Problem1:

Q: What are the states?

A: The set of all reachable states when searching the goal state. For instance,

Q: What are the operators?

A: Operator is the action from one state to another state. For instance, add task or delete task.

Q: What is the branching factor?

A: Branching factor is the number of successors (child nodes) generated by a node (parent node). As Fig.1 shown, the branching factor of A is 2 (because A has two successor).

Q: Is the depth of the goal node known initially?

A: No, the depth of the goal node isn’t known initially. What we know is the upper limit of the depth of the goal node. For instance, some search tree G has n nodes in total. So the depth of the goal node is no greater than n.

The upper limit of the depth of the goal node is base on the assumption that the task of each level requires X = min T.V value, and Y = min T.L time. Since the requirement of Problem1 is an overall deadline D and a target value is M, the overall time of a correct schedule must be no more than D and overall value of a correct schedule must be more than M. The depth must satisfy that it is no greater than ceil (M / X) , floor(D / Y) and N(as mentioned above). So the depth must be no greater than the smallest value of these three values.

Problem 2:

A. Fig.2 is the graph of Problem 1 using DFS. So the portion of the state space is the total number of nodes. As Fig.2 shown, it is 6 (A, C, E, AC, AE, AEB).

B. Fig.3 is the graph of Problem 2 using BFS. So the portion of the state space is the total number of nodes. As Fig.2 shown, it is 10 (A, C, E, AC, AE, CA, CE, EA, EC, AEB).

Problem 3:

A: Fig.3 is the graph of Problem 1 using BFS, assuming that those repeated states are eliminated. So we don’t touch the CA, EA, EC. The portion of states space becomes 7.

B: We can use hash table as explored set. The nodes that have been visited can be put into the explored set (hash table). Hash key is the frontier node that is currently visited. Under some order form, hash value is the node formed by the ordered task combination. For Instance, we can hash the task combination (node) by Alphabetical order. Firstly, AC has been visited and put into hash table. When CA is visiting, the hash key is CA and hash value is AC. We check the hash table to find whether AC is in the hash table or not. If it is there, we discard CA. If not, we put it in hash table.

C: Since BFS stores the visited state in the explored set; the solution in B can eliminated the repeated state in the explored set. However, DFS and ID will not store the previous states until we find the goal node. So there is no need to lookup the explored set.