket v_n, [e_n] \rrbracket])$
$\llbracket NewArray(\tau, [e_0, \dots, e_n]) \rrbracket$	$= List(\llbracket \tau \rrbracket, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket])$
$\llbracket TupleGet(., i, e) \rrbracket$	$= TupleGet(\llbracket e \rrbracket, i)$
$\llbracket FSharpFieldGet(e, ., field) \rrbracket$	$= RecordGet(field, \llbracket e \rrbracket)$
$\llbracket Call(., GetArray, ., nil, [e_0, e_1]) \rrbracket$	$= ArrayIndex(\llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket)$
$\llbracket Call(., name, ., nil, [e_0, \dots, e_n]) \rrbracket$	$= Call(name, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket])$
$\llbracket Call(., name, ., \tau, [e_0, \dots, e_n]) \rrbracket$	$= TypedCall(\llbracket \tau \rrbracket, name, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket])$
$\llbracket Call(., infixOp, ., \tau, [e_0, e_1]) \rrbracket$	$= InfixOp(infixOp, \llbracket \tau \rrbracket, \llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket)$
$\llbracket Call(., unaryOp, ., \tau, [e_0]) \rrbracket$	$= UnaryOp(unaryOp, \llbracket \tau \rrbracket, \llbracket e_0 \rrbracket)$
$\llbracket Let(v, e_0, e_1) \rrbracket$	$= LetIn(v, \llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket)$
$\llbracket IfThenElse(e_0, e_1, e_2) \rrbracket$	$= If(\llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket, \llbracket e_2 \rrbracket)$
$\llbracket Lambda((v : \tau), e) \rrbracket$	$= Lambda(v, \llbracket \tau \rrbracket, \llbracket e \rrbracket)$
$\llbracket Application(func, ., [e_0, \dots, e_n]) \rrbracket$	$= Application(\llbracket func \rrbracket, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket])$
$\llbracket TypeLambda(e) \rrbracket$	$= \llbracket e \rrbracket$
$\llbracket DecisionTree(., .) \rrbracket$	$= Pass$
$\llbracket DecisionTreeSuccess(., .) \rrbracket$	$= Pass$

Figure 5.8: Translation rules for FSharp expressions to FSharkIL expressions

$\llbracket Const(obj, \tau) \rrbracket$	$= obj[\llbracket \tau \rrbracket]$
$\llbracket Var(v) \rrbracket$	$= v$
$\llbracket Tuple([e_0, \dots, e_n]) \rrbracket$	$= ([\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket])$
$\llbracket Record([\llbracket (v_0, e_0), \dots, (v_n, e_n) \rrbracket]) \rrbracket$	$= \{v_0 = \llbracket e_0 \rrbracket, \dots, v_n = \llbracket e_n \rrbracket\}$
$\llbracket List(\llbracket \tau \rrbracket, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket]) \rrbracket$	$= [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket]$
$\llbracket TupleGet(\llbracket e \rrbracket, i) \rrbracket$	$= \llbracket e \rrbracket.i$
$\llbracket RecordGet(field, e) \rrbracket$	$= \llbracket e \rrbracket.field$
$\llbracket ArrayIndex(e_{arr}, [e_0, \dots, e_n]) \rrbracket$	$= \llbracket e_{arr} \rrbracket[\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket]$
$\llbracket Call(name, [e_0, \dots, e_n]) \rrbracket$	$= name(\llbracket e_0 \rrbracket) \dots (\llbracket e_n \rrbracket)$
$\llbracket TypedCall(\llbracket \tau \rrbracket, name, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket]) \rrbracket$	$= \llbracket \tau \rrbracket.name(\llbracket e_0 \rrbracket) \dots (\llbracket e_n \rrbracket)$
$\llbracket InfixOp(infixOp, \llbracket \tau \rrbracket, \llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket) \rrbracket$	$= (\llbracket e_0 \rrbracket) \llbracket \tau \rrbracket.infixOp(\llbracket e_1 \rrbracket)$
$\llbracket UnaryOp(unaryOp, \llbracket \tau \rrbracket, \llbracket e_0 \rrbracket) \rrbracket$	$= \llbracket \tau \rrbracket.unaryOp(\llbracket e_0 \rrbracket)$
$\llbracket LetIn(v, \llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket) \rrbracket$	$= \text{let } v = \llbracket e_0 \rrbracket \text{ in } \llbracket e_1 \rrbracket$
$\llbracket If(\llbracket e_0 \rrbracket, \llbracket e_1 \rrbracket, \llbracket e_2 \rrbracket) \rrbracket$	$= \text{if } \llbracket e_0 \rrbracket \text{ then } \llbracket e_1 \rrbracket \text{ else } \llbracket e_2 \rrbracket$
$\llbracket Lambda(v, \llbracket \tau \rrbracket, \llbracket e \rrbracket) \rrbracket$	$= \lambda(v : \llbracket \tau \rrbracket) \rightarrow \llbracket e \rrbracket$
$\llbracket Application(\llbracket func \rrbracket, [\llbracket e_0 \rrbracket, \dots, \llbracket e_n \rrbracket]) \rrbracket$	$= (\llbracket func \rrbracket)(\llbracket e_0 \rrbracket) \dots (\llbracket e_n \rrbracket)$
$\llbracket Pass \rrbracket$	$= \epsilon$

Figure 5.9: FSharkIL expressions to Futhark

## Chapter 6

# The Futhark C# backend

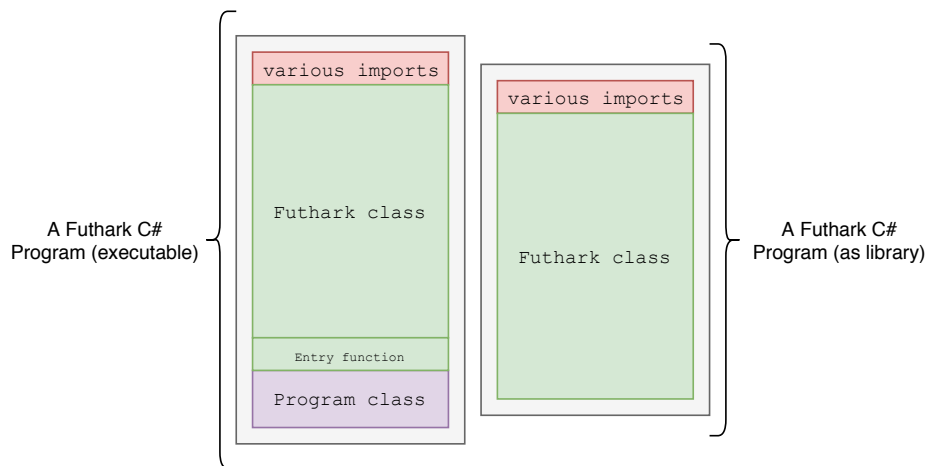


Figure 6.1: The two possible types Futhark C# programs

To be able to use Futhark with F# programs, it was necessary to compile Futhark programs to a language that F# could work with. Although the difference between running a compiled Futhark C- and C# executable from the command line is negligible, a Futhark C# backend would allow .NET projects to use Futhark libraries natively, instead of running their Futhark calculations through separate C or Python modules.

Because F# has almost frictionless interoperability with C#, and C#'s imperative constructs are very close to the intermediate code that Futhark generates for its code generation, it was an easy decision to implement a C# generating backend for Futhark, to accompany the already existing C- and Python backends.

A Futhark backend must be able to do two different programs from a given Futhark program:

First, it must be able to generate standalone executables which can take input data from the `stdin` stream, and send the results to the `stdout` stream. Although a Futhark C,

-Python or C# executable should have equivalent functionality, their performance may vary, and the users may alter between the versions depending on which platforms that are available on their systems.

Second, and more interesting, it must be able to generate single file libraries which can then be imported and used in other C, Python or C# projects, in the same manner as any other library.

---

```
let main (xs : []i32) : []i32 = map (+2) xs
```

---

Figure 6.2: A very small Futhark program `map2.fut`

In example, if we compile the Futhark program in figure 6.2 as a Python library, we will be able to use it in a Python program, as showed in figure 6.3. Likewise, we would like to be able to do the same thing in a C# or an F# context.

---

```
import numpy as np
from map2 import map2

xs = np.array([1,2,3])
map2object = map2()
xs_res = map2object.map2(xs)
print xs_res # prints [3,4,5]
```

---

Figure 6.3: A very small Python program

## The anatomy of a Futhark C# program

In figure 6.1, we see the two different ways we can compile a Futhark program to C#. They're largely the same, except for that the executable Futhark program must have a `Program` class with a `Main` method defined, so that there is an entrypoint defined for the compiled executable. Furthermore, the Futhark class in the executable version contains an entry function which chooses what Futhark function to run (in cases where the Futhark program has more than one entry function defined.)

The `Program` class itself (as seen in figure 6.4) is not especially interesting, and does only contain a `main` method which initialises the Futhark class, and calls the entry function inside the Futhark class. For both Futhark programs, the top consists of the various imports needed for the program.

This leaves us with the Futhark class itself. Figure 6.5 shows the different parts that make up the generated Futhark C# class. In the following sections we will walk through the individual parts.

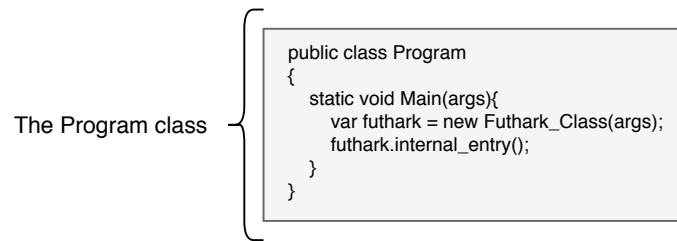


Figure 6.4: The FShark compilation pipeline

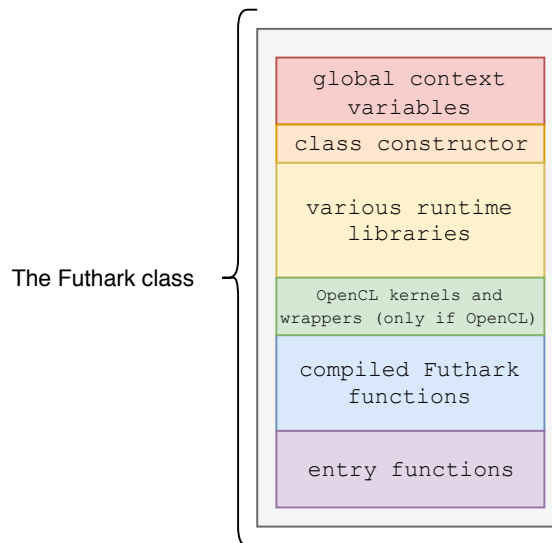


Figure 6.5: The layout of the C# Futhark class

## Global context variables

Compiled Futhark programs need to keep track of several variables. Both normal and OpenCL-enabled Futhark C# programs can take several options when they're launched from the command line. In example, `num_runs` tells the Futhark runtime how many times the chosen entry function should be executed, and the variable `runtime_file` tells the Futhark runtime where it should write timing information to, for example for benchmarking purposes.

Instead of passing an argument array along throughout all the functions in the Futhark class, like we usually would if we were writing purely functional programs, we instead set these arguments as class variables at class initialization, so we can refer to them everywhere throughout the rest of the class.

For non-OpenCL programs, the variables are exclusively for benchmarking and debugging purposes. For OpenCL programs however, the global variables are vital for the program's execution. In an OpenCL program, the Futhark class must keep track of two extra variables.

The struct `futhark_context ctx` is the struct that contains the global state of

the current program's execution. Contained in the global state there is the current list of unused but allocated OpenCL buffers on the device, kernel handles for all the OpenCL kernels used in the Futhark program, and a counter for the total running time of the program. There is even another context contained in the `futhark_context`, namely the `opencl_context`, which contains the current state of the device, and also information about its platform, its queue and so forth.

The struct `futhark_context_config cfg` is similar to the `futhark_context`, but is only used for constructing the actual `futhark_context`.

## The class constructor

The class constructor is necessary to setup the global variables needed throughout the Futhark class. When the Futhark program is compiled as an executable, the command line arguments are passed to the class constructor by the `Program` class. If the Futhark program is compiled as a library, the programmer can pass a string array of arguments to this constructor manually.

Besides setting class variables, OpenCL-enabled versions will initialize (and set) first the `futhark_context_config cfg` variable, and afterwards the `futhark_context` itself.

## The various runtime libraries

The runtime libraries are a set of separate C# files that are written and distributed through the Futhark compiler. When a Futhark program is compiled, these library files are concatenated and embedded directly into the rest of the generated code. They contain functionality which the generated Futhark programs depend on. The runtime libraries are the following:

### **memory.cs**

As Futhark's stores all array values (no matter the dimensionality) as a flat one-dimensional byte arrays (with an accompanying array of 64-integers which denote the dimensions of the flat array), it was necessary to define a set of functions to interact with these byte arrays. I.e., `memory.cs` contains the `writeScalarArray` functions, which writes a scalar value to a byte array. The function is overloaded so it works with scalars of any integer or floating point primitive. See figure 6.6 for an example:

### **scalar.cs**

This library contains all the scalar functions necessary for Futhark C# programs. In Futhark, arithmetic operators are defined for integers and floats of all sizes, and bitwise operators are defined for all integers. However, this is not the case in C#, where many arithmetic operators are only defined for 32- and 64 bit integers.

If these operators are used with 8- or 16 bit operands, the operands are implicitly casted to 32 bit integers at compile time, which also means that the final result of the operation is a 32 bit integer, which doesn't have the right type.

Therefore, wrapper functions must be defined for even the simplest arithmetic functions. I.e., integer addition in C# Futhark is actually four different functions:

---

```

void writeScalarArray(byte[] dest, int offset, double value)
{
    unsafe
    {
        fixed (byte* dest_ptr = &dest[offset])
        {
            *(double*) dest_ptr = value;
        }
    }
}

```

---

Figure 6.6: `writeScalarArray` writes a value at the specified offset in some byte array.

---

```

static sbyte add8(sbyte x, sbyte y){ return Convert.ToSByte(x + y); }
static short add16(short x, short y){ return Convert.ToInt16(x + y); }
static int add32(int x, int y){ return x + y; }
static long add64(long x, long y){ return x + y; }

```

---

Besides, `scalar.cs` also contains the C# definitions for the various mathematical functions from Futhark's `math.futlibrary`, such as `exp`, `sin` and `cos`.

#### **reader.cs**

The reader contains the entire functionality for receiving function parameters through `stdin`. The reader reads scalars of any of the Futhark-supported primitives, and also arrays and multidimensional arrays of scalars. The reader also supports reading streams of binary data. It is only necessary for Futhark executables.

#### **opengl.cs**

MAYBE WRITE ALL OF THIS ALSO? ALRIGHT THANKS

## **The compiled Futhark functions**

The compiled Futhark functions are the Futhark Intermediate Code functions, expressed in the target language, and corresponds to the entry functions found in the entry functions-section of the Futhark class. Only the Futhark `entry` functions are compiled to individual functions, and remaining helper functions are inlined here.

In OpenCL programs, all array functions and SOAC calls are compiled as individual (or fused) OpenCL kernels. Therefore, the compiled Futhark functions in these programs consists of mainly some scalar operations and memory allocations, and calls to Futhark-generated kernel wrapper functions.

In non-OpenCL programs, the array functions and SOAC calls are not stored in separate wrapper functions, but inlined in the Futhark functions.

## OpenCL kernels and wrappers

If the Futhark program is compiled for OpenCL, all array handling function- and SOAC calls are compiled as OpenCL kernels. This part of the Futhark class has two parts:

1. The string (actually a single string in an array) `opencl_prog`, which contains the entire Futhark-generated OpenCL source code for the Futhark program in question. This source code contains all the OpenCL kernels for the program, and is passed to the OpenCL device, compiled and loaded, when the Futhark class is initialized. Handles to the individual kernels are then stored in the `futhark_context`.
2. For each kernel in the `opencl_prog`, the Futhark compiler generates a kernel wrapper function. These wrapper functions takes the kernel arguments (such as scalar values, array values and indexes) as input, and performs all the OpenCL specific work necessary for the actual kernel launch; in example setting the kernel arguments on the device, and copying data back and forth between host and device buffers.

### 6.0.1 Entry functions

Futhark’s internal representation of array values are one dimensional byte arrays (which can represent arrays of any type and dimensionality), and an accompanying list of integers denoting the lengths of the array’s dimensions. However, Futhark does not expect it’s users to pass this form of arrays as function arguments, which is why each Futhark `entry` function has a corresponding entry function in the final compiled code.

To discern between Futhark functions and entry functions, the Futhark function’s name is prefixed with “`futhark_`”, as in for example “`futhark_foo`”. Depending on whether the Futhark program is compiled as an executable or a library, the entry function itself is then named “`entry_foo`” or just “`foo`”.

For executables, “`entry_foo`” is a function that doesn’t take any arguments. Instead, it uses the reader functions from `reader.cs` to parse the arguments for “`foo`” from `stdin`, and passes them to the Futhark function. For all array values in the arguments, the array values are converted into Futhark representations of them. When the Futhark function returns the result, the result is then printed to `stdout`.

For libraries, the “`entry_`” prefix is dropped, and the function just takes care of converting array arguments into and back-from their Futhark representations.

MENTION THE INPUT ARGUMENTS SPECIFICALLY ARRAYS THAT HAVE TO BE FLATTENED FIRST

## The C# backend, compared to the C- and Python counterparts

THE PYTHON BACKEND HAS MUCH FUNCTIONALITY ENCAPSULATED IN PYOPENCL, AND DOESN'T NEED TO DECLARE VARIABLES BEFORE SETTING THEM LESS COMPLEX GENERATOR NEEDED AS VARIOUS OPENCL STATEMENTS ARE HANDLED AUTOMATICALLY BY LIBRARY

C BACKEND MUST BE AWARE OF ALL SIZES AND EVERYTHING AT COMPILE TIME, WHICH MEANS STATES MUST BE ALLOCATED THROUGH COMPLEX STRUCTS AT COMPILE TIME, AND STRUCTS MUST BE DEFINED AT COMPILE TIME AS WELL

C ALLOWS NULL POINTERS, CS DOES NOT WHICH MEANS WE NEED PLACEHOLDER VARIABLES

CSHARP GENERATOR IS SOMEWHERE INBETWEEN AS IT IS CAN HANDLE OBJECTS WHICH CAN CARRY STATE, FURTHERMORE DYNAMIC MEMORY ALLOCATION

## Memory management in Futhark C#

As Futhark stores array values around in byte arrays, it is relevant to compare the difference between how the array handling differs between Futhark's C backend, and this C# backend. For OpenCL programs, the memory management of C# and C is largely the same, as the OpenCL side of these programs are the same. C# does after all just use C bindings for its OpenCL interactions.

However, for non-OpenCL C# programs, we have to take C#'s memory model into consideration

C implicitly allows unsafe programming. In this case, it means interacting with system memory by reading and writing arbitrary values from/to arbitrary locations, designating the values and destinations as whatever type we want. In figure 6.7, we see a `for`-loop that performs a summing scan on an array of integers. On line 6, reading from right to left, we are first creating reference to a location in the byte array `xs_mem_4223`. However, as the reference is a pointer to a byte in the array, we must recast it as an `int32_t` pointer. After we do this, we can finally derefer the pointer to retrieve a four byte integer from the byte array.

We add the retrieved integer to our accumulating variable `scanacc_4187`, before we cast a reference in our destination byte array as an integer pointer, and store the result there.

WHY IS THIS NOT ALLOWED?

DESCRIBE TWO DIFFERENT WAYS OF DOING IT ANYHOW 1) MARSHAL 2) unsafe and fixed

what was chosen and why



---

```

1 memblock mem_4226;
2 memblock_alloc(&mem_4226, bytes_4224);
3 int32_t scanacc_4187 = 0;
4
5 for (int32_t i_4189 = 0; i_4189 < sizze_4135; i_4189++) {
6     int32_t x_4147 = *(int32_t *) &xs_mem_4223[i_4189 * 4];
7
8     scanacc_4187 += x_4147;
9
10    *(int32_t *) &mem_4226[i_4189 * 4] = scanacc_4187;
11 }

```

---

Figure 6.7: A short snippet from a Futhark C program

## Selection an OpenCL interface for C#

OpenCL interaction is not a part of the .NET standard library, but several libraries do exist for .NET/OpenCL interactions. For this thesis, I researched a selection of these libraries, to determine which one that would fit the best for my purposes. As Futhark depends on being able to interface with the OpenCL platform directly, it was necessary to find an OpenCL library for .NET which had direct bindings to the OpenCL developer library.

The .NET libraries I took into consideration was `NOpenCL`, `OpenCL.NET` and `Cloo`. All three libraries have been designed to aide OpenCL usage in C# programs, by simplifying OpenCL calls behind methods GØR BEDRE.

### **NOpenCL**

`NOpenCL` was the first candidate for the C# backend, and had several advantages to the other two: As per February 2018, it had been updated within the last year, and was therefore the least deprecated library. Second, the `NOpenCL` repository on Github contains both unit tests and example programs.

However, `NOpenCL` is also tailored for Windows use, and therefore not a good fit for Futhark, as Futhark is available on both Windows, Linux and Mac OS. Furthermore, the library is not available through the NuGet package manager, and the OpenCL API calls are needlessly complex to work with through the library.

### **OpenCL.NET**

`OpenCL.NET` also has a test suite, is available through NuGet, and is used as the backend for other libraries, such as the F# GPU library `Brahma.FSharp??`.

However, this library hardcoded to work on a in a Windows context, and has not been updated for more than five years.

### **Cloo**

`Cloo` is usable on all three platforms, and it is available on NuGet. Furthermore, as opposed to the other two libraries, the `Cloo` library contains a class with static functions that does nothing but passing arguments on to the OpenCL library,

using C#'s `DllImport` attribute. It is immediately possible to skip most of `Cloos` features, and just use the library for its OpenCL bindings.

Even then, `Cloos` has not been updated within the last five years, and probably won't be in the future either.

Given these three candidates, I chose to work with `Cloos`: It was the only one that had the necessary OpenCL bindings readily available, and the only one that was platform agnostic.

### **Writing a custom OpenCL bindings library**

Though `Cloos` is a good fit for Futhark C#, it is also slightly risky to depend on a five year old unmaintained library in a modern project. Therefore, it could be a good idea to write a smaller library similar to `Cloos`, specifically for Futhark - or maybe even just include it with Futhark as one of the C# runtime libraries.

## Chapter 7

# Benchmarks and evaluation

### **FShark generated Futhark compared to original Futhark code**

Appendices show

#### **The LocVolCalib benchmark**

small.in: FShark (openCL) took 211882 microseconds. Average invocation (fshark non openCL) time was 81194767 microseconds. Native took 438 929 311 microseconds.

medium.in: (Fshark opencl)invokation time was 310833 microseconds Fshark nonopencl Average invocation time was 154 141 321 ms Native took 900 643005 microseconds.

large.in:

fshark with opencl Memory Allocation Error fshark sans opencl 2450 637 053 microseconds Native took 24757 874 577 microseconds.

for all three datasets

#### **The nbody benchmark**

for all three datasets

#### **Specifications for benchmark**

We have run the benchmarks on a system with these attributes:

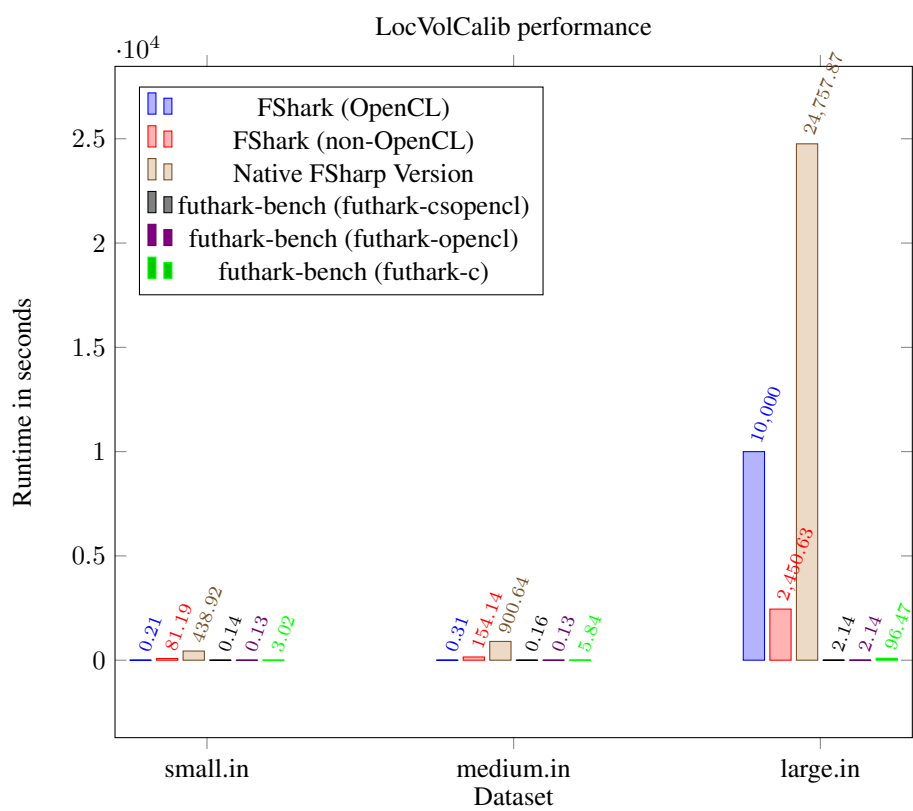


Figure 7.1: Comparison between Python and Futhark performance for simple model

- CPU: 4 cores of Intel Core i5-6500 at 3.20GHz
  - L1 cache: 128 KiB
  - L2 cache: 1024 KiB
  - L3 cache: 6144 KiB
- GPU: GeForce GTX 970

Introduction for the two benchmarks LocVolCalib and nbody  
why are they faster in general

## **Chapter 8**

# **Current limitations**

## **Chapter 9**

### **Related work**

## **Chapter 10**

# **Conclusion and future work**



## Chapter 11

# FSharks interoperability between F# and Futhark (C#)

FShark stands on three legs: The FShark compiler, the Futhark C# code generator, and the FSharkWrapper. The compiler is responsible for compiling FShark source code into Futhark source code, and the C# code generator takes the result Futhark source code, and compiles OpenCL powered C# libraries, which can be imported directly back into F#.

It is of course possible to use the compiler and the code generator as individual modules, but for this project, the FSharkWrapper has been designed to let users use FShark without having to understand any of the underlying pipeline.

To illustrate this; take a look at figure ???. In the first line, the user initializes the FSharkWrapper with the arguments necessary to use the wrapper itself. In the second line, the user adds a source file to the wrapper by its path. In the third line, the user tells the wrapper to run the compilation pipeline. Assuming that the compilation goes well, the user can then invoke some function from the FShark program in line four.

Here, calling the `CompileAndLoad()` function triggers the entire FShark pipeline as described in ??, and does then have a function available for the user to call afterwards.

However, as this is the default way of using FShark, we are currently calling `CompileAndLoad()` every time we use the FShark program. This is happening even though we only need the final compiled C# assembly to load back into F# at run-time.

Everytime we run the FShark compiler pipeline, we are therefore also

1. parsing, typechecking and generating a TAST from the FShark code, using FSharp's compiler.
2. generating Futhark source code from the FSharp TAST
3. Writing the Futhark source code to disk
4. running the Futhark compiler and C# code generator on the Futhark source code

5. running the mono C# compiler on the resulting C# source code

For two selected benchmarks we have the following times

### 11.0.1 Pros and cons of the current design

As there are demonstratively great performance gains to be won by only using the compiler part of the FShark pipeline, it is worth discussing whether the rest of the FShark pipeline should remain.

Besides eliminating the entire compilation operation at every FShark execution, a compiler-only approach to the FShark compiler would give us the following advantages:

- **Standalone-modules first:** As the compiler is now only used once, the resulting Futhark assembly is readily importable in any .NET project, as long as the required Mono libraries are also available. This goes not only for the user who just compiled the assembly, but also for any other user who has acquired the necessary Mono libraries. This means that the FShark developer can use and share the FShark assemblies with colleagues and coworkers like any other sharable .NET library, as this is indeed what a compiled Futhark C# library amounts to.

Corollarily; this FShark design would make FShark is useful for generating high-performing .NET libraries. (Although one could write such libraries in Futhark instead of FShark.)

- **Static typing of FShark module:** The current runtime-only approach means that the user must rely on reflection to call FShark functions. In this situation, all modern IDE comforts like autocompletion, and especially static type checking and inference falls away. For the following example , the current design demands that we first wrap our arguments in an `obj` array, before invoking the function `foo`. Furthermore, we must also downcast the result using F#'s downcasting operator `:?>`. Because we are upcasting our arguments to an `obj` array, we can actually pass any (correctly casted) array as an argument to our reflection-invoked function, without triggering any type errors at compile time. The same goes for the downcasted result from the function. We can cast the result as whatever type we like, and not run into any trouble until we finally run the compiled program. However, if we use FShark to generate assemblies instead, we now have all the type information available at compile time. Our compiler will block us from compiling the program by giving us useful type errors. Last but not least, we can remove all the casting operations that are littering the program.

- **Getting rid of, or at least trimming down, the FSharkWrapper:**

However, the current design of FShark also has some advantages that follows automatically from the design.

- **Rapid FShark code development:** Currently, it is recommended that any FShark code for a project is also built as part of the project. By including the FShark file in the original project's source list, we can call the FShark module natively, without running the FShark compiler, to prototype and debug the FShark

code directly in our IDE, before we switch to using the compiled version of the FShark code.

- **En mere**

## 11.1 The future of FShark interoperability

With these considerations in mind, my future work on FSharks interoperability consists of reducing FSharkWrapper in size, so it only takes an FShark source path and a .NET assembly outpath as inputs, and does nothing more than orchestrating the FShark-, the Futhark- and the C#compiler. The current design is too complex, largely from supporting too many superflous features like concatenating multiple sources, and so on.

I will also be researching the optimal way to keeping the FShark module development as close to the rest of the FShark-using project as possible, without

The current design enables direct prototyping, which must of course be kept in later versions of FShark.