



Brute-force hash function

## **CR - Collision Resistance**

➡ given  $H$ , hard to find  $m$  and  $m'$  such that  $H(m) = H(m') = x$

Given a hash function  $H$  of  $n$  bits output

- Reaching all possibilities
- On average, an attacker should try half of them

m



X

2n cases

$2n-1$  cases



Brute-forcing a hash function



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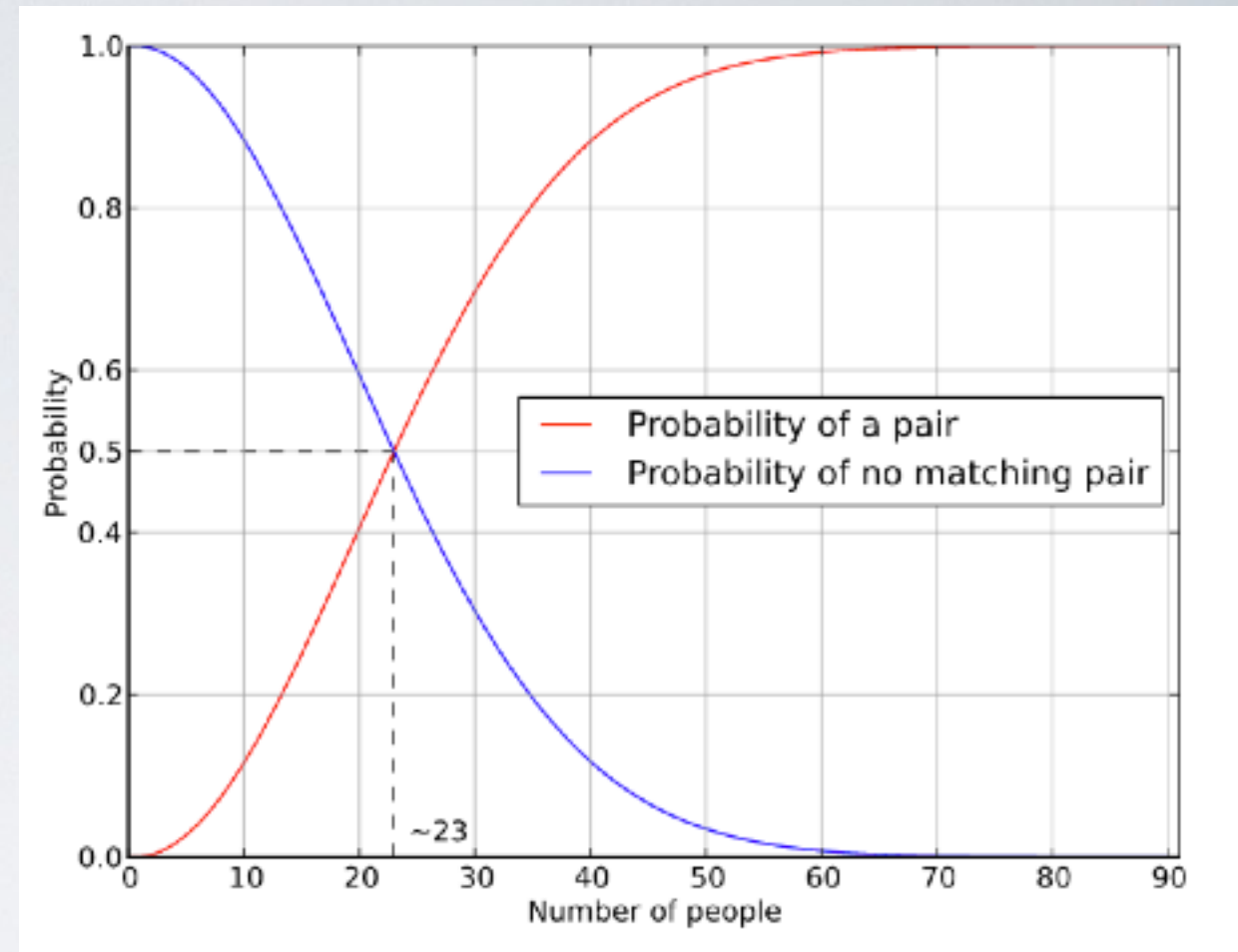
$2^n$  cases

~~$2^{n-1}$~~  cases



# Birthday Paradox

*“There are 50% chance that 2 people have the same birthday in a room of 23 people”*



## N-bits security

- ➡ Given a hash function **H** of **n** bits output, a collision can be found in around  **$2^{n/2}$**  evaluations  
e.g SHA-256 is 128 bits security