

Since the input array is sorted in **ascending** order,
if it's rotated at some pivot, we surely would find a **gap**,
or an **inflection point** in the rotated array.

The index of this **gap (or inflection point)** is the pivot position.

Version 1: Iterate the input array to see if the target exists.

Time Complexity: $O(n)$; $n == \#$ of elements in array, too slow.

Version 2: 2 times binary search

- ① find the rotated position;
- ② find the target according to two ascending parts of the array.

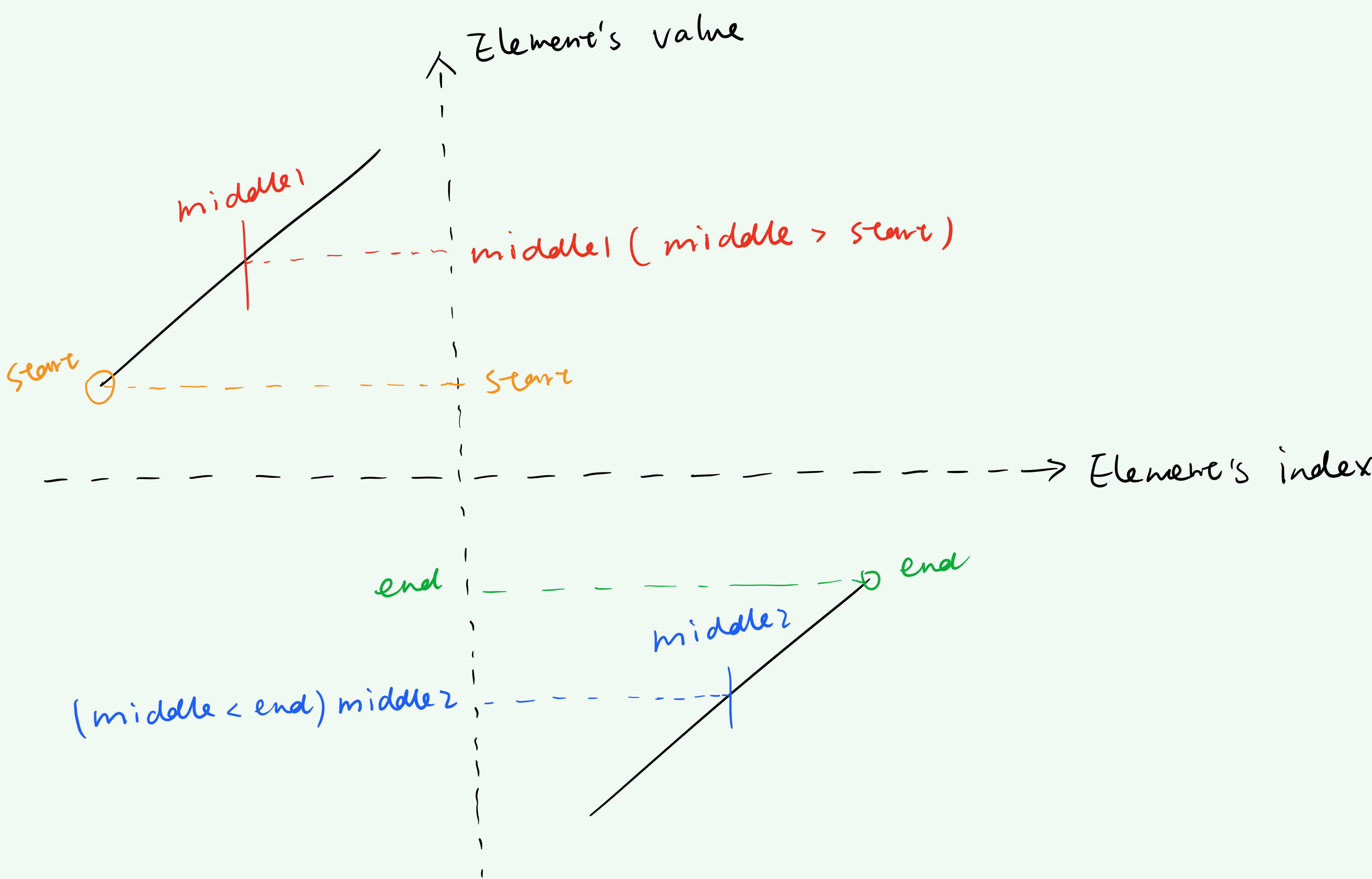
Time Complexity: $O(\log n + \log n) \sim O(\log n)$, $n == \#$ of elements in array

Version 3: One time binary search.

There're two possibilities of the middle point

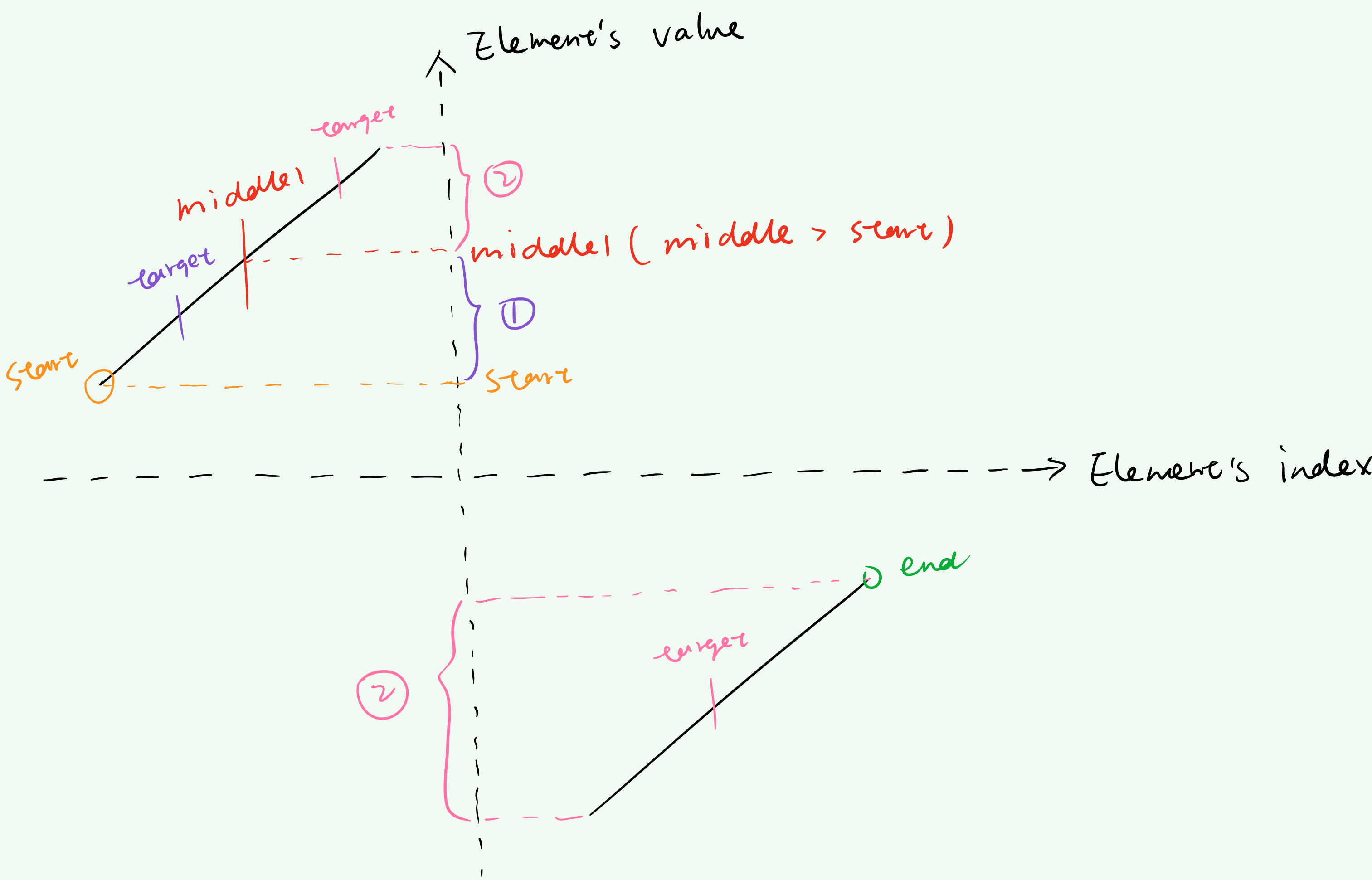
- \swarrow middle > start \Rightarrow **middle 1**
- \searrow middle < end \Rightarrow **middle 2**

4 5 **6** 7 0 **1** 2



1. For **middle 1**

- ① if **start** <= **target** <= **middle 1**
if \updownarrow else \downarrow let binary search's end = middle 1
- ② if **target** > **middle 1**
or **target** < **start**
let binary search's start = middle 1



2. For **middle 2**

- ① if **middle 2** <= **target** <= **end**
if \updownarrow else \downarrow let binary search's start = middle 2
- ② if **target** < **middle 2**
or **target** >= **start**
let binary search's end = middle 2

