HW#1 Advanced Operating Systems, Fall 2024

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1. Write a program that calls fork(). Before calling fork(), have the main process access a variable (e.g., x) and set its value to something (e.g., 100). What value is the variable in the child process? What happens to the variable when both the child and parent change the value of x?

Solution: Please refer to List 1 (q1.c):

```
Listing 1: q1.c
1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <unistd.h>
5 int
6 main(int argc, char *argv[])
7 {
       int x=100;
8
       int rc = fork();
9
       if (rc < 0) {
10
           // fork failed; exit
11
           fprintf(stderr, "fork_failed\n");
12
           exit(1);
13
       } else if (rc == 0) {
14
           // child (new process)
15
           printf("hello,_I_am_child_(pid:%d)\n", (int) getpid());
16
           printf("x_was_%d_", x);
17
           x = 222;
18
           printf(",but_I_have_changed_it_to_222... \n");\\
19
           printf("Now, _x_is_%d! \n", x);
20
       } else {
^{21}
           // parent goes down this path (original process)
22
23
           printf("hello, _I_am_parent_of_%d_(pid:%d) \n",
                   rc , (int) getpid());
24
           printf("x_was_%d_w", x);
25
26
           x = 1111;
           printf(",but_I_have_changed_it_to_111...\n");
27
28
           printf("Now, _x_is_%d! \setminus n", x);
29
30
       return 0;
31 }
```

Its execution results are as follows:

As we can see in Line 8 of Listing 1, we have declared a variable x with a value of 100. Lines 14-20 and 22-28 show the code for the child process and the parent process, respectively. In Lines 17 and 25, we print out the value of x for the child and the parent process, which both are 100 (as shown in lines 4 and 7 of the execution results). It is because the child process is created with the same content as the parent

- including the data segment as well as the stack. Later, in Lines 18 and 26, we changed the values of x in the child and the parent to 222 and 111, respectively. We also print out the changed values in Lines 20 and 28, which are 222 and 111. This shows that the forked child process has the same contents as its parent when it was created. However, after that, the child and the parent are two independent processes.
- 2. Write a program that opens a file (with the open() system call) and then calls fork() to create a new process. Can both the child and parent access the file descriptor returned by open()? What happens when they are writing to the file concurrently, i.e., at the same time?

Solution:

Consider the following program:

```
1 #include <stdio.h>
2 #include < stdlib.h>
3 #include <unistd.h>
4 #include <fcntl.h>
5 #include < string.h>
6 #include <sys/types.h>
7 \#include <sys/wait.h>
9 int main() {
       //fd : file descripter
10
       int fd = open("output.txt", O_CREAT | O_WRONLY | O_TRUNC, 0644);
11
       if (fd < 0) {
    perror("open_failed");</pre>
12
13
           exit (EXIT_FAILURE);
14
15
16
       int rc = fork();
17
       if (rc < 0) {
18
           perror("Fork failed");
19
           exit (EXIT_FAILURE);
20
       } else if (rc == 0) {
21
           const char *child_message = "Hello_from_child!\n";
22
23
           write(fd, child_message, strlen(child_message));
           printf("Child_wrote_to_file.\n");
24
25
           const char *parent_message = "Hello_from_parent!\n";
26
           write(fd, parent_message, strlen(parent_message));
27
           printf("Parent_wrote_to_file.\n");
28
           wait (NULL);
29
30
31
       close (fd);
32
       return 0;
33
34 }
```

Its execution results are as follows:

```
1 $ cc q2.c
2 $ ./a.out
```

```
3 Parent wrote to file.
4 Child wrote to file.
5 $
```

The output.txt content:

1 Hello from parent!
2 Hello from child!

As shown in Listing 2, the program's main idea is to create a parent and child process to write messages into a file named output.txt. In Line 10, the open() function is used to open or create the file output.txt. The flags O_CREAT indicate that the file should be created if it does not already exist, O_WRONLY opens the file in write-only mode, and O_TRUNC clears the contents of the file if it already exists. The permission mode 0644 allows the owner to read and write the file, while others can only read it. Lines 17-20 handle error messages. If fork() fails, the program will display an error message and exit. Lines 20-24 represent the child process segment, which writes a message to the file and prints a success message. Lines 24-29 represent the parent process segment, which functions similarly to the child process. It writes a message to the file and prints a success message as well. As we can see in output.txt, the messages are written successfully.

3. Write a program that calls fork(). Before calling fork(), have the main process access a variable (e.g., x) and set its value to something (e.g., 100). What value is the variable in the child process? What happens to the variable when both the child and parent change the value of x?

Solution:

23 }

Listing 3: q3.c 1 #include <stdio.h> 2 #include <stdlib.h> 3 #include <unistd.h> int main() { 5 int pipefd[2]; 6 pipe (pipefd); int rc = fork();if (rc < 0){ 9 10 fprintf(stderr, "fork_failed\n"); exit(1); 11 $}$ else if (rc == 0) { 12 printf("hello\n"); 13 write(pipefd[1], "finished", 8);14 15 close (pipefd[1]); } else { 16 17 char buffer [4]; $\verb"read" (\verb"pipefd" [0]", buffer", 8);$ 18 printf("goodbye\n"); 19 close (pipefd [0]); 20 21 22 return 0;

```
1 $ cc q3.c
2 $ ./a.out
3 hello
4 goodbye
5 $
```

As we can see in Line 7 of the code, we use pipe() to create a pipeline for inter-process communication, where pipefd[0] is the read end and pipefd[1] is the write end of the pipe. In Line 8, we call fork() to create a child process.

Lines 12-16 handle the child process. In Line 13, the child prints "hello", and in Line 14, it writes "finished" to the pipe using write(pipefd[1], "finished", 8), which sends the message to the parent process through the pipe. The child then closes the write end of the pipe in Line 15.

Lines 16-20 deal with the parent process. In Line 18, the parent waits for message from the child by calling read(pipefd[0], buffer, 8), which blocks until the child sends its message. After receiving the message, the parent prints "goodbye" in Line 19, representing the message was received.

From upon example, we can know how the parent and child process switch the process by "finished" message. :)

4. Write a program that calls fork() and then calls some form of exec() to run the program /bin/ls. See if you can try all of the variants of exec(), including (on Linux) execl(), execle(), execlp(), execvp(), and execvpe(). Why do you think there are so many variants of the same basic call?

Solution:

Listing 4: q4.c

```
1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <unistd.h>
  int main() {
       printf("Child_process_is_going_to_run_/bin/ls\n");
6
       int rc = fork();
       if (rc = 0) {
8
           printf("Running_execl\n");
9
           execl("/bin/ls", "ls", NULL);
10
           perror("excel_error_QAQ.");
11
           exit(1);
12
13
       rc = fork();
14
       if (rc = 0) {
15
           printf("Running_execle\n");
16
           char * env[] = {"User=MingfaUser", NULL};
17
           execle("/bin/ls", "ls", NULL, env);
18
19
           perror("excel_error_QAQ.");
           exit(1);
20
21
      rc = fork();
22
       if (rc = 0) {
23
           printf("Running_execlp\n");
24
           execlp("ls", "ls", "-1", NULL);
25
```

```
26
             perror("excelp_error_QAQ.");
             exit(1);
27
28
        rc = fork();
29
30
        if (rc = 0)  {
             printf("{\tt Running\_execv} \verb|\| n");
31
             \mathbf{char} \ * \mathrm{args} \left[\right] \ = \ \overline{\{" \text{ls", "-l", NULL}\}};
32
             execv("/bin/ls", args);
             perror ("excev_error_QAQ.");
34
             exit(1);
35
36
        rc = fork();
37
38
        if (rc = 0) {
             printf("Running_execvp\n");
39
             char *args[] = {"ls", "-1", NULL};
40
             execvp("ls", args);
41
             perror("execvp_error_QAQ.");
42
43
             exit(1);
44
45
        return 0;
46
        printf("All_Processes_finished.");
47
48 }
```

```
_{1} $ cc q4.c
2 $ ./a.out
3 Child process is going to run /bin/ls
4 Running execl
5 Running execle
6 Running execlp
7 Running execv
s Running execvp
                                q3.c
                                               q3_explain.c
9 a.out
                 q1.c
                                                               q4.c
                                q3.c
                                               q3_explain.c
10 a.out
                 q1.c
                                                              q4.c
11 total 104
12 total 104
13 total 104
_{14} -rwxr-xr-x 1 limingfa
                               33488\ 10\ 1\ 15:31\ a.out
                         staff
staff
                                 811
                                     3 10
                                           2022 q1.c
_{16} -rw-r-r-@ 1 limingfa
                                        1 14:39 q3.c
                         staff
                                 487 10
17 -rw-r-r- 1 limingfa
                         staff
                                 854 10
                                        1 14:01 q3_explain.c
_{18} -rw-r-r-- 1 limingfa
                         staff
                                1179 10
                                        1 15:28 q4.c
                         staff
                               33488 10
19 -rwxr-xr-x 1 limingfa
                                        1 15:31 a.out
staff
                                     3 10
                                           2022 q1.c
                                 811
                                        1 15:31 a.out
21 -rwxr-xr-x 1 limingfa
                               33488 10
                         staff
staff
                                 487 10
                                        1 14:39 q3.c
23 -rw-r-r- 1 limingfa
                         staff
                                 854 10
                                        1 14:01 q3_explain.c
                                 811 3 10 2022 q1.c
staff
_{25} -rw-r-r-- 1 limingfa
                         staff
                                1179 10
                                         1 15:28 q4.c
26 — rw— r—— @ 1 limingfa
                                 487 10
                         staff
                                        1 14:39 q3.c
27 -rw-r-r- 1 limingfa
                         staff
                                 854 10
                                        1 14:01 q3_explain.c
                                1179 10
28 -rw-r-r- 1 limingfa
                         staff
                                        1 15:28 q4.c
```

As we can see from the example, there are many different styles of exec. I think the basic functions are:

- Parameter Transfer: Some functions, like execl() and execle(), allow you to directly list parameters, which is suitable for a small number of known parameters. Others, like execv() and execvp(), use an array, making it easier to handle an uncertain number of parameters.
- Control of Environment Variables: Functions such as execle() and execvpe() allow you to set environment variables, which is really useful when you need to change the environment during execution.
- Path Searching: Functions like execvp() and execvpe() can automatically search for executable files in the PATH environment variable, so you don't need to write the full path.

With those functions, we can use the command more wisely. :)

5. Now write a program that uses wait() to wait for the child process to finish in the parent. What does wait() return? What happens if you use wait() in the child?

Solution:

First Example will show the parent process with wait() function, waiting for child process.

```
Listing 5: q5-1.c
1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <unistd.h>
  #include <sys/wait.h>
  int main() {
6
      int rc = fork();
8
      if (rc < 0) {
9
           fprintf(stderr, "Fork_failed\n");
10
           exit(1);
11
      else if (rc == 0) 
12
           printf("Child_process_(PID:_%d)_is_running.\n", getpid());
13
           printf("Child_process_(PID:_%d)_is_done.\n", getpid());
15
           exit(0);
16
      } else {
17
           printf("Parent_process_(PID:_%d)_is_waiting_for_child_process_(PID:_%d)_to_
18
               finish.\n", getpid(), rc);
           int status;
19
           printf("wait_the_child_process_PID:_%d\n", wait(&status));
20
           if (wait(&status)){
21
                printf("Child, process, exited, with, wait(&status), returned: ..%d\n", wait(&
22
                    status));
23
           printf("Parent_process_(PID:_%d)_is_done._:_) \n", getpid());
25
26
      \mathbf{return} \ \ 0\,;
27
28 }
```

Its execution results are as follows:

```
1 $ cc q5_1.c
2 $ ./a.out
3 Parent process (PID: 27404) is waiting for child process (PID: 27405) to finish.
4 Child process (PID: 27405) is running.
5 Child process (PID: 27405) is done.
6 wait the child process PID: 27405
7 Child process exited with wait(&status) returned: -1
8 Parent process (PID: 27404) is done. : )
```

As we can see the first example has printed out the return value waiting for child process (on the line 6). It actually return the Pid from the child process. After that it will return the -1 as the child process is finished (on the line 7).

Second Example will show the child process with wait() function, waiting for child's child process (not exist).

Listing 6: q5-2.c

```
1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <unistd.h>
4 #include <sys/wait.h>
5 #include <errno.h>
  int main() {
       int rc = fork();
8
       if (rc < 0) {
9
           fprintf(stderr, "fork_failed\n");
10
11
           exit(1);
       } else if (rc == 0) {
12
           printf("Child_process_is_running\n");
13
           int status;
14
           printf("wait(&status):_%d\n", wait(&status));
15
16
           if (wait(&status)){
               printf("No_more_child_process_here.\n");
17
18
19
        else {
           wait (NULL);
20
21
      return 0;
22
23 }
```

Its execution results are as follows:

```
1 $ cc q5.2.c
2 $ ./a.out
3 Child process is running
4 wait(&status): -1
5 No more child process here.
```

As we can see in lines 12-19, the child segment is designed to wait for its child process to finish. However, since there are no child processes created after the child process, the wait (&status) function will return -1, indicating that there are no child processes left to wait for.

6. Write a slight modification of the previous program, this time using waitpid() instead of wait(). When

would waitpid() be useful?

Solution:

```
Listing 7: q6.c
 1 #include <stdio.h>
 2 #include <stdlib.h>
 3 #include <unistd.h>
   int main() {
 5
        int rc = fork();
        \begin{array}{ll} \textbf{if} & (\,\mathrm{rc}\,<\,0\,) \; \{ \\ & \mathtt{perror}\,(\,\texttt{"Fork\_failed"}\,)\,; \end{array}
 8
 9
              exit(1);
10
        } else if (rc == 0) {
11
              printf("Child_process_is_running._PID=%d\n", getpid());
12
13
              sleep (2);
        } else {
14
              int status;
15
              int waited_pid = waitpid(rc, &status, 0);
16
              printf("\texttt{Parent\_process\_is\_running.\_PID=\$d\n"}\,,\ getpid());
17
              printf("waitpid()_returned:_%d\n", waited_pid);
18
19
20
        return 0;
21
22 }
```

Its execution results are as follows:

```
1 $ cc q6.c
2 $ ./a.out
3 Child process is running. PID=36153
4 Parent process is running. PID=36149
5 waitpid() returned: 36153
```

The getpid() function will be useful for identifying a specific process. If we want to know the process ID of the currently running process, we can call getpid(), which returns its own process ID. This can help in debugging or when we need to manage processes, as we can keep track of which process is executing certain tasks. For example, if a parent process needs to wait for a particular child process, it can use getpid() to know its own ID and ensure it's managing the right child process effectively.

7. Write a program that creates a child process, and then in the child closes standard output (STDOUT FILENO). What happens if the child calls printf() to print some output after closing the descriptor?

Solution:

Listing 8: q7.c 1 #include <stdio.h> 2 #include <stdlib.h> 3 #include <unistd.h>

```
5 int main() {
      int rc = fork();
6
       if (rc < 0) {
8
9
           perror("Fork_failed");
           exit(1);
10
      } else if (rc = 0) {
11
           close(STDOUT_FILENO);
12
           printf("This_message_will_not_be_printed_to_the_terminal.\n");
13
           exit(0);
14
       } else {
15
           int status;
16
           waitpid (rc, &status, 0);
17
           printf("Child_process_has_finished._:)\n");
18
19
20
      return 0;
21
22 }
```

```
1 $ cc q7.c
2 $ ./a.out
3 Child process has finished. :)
```

In the child process (when rc == 0), nothing will be printed out after closing the standard output until the process finishes and control switches back to the parent process. (when rc == 1)

8. Write a program that creates two children, and connects the standard output of one to the standard input of the other, using the pipe() system call.

Solution:

1 #include <stdio.h> 2 #include <stdlib.h> 3 #include <unistd.h> 4 #include <string.h> 6 **int** main() { int pipeEnd[2];if (pipe(pipeEnd) = -1) { perror("pipe_failed"); 9 10 exit(1); 11 12 13 int rc1 = fork();14 15 if (rc1 < 0){ perror("Fork_failed_for_first_child"); 16 17 exit(1); $\}$ else if (rc1 = 0) { 18 close (pipeEnd[0]); 19 const char *message = "Hello_from_child_1"; write(pipeEnd[1], message, strlen(message) + 1); 21

```
22
           close (pipeEnd[1]);
23
24
       int rc2 = fork();
25
26
       if (rc2 < 0) {
27
           perror("Fork_failed_for_second_child");
28
            exit (1);
       } else if (rc2 == 0) {
30
           close (pipeEnd[1]);
31
           char buffer [100];
32
           read(pipeEnd[0], buffer, sizeof(buffer));
33
           printf("Child_2_received:_%s\n", buffer);
34
           close(pipeEnd[0]);
35
36
37
       close (pipeEnd[0]);
38
       close(pipeEnd[1]);
39
40
41
       return 0;
42 }
```

```
1 $ cc q8.c
2 $ ./a.out
3 Child 2 received: Hello from child 1
```

In the child process (when rc1 == 0), it closes the read end of the pipe (pipeEnd[0]) to ensure it only writes. It then prepares a message, "Hello from child 1", and writes this message to the write end of the pipe (pipeEnd[1]). After writing, it closes the write end of the pipe and exits.

A second child process is created with another fork() call (rc2). If the fork fails, an error message is printed. In the second child process, it closes the write end of the pipe (pipeEnd[1]) to ensure it only reads. It then declares a buffer to hold the incoming message and reads from the read end of the pipe (pipeEnd[0]). After successfully reading the message, it prints the received message to the console and closes the read end before exiting.

After all the process ended, close the writing and reading ends. :)