

inorder: left - root - right

for ①: It's inorder successor is in / ①'s right subtree  
  \ ①'s node

for ② : it's inorder successor is in — ②'s right subtree  
 \ null if ② has no right child

for ③: 'His inorder successor is in / ③'s right subtree  
 \ null if ③ has no right child

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Node's Location	It's Inorder Successor in BST
Left subtree	It's right subtree It's root
Root	It's right subtree null if it has no right child
Right subtree	It's right subtree null if it has no right child

-if we use divide and conquer.

## // Corner Case

```
if (root == null) {  
    return root;  
}
```

// p is the root or p is in the right subtree

// find p's inorder successor in the right subtree

if (proc.val <= p.val) {

return inorderSuccessor ( root.right, p );

// p is in the left subtree

```

// find p's inorder successor in the left subtree, if not found, return root;

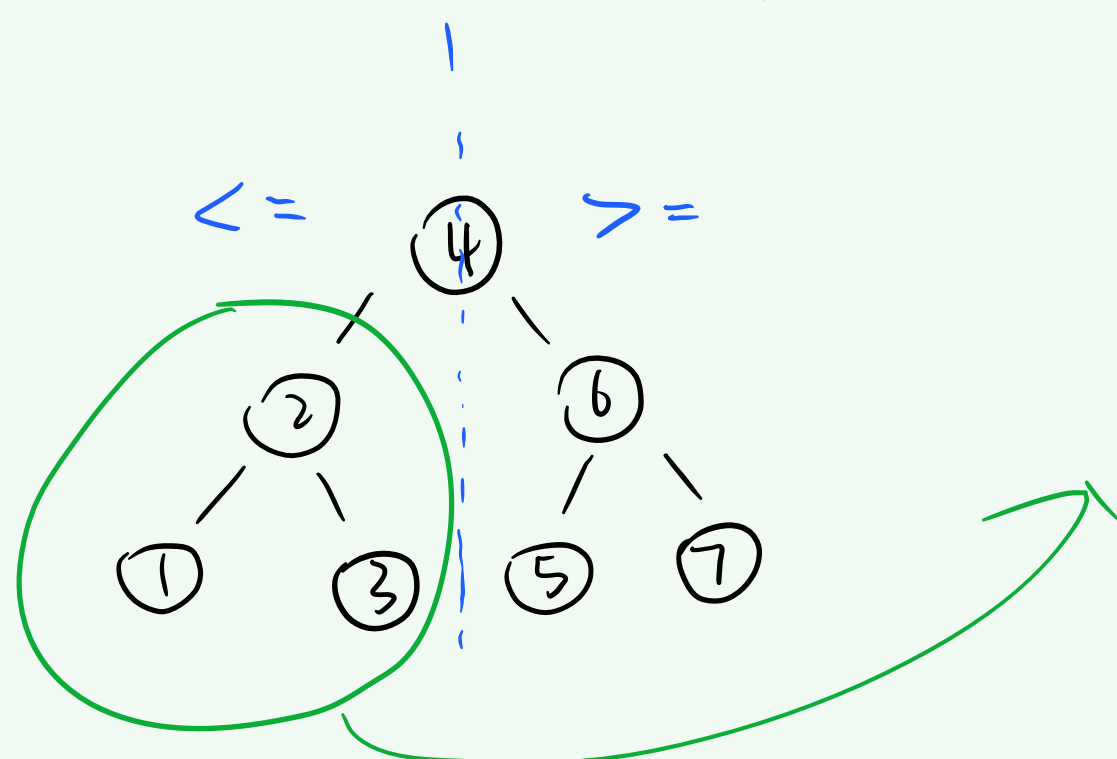
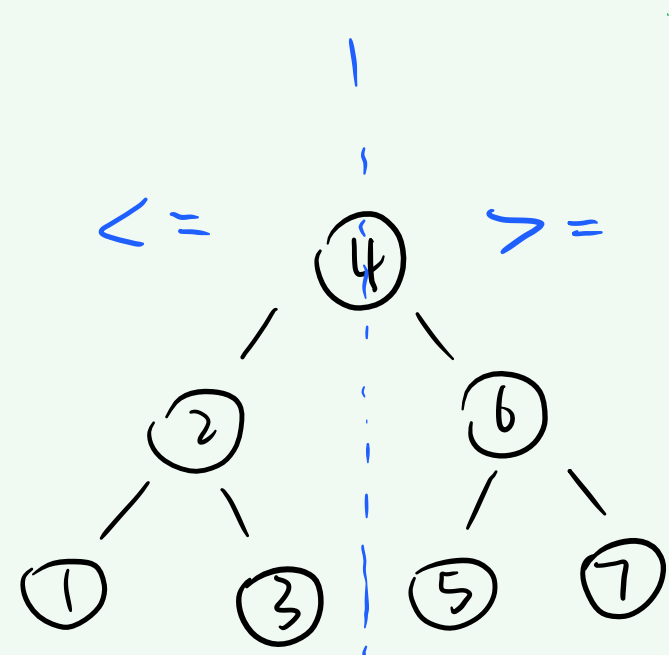
```

```
} else {
```

```
TreeNode left = inorderSuccessor ( root.left, P );
```

```
return (left != null) ? left : root;
```

}



Left subtree has 3 nodes

if  $k \leq 3$ , we find  $k$  in left subtree

if  $k > 3 + 1$ , we find  $k$  in right subtree

if  $k = 3+1$ ,  $k$  is the root

$\therefore$  BST

$\therefore$  Left tree's nodes' value  $\leq$  root.val  $\leq$  Right tree's node's value

The smallest elements distribution should be: left tree  $\rightarrow$  root  $\rightarrow$  right tree

We need to count total # of nodes in left subtree, let's call it  $l$ .

If  $k \leq l$ , we do not need to find  $k$ th smallest element in the root and right sub tree.

2f  $k > l + 1$ , that means  $k$ th smallest element is in the right subtree.

Otherwise,  $k = l+1$ ,  $k$ -th smallest element is the root.