



TCP and OSPF Assignment ELC 3080

Name	ID
كلارا عيسى اسحاق عبدالمسيح	9203081
مينا صبحي كامل السيد	9203591

Submitted To: Prof. Dr. Khaled Fouad 3rd Year – Spring 2023

Contents

1. 7	ГСР Lab)	4
1.1	. Effe	ect of TCP Window Size	4
1 f v t t	function window hroughpouttps://ch	Vary iperf3 window size from 1Kbytes to 6 Kbytes in increments of 1 Kbyte, then set it 4, 32 Kbytes. Plot the average throughput and the average number of retransmissions as of window size. Comment on the results and explain the zigzag behaviour noticed for large sizes. Note that number of retransmissions is the 5th column in iperf3 default output & but is the 4th. You can plot data using matlab/python/any online graphing tool (Ex. nart-studio.plotly.com/create/#/). The window size is the last option in the iperf command of INTRO Lab).	ger
f		Click on any of the TCP a data segment whose source is node n7, dissect the segment by g different protocol layer headers from TCP->IP->Ethernet identifying how many header added by each layer, identify the TCP options used	
1	1.1.3.	Repeat for an ACK packet sent from node n11.	10
1.2	. TCI	P short VS long paths	11
	1.2.1. Why thro	Compare the result of throughput with the case when connection was made to node n11. oughput drops when connecting to n8 although capacities on the two paths are the same	
1.3	. Hig	her link Capacity with Drops VS Reliable Lower Capacity	11
1	1.3.1.	Compare throughputs in cases a, b, c. Why b is better than c?	13
1	1.3.2.	Compare throughputs in cases b, c and d. Which is better? Why?	13
1	1.3.3.	Compare throughputs in cases e and f? Which is better? Why?	13
2. (OSPF La	ıb	14
2.1	. OSI	PF Link Cost Change	14
_	2.1.1. what hap	Check what happens to the path between n7 and n11 (as seen after steps 3 and 6)? Explanation Explanation Check what happens to the path choose toolbar->widgets-> Throughput	
r i	nterface	Set the cost of eth1 at node n5 back to 10. Establish two iperf3 connections: one from n7 the second from n11 to n7 both for duration of 500 seconds. Now go to node n4 and set cost for interface connecting n4 with n5 to 40. What happens in the paths of the two ons? Explain what happens. What do you conclude?	
2.2	OSI	PF Database Updates	17
	2.2.1 nore tha	Capture and explain the outputs due to execution of step 2. Why some destinations have n route in the routing tables?	
ŗ		After executing step 3, determine how long it took the network to exchange link state and adjust routing tables. (Hint: you can calculate the required time by observing the time PF update message and the last ACK from Wireshark).	
	2.2.3 Explain t	After execution of step 4, identify the new routing table and router database at router n2. the updates in the new routing table and the new database.	
2.3	OSI	PF Link State Advertisement Periodicity	20

2.3.1	Bring router n4 down and then up again and determine how long it took the network to
recognize	the router is down/up
2.3.2	Now, the following is a bonus question (extra marks). Explain the noticed behavior of the
link state	update packet storms (flooding) where it generally takes longer times for the LSP exchange
to die out	(i.e. LSP exchange stops) when a router is down as compared with the case when the
router is u	ıp again22

1. TCP Lab

Each of the following sections describes an experiment based on the network topology provided by the Project_OSPF_TCP.imn file. In your report, divide the answers to correspond to one of these sections in exactly the same sequence and same name, for example use Question 2.1, Question 3.2 and so on. The Introductory Lab explained and introduced most of the needed tools like iperf3, vtysh, wireshark, etc. If you have not gone through it, do not start working on the assignment below. It is better to start the network emulation once and run Wireshark captures. You can reset Wireshark captures from experiment to another.

1.1. Effect of TCP Window Size

- Set wireshark filter to display TCP packets only.
- Start an iperf3 server on node n11.
- Start an iperf3 client on node n7 connecting to server on node n11 for a duration of 40 seconds and reporting interval of 10. Note that in iperf, the client is the node sending the traffic and the server simply receives and sends an ACK.
- 1.1.1. Vary iperf3 window size from 1Kbytes to 6 Kbytes in increments of 1 Kbyte, then set it to 12, 16, 24, 32 Kbytes. Plot the average throughput and the average number of retransmissions as function of window size. Comment on the results and explain the zigzag behaviour noticed for larger window sizes. Note that number of retransmissions is the 5th column in iperf3 default output & throughput is the 4th. You can plot data using matlab/python/any online graphing tool (Ex. https://chartstudio.plotly.com/create/#/). The window size is the last option in the iperf command (return to INTRO Lab).

Window size (Kbytes)	Average number of transmissions	Average throughput (Mbytes/sec)
1k	0	0.123375
2k root@n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 2k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40781 connected to 10.0.13.20 port 5201 [ID] Interval	0	0.3175
3k root@n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 3k Connecting to host 10.0.13.20. port 5201 [4] local 10.0.12.20 port 40783 connected to 10.0.13.20 port 5201 [III] Interval	0	0.6375
### Toot@n7:/tmp/pycore.51112/n7.conf# iperf3 =c 10.0.13.20 =t 40 =i 10 =w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40785 connected to 10.0.13.20 port 5201 [III] Interval	0	0.64

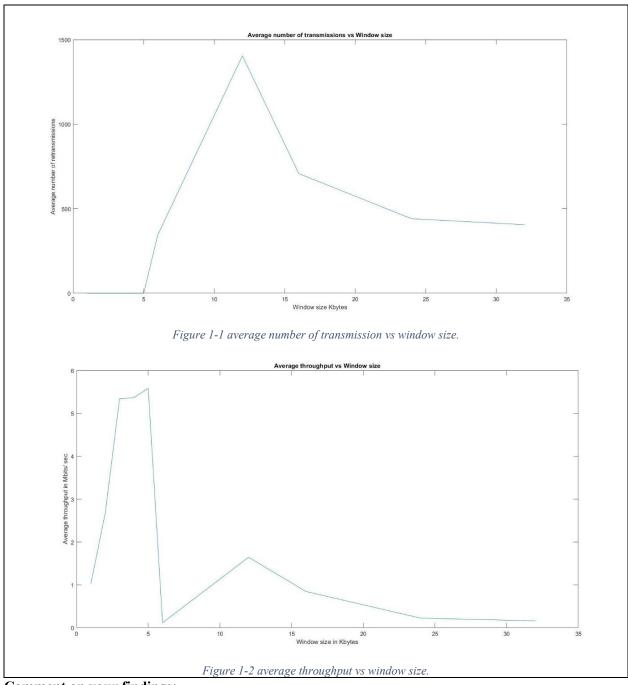
5k		
root@n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 5k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40787 connected to 10.0.13.20 port 5201 [ID] Interval Transfer Bandwidth Retr Cwnd [4] 0.00-10.00 sec 5.99 MBytes 5.80 Mbits/sec 0 14.1 KBytes [4] 10.00-20.00 sec 6.92 MBytes 5.80 Mbits/sec 0 14.1 KBytes [4] 20.00-30.00 sec 6.81 MBytes 5.72 Mbits/sec 0 14.1 KBytes [4] 30.00-40.00 sec 6.91 MBytes 5.80 Mbits/sec 0 14.1 KBytes [4] 30.00-40.00 sec 6.91 MBytes 5.80 Mbits/sec 0 sender [4] 0.00-40.00 sec 26.6 MBytes 5.59 Mbits/sec 0 sender [4] 0.00-40.00 sec 26.6 MBytes 5.58 Mbits/sec receiver iperf Done.	0	0.665
6k root@n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 6k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40789 connected to 10.0.13.20 port 5201 [IID] Interval	345	0.0148038
12k	1405	0.196
16k root@n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 16k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40793 connected to 10.0.13.20 port 5201 [III] Interval	708	0.100875

```
f3 -c 10.0.13.20 -t 40 -i 10 -w 24k
                          port 40795 connected to 10.0.13.20 port 5201
                                                                   Cwnd
7.07 KBytes
5.66 KBytes
                                            254 Kbits/sec
                                                Kbits/sec
                                                                   8.48 KBytes
7.07 KBytes
                                                Kbits/sec
                                                                                                        441
                                                                                                                                  0.027
        30,00-40,00
                             277 KBytes
                                            227 Kbits/sec
      Interval
                                           Bandwidth
         0,00-40,00 sec
                            1.10 MBytes
1.06 MBytes
                                            230 Kbits/sec
                                                                               sender
                                            221 Kbits/sec
         0.00-40.00 sec
                                                                               receiver
iperf Done.
                                           32k
      ting to host 10.0.13.20, port
       local 10.0.12.20 port 40797 connected to 10.0.13.20 port 5201
                                                                    Cwnd
7.07 KBytes
12.7 KBytes
      Interval
                                           Bandwidth
                                                             Retr
                                            202 Kbits/sec
155 Kbits/sec
                             246 KBytes
         0,00-10,00
                                                               96
        10,00-20,00
                             189 KBytes
                                            132 Kbits/sec
175 Kbits/sec
                                                                         KBytes
                                  KBytes
                                                                                                        406
                                                                                                                                0.01915
                                                             112
                             214 KBytes
                                           Bandwidth
                                                             Retr
                             810 KButes
                                            166 Kbits/sec
                      sec
                                                                                sender
         0,00-40,00 sec
                             759 KButes
                                                                                receiver
```

From the above data the following code has written to plot the average number of retransmission and throughput vs the window size

```
%setting the window size values
window size = [1 2 3 4 5 6 12 16 24 32];
%setting the average retransmissions values
avg retr = [0 \ 0 \ 0 \ 0 \ 345 \ 1405 \ 708 \ 441 \ 406];
%setting the average throughput values
avg throughput = [1.03 2.67 5.34 5.37 5.585 0.118 1.645 ...
    0.8465 0.2255 0.161];
%plotting average retransmissions vs window size
figure
plot(window size, avg retr)
xlabel("Window size Kbytes");
ylabel("Average number of retransmissions");
title("Average number of transmissions vs Window size");
%plotting average throughput vs window size
figure
plot(window size, avg throughput)
xlabel("Window size in Kbytes");
ylabel("Average throughput in Mbits/ sec");
title("Average throughput vs Window size");
```

This code generates the following two figures:



Comment on your findings:

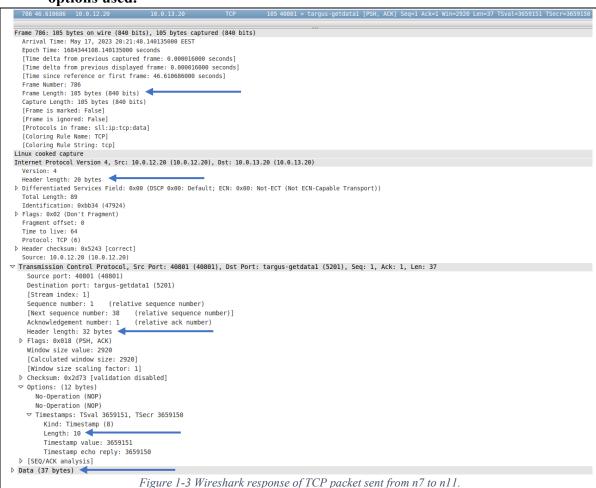
Firstly, when the window size is from 1 Kbytes to 6 Kbytes, the retransmission packets are equal to zero, meaning that the receiver buffer is not filled yet, therefore by increasing the window size the throughput increases until reaching 6 Kbytes window.

At 6 Kbytes window, the retransmission packets increase as the buffer is filled and the packets are lost as the TCP uses Go Back N protocol, therefore it retransmits the whole window which increases the congestion and decreases the throughput, then the receiver must upgrade the buffer size to handle the new window size.

By increasing the window size to 12 Kbytes, the retransmission packets reduce, and the throughput increases until the buffer is filled so the retransmission increases, and the throughput decreases again which form the zigzag shape.

As the window size increases, the chance to retransmit a packet increases as the window size has more packets, then the chance to retransmit the window again will increase which make the throughput decreases at the end of the graph comparing to the initial values.

1.1.2. Click on any of the TCP a data segment whose source is node n7, dissect the segment by following different protocol layer headers from TCP->IP->Ethernet identifying how many header bytes are added by each layer, identify the TCP options used.



Comment on Your findings:

- TCP header length = 32 bytes (including options: timestamp (10 bytes) + 2 bytes No Operation for padding the options and making the total header multiple of 32 bits)
- IP header length = 20 bytes
- Data length = 37 bytes
- Ethernet length = Frame length − TCP header length − IP header length Data = 105 − 32 − 20 − 37 = 16 bytes

1.1.3. Repeat for an ACK packet sent from node n11.

```
▽ Frame 705: 68 bytes on wire (544 bits), 68 bytes captured (544 bits)
     Arrival Time: May 17, 2023 20:21:48.087658000 EEST 
Epoch Time: 1684344108.087658000 seconds
      [Time delta from previous captured frame: 0.000025000 seconds]
      [Time delta from previous displayed frame: 0.000025000 seconds]
      [Time since reference or first frame: 46.558209000 seconds]
     Frame Number: 705
Frame Length: 68 bytes (544 bits)
      Capture Length: 68 bytes (544 bits)
      [Frame is marked: False]
      [Frame is ignored: False]
[Protocols in frame: sll:ip:tcp]
      [Coloring Rule Name: TCP]
[Coloring Rule String: tcp]
Dinux cooked capture
 ▼ Internet Protocol Version 4, Src: 10.0.13.20 (10.0.13.20), Dst: 10.0.12.20 (10.0.12.20)
      Version: 4
   Differentiated Services Field: 0x00 (DSCP 0x00: Default; ECN: 0x00: Not-ECT (Not ECN-Capable Transport))
      Total Length: 52
      Identification: 0x348c (13452)
   ▶ Flags: 0x02 (Don't Fragment)
Fragment offset: 0
     Time to live: 64
Protocol: TCP (6)
   ▶ Header checksum: 0xd910 [correct]
      Source: 10.0.13.20 (10.0.13.20)
     Destination: 10.0.12.20 (10.0.12.20)
▽ Transmission Control Protocol, Src Port: targus-getdata1 (5201), Dst Port: 40000 (4000), Seq: 1, Ack: 38, Len: 0
    Source port: targus-getdata1 (5201)
Destination port: 40800 (40800)
    [Stream index: 0]
     Sequence number: 1
                             (relative sequence number)
    Acknowledgement number: 38
                                     (relative ack number)
    Header length: 32 bytes 🔫
  ▶ Flags: 0x010 (ACK)
    Window size value: 905
[Calculated window size: 14480]
    [Window size scaling factor: 16]
Checksum: 0x2d4e [validation disabled]
  ▽ Options: (12 bytes)
       No-Operation (NOP)
    No-Operation (NOP)

▼ Timestamps: TSval 3659138, TSecr 3659137
         Kind: Timestamp (8)
          Length: 10
         Timestamp value: 3659138
  ▷ [SEQ/ACK analysis]
                                           Figure 1-4 Wireshark response of TCP packet sent from n11 to n7.
```

Comment on your findings here:

- TCP header length = 32 bytes (including options: timestamp (10 bytes) + 2 bytes No Operation for padding the options and making the total header multiple of 32 bits)
- IP header length = 20 bytes
- Data length = 0 bytes (just an ACK)
- Ethernet length = Frame length TCP header length IP header length = 68 32 20 = 16 bytes

1.2. TCP short VS long paths

- Run an iperf3 server on node n8.
- Start an iperf3 client on node n7 connecting to server on node n8 for a duration of 40 seconds and reporting interval of 10 with window size 4K.
- 1.2.1. Compare the result of throughput with the case when connection was made to node n11. Why throughput drops when connecting to n8 although capacities on the two paths are the same.



Comment in your findings here:

The throughput of sending from n7 to n8 (2.18 Mbits/sec) is much smaller than the throughput of sending from n7 to n11 (5.37 Mbits/sec). The path from n7 to n8 is larger than the path from n7 to n11, hence there is more propagation delays with larger RTT and more cost on the path which leads to less throughput.

1.3. Higher link Capacity with Drops VS Reliable Lower Capacity

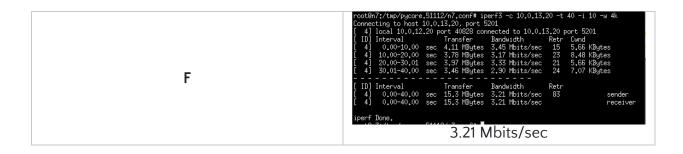
This guestions for this part are based on the path between n7 and n11.

For each case of the following, run iperf3 client from n7 to n11 with window size 4K.

- A Select link between n4 and n5, configure it to have capacity of 10 Mbps with zero loss in both directions.
- B Select link between n4 and n5, configure it to have capacity of 3 Mbps with zero loss in both directions.
- C Select link between n4 and n5, configure it to have capacity of 10 Mbps with 5% loss in both directions.
- D Select link between n4 and n5, configure it to have capacity of 100 Mbps with 10% loss in both directions.
- E Select link between n4 and n5, configure it to have capacity of 10 Mbps with 1% loss in direction from n4 to n5 and 0% loss in the other direction.
- F Select link between n4 and n5, configure it to have capacity of 10 Mbps with 0% loss in direction from n4 to n5 and 1% loss in the other direction.

Please fill the following table with the AVG throughput of each case:

Case	Throughput
A	root8n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40811 connected to 10.0.13.20 port 5201 [III] Interval [4] 0.00-10.00 sec 5.94 MBytes 4.98 Mbits/sec 0 14.1 KBytes [4] 10.00-20.00 sec 6.48 MBytes 5.48 Mbits/sec 0 14.1 KBytes [4] 10.00-30.00 sec 7.35 MBytes 5.16 Mbits/sec 0 14.1 KBytes [4] 20.00-30.00 sec 7.58 MBytes 6.18 Mbits/sec 0 14.1 KBytes [4] 30.00-40.00 sec 7.58 MBytes 6.37 Mbits/sec 0 14.1 KBytes [4] 30.00-40.00 sec 7.58 MBytes 6.37 Mbits/sec 0 14.1 KBytes [1II] Interval
В	root8n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40014 connected to 10.0.13.20 port 5201 [III] Interval
C	root@n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40816 connected to 10.0.13.20 port 5201 [III] Interval
D	rootBm7:/tmp/pycore.51112/n7.conf# iperf3 -c 10,0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10,0.12.20 port 40618 connected to 10.0.13.20 port 5201 [ID] Interval Transfer Bandwidth Retr Cund [4] 0.00-10.00 sec 198 KBytes 162 Kbits/sec 24 4.24 KBytes [4] 10.00-20.00 sec 138 KBytes 114 Kbits/sec 15 4.24 KBytes [4] 20.00-30.00 sec 106 KBytes 86.9 Kbits/sec 16 5.66 KBytes [4] 30.00-40.00 sec 218 KBytes 178 Kbits/sec 24 2.83 KBytes [1D] Interval Transfer Bandwidth Retr [4] 0.00-40.00 sec 660 KBytes 135 Kbits/sec 79 sender [4] 0.00-40.00 sec 659 KBytes 135 Kbits/sec receiver iperf Done. 135 Kbits/sec
E	root8n7:/tmp/pycore.51112/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 40830 connected to 10.0.13.20 port 5201 [ID] Interval



1.3.1. Compare throughputs in cases a, b, c. Why b is better than c?

	<i>O</i> 1	•	
	A	В	C
Throughput	5.74 Mbits/sec	2.73 Mbits/sec	449 Kbits/sec

Comment in your findings here:

- The throughput of case A is better than case B as the capacity of case A is 10 Mbps while case B is 3 Mbps.
- Although, the capacity of C is better than the capacity of B (10 Mbps to 3 Mbps) but case C has losses which will make retransmission in our system and decreases its throughput making the throughput of B is better than the throughput of C.

1.3.2. Compare throughputs in cases b, c and d. Which is better? Why?

	В	\mathbf{C}	D
Throughput	2.73 Mbits/sec	449 Kbits/sec	135 Kbits/sec

Comment in your findings here:

- Comparing the throughput of the 3 cases, we can observe that case B is better as it has no losses so there is no retransmission, while case C and D have losses which will make their throughput affected by the retransmission.
- Although, the capacity of D is better than the capacity of C (100 Mbps to 10 Mbps) but case D has
 more losses which will make retransmission in our system much worse and decreases its throughput
 more making the throughput of C is better than the throughput of D.
 Overall, B > C > D.

1.3.3. Compare throughputs in cases e and f? Which is better? Why?

	E	${f F}$
Throughput	3.94 Mbits/sec	3.21 Mbits/sec

Comment in your findings here:

The transmission is from n7 to n11, using the data path from n5 to n4.

- In case E, the losses in the direction of n4 to n5, so the losses will be only in the ACK sent from n11 to n7 to retransmit the data if doesn't arrive to n7 and as TCP using Go Back N protocol with accumulative ACK, it won't affect the losses as last ACK received at transmitter means that all the packets with sequence number less than the one received with ACK has been received at receiver while all the transmitted and retransmitted data use path from n5 to n4 which has no losses.
- In case F, the losses in the direction of n5 to n4, so the losses will affect all the transmitted and retransmitted data sent from n7 to n11 while the ACK use path from n4 to n5 which has no losses.

From case E and F, we can conclude that the throughput of E will be better as all the data sent doesn't affect by losses.

2. OSPF Lab

Each of the following sections describes an experiment based on the network topology provided by the **Project OSPF TCP.imn** file.

The <u>Introductory Lab</u> explained and introduced most of the needed tools like iperf3, vtysh, wireshark, etc. If you have not gone through it, do not start working on the assignment below.

It is better to start the network emulation once and run Wireshark captures. You can reset Wireshark captures from experiment to another.

2.1. OSPF Link Cost Change

- 1. Stop any running iperf3 clients.
- 2. Set all links to have zero loss in the two directions with 100 Mbps speed.
- 3. Run iperf3 between n7 and n11 for a duration of 500 seconds or longer. Identify the path between n7 and n11.
- 4. Open vtysh on node n5 by opening a bash terminal and typing vtysh. Now we can configure router and link costs.
- 5. Type the following in vtysh: show ip ospf interface eth1 (To know the ethernet number of each interface you can choose toolbar->view->show->interface names) This displays information about interface eth1. Note its ospf cost.
- 6. Type the following configure terminal
 - interface eth1 (the interface for the link between n5 and n4)
 - ospf cost 40 (set cost to 40)

2.1.1. Check what happens to the path between n7 and n11 (as seen after steps 3 and 6)? Explain what happens. To help visualize the path choose toolbar->widgets-> Throughput.

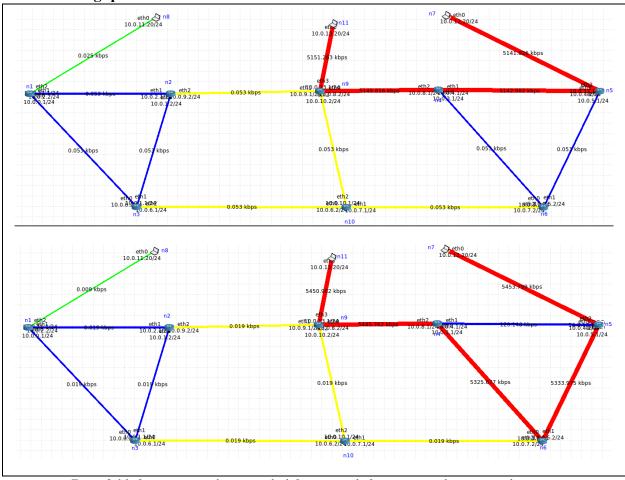
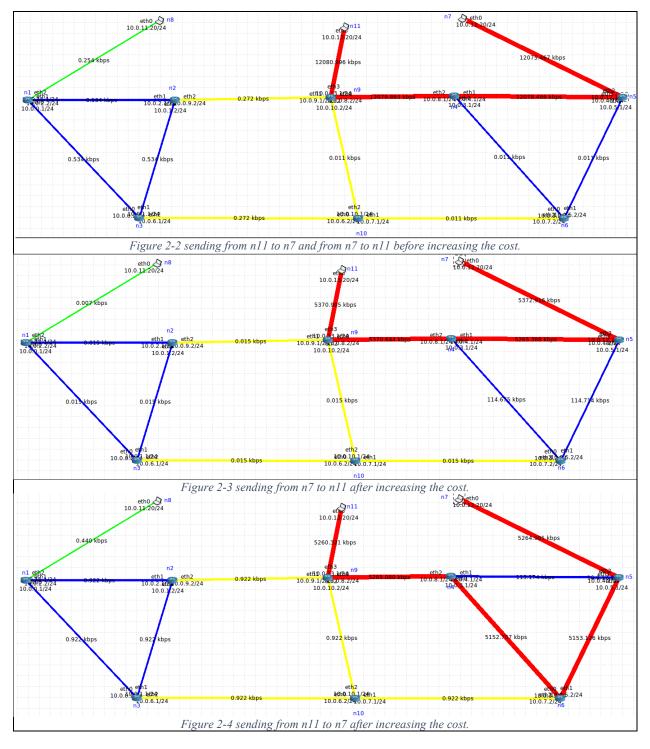


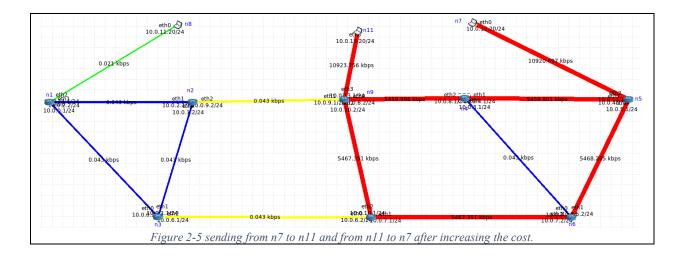
Figure 2-1 before increasing the cost on the left picture and after increasing the cost on right picture.

Comment on your findings:

Before increasing the cost of ethernet 1, OSPF protocol choose the shortest path according to the cost where it moves from n5 to n4 through eth1 directly which has a cost = 10 that represents the minimum path cost, while after increasing the cost of this link to 40, now, the cost of moving from n5 to n6 through ethernet 0 has cost = 10 and then from n6 to n4 has cost = 10 which has a total cost from n5 to n4 equal to 20 instead of moving directly by ethernet 1 from n5 to n4 with cost 40. So, most of the data will move in the new shortest path.

2.1.2 Set the cost of eth1 at node n5 back to 10. Establish two iperf3 connections: one from n7 to n11 and the second from n11 to n7 both for duration of 500 seconds. Now go to node n4 and set interface cost for interface connecting n4 with n5 to 40. What happens in the paths of the two connections? Explain what happens. What do you conclude?





Comment on your findings:

Before increasing the cost, sending in both direction is the same as shown in figure 2-2 as it is the lowest cost path for both directions, while after increasing the cost, sending from n7 to n11 will not be affected as the connection from n5 to n4 (eth1) didn't change as shown in figure 2-3, but sending from n11 to n7 will go from n4 to n6 then n5 instead of n4 to n5 directly as the cost of n4 has been increased to 40 as shown in figure 2-4. If both directions are active at the same time, sending from n7 to n11 will be the same as illustrated before, while sending from n11 to n7 will use the node n10 to avoid using the same path of sending from n7 to n11 as these two paths have the same cost as shown in figure 2-5.

2.2 OSPF Database Updates

- 1. Start wireshark and have its filter to capture only OSPF related packets.
- 2. Go to another router say router n2, open vtysh and issue the commands:
 - config terminal
 - show ip ospf database (check this link to understand more)
 - show ip ospf route
- 3. On router n4, set link cost of eth1 to 20. Capture the link state packets advertised in Wireshark.
- 4. Go to router n4, go out of vtysh, or open new bash terminal and issue the following commands to disconnect router n4 from the network:
 - ifconfig eth0 down
 - ifconfig eth1 down
 - ifconfig eth2 down

2.2.1 Capture and explain the outputs due to execution of step 2. Why some destinations have more than route in the routing tables?

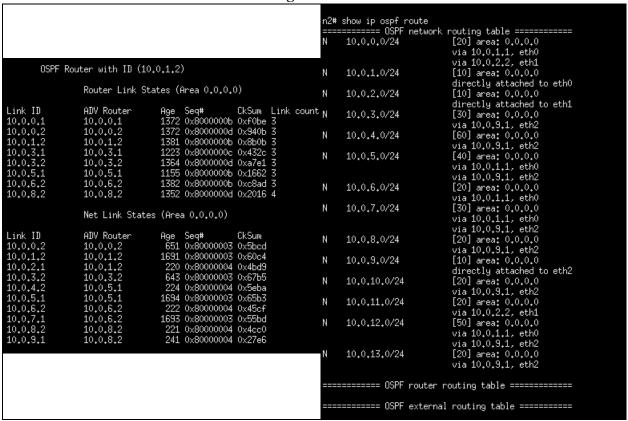


Figure 2-6 left picture represents database and right picture represents routing table.

Comment on your findings:

From the data shown in the database and routing table as shown in figure 2-6, we can observe that the routing table gets the shortest path to the destinations. Knowing that the cost of all links is the same which is equal to 10 except eth1 of node 4 has cost of 40, hence some destinations can have more than one path which have the same total cost, so the router can send the packets in either path of them.

2.2.2 After executing step 3, determine how long it took the network to exchange link state packets and adjust routing tables. (Hint: you can calculate the required time by observing the time of first OSPF update message and the last ACK from Wireshark).

	0				
683764 48.60068	0 10.0.12.1	224.0.0.5	0SPF	80 Hello Packet	
683765 48.60068	3 10.0.12.1	224.0.0.5	0SPF	80 Hello Packet	
683768 48.60134	2 10.0.5.1	224.0.0.5	0SPF	84 Hello Packet	
683769 48.60134	4 10.0.4.2	224.0.0.5	0SPF	84 Hello Packet	
683770 48.60134	6 10.0.12.1	224.0.0.5	0SPF	80 Hello Packet	
704896 50.11212	2 10.0.3.1	224.0.0.5	0SPF	124 LS Update	
704897 50.11212	25 10.0.3.1	224.0.0.5	0SPF	124 LS Update	
704898 50.11213	35 10.0.4.1	224.0.0.5	0SPF	124 LS Update	
704899 50.11213	88 10.0.4.1	224.0.0.5	0SPF	124 LS Update	
704900 50.11214	5 10.0.8.1	224.0.0.5	0SPF	124 LS Update	
704901 50.11214	6 10.0.8.1	224.0.0.5	0SPF	124 LS Update	
704903 50.11227	9 10.0.8.1	224.0.0.5	0SPF	124 LS Update	
704909 50.11233	8 10.0.9.1	224.0.0.5	0SPF	124 LS Update	
704910 50.11234	1 10.0.9.1	224.0.0.5	0SPF	124 LS Update	
704911 50.11234	8 10.0.10.2	224.0.0.5	0SPF	124 LS Update	
704912 50.11235	0 10.0.10.2	224.0.0.5	0SPF	124 LS Update	
709143 50.460007	10 0 10 1	224 0 0 5	OSPF 80 LS	A cknowledge	
709145 50.460293		224.0.0.5 224.0.0.5		Acknowledge Acknowledge	
709146 50.460294		224.0.0.5		Acknowledge	
709147 50.460295		224.0.0.5		Acknowledge	
709149 50.460323		224.0.0.5		Acknowledge	
709190 50.464107	10.0.8.2	224.0.0.5		Acknowledge	
709191 50.464118	10.0.8.2	224.0.0.5	OSPF 80 LS	Acknowledge	
709198 50.464292	10.0.8.2	224.0.0.5	OSPF 80 LS	Acknowledge	
758796 53.882216	fe80::200:ff:feaa:d	ff02::5	OSPF 96 Hel	lo Packet	
	fe80::200:ff:feaa:d			lo Packet	
	fe80::200:ff:feaa:e			lo Packet	
758799 53.882236	fe80::200:ff:feaa:e	ff02::5	OSPF 96 Hel	lo Packet	

Figure 2-7 Wireshark response of OSPF packets after a changing the cost in the network.

Comment on your findings:

Initially, the cost of eth1 of node 4 equals to 40, the routers only send hello packets until any change happens. After changing the cost to 20, the routers must send the new link state update to let the network know that there is some change happened and adjust its routing table according to the new shortest path for every destination. The time taken to spread the update = time of last ACK – time of first update = 50.464292 - 50.112122 = 0.35217 sec.

2.2.3 After execution of step 4, identify the new routing table and router database at router n2. Explain the updates in the new routing table and the new database.

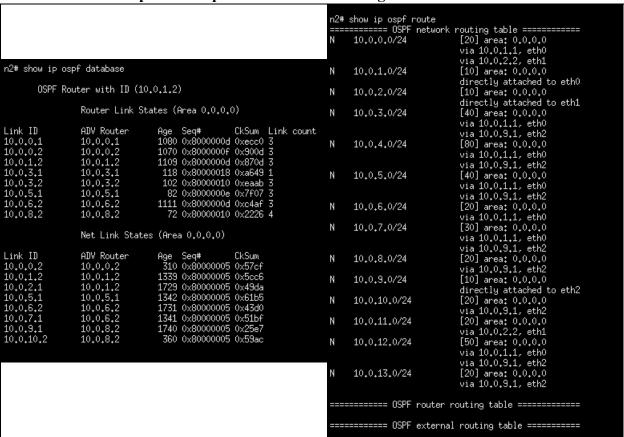


Figure 2-8 left picture represents database and right picture represents routing table.

Comment on your findings:

- In the database table, the number of links connected to 10.0.3.1 reduced from 3 to 1 as n4 is disconnected, therefore 10.0.3.1 only has link to 10.0.3.2 instead of 10.0.3.2 & 10.0.4.1 & 10.0.8.1.
- In the routing table, the cost of 10.0.3.0 increased from 30 to 40 & the cost of 10.0.4.0 increased from 40 to 60, as after disconnecting n4 some links will get another shortest path which will be with higher cost than with connecting n4.

2.3 OSPF Link State Advertisement Periodicity

This part is on your own. No instructions are given. Bring router n4 ethernet interfaces back up.

Stop any previous iperf3 connections. Let Wireshark capture focus on Ospfv2 only (IPv4).

Apply the filter (ospf.msg.lsupdate || ospf.msg.lsack) && ip.addr == 224.0.0.5

2.3.1 Bring router n4 down and then up again and determine how long it took the network to recognize the router is down/up.

Filter:	msg.lsupdat	e ospf.msg.lsack) &	& ip.addr == 224.0.0.5	Expression	Clear Apply
No.	Time	Source	Destination	Proto	ocol Length Info
2888	3 100.460902	10.0.8.1	224.0.0.5	OSPF	172 LS Update
2889	9 100.460911	10.0.8.1	224.0.0.5	0SPF	172 LS Update
2898	3 100.461842	10.0.9.1	224.0.0.5	0SPF	100 LS Update
2899	9 100.461844	10.0.10.2	224.0.0.5	0SPF	100 LS Update
2900	0 100.461926	10.0.1.2	224.0.0.5	0SPF	100 LS Update
2902	2 100.461932	10.0.1.2	224.0.0.5	0SPF	100 LS Update
2904	4 100.461945	10.0.2.1	224.0.0.5	0SPF	100 LS Update
2905	5 100.461948	10.0.2.1	224.0.0.5	0SPF	100 LS Update
2906	5 100.462008	10.0.6.2	224.0.0.5	0SPF	100 LS Update
					,
	139.886829 1		224.0.0.5	0SPF	120 LS Acknowledge
	139.886916 1		224.0.0.5	0SPF	100 LS Acknowledge
	139.886920 1		224.0.0.5	0SPF	100 LS Acknowledge
3739	139.886922 1	0.0.1.1	224.0.0.5	0SPF	80 LS Acknowledge
3740	139.886926 1	0.0.6.1	224.0.0.5	0SPF	120 LS Acknowledge
3741	139.887094 1	0.0.9.2	224.0.0.5	0SPF	100 LS Acknowledge
3742	139.887108 1	0.0.9.2	224.0.0.5	0SPF	100 LS Acknowledge
3743	139.887269 1	0.0.1.2	224.0.0.5	0SPF	100 LS Acknowledge
3744	139.887280 1	0.0.9.2	224.0.0.5	0SPF	100 LS Acknowledge

Figure 2-9 Wireshark response of OSPF packets after changing the state of n4 from up to down.

Filter:	msg.lsupda	ate ospf.msg.lsa	ck) && ip.addr == 224.0.0.5	Expression Cl	lear Apply
No.	Time	Source	Destination	Protoco	l Length Info
25	9 15.828157	10.0.8.1	224.0.0.5	0SPF	112 LS Update
26	0 15.828164	10.0.8.1	224.0.0.5	0SPF	112 LS Update
26	2 15.828387	10.0.8.1	224.0.0.5	0SPF	112 LS Update
26	5 15.828569	10.0.9.1	224.0.0.5	0SPF	112 LS Update
26	6 15.828578	10.0.9.1	224.0.0.5	0SPF	112 LS Update
26	7 15.828593	10.0.10.2	224.0.0.5	0SPF	112 LS Update
26	8 15.828597	10.0.10.2	224.0.0.5	0SPF	112 LS Update
27	0 15.828917	10.0.9.1	224.0.0.5	0SPF	112 LS Update
27	2 15.828920	10.0.10.2	224.0.0.5	0SPF	112 LS Update
27	3 15.829051	10.0.1.2	224.0.0.5	0SPF	112 LS Update
27	4 15.829060	10.0.1.2	224.0.0.5	0SPF	112 LS Update
27	5 15.829076	10.0.2.1	224.0.0.5	0SPF	112 LS Update
27	6 15.829082	10.0.2.1	224.0.0.5	0SPF	112 LS Update
82	1 26.066780	10.0.7.2	224.0.0.5	0SPF	80 LS Acknowledge
82	2 26.100484	10.0.8.2	224.0.0.5	0SPF	80 LS Acknowledge
82	3 26.100530	10.0.8.2	224.0.0.5	0SPF	80 LS Acknowledge
82	4 26.100793	10.0.8.2	224.0.0.5	0SPF	80 LS Acknowledge
82	5 26.160029	10.0.5.1	224.0.0.5	0SPF	80 LS Acknowledge
	6 26.160085	10.0.5.1	224.0.0.5	0SPF	80 LS Acknowledge
	7 26.160338	10.0.5.1	224.0.0.5	0SPF	80 LS Acknowledge
	8 26.269086	10.0.0.1	224.0.0.5	0SPF	80 LS Acknowledge
	9 26.269116	10.0.0.1	224.0.0.5	0SPF	80 LS Acknowledge
83	0 26.269729	10.0.0.1	224.0.0.5	0SPF	80 LS Acknowledge

Figure 2-10 Wireshark response of OSPF packets after changing the state of n4 from down to up.

Comment on your findings:

Initially router n4 is up, then after disconnected all its interfaces and bring them down, the network sends the new link state packets to update the routing table, the time taken to spread the update = 139.887280 - 100.460902 = 39.426378 as shown in figure 2-9.

Bringing back router n4 by reconnecting all its interface, the time taken to spread the update = 26.269729 - 15.828157 = 10.441572 sec.

2.3.2 Now, the following is a bonus question (extra marks). Explain the noticed behavior of the link state update packet storms (flooding) where it generally takes longer times for the LSP exchange to die out (i.e. LSP exchange stops) when a router is down as compared with the case when the router is up again.

```
root@n5:/tmp/pycore.51567/n5.conf# vtysh

Hello, this is Quagga (version 0.99,20.1).
Copyright 1996-2005 Kunihiro Ishiguro, et al.

n5# show ip ospf interface eth1
eth1 is up
ifindex 123, MTU 1500 bytes, BW 0 Kbit <UP,BROADCAST,RUNNING,MULTICAST>
Internet Address 10.0,4,2/24, Area 0.0,0.0

MTU mismatch detection;enabled
Router ID 10.0,5.1, Network Type BROADCAST, Cost: 10
Transmit Delay is 1 sec, State DR, Priority 1
Designated Router (ID) 10.0,5.1, Interface Address 10.0,4.2
Backup Designated Router (ID) 10.0,3.1, Interface Address 10.0,4.1
Saved Network-LSA sequence number 0x80000007
Multicast group memberships: OSPFAlRouters OSPFDesignatedRouters
Timer intervals configured, Hello 10s, Dead 40s, Wait 40s, Retransmit 5
Hello due in 4,603s
Neighbor Count is 1, Adjacent neighbor count is 1
```

Figure 2-11 eth1 interface of router n5.

Comment on your findings:

From figure 2-11, the dead time equals to 40 sec, means that to know that the link is disconnected, the node sends packets and wait for the ACK from the other node for 40 sec, if it didn't get anything from that link in this interval of time, the node understands that this link is disconnected, and this can be proved when disconnecting router n4 from the previous question.

Also, the hello time equals to 10 sec, means that when router enters the network, it will send Hello packets for 10 sec to let the network knows its existence and update its routing table, and this can be proved when reconnecting router n4 and bringing it back to network from the previous question.