Converting an NFA to a DFA

Given:

A *non-deterministic* finite state machine (NFA)

Goal:

Convert to an equivalent <u>deterministic</u> finite state machine (DFA)

Why?

Faster recognizer!

Approach:

Consider simulating a NFA.

Work with sets of states.

IDEA: Each *state* in the DFA will correspond to a *set of* NFA states.

Worst-case:

There can be an exponential number $O(2^N)$ of sets of states.

The DFA can have exponentially many more states than the NFA ... but this is rare.

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Lexical Analysis - Part 3

NFA to DFA

<u>Input:</u> A NFA $S = States = \{ s_0, s_1, ..., s_N \} = S_{NFA}$

 $\delta = Move function = Move_{NFA}$

Move'(S, a) \rightarrow Set of states

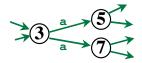
<u>Output:</u> A DFA $S = States = \{?, ?, ..., ?\} = S_{DFA}$

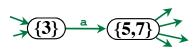
 $\delta = Move function = Move_{DFA}$

Move(s, a) \rightarrow Single state from S_{DFA}

Main Idea: Each state in S_{DFA} will be a set of states from the NFA

 $S_{DFA} = \{ \{...\}, \{...\}, ..., \{...\} \}$

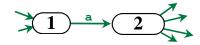




NFA

DFA

(The names of the states is arbitrary and can be changed later, if desired.)



Algorithm: Convert NFA to DFA

We'll use...

the transition function from NFA Move_{NFA} (S, a) where s is a single state from NFA ε-Closure(s) where S is a set of states from NFA ε-Closure(S)

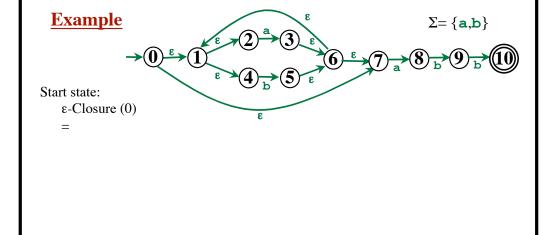
We'll construct...

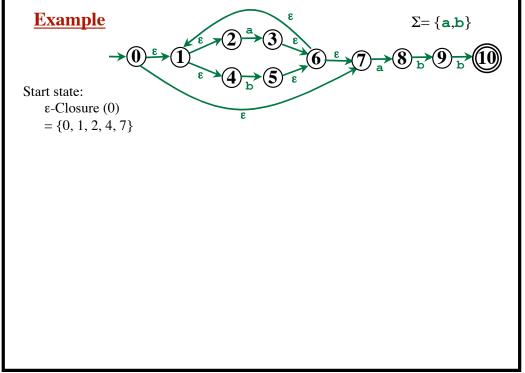
the set of states in the DFA S_{DFA} Initially, we'll set S_{DFA} to $\{\}$ Add \mathbf{X} to \mathbf{S}_{DFA} where \mathbf{X} is some set of NFA states Example: "Add $(\{3,5,7\})$ to S_{DFA} " We'll "mark" some of the states in the DFA. Marked = "We've done this one" ($\sqrt{}$) Unmarked = "Still need to do this one" The transition function from DFA $Move_{DFA}(T,b)$ To add an edge to the growing DFA... Set Move_{DFA}(T,b) to S

...where **S** and **T** are sets of NFA states

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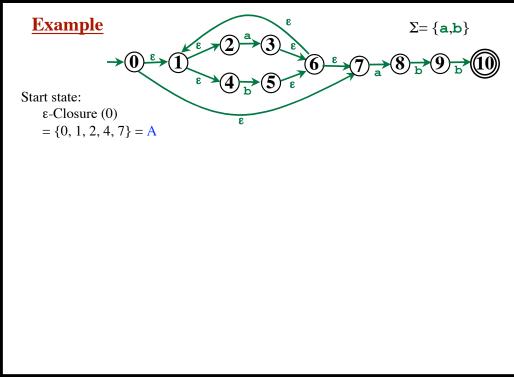
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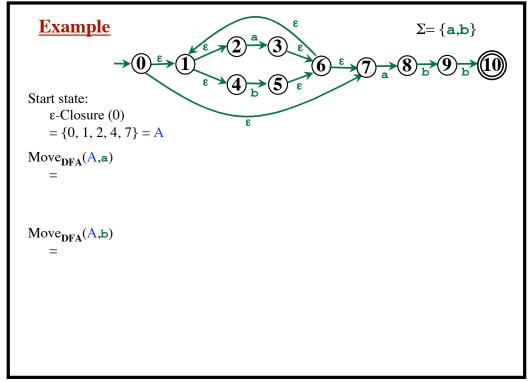




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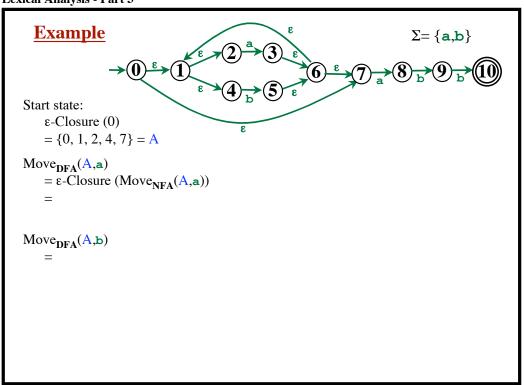
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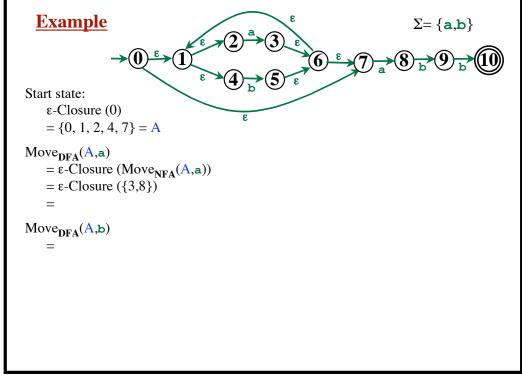




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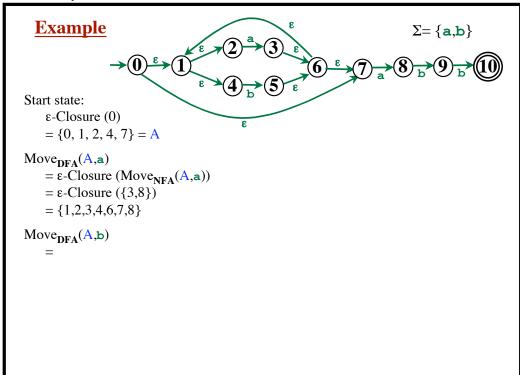
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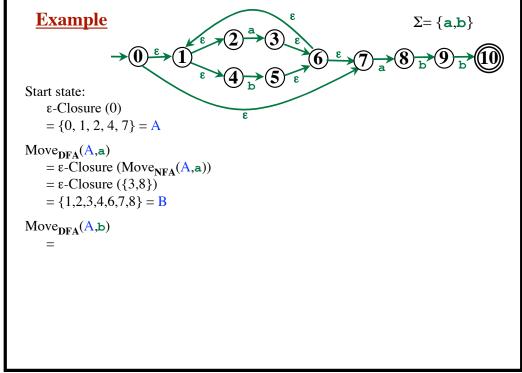




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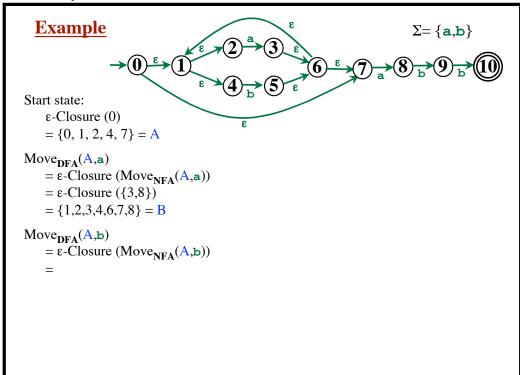
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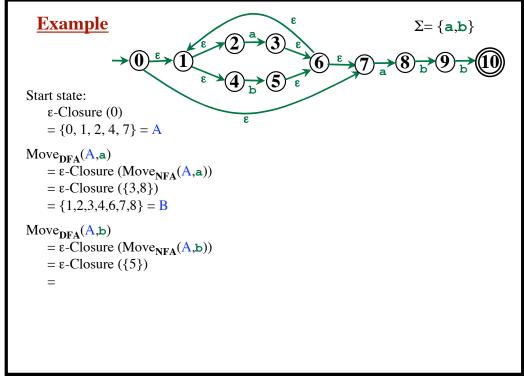




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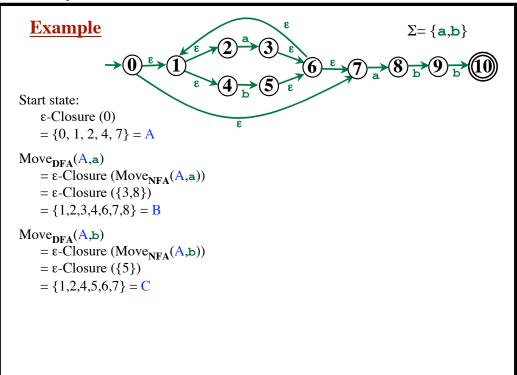


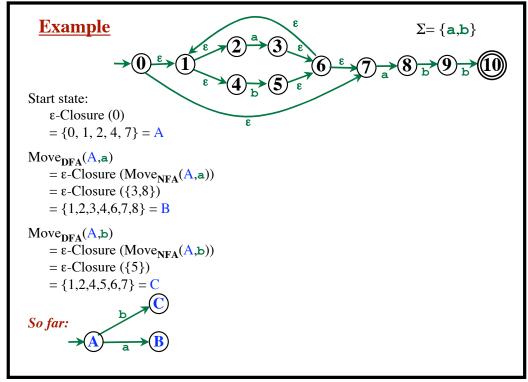


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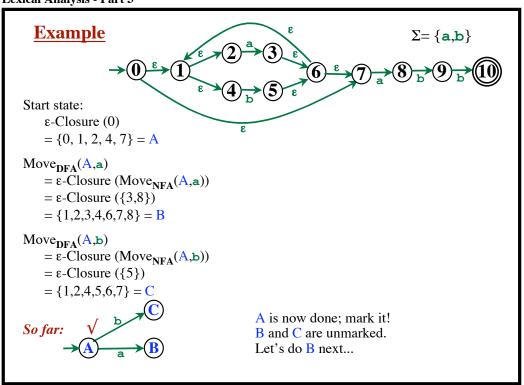
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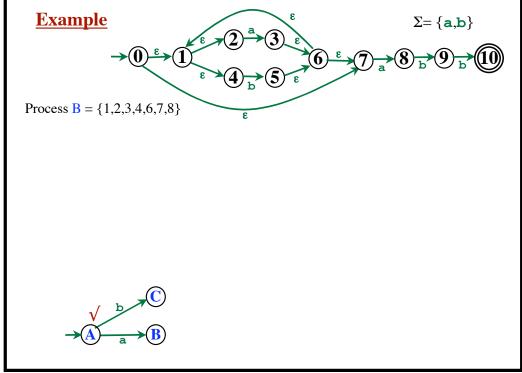




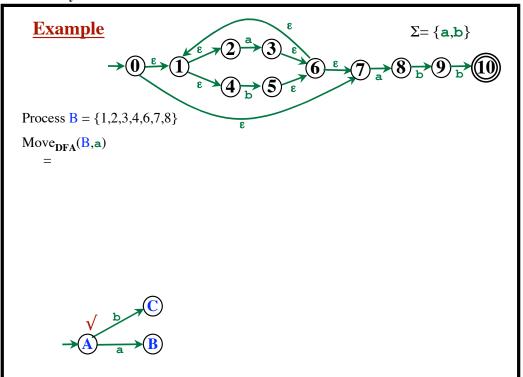
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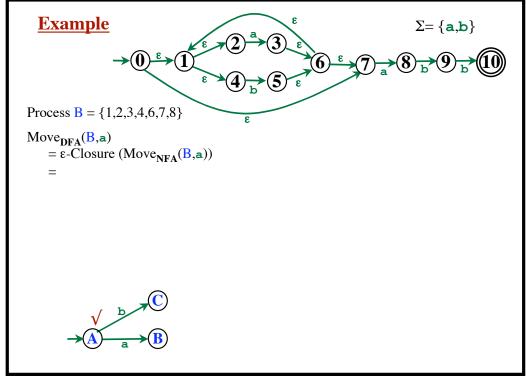
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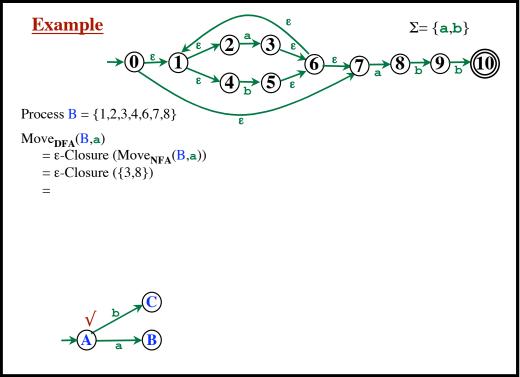
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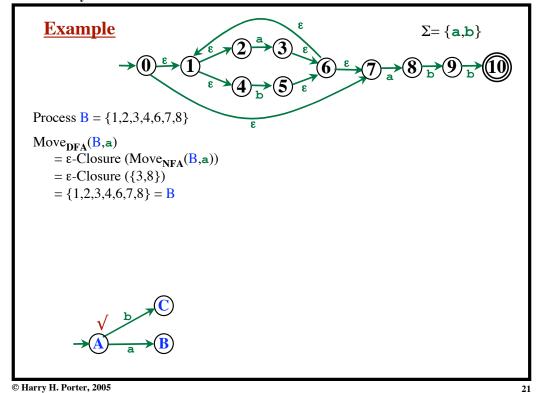


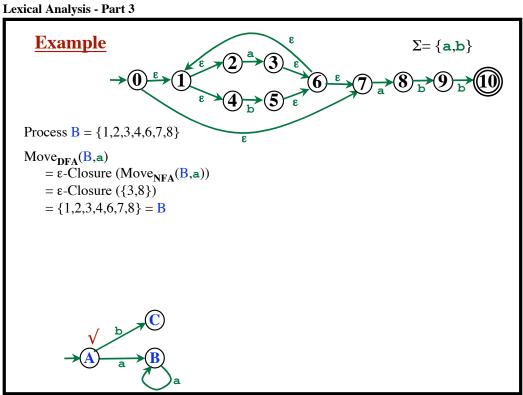


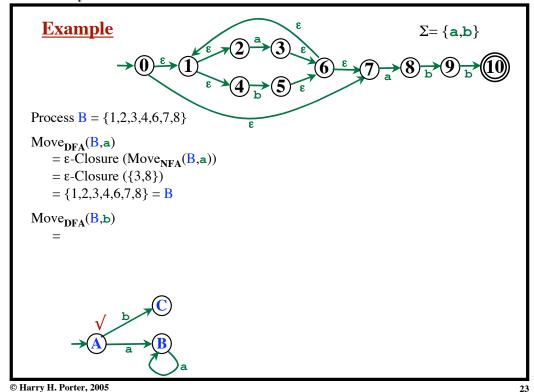
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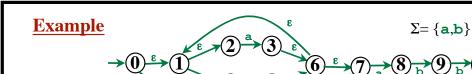
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Process $\mathbf{B} = \{1,2,3,4,6,7,8\}$

 $\mathsf{Move}_{\mathbf{DFA}}({\color{red}\mathbf{B}},\!{\color{blue}\mathbf{a}})$

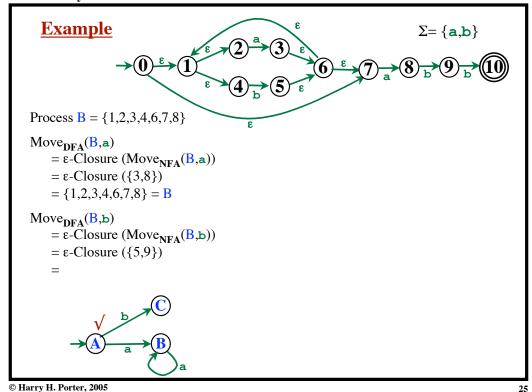
= ε -Closure (Move_{NFA}(B,a))

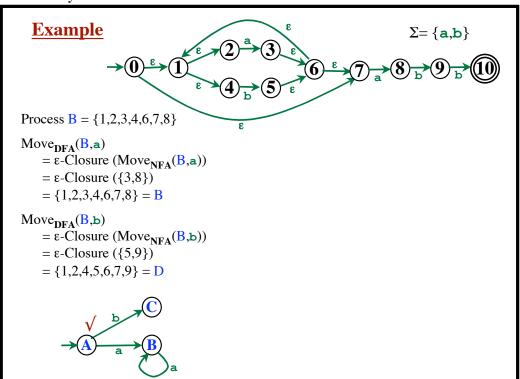
= ε -Closure ({3,8})

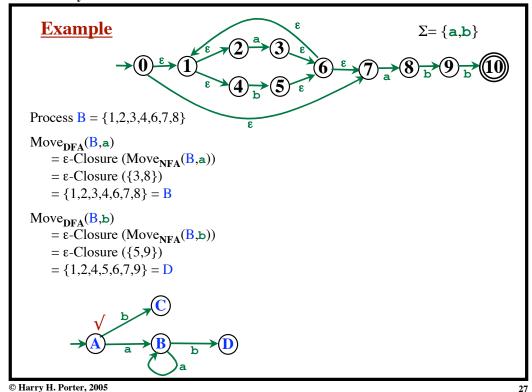
 $= \{1,2,3,4,6,7,8\} = \mathbf{B}$

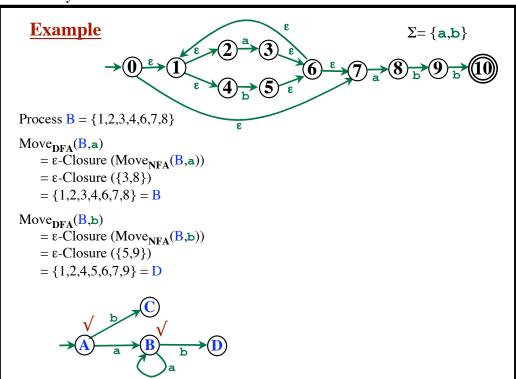
 $Move_{\mathbf{DFA}}(\mathbf{B},\mathbf{b})$

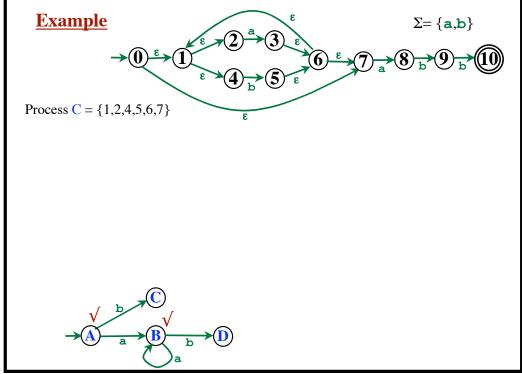
= ϵ -Closure (Move_{NFA}(B, \mathfrak{b}))



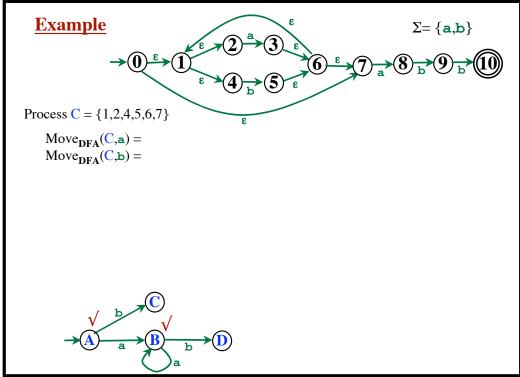


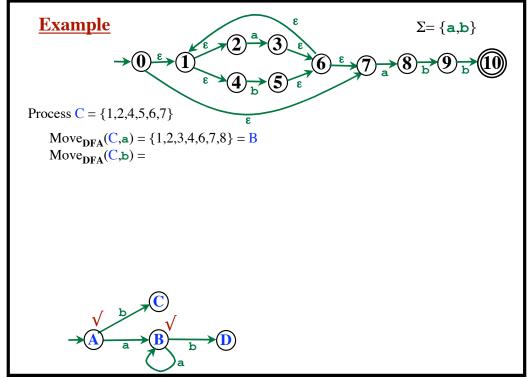




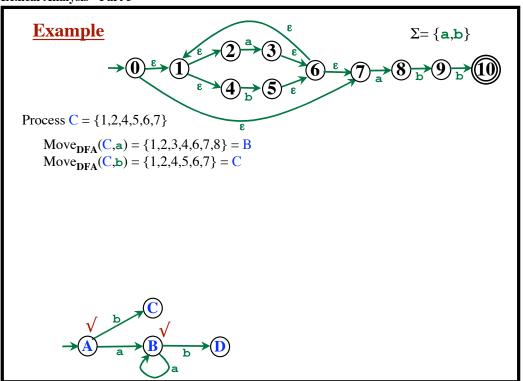


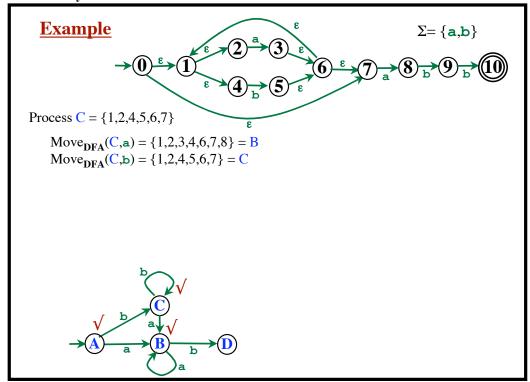
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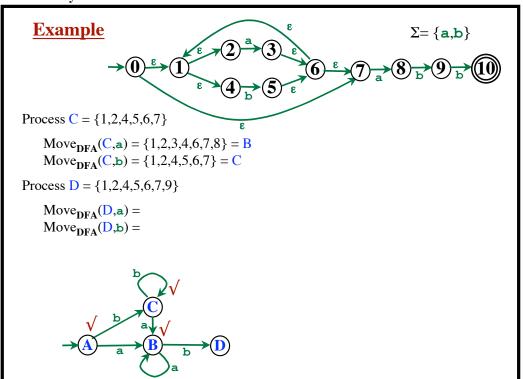
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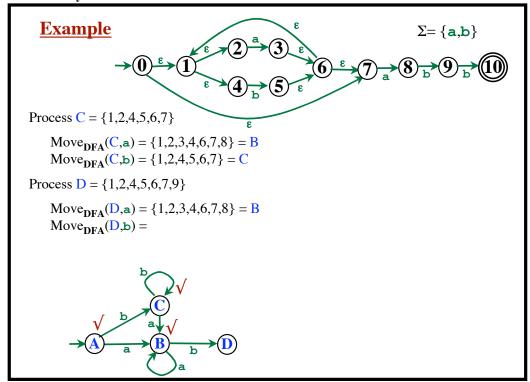




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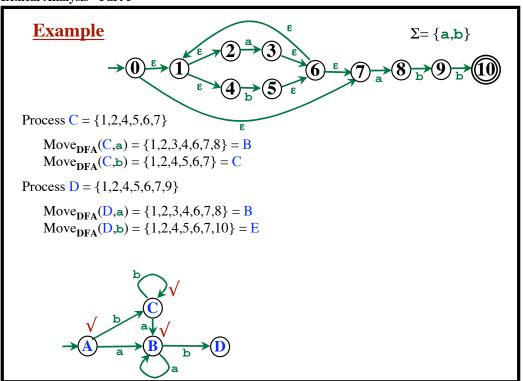
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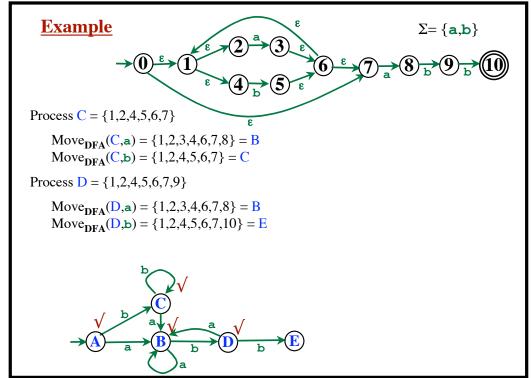




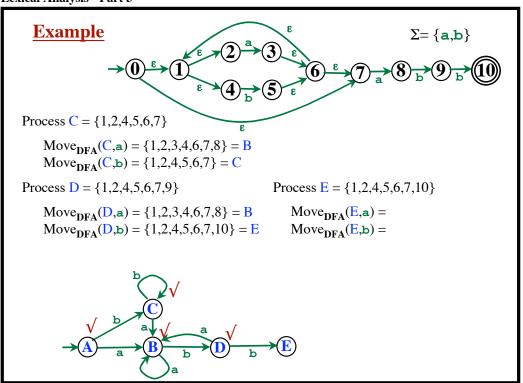
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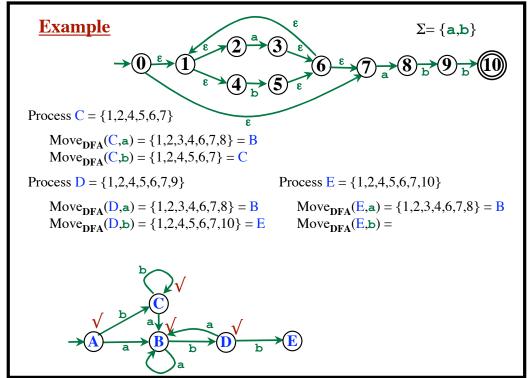
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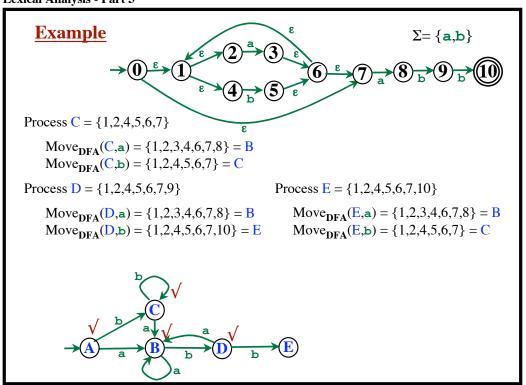


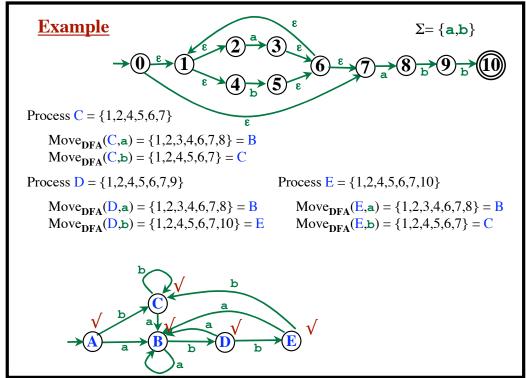
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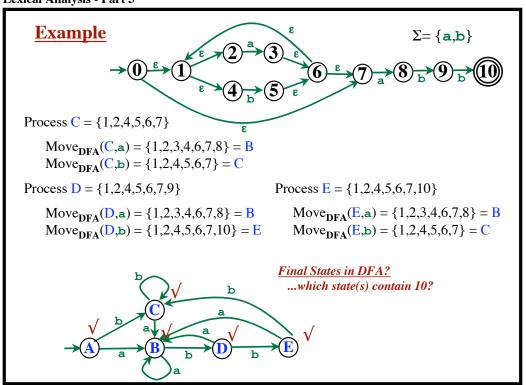


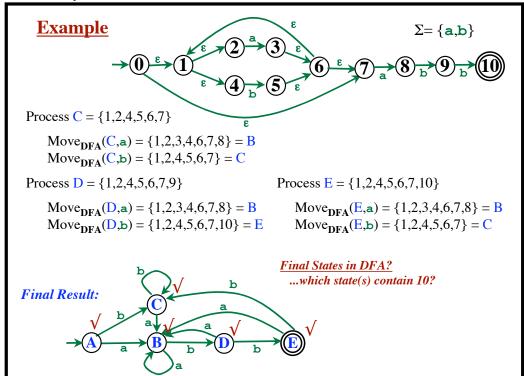
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Lexical Analysis - Part 3

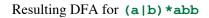


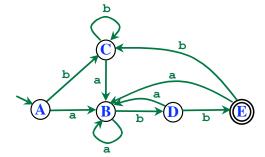


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Lexical Analysis - Part 3

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Algorithm: Convert NFA to DFA
Add \epsilon-Closure(s<sub>0</sub>) to S<sub>DFA</sub> as the start state
Set the only state in \mathbf{S}_{\mathrm{DFA}} to "unmarked"
\underline{\text{while}} \ S_{\text{DFA}} \ \text{contains an unmarked state} \ \underline{\text{do}}
  Let {\bf T} be that unmarked state
                                                         A set of NFA states
  Mark T
   for each a in Σ do
                                                             Everywhere you could
     S = \varepsilon-Closure (Move<sub>NFA</sub> (T, a))
     \underline{\underline{\text{if}}} S is not in S_{DFA} already \underline{\text{then}}
                                                             possibly get to on an a
        Add S to S<sub>DFA</sub> (as an "unmarked" state)
     Set Move<sub>DFA</sub>(T,a) to S
                                                   i.e, add an edge to the DFA...
  endFor
<u>endWhi</u>le
for each S in SDFA do
  if any s∈S is a final state in the NFA then
     Mark S an a final state in the DFA
  <u>end</u>If
endFor
```





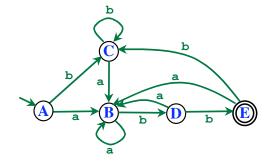
Is it minimal?

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Lexical Analysis - Part 3

Resulting DFA for (a|b) *abb



Is it minimal?



