### **Wave Algorithms**

- ning algorithm
- tree algorithm
- echo algorithm
- polling algorithm

# Ring algorithm

```
For the initiator: 
 \mathbf{begin} \  \, \mathbf{send} \  \, \langle \mathbf{tok} \rangle \  \, \mathbf{to} \  \, Next_p \  \, ; \  \, \mathbf{receive} \  \, \langle \mathbf{tok} \rangle \  \, ; \  \, decide \  \, \mathbf{end} 
 For non-initiators: 
 \mathbf{begin} \  \, \mathbf{receive} \  \, \langle \mathbf{tok} \rangle \  \, ; \  \, \mathbf{send} \  \, \langle \mathbf{tok} \rangle \  \, \mathbf{to} \  \, Next_p \  \, \mathbf{end}
```

- Processes are arranged in a unidirectional ring (each process has a sense of direction or knowledge of one dedicated neighbor)
- initiator send message (tok) along the cycle
- each process passes it on
- when it returns to initiator, initiator decides
- Theorem: ring algorithm is a wave algorithm

3

# Polling algorithm

```
\begin{array}{ll} \mathbf{var}\; rec_p &: \mathbf{integer} & \mathbf{inte}\; 0 \;; \; (^\bullet \; \mathbf{for}\; \mathbf{initiator}\; \mathbf{only}\; ^\bullet) \\ \\ \text{For the initiator:} & \mathbf{begin}\; \mathbf{forall}\; q \in \mathit{Neigh}_p \; \mathbf{do} \; \mathbf{send} \; \langle \; \mathbf{tok} \rangle \; \mathbf{to} \; \mathbf{q} \; ; \\ \mathbf{begin}\; \mathbf{receive} \; \langle \; \mathbf{tok} \rangle \; ; \; rec_p := rec_p + 1 \; \mathbf{end} \; ; \\ \mathbf{docide} & \mathbf{end} \\ \\ \text{For non-initiators:} & \mathbf{begin}\; \mathbf{receive} \; \langle \; \mathbf{tok} \rangle \; \mathbf{from} \; q \; ; \; \mathbf{send} \; \langle \; \mathbf{tok} \rangle \; \mathbf{to} \; \mathbf{q} \; \mathbf{end} \\ \end{array}
```

- Works on cliques (complete networks)
- one initiator (centralized)
- how many processes decide?
- How many messages are exchanged in the algorithm
- what other topology can this algorithm be used in
- Is this a wave algorithm?

#### **Wave Algorithms**

- Wave algorithm satisfies the following three properties
  - \_ termination: each computation is finite
  - decision: each computation contains at least one decide event
  - dependence: in each computation each decide event is causally preceded by an event in each process
- initiator(starter) process that execution of its actions spontaneously
- non-initiator(follower) starts execution only when receives a message
- wave algorithms differ in many respects, some features:
  - centralized (single-source) one initiator; decentralized (multi-source) multiple initiators
  - topology ring, tree, clique, etc.
  - initial knowledge:
    - each process knows its own unique name
    - each process knows the names of its neighbors
  - number of decisions to occur in each process
- usually wave algorithms exchange messages with no content tokens

# Tree algorithm

Operates on tree network (can work on spanning tree of arbitrary network) - no root, edges are undirected (bi-directional)

```
var rec_p[q] for each q \in Neigh_p; boolean init false; (* rec_p[q] is true if p has received a message from q *) begin while \#\{q: rec_p[q] \text{ is false}\} > 1 do begin receive \{\text{tok}\} from q; rec_p[q] := \text{true} end (* Now there is one q_p with rec_p[q_p] is false *) send \{\text{tok}\} to q_p with rec_p[q_p] is false ; p receive p to p with p receive p re
```

- leaves of tree initiate the algorithm
- if a process has received a message from all neighbors but one (initially true for leaves), the process sends a message to the remaining neighbor.
- If process gets messages from all neighbors it decides
- Excluding the **forall** statement how many processes can decide? What are these processes?
- Why do we need **forall** statement?
- How many messages are sent in the algorithm?
- ls this a wave algorithm?

## Chang's Echo algorithm

```
var rec<sub>p</sub>: briege: list 0; (* Counts number of received messages *)
fistler<sub>p</sub>: P init udef.

For the initiator:
begin forall q c Neigh, do send (tok) to q;
while rec<sub>p</sub> < q/Knigh, do
begin receive (tok); rec<sub>p</sub> := rec<sub>p</sub> + 1 and;
decide
end

For non-initiators:
begin receive (tok) from neighbor q: failer<sub>p</sub>:= q: rec<sub>p</sub>:= rec<sub>p</sub> + 1;
while rec<sub>p</sub> < q/Knigh, q ≠ failer<sub>p</sub> do send (tok) to q:
while rec<sub>p</sub> < q/Knigh, q ≠ failer<sub>p</sub> do send (tok) to q:
begin receive (tok); rec<sub>p</sub>:= rec<sub>p</sub> + 1 end;
end (tok) to failer<sub>p</sub>
```

- Works on networks of arbitrary topology
- one initiator (centralized)
- initiator sends messages to all neighbors
- when non-initiator receives the first message it forwards it to all other neighbors, when it gets tokens from all other neighbors it replies back
- how many processes decide?
- How many messages are exchanged in the algorithm
- ls this a wave algorithm?

6