

C0: Basic Topology (MA2108S Recap)

Rudin’s Point Terminologies (E° = Interior(E), ∂E = Boundary(E), Ē = Closure(E) $\overset{\text{def}}{=} E^\circ \oplus \partial E \overset{\text{def}}{=} E + E'$)

- Limit Point $p \in E' : \forall r > 0 (N_r(p) \setminus \{p\} = \emptyset)$ (‘deleted neighbourhood)
- Isolated Point: $p \in E - E' . \therefore p \in E \Leftrightarrow p$ limit point xor isolated point.
- Interior point $p \in E^\circ : \exists r > 0 (N_r(p) \subseteq E \Leftrightarrow N_r(p) \cap E^c = \emptyset)$

Boundary point $p \in \partial E : \forall r > 0 (N_r(p) \cap E \neq \emptyset \wedge N_r(p) \cap E^c \neq \emptyset) \overset{\text{def}}{\Leftrightarrow} p \in \bar{E} \cap \bar{E^c} \overset{\text{def}}{\Leftrightarrow} p \in \bar{E} - E^\circ . \therefore p \in \partial E \nRightarrow p \in E$.
Exterior Point: $\exists r > 0 (N_r(p) \cap E = \emptyset \Leftrightarrow N_r(p) \subseteq E^c) . \therefore p \in E \Leftrightarrow p$ interior, boundary, xor exterior point.

- Corollaries: 1. Limit point is interior xor boundary, and provably not exterior.
Purely-interior eg: $p = 0 \in \mathbb{R}$. Purely-boundary eg: $p = 0 \in \{0\} \cup \{\frac{1}{n} : n \in \mathbb{N}^+\}$.
2. Isolated point is interior xor boundary.
Purely-interior eg: $p = 0 \in \mathbb{Z}$. Purely-boundary eg: $p = 0 \in \{0\}$
3. If metric space X connected, then isolated \Rightarrow boundary $\wedge E^\circ \subseteq E' \wedge (E' \neq \emptyset \Rightarrow E$ infinite)
4. $\forall p \in E' (\forall r > 0) (N_r(p) \cap E$ infinite)

Rudin’s Set Terminologies [LUB/GCB: $\exists \min\{E\}, \max\{E\}$ of bdd E in \mathbb{R}^k]

Open set: $E = E^\circ$ | Closed set: $E' \subseteq E \overset{\text{def}}{\Leftrightarrow} E = \bar{E}$ | Crowded set: $E' \supseteq E$ (eg: \mathbb{Q}^c) | Perfect (closed crowded) set: $E' = E$
E bounded in X : $\exists r > 0 (\exists \bar{q} \in X) (N_r(\bar{q}) \supseteq E) \Leftrightarrow \exists r > 0 (\forall \bar{q} \in E) (||\bar{q}|| \leq r)$ |

- Corollaries: 1. de-Morgan’s Law: $\forall A \succeq \mathbb{N} (\forall a \in A) ((\bigcup_{a \in A} E_a)^c = \bigcap_{a \in A} E_a^c) . \therefore E$ open $\Leftrightarrow E$ closed.
2. Any set E is open and perfect in itself. $\bar{E} = \min$ closed superset, $E^\circ = \max$ open subset.
3. (Countable \bigcup /Finite \bigcap) keeps openness. (Countable \bigcap /Finite \bigcup) keeps closure
Countable open \bigcap eg: $\bigcap_{n=1}^\infty (-\frac{1}{n}, \frac{1}{n}) = \{0\}$ | Countable closed \bigcup eg: $\bigcup_{n=1}^\infty [\frac{1}{n}, \frac{n-1}{n}] = (-1, 1)$

Relative Openness/Relative Closeness

$\forall E \subseteq$ subspace $Y \subseteq$ metric space $X (E$ open/closed in $Y \Leftrightarrow \exists$ open/closed $G \subseteq X (E = Y \cap G))$

E open in $X \Rightarrow E$ open in $Y \subseteq X$. E closed in $Y \Rightarrow E$ closed in $X \supseteq Y$.

Compactness

E compact in X: \exists open $U_i \subseteq X \Rightarrow \exists n \in \mathbb{N}^+ (\bigcup_{i=1}^n U_{ij} \supseteq E)$ i.e, \exists finite sub-cover per open cover.

$\forall E \subseteq$ subspace $Y \subseteq$ metric space $X (E$ compact in $Y \Leftrightarrow E$ compact in $X) | \forall$ closed $F \subseteq$ compact $E (F$ compact in $X)$

k-cell compact in $\mathbb{R}^k | \forall$ infinite $G \subseteq$ compact $E (E \cap G' \neq \emptyset)$

E compact $\Rightarrow E$ closed bdd in X. Converse eg: E closed bdd in $X = \mathbb{Q} \cap [0, 1]$ not complete (‘ $\exists \{q_n\} \rightarrow \frac{\sqrt{2}}{2}$ divergent)

Heine-Borel Theorem: If $X = \mathbb{R}^k$, E compact $\Leftrightarrow E$ closed bounded

Perfectness

\forall perfect $E \subseteq \mathbb{R}^k (E = \emptyset \vee E \succ \mathbb{N}) | \exists$ perfect, disconnected E eg: Cantor’s set = $\bigcap_{n=1}^\infty \bigcup_{k=0}^{3^n-1-1} ((\frac{3k}{3^n}, \frac{3k+1}{3^n}) \cup [\frac{3k+2}{3^n}, \frac{3k+3}{3^n}])$

Connectedness $[\forall A, B \subseteq X (A, B$ separated in $X \overset{\text{def}}{\Leftrightarrow} A, B \neq \emptyset \wedge A \cup B = X, \bar{A} \cap B = A \cap \bar{B} = \emptyset)$, so $A \cap B = \emptyset$]

$\forall E \subseteq X (E$ disconnected $\overset{\text{def}}{\Leftrightarrow} \exists$ separated partition (A, B) of $E) | E$ connected $\Rightarrow E$ connected (‘take $(A \cap E, B \cap E)$ of E)

E connected $\nRightarrow E^\circ$ connected. Eg: $\bigcup 2$ open tangential disc. If $X = \mathbb{R}^1$, E° connected | E connected in $\mathbb{R}^k \Leftrightarrow E$ interval.

C1: Terminologies

Constrained NLP format: $\min/\max f(\vec{x})$ s.t. $\begin{cases} \forall i (g_i(\vec{x}) = 0) & \text{equality constraints} \\ \forall j (h_j(\vec{x}) \leq 0) & \text{inequality constraints} \end{cases}$

Unconstrained NLP format: $\min/\max f(\vec{x})$ s.t. $\vec{x} \in$ open convex D , f convex on D
Feasible region $S = \{\vec{x} \in \mathbb{R}^n : \forall i \forall j (g_i(\vec{x}) = 0 \wedge h_j(\vec{x}) \leq 0)\}$ closed set of feasible solutions (Generalisation)

Optimal $\vec{x}^* \in \mathbb{R}^k$ (or ∞ if max or $-\infty$ if min) $\overset{\text{def}}{\Leftrightarrow} f(\vec{x}^*)$ optimal objective value (or unbounded, so NLP unbounded).

C2: Weierstrass EVT, Graphical Method, Optimality Existence

Inner Product $\langle \vec{x}, \vec{y} \rangle = \vec{x}^T \vec{y} = \sum_{i=1}^n x_i y_i = ||\vec{x}|| ||\vec{y}|| \cos \theta_{xy}$ (Cosine Rule) Properties

1. $\langle \vec{u} + \vec{v}, \vec{w} \rangle = \langle \vec{u}, \vec{w} \rangle + \langle \vec{v}, \vec{w} \rangle$ [Additivity] | 2. $\langle k\vec{u}, \vec{v} \rangle = k \langle \vec{u}, \vec{v} \rangle$ [Homogeneity] | 3. $\langle \vec{u}, \vec{v} \rangle = \langle \vec{v}, \vec{u} \rangle$ [Symmetry]

norm $||\vec{u}|| = \langle \vec{u}, \vec{u} \rangle$ properties: 1. $||\vec{u}|| \geq 0$, equality iff $\vec{u} = \vec{0}$ [Positive-Definiteness]
2. $||\vec{u}|| - ||\vec{v}|| \leq ||\vec{u} + \vec{v}|| \leq ||\vec{u}|| + ||\vec{v}||$ (‘expand abs sign) [(Reverse) Triangle Inequality]
3. $||\langle \vec{x}, \vec{y} \rangle| \leq ||\vec{x}|| ||\vec{y}|| \Leftrightarrow \sum x_i y_i \leq \sum x_i^2 \sum y_i^2$, equality iff \vec{x}/\vec{y} [Cauchy-Schwarz Inequality]

Vector Geometry [MA2104 Recap]

Scalar projection of \vec{b} onto \vec{a} , $\text{comp}_{\vec{a}} \vec{b} = ||\vec{b}|| \cos \theta_{ab} = \frac{\vec{a} \cdot \vec{b}}{||\vec{a}||}$ [Vector projection of \vec{b} onto \vec{a} , $\text{proj}_{\vec{a}} \vec{b} = \frac{\vec{a} \cdot \vec{b}}{||\vec{a}||^2} \vec{a}$

Cross-product: $||\vec{a} \times \vec{b}|| = ||\vec{a}|| ||\vec{b}|| \sin \theta_{ab}$ | (Right-Hand Grip Rule) \Rightarrow Perpendicular of \vec{b} onto $\vec{a} = ||\vec{b}|| \sin \theta_{ab}$ | $= \frac{||\vec{a} \times \vec{b}||}{||\vec{a}||}$

Angle between planes=Angle between normals = $\cos^{-1}(\frac{\vec{n}_1 \cdot \vec{n}_2}{||\vec{n}_1|| ||\vec{n}_2||}) \in [0, \pi]$ | Arc length of curve $C = \vec{r}(t)$, $s = \int_a^b ||\vec{r}'(t)|| dt$

Extremisers [Maximisers’ Definitions follow analogously]

Local Min: $\exists r > 0 (\forall \vec{x} \in N_r(\vec{x}^*) \cap S) (f(\vec{x}) \geq f(\vec{x}^*))$ | Strict Local Min: $\exists r > 0 (\forall \vec{x} \in N_r(\vec{x}^*) \cap S \setminus \{\vec{x}^*\}) (f(\vec{x}) > f(\vec{x}^*))$

Global Minimiser: $\forall \vec{x} \in S (f(\vec{x}) \geq f(\vec{x}^*))$ | [Strict Global Min: $\forall \vec{x} \in S \setminus \{\vec{x}^*\} (f(\vec{x}) > f(\vec{x}^*))$]

(strict) Global maximisers \Rightarrow (strict) Local maximisers. $\vec{x}^* \in E^\circ \Rightarrow \vec{x}^*$ stationary (not necessarily if $\in \partial E$)

Weierstrass Extreme Value Theorem (EVT), IVT, MVT

Weierstrass EVT: If f cont on nonempty closed bdd $S \subseteq \mathbb{R}^k$, \exists global minimiser \wedge maximiser in S not-necessarily unique.

Boundedness Theorem: If f continuous on nonempty closed bounded $[a, b]$, f bounded in $[a, b]$.

‘. IVT: If real-valued f cont on closed bounded interval $[a, b] \neq \emptyset, \forall c \in (a, b), \exists f^{-1}(c) \in (\min\{f(a), f(b)\}, \max\{f(a), f(b)\})$

‘. Bolzano’s Theorem: If real-valued f continuous on closed bounded $[a, b] \neq \emptyset$ s.t. $f(a)f(b) < 0, \exists c \in (a, b)$

‘. MVT: If real-valued f cont on closed bounded interval $[a, b]$, diff on (a, b) , $\exists c \in (a, b), f'(c) = \frac{f(b)-f(a)}{b-a}$

‘. Rolle’s Theorem: If real-valued f cont on $[a, b]$, diff on (a, b) s.t. $f(a) = f(b), \exists c \in (a, b) f'(c) = 0$

C3: Convex/Non-Convex Set & Function

Convex Set: $\forall \vec{x}, \vec{y} \in S (\forall \lambda \in [0, 1]) (\lambda \vec{x} + (1 - \lambda) \vec{y} \in S)$ i.e, closure under affine combination

Algebraic proofs oft-applies Traingle/Cauchy-Schwarz Inequalities for set/function. \nexists ”concave” analogue. Properties:

- If $X = \mathbb{R}$ or \mathbb{C} , S convex \Rightarrow S path-connected \Rightarrow S connected. Converse fails. Eg: torus | 2. ∂S convex curve/function.
- \emptyset, X convex set. $\forall S_1, S_2$ convex, countable intersection convex, but union convex iff chain (totally-ordered set).
- Non-eg: Venn Diagram. | 4. Strictly convex set: $\forall \vec{x}, \vec{y} \in S (\forall \lambda \in (0, 1)) (\lambda \vec{x} + (1 - \lambda) \vec{y} \in \circ)$.

Convex-downward Function: $\forall \vec{x}, \vec{y} \in D_f (\forall \lambda \in [0, 1]) (f(\lambda \vec{x} + (1 - \lambda) \vec{y}) \leq \lambda f(\vec{x}) + (1 - \lambda) f(\vec{y}))$ [\exists concave analogue]

Strictly convex-downward function: $\forall \vec{x}, \vec{y} \in D_f (\forall \lambda \in (0, 1)) (f(\lambda \vec{x} + (1 - \lambda) \vec{y}) < \lambda f(\vec{x}) + (1 - \lambda) f(\vec{y}))$. Properties:

- $\forall f_1, \dots, f_n$ convex, $\forall k_1, \dots, k_n \geq 0, \sum_{i=1}^n k_i f_i$ convex [Linearity]. | 2. $\forall f_1$ convex, $\forall k < 0, k f_1$ concave.
- $\forall f_1, f_2$ convex, $\max(\min)\{f_1, f_2\}$ convex (concave) | 4. h convex (concave) $\wedge g \begin{cases} \text{non-dec} \\ \text{non-inc} \end{cases} \Rightarrow g \circ h \begin{cases} \text{convex (concave)} \\ \text{concave (convex)} \end{cases}$

C4: Gradient Vector, Tangent Plane

Convex Function f on Convex Set $S \subseteq \mathbb{R}^n : \forall k \in \mathbb{R} (S_k = \{\vec{x} \in D_f : f(\vec{x}) \leq k\})$ convex, epigraph E_f convex

$\forall \vec{x} = \sum \lambda_i \vec{x}_i \in D_f, \sum \lambda_i = 1, f(\vec{x}) \leq \sum \lambda_i f(\vec{x}_i)$ (‘Convex f’s definition generalised by well-defined $f(\vec{x}_i)$ in convex S)

Gradient Vector $\nabla f(\vec{x}) = \begin{pmatrix} f_{x_1}(\vec{x}) \\ \vdots \\ f_{x_n}(\vec{x}) \end{pmatrix} : \forall \vec{d} \in \mathbb{R}^n (\forall \lambda \in \mathbb{R}) (f'(\vec{x} + \lambda \vec{d}) = (\nabla f(\vec{x} + \lambda \vec{d}))^T \vec{d})$

1. If $\lambda = 0, f'(\vec{x}) = \lim_{\lambda \rightarrow 0} \frac{f(\vec{x} + \lambda \vec{d}) - f(\vec{x})}{\lambda} = (\nabla f(\vec{x}))^T \vec{d} =$ signed rate of change from \vec{x} in direction \vec{d}

‘. $f \downarrow$ most rapidly from \vec{x} in direction $\vec{d} = -\nabla f(\vec{x})$ and \uparrow most rapidly from \vec{x} in direction $\vec{d} = \nabla f(\vec{x})$.

2. \forall local min/maximiser \vec{x}^* of f s.t. $g(\vec{x}) = c$ eq constraint, $\exists \lambda$ s.t. $\nabla f(\vec{x}^*) = \lambda \nabla g(\vec{x}^*)$ [Lagrange Multiplier Method]

3. If f:convex $D_f \rightarrow \mathbb{R}^n$ convex, $\begin{cases} \vec{x}^* \text{ global minimiser} \\ \vec{x}^* \text{ global maximiser} \end{cases} \Leftrightarrow \forall \vec{x} \in D_f ((\nabla f(\vec{x}^*)^T)(\vec{x} - \vec{x}^*) \geq 0) \text{ (‘} f \uparrow \text{ in all dir from } \vec{x}^*)$
 $\Leftrightarrow \forall \vec{x} \in D_f ((\nabla f(\vec{x}^*)^T)(\vec{x} - \vec{x}^*) \leq 0) \text{ (‘} f \downarrow \text{ in all dir from } \vec{x}^*)$

Tangent-plane below f Definition of Function Convexity

f convex $\Leftrightarrow \forall \vec{x}, \vec{y} \in D_f (f(\vec{x}) + (\nabla f(\vec{x}))^T (\vec{y} - \vec{x}) \leq f(\vec{y}))$ [Strictly convex, concave, strictly concave apply analogously]

C5: Hessian Matrix, Definiteness Tests

Hessian $H_f(\vec{x}) = \begin{pmatrix} f_{x_1 x_1} & f_{x_2 x_1} & \dots & f_{x_n x_1} \\ f_{x_1 x_2} & f_{x_2 x_2} & \dots & f_{x_n x_2} \\ \vdots & \vdots & \ddots & \vdots \\ f_{x_1 x_n} & f_{x_2 x_n} & \dots & f_{x_n x_n} \end{pmatrix} \forall \vec{x} \in D_f^\circ, (f_{x_i})_{x_j} = \frac{\partial^2 f}{\partial x_j \partial x_i}, \Delta_k = \det(\begin{pmatrix} f_{x_1 x_1} & f_{x_2 x_1} & \dots & f_{x_k x_1} \\ f_{x_1 x_2} & f_{x_2 x_2} & \dots & f_{x_k x_2} \\ \vdots & \vdots & \ddots & \vdots \\ f_{x_1 x_k} & f_{x_2 x_k} & \dots & f_{x_k x_k} \end{pmatrix}))$

Clairaut’s Theorem: f has continuous 2nd-order partial derivatives $\Rightarrow H_f(\vec{x})$ symmetric. Else, non-eg is pathological.

3 Definiteness Tests

$\begin{cases} \geq 0 \\ > 0 \end{cases} \overset{\text{def}}{\Leftrightarrow} \begin{cases} \text{A p.s.d.} \\ \text{A p.d.} \end{cases}$
1. $\forall \vec{x} \in \mathbb{R}^n (\vec{x}^T A \vec{x} = \begin{cases} \leq 0 \\ < 0 \end{cases} \overset{\text{def}}{\Leftrightarrow} \begin{cases} \text{A n.s.d.} \\ \text{A n.d.} \end{cases})$ [Definition]
 $\begin{cases} \leq 0 \\ < 0 \end{cases} \overset{\text{def}}{\Leftrightarrow} \begin{cases} \text{A n.d.} \\ \text{A indefinite} \end{cases}$
($\exists x_1 \neq x_2 (x_1^T A x_1 > 0 \wedge x_2^T A x_2 < 0)$)

2. \forall eigenvalue λ of symmetric $A, \lambda = \begin{cases} \text{all } > 0 \\ \text{all } \geq 0 \\ \text{all } \leq 0 \\ \text{all } < 0 \end{cases} \Leftrightarrow \begin{cases} \text{A p.s.d.} \\ \text{A p.d.} \\ \text{A n.s.d.} \\ \text{A n.d.} \end{cases}$ [Eigenvalue Test, $p_A(\lambda) = \det(A - \lambda I_n)$]
 $\begin{cases} \text{all } > 0 \\ \text{all } \leq 0 \end{cases} \Leftrightarrow \begin{cases} \text{A p.d.} \\ \text{A indefinite} \end{cases}$
 $\exists \lambda_1 < 0 \wedge \lambda_2 > 0$

3. For symmetric $A, \begin{cases} \forall k \in [1, n] \Delta_k > 0 \\ \forall k \in [1, n] (-1)^k \Delta_k > 0 \end{cases} \Leftrightarrow \begin{cases} \text{A p.d.} \\ \text{A n.d.} \end{cases}$ [Principal Minor Test on Definiteness Only]

$\forall k \in [1, n] \Delta_k \geq 0 \nRightarrow$ A p.s.d. eg: n.s.d. $A = \begin{pmatrix} 0 & 0 \\ 0 & -1 \end{pmatrix}$. $\forall k \in [1, n] (-1)^k \Delta_k \geq 0 \nRightarrow$ A n.s.d. eg: p.s.d. $A = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$

2nd Partial Derivative Test

$f : \text{open } D_f \rightarrow \mathbb{R}^n \begin{cases} \text{convex} & \Leftrightarrow \forall \vec{x} (H_f(\vec{x}) \text{ p.s.d.}) \\ \text{concave} & \Leftrightarrow \forall \vec{x} (H_f(\vec{x}) \text{ n.s.d.}) \\ \text{strictly convex} & \Leftrightarrow \forall \vec{x} (H_f(\vec{x}) \text{ p.d.}) \text{ [Non-eg: } f(x) = x^4 \Rightarrow H_f(x) = 12x^2 \text{ p.s.d. only}] \\ \text{strictly concave} & \Leftrightarrow \forall \vec{x} (H_f(\vec{x}) \text{ n.d.}) \text{ [Non-eg: } f(x) = -x^4 \Rightarrow H_f(x) = -12x^2 \text{ n.s.d. only}] \\ \text{neither} & \Leftrightarrow \exists \vec{x} (H_f(\vec{x}) \text{ indef}) \text{ [Non-eg: } f(x) = x^3 \text{ neither by } (\nabla f)^T \Rightarrow H_f(x) = 6x \text{ p/n.s.d. } \forall x] \end{cases}$

Taylor’s Theorem on f : $\mathbb{R}^n \rightarrow \mathbb{R}$ with continuous 2nd partial derivatives

$\forall \vec{x}, \vec{y} \in D_f (||\vec{x}, \vec{y}|| \subseteq D_f^\circ \Rightarrow \exists \lambda_w \in [0, 1] (f(\vec{y}) = f(\vec{x}) + (\nabla f(\vec{x}))^T (\vec{y} - \vec{x}) + \frac{1}{2} (\vec{y} - \vec{x})^T H_f(\lambda_w \vec{x} + (1 - \lambda_w) \vec{y}) (\vec{y} - \vec{x})))$

C6-C7: Function Coercivity, Critical Point Test of Semi-Definiteness

Coercive Function $f : \mathbb{R}^n \rightarrow \mathbb{R}$

Definition: $\lim_{||\vec{x}|| \rightarrow \infty} f(\vec{x}) \rightarrow +\infty \overset{\text{def}}{\Leftrightarrow} \forall M > 0 (\exists r > 0) (\forall \vec{x} \in N_r(\vec{0})^c) (f(\vec{x}) > M)$ i.e, all paths to $\pm\infty$ from $\vec{0}$ tends to $+\infty$

‘. \forall finite $p \in \mathbb{N}^+ (||\vec{x}||_p = (\sum_{k=1}^n |x_k|^p)^{\frac{1}{p}} \Rightarrow ||\vec{x}||_\infty = \lim_{p \rightarrow \infty} ||\vec{x}||_p \cong \lim_{p \rightarrow \infty} (\max_{|x_i|} \{|x_i|^p\})^{\frac{1}{p}} = \max\{|x_i| : i \in [1, n]\}$

‘. $||\vec{x}|| \rightarrow \infty \Leftrightarrow ||\vec{x}||_\infty \rightarrow \infty$. Besides, $\forall \vec{x} \in \mathbb{R}^n (||\vec{x}||_\infty \leq ||\vec{x}|| \leq n ||\vec{x}||_\infty)$.

f coercive $\Rightarrow \exists$ finite global minimiser (‘Weierstrass EVT on closed ball $\overline{N_r(\vec{0})}$). Converse non-eg: $f(x) = \frac{x^2}{1+x^2}$

‘. If coercive f has local minimisers, global \vec{x}^* must be within.

Interior Critical Point Tests

$\vec{x}^* \begin{cases} \text{local minimiser} & \Rightarrow \nabla f(\vec{x}^*) = \vec{0} \wedge H_f(\vec{x}^*) \text{ p.s.d.} \text{ [Non-eg: } x = 0, f(x) = -x^4] \\ \text{local maximiser} & \Rightarrow \nabla f(\vec{x}^*) = \vec{0} \wedge H_f(\vec{x}^*) \text{ n.s.d.} \text{ [Non-eg: } x = 0, f(x) = x^4] \\ \text{strict local minimiser} & \Leftarrow \nabla f(\vec{x}^*) = \vec{0} \wedge H_f(\vec{x}^*) \text{ p.d.} \text{ [Non-eg: } x = 0, f(x) = x^4 \text{ with } H_f(0) = 0 \text{ p.s.d. only}] \\ \text{strict local maximiser} & \Leftarrow \nabla f(\vec{x}^*) = \vec{0} \wedge H_f(\vec{x}^*) \text{ n.d.} \text{ [Non-eg: } x = 0, f(x) = -x^4 \text{ with } H_f(0) = 0 \text{ n.s.d. only}] \\ \text{saddle point} & \Leftarrow \nabla f(\vec{x}^*) = \vec{0} \wedge H_f(\vec{x}^*) \text{ indef [Non-eg: } f(x, y) = x^3 - y^3, H_f(0, 0) = (0) \text{ singular p/n.s.d.}] \end{cases}$

f coercive $\wedge \partial D_f = \emptyset$ (‘unconstrained) \Rightarrow any global minimiser is interior critical point, saving sufficient ∂D_f analysis.

Convex NLP format: min/max f(x̄) s.t. ⌐convex D ≠ ∅, f convex D ⇒ ∅, f convex on D **non-convex if constrained**
x̄* local minimiser (maximiser) ∧ f convex (concave) ⇒ x̄* global minimiser (maximiser) [x̄* strict ⇔ f strict]
Unconstrained Convex Quadratic Programming format

- min/max q(x̄) = ½ x̄^T Q x̄ + c^T x̄, Q **symmetric s.t.** x̄ ∈ **open convex** D ≠ ∅, q **convex on** D. Properties:
- q(x̄) convex on D ⇔ Q p.s.d. (∴ ∇q(x̄) = Q x̄ + c̄ ⇒ H_q(x̄) = Q, 2nd Partial Derivative Test)
 - q(x̄) convex (concave) on D ⇒ (x̄* local (hence global) minimiser (maximiser) ⇔ Q x̄* = -c)
 - inf_{x̄ ∈ ℝⁿ} {½ x̄^T Q x̄ - c^T x̄} = $\begin{cases} -\frac{1}{2} \bar{w}^T Q \bar{w} & \forall \bar{c} \in \text{Col}(Q) \stackrel{\text{d.f.}}{\Leftrightarrow} \exists \bar{w} \in \mathbb{R}^n (\bar{c} = Q \bar{w}) [\text{Non-unique } \bar{w} \Leftrightarrow \text{Nullity}(Q) > 0] \\ -\infty & \forall \bar{c} \notin \text{Col}(Q) \text{ (}\therefore Q \text{ deficient-rank)} \end{cases}$

Above infinum takes Q symmetric p.s.d so LHS convex, -½ w̄^T Q w̄ unique.
C8: 3 Numerical Methods to 1-variable Unconstrained NLP Bisection (aka Binary Search) Method (∴ Bolzano's Theorem)

- Necessary conditions: f continuous, 1-differentiable
- Pick interval [a₁, b₁] s.t. f'(a₁)f'(b₁) ≤ 0, and width tolerance ε > 0.
 - if f'(a₁) = 0: **return** a₁. if f'(b₁) = 0: **return** b₁.
 - while b_k - a_k > 2ε:
 - med = a_k + $\frac{b_k - a_k}{2}$
 - if f'(med) = 0: return med.
 - if f'(med)f'(b_k) < 0: a_{k+1} = med; b_{k+1} = b_k
 - else b_{k+1} = med; a_{k+1} = a_k
 - return** a_f + $\frac{b_f - a_f}{2}$

Convergence: b_k - a_k = $\frac{b_1 - a_1}{2^{k-1}} \Rightarrow f \leq \lceil \log_2(\frac{b_1 - a_1}{\epsilon}) \rceil \Leftrightarrow$ max steps = $\lceil \log_2(\frac{b_1 - a_1}{\epsilon}) \rceil$ (∴ 1-indexing)

Newton's Method (∴ Truncated Taylor's Expansion)

Necessary conditions: f continuous, 2-differentiable

∴ f(x) ≈ q(x) = f(x_k) + f'(x_k)(x - x_k) + ½ f''(x_k)(x - x_k)² = ½ f''(x_k)[x - x_k + $\frac{f'(x_k)}{f''(x_k)}$]² + [f(x_k) - $\frac{(f'(x_k))^2}{2f''(x_k)}$]

∴ Newton's univariate recurrence: x* ≈ x_{k+1} ≈ x_k - $\frac{f'(x_k)}{f''(x_k)}$

- Pick x₀, derivative tolerance ε > 0.
 - while |f'(x_k)| > ε : x_{k+1} = x_k - (f''(x_k))⁻¹ f'(x_k)
 - return** x_f
- Convergence for close x₀ to x*: quadratic.

Golden Section Method (∴ Unimodality)

Necessary conditions: f continuous, unimodal

- Pick interval [a₁, b₁] s.t. ∃! local minimiser x* ∈ [a₁, b₁], and width tolerance ε > 0. Cache golden ratio φ = $\frac{\sqrt{5}-1}{2}$
- ℓ₁ = φa₁ + (1 - φ)b₁; h₁ = (1 - φ)a₁ + φb₁
- while b_k - a_k > ε:
- if f(ℓ_k) > f(h_k) : a_{k+1} = ℓ_k; b_{k+1} = b_k; h_{k+1} = ℓ_k + φ(b_k - ℓ_k); ℓ_{k+1} = h_k (∴ 4th assignment: cost-saving narrowing)
- else: b_{k+1} = h_k; a_{k+1} = a_k; ℓ_{k+1} = h_k - φ(h_k - a_k); h_{k+1} = ℓ_k
- return** $\frac{a_f + b_f}{2}$

Convergence: b_k - a_k = φ^{k-1}(b₁ - a₁) ⇒ f ≤ $\lceil \log_{\phi^{-1}}(\frac{b_1 - a_1}{\epsilon}) \rceil \Leftrightarrow$ max steps = $\lceil \log_{\phi^{-1}}(\frac{b_1 - a_1}{\epsilon}) \rceil$ (∴ 1-indexing)

C9-C11: 3 Numerical Methods to 1-variable Unconstrained NLP

Multivariable Newton's Method

4 Necessary conditions: f continuous, 2-differentiable, ∀k ∈ ℕ₀ (H_f(x̄^(k)) invertible), x̄⁽⁰⁾ close to unknown x*
H_f(x̄^(k)) Lipschitz-continuous: (∃r, L > 0)(∀x̄, ȳ ∈ N_r(x̄*))((||H_f(x̄) - H_f(ȳ)|| ≤ L||x̄ - ȳ||)

- Pick x̄⁽⁰⁾, and derivative tolerance ε > 0.
- while ||∇f(x̄^(k))|| ≥ ε: x̄^(k+1) = x̄^(k) - H_f(x̄^(k))⁻¹ ∇f(x̄^(k))
- return** x̄^(f)

Convergence for close x̄⁽⁰⁾ to x*: quadratic. Optionally improve via Armijo line search:

- ε_{arm} ∈ (0, 0.5); r ∈ (0, 1); t = 1 usually
- while f(x̄^(k) + t d̄^(k)) > f(x̄^(k)) + ε_{arm} t ∇f(x̄^(k))^T d̄^(k): t = t * r [Generalisable to other methods' d̄^(k)]

Steepest Descent Method

Necessary conditions: f continuous, 1-differentiable

- Pick x̄⁽⁰⁾, and derivative tolerance ε > 0
- while ||∇f(x̄^(k))|| ≥ ε: x̄^(k+1) = x̄^(k) - t_k ∇f(x̄^(k)), t_k = argmin_{t ≥ 0} g(x̄^(k) - t_k ∇f(x̄^(k))) default line search.
- return** x̄^(f)

Convergence for quadratic q(x̄) := ½ x̄^T Q x̄ - c^T x̄, ρ(Q) := ($\frac{\kappa(Q)-1}{\kappa(Q)+1}$)² = 1 - $\frac{4}{\kappa(Q)+1} + \frac{4}{(\kappa(Q)+1)^2} \approx 1 - \frac{4}{\kappa(Q)}$
Matrix condition κ(Q) := $\frac{\lambda_{\max}(Q)}{\lambda_{\min}(Q)}$ = ($\frac{\text{major axis}}{\text{minor axis}}$)² ⇒ step = $\lfloor \frac{\ln \epsilon}{\ln \rho(Q)} \rfloor + 1$. Step length t_k = $\frac{\|\nabla f(\bar{x}^{(k)})\|^2}{\langle \nabla f(\bar{x}^{(k)}), Q \nabla f(\bar{x}^{(k)}) \rangle}$

Conjugate Descent Method only for convex quadratic programming min f(x̄) := ½ x̄^T Q x̄ - c^T x̄

Necessary conditions: Q symmetric p.d (∴ invertible).

- Pick x̄⁽⁰⁾. Exact method lacks ε. Compute initial residual r̄⁽⁰⁾ := -∇f(x̄⁽⁰⁾) := Q x̄⁽⁰⁾ - c̄; p̄⁽⁰⁾ = -r̄⁽⁰⁾.
- while r̄⁽⁰⁾ ≠ 0̄: x̄^(k+1) = x̄^(k) + t_k p̄^(k) = x̄^(k) - $\frac{\langle \bar{r}^{(k)}, \bar{p}^{(k)} \rangle}{\langle \bar{p}^{(k)}, Q \bar{p}^{(k)} \rangle}$ p̄^(k)
- r̄^(k+1) = ∇f(x̄^(k+1)) = Q x̄^(k+1) - c̄ = r̄^(k) + t_k Q p̄^(k) = r̄^(k) - $\frac{\langle \bar{r}^{(k)}, \bar{p}^{(k)} \rangle}{\langle \bar{p}^{(k)}, Q \bar{p}^{(k)} \rangle}$ Q p̄^(k)
- p̄^(k+1) = -r̄^(k+1) + β_{k+1} p̄^(k) = -r̄^(k+1) + $\frac{\langle \bar{r}^{(k+1)}, Q \bar{p}^{(k)} \rangle}{\langle \bar{p}^{(k)}, Q \bar{p}^{(k)} \rangle}$ p̄^(k) = -r̄^(k+1) + $\frac{\langle \bar{r}^{(k+1)}, \bar{r}^{(k+1)} \rangle}{\langle \bar{r}^{(k)}, \bar{r}^{(k)} \rangle}$ p̄^(k)

Properties: {p̄⁽⁰⁾ ... p̄⁽ⁿ⁾} conjugate wrt Q $\stackrel{\text{d.f.}}{\Leftrightarrow}$ (∀i ≠ j)(⟨p̄⁽ⁱ⁾, Q p̄^(j)⟩ = 0). Conjugacy non-transitive, symmetric iff Q is.
r̄^(k+1) = r̄⁽⁰⁾ + ∑_{i=0}^k t_i p̄⁽ⁱ⁾ inductively ⇒ (∀i < k)(⟨r̄^(k), p̄⁽ⁱ⁾⟩ = 0) given p̄⁽⁰⁾ = -r̄⁽⁰⁾ base case.

Convergence: x̄^(k) = x̄⁽⁰⁾ + ∑_{i=0}^{k-1} t_i p̄⁽ⁱ⁾ ∈ x̄⁽⁰⁾ + Span{p̄⁽⁰⁾ ... p̄^(k-1)} ∧ p̄⁽ⁱ⁾'s L.I. ⇒ x̄⁽ⁿ⁾ min f in ℝⁿ ⇒ exact steps ≤ n

C12-C13: Regular & KKT Points

Regular Points x̄* := feasible points fulfilling Linear Independence Constraint Qualification (LICQ)

LICQ: Active constraint gradient set (Constraint Jacobian) A(x̄*) = {∇g_i(x̄*)} ∪ {∇h_j(x̄*) : j ∈ J(x̄*) active indices} LI
Corollary: βg_i ⇒ all interior points are vacuously-regular (J(x̄*) = ∅ if all inequalities inactive/slack)

Karush-Kuhn-Tucker (KKT) points x̄* Defintion

- Stationarity: (∃ Lagrangian multiplier λ_i^{*}, μ_j^{*}'s)(Lagrangian L_{x̄}(x̄*, λ̄*, μ̄*) := f(x̄*) + ∑ λ_i^{*} g_i(x̄*) + ∑ μ_j^{*} h_j(x̄*) = 0̄)
 - Primal feasibility: (∀i)(g_i(x̄*) = 0̄) ∧ (∀j)(h_j(x̄*) = 0̄)| 3. Dual feasibility: (∀j)(μ_j^{*} ≥ 0) from stationarity
 - Complementary slackness: (∀j)(μ_j^{*} h_j(x̄*) = 0̄ ⇔ h_j(x̄*) active ∨ μ_j^{*} = 0 ∨ h_j(x̄*) inactive)
- Strict complementarity (not needed): (∀h_j(x̄*) active)(μ_j^{*} > 0)

KKT 1ONC: x̄* regular ∧ local min ⇒ x̄* KKT

Analogous unconstrained 1st-Order Partial Derivative Test (∴ x̄* interior): x̄* local minimiser ⇒ x̄* critical point

Critical point x̄* : ∇f(x̄*) = 0̄ ∨ undefined | β local constrained general tests ⇒ KKT 1ONC hard to apply.

KKT 2ONC: x̄* regular ∧ local min ∧ f, g_i, h_j 2-diff ⇒ x̄* KKT ∧ (∀ȳ ∈ C(x̄*, λ̄*, μ̄*)) (H_L(x̄*) p.s.d.)

Critical cone C(x̄*, λ̄*, μ̄*) := {ȳ ∈ ℝⁿ : $\begin{cases} \langle \nabla g_i(\bar{x}^*), \bar{y} \rangle = 0 \\ \langle \nabla h_j(\bar{x}^*), \bar{y} \rangle = 0 & \forall \text{ active } h_j(\bar{x}^*), \mu_j^* > 0 \\ \langle \nabla h_j(\bar{x}^*), \bar{y} \rangle \leq 0 & \forall \text{ active } h_j(\bar{x}^*), \mu_j^* = 0 \end{cases}$ independent of inactive h_j(x̄*)}

C(x̄*, λ̄*, μ̄*) ⊆ Linearised feasible direction set F(x̄*) := {ȳ ∈ ℝⁿ : ⟨∇g_i(x̄*), ȳ⟩ = 0 ∧ ⟨∇h_j(x̄*), ȳ⟩ ≤ 0 ∀ active h_j(x̄*)}
Analogous unconstrained 2nd-Order Partial Derivative Test (∴ x̄* interior): x̄* local minimiser ∧ f 2-diff ⇒ H_f(x̄*) p.s.d.
Corollary: X regular throughout, or not for finite points ⇒ global (∴ local) minimiser x̄* KKT or in finite points.

KKT 2OSC: x̄* KKT (∀ȳ ∈ C(x̄*, λ̄*, μ̄*)) (H_L(x̄*) p.d.) ⇒ x̄* strict local minimiser, H_L(x̄*) = ∇_{x̄} L(x̄*, λ̄*, μ̄*)

Analogue unconstrained 2nd-Order Partial Derivative Test (∴ x̄* interior): ∇f(x̄*) = 0̄ ∧ H_f(x̄*) pd ⇒ x̄* strict local min

Farka's Lemma: (∀A ∈ M_{m × n}(ℝ))(∀b̄ ∈ ℝ^m)(∃x̄ ∈ ℝⁿ)(A x̄ = b̄ ∧ x̄ ≥ 0̄ **item-wise**) ⊕ (∃ȳ ∈ ℝ^m)(b̄^T ȳ < 0̄ ≤ A^T ȳ))

C14-C15: Constrained Convex NLP

Convex Implications and Slater's Condition

Constrained convex NLP format: min_{x̄ ∈ convex X} convex f(x̄) s.t. $\begin{cases} g_i(\bar{x}) := \bar{a}_i^T \bar{x} - b_i = 0 \text{ affine/linear} \\ h_j(\bar{x}) \leq 0 \text{ convex} \\ \text{Feasible set } X \text{ convex} \end{cases}$
Slater's Condition: (∃ strictly-interior x̄ ∈ X)(g_i(x̄) = 0 ∧ h_j(x̄) < 0)
Corollaries: 1. x̄* KKT in convex program ⇒ x̄* global minimiser

- x̄* global-minimiser in convex program ∧ Slater's ⇒ x̄* KKT. ∴ If holds
- Linear convex ECP: J(x̄*) = ∅ ⇒ vacuous LICQ ∧ Slater's ⇒ (x̄* local min ⇔ x̄* KKT ⇔ x̄* global min)
- ∴ ⇒: KKT 1ONC, convexity, ⇐: Global, Slater's
- Linear general ECP (X ≠ ∅), x̄ KKT ⇔ x̄* global min. ∴ Suffice to check stationarity.

Orthogonality

ū ⊥ v $\stackrel{\text{d.f.}}{\Leftrightarrow}$ Standard (ū, v) = ū^T v = 0| subspaces U ⊥ V $\stackrel{\text{d.f.}}{\Leftrightarrow}$ (∀ū ∈ U)(∀v̄ ∈ V)(ū ⊥ v̄). Eg: Im(A^T) ⊥ Ker(A) provably.

Constrained (convex non-ECP) quadratic programming: min q(x̄) := ½ x̄^T Q x̄ + c^T x̄ **s.t.** A x̄ - b̄ ≤ 0̄ ∧ x̄ ≥ 0̄

Support Vector Machine (SVM): max $\frac{2}{\|\bar{w}\|} \equiv \min f(\bar{w}) = \frac{1}{2} \|\bar{w}\|^2$ s.t. h_j(w̄) := -y_j(x_j^T w̄ + b) + 1 ≤ 0, y_j ∈ {±1}

∴ Q = I_n symmetric p.d. ∧ c̄ = 0̄. x̄* KKT in convex QP SVM ⇒ x̄* global-minimiser, hence find global minimum.

Hinge-loss SVM: min f(w̄) = ½ ||w̄||² + λ ∑ ζ_i s.t. h_j(w̄) := -y_j(x_j^T w̄ + b) - ζ_j + 1 ≤ 0 ∧ h_k(w̄) := -ζ_k ≤ 0, y_j ∈ {±1}

Constrained (convex) LP: min f(x̄) := c^T x̄ **s.t.** g_i(x̄) := b_i - a_i^T x_i = 0 ∧ h_j(x̄) := -e_j^T x_j ≤ 0 (∴ x̄ ≥ 0)

C16-C17: Lagrangian Dual Problem

Lagrangian Dual Formulation

∴ Primal minimisation format: min_{x̄ ∈ X} {max_{x̄ ∈ ℝ^m, μ̄ ≥ 0̄_n} {L(x̄, λ̄, μ̄) := f(x̄) + ∑ λ_i g_i(x̄) + ∑ μ_j h_j(x̄)}}

∴ Swap: Lagrangian dual function θ(λ̄, μ̄) := inf_{x̄ ∈ X} L(x̄, λ̄, μ̄) ⇒ Lagrangian dual problem D := max_{λ̄ ∈ ℝ^m, μ̄ ≥ 0̄_n} θ(λ̄, μ̄)

- Corollaries: 1. θ concave independent of f, g_i, h_j's convexity.
- If single IECP, graphically, primal := min_{x ≤ 0} y ≥ dual := min y-intercept of downward-sloping tangent

Duality Theorems

Weak Duality Theorem: (∀ primal feasible x̄)(∀ dual feasible (λ̄, μ̄))(f(x̄) ≥ θ(λ̄, μ̄) ⇔ duality gap ≥ 0)

Strong Duality Theorem: X convex set ∧ f, h_j convex ∧ g_i affine/linear ∧ Slater's condition ⇒ P - D := duality gap = 0

∴ If f(x̄*) = θ(λ̄*, μ̄*), then L(x̄, λ̄*, μ̄*) = inf_{x̄ ∈ X} L(x̄, λ̄*, μ̄*) := θ(λ̄*, μ̄*) = f(x̄*) = L(x̄*, λ̄*, μ̄*) ⇒ (μ̄*)^T h(x̄*) = 0.

∴ P = D ⇒ complementary slackness.

Convex proximal-mapping format P: min_{B x̄ = ū} f(x̄) := ½ ||x̄ - c̄||² + ρ ||ū||_∞

∴ D := max_{λ̄ ∈ ℝ^m, ∑ |λ_i| ≤ ρ} θ(λ̄) := -½ ||B^T λ̄||² + ⟨B c̄, λ̄⟩

C18: Saddle Point of Lagrangian function L(x̄, λ̄, μ̄)

Definition: (x̄*, λ̄*, μ̄*) saddle-point $\stackrel{\text{d.f.}}{\Leftrightarrow}$ (∀x̄ ∈ X)(∀(λ̄, μ̄) ∈ ℝ^m × ℝⁿ₊)(L(x̄*, λ̄, μ̄) ≤ L(x̄*, λ̄*, μ̄*) ≤ L(x̄, λ̄*, μ̄*))

Corollaries: 1. (x̄*, λ̄*, μ̄*) saddle-point ⇒ x̄* primal global minimiser ∧ (λ̄*, μ̄*) dual global maximiser.

- (x̄*, λ̄*, μ̄*) saddle-point ∧ x̄* ∈ X^o ⇒ (x̄*, λ̄*, μ̄*) KKT (∴ ∇_{x̄} L(x̄*, λ̄*, μ̄*) = 0̄)
- x̄* KKT-point of convex program with Lagrangian λ̄*, μ̄* ⇒ (x̄*, λ̄*, μ̄*) saddle-point.
- f(x̄*) > θ(λ̄*, μ̄*) ⇒ β saddle.

C19-C22: 5 Numerical Methods to ≥ 1-variable Constrained NLP
Subdifferential Terminologies & Properties

Subdifferential $\partial f(\vec{x}^*) := \{\text{subgradient } \zeta : (\forall \vec{x} \in X)(\begin{cases} f(\vec{x}) \leq f(\vec{x}^*) + \zeta^T(\vec{x} - \vec{x}^*) & \forall \text{concave } f(\text{eg: } \theta) \\ f(\vec{x}) \geq f(\vec{x}^*) + \zeta^T(\vec{x} - \vec{x}^*) & \forall \text{convex } f \end{cases})\}$ convex set

∴ Constraint vector $\beta(\vec{x}) := (g(\vec{x}), h(\vec{x}))$, Lagrangian multiplier vector $\vec{w} := (\vec{\lambda}, \vec{\mu})$

∴ (∀ Lagrangian minimiser $\vec{x}^* \in X(\vec{w})$)($\theta(\vec{w}) = L(\vec{x}^*, \vec{w}^T \beta(\vec{x}^*))$). Corollaries:

1. If $X(\vec{w}) = \{\vec{x}^*\}$ singleton, then $\nabla \theta(\vec{w}) = \beta(\vec{x}^*)$ well-defined at \vec{w} | 2. $\partial \theta(\vec{w}) = \text{conv}\{\beta(\vec{x}^*) : \vec{x}^* \in X(\vec{w})\}$
3. $(\exists \vec{x}^* \in X(\vec{w}))(\vec{d}^T \beta(\vec{x}^*) \leq \text{r.o.c. of } \theta \text{ in direction } \vec{d} = \text{directional derivative } \theta'(\vec{w}; \vec{d}) := \lim_{\gamma \rightarrow 0^+} \frac{\theta(\vec{w} + \gamma \vec{d}) - \theta(\vec{w})}{\gamma})$

(3)'s continuous analogy: $\vec{d}^T \nabla f(\vec{x}) = \text{r.o.c. of } f \text{ in direction } \vec{d}$ | ∴ $\theta'(\vec{w}; \vec{d}) = \inf_{\zeta \in \partial \theta(\vec{w})} \{\vec{d}^T \zeta\}$

∴ \vec{d} ascent direction of θ at $\vec{w} \stackrel{\text{def}}{\Leftrightarrow} (\exists r > 0)(\forall t \in (0, r))(\theta(\vec{w} + t\vec{d}) > \theta(\vec{w})) \stackrel{\text{def}}{\Leftrightarrow} (\exists r > 0)(\forall \zeta \in \partial \theta(\vec{w}))(\vec{d}^T \zeta \geq r)$

∴ Normalised steepest ascent direction of θ at \vec{w} , $\vec{d} := \frac{\hat{\zeta}}{\|\hat{\zeta}\|}, \hat{\zeta} := \min\{\zeta \in \partial \theta(\vec{w}) = \text{conv}\{\beta(\vec{x}^*) : \vec{x}^* \in X(\vec{w})\}\}$

∴ Ascent directions' cone $C(\vec{d}) := \{\vec{d} : \inf_{\zeta \in \partial \theta(\vec{w})} \{\vec{d}^T \zeta\} \geq 0\}$. Graphically, $\forall \zeta \in \partial \theta(\vec{w})$, angle between ζ and $\theta \leq \frac{\pi}{2}$.

Steepest Ascent Direction Method for $\max_{\vec{w} \in \mathbb{R}^m \times \mathbb{R}^p_+} \theta(\vec{w}) := \inf_{\vec{x} \in X} \{f(\vec{x}) + \vec{w}^T \beta(\vec{x})\}$ dual maximisation

1. Pick $\vec{w}^{(0)}$, and min-norm tolerance $\epsilon > 0$. Compute min-norm $\hat{\zeta}$ as below.
2. while $\frac{1}{2} \|\hat{\zeta}\|^2 > \epsilon$: subdifferential $\partial \theta(\vec{w}^{(k)}) = \text{conv}\{\beta(\vec{x}^*) : \vec{x}^* \in X(\vec{w}^{(k)})\}$
3. normalised steepest ascent direction $\vec{d}^{(k)} = \frac{\hat{\zeta}}{\|\hat{\zeta}\|}, \hat{\zeta} = \text{argmin}_{\zeta \in \partial \theta(\vec{w}^{(k)})} \{\frac{1}{2} \|\zeta\|^2\}$
4. $\vec{w}^{(k+1)} = \vec{w}^{(k)} + t_k \vec{d}^{(k)}, t_k = \text{argmin}_{t \geq 0} \{\vec{w}^{(k)} + t \vec{d}^{(k)}\}$
5. **return** $\vec{w}^{(f)}$. (∴ if steepest descent, only edit Step 3 to $-\vec{d}^{(k)}$)

Frank-Wolfe Algorithm for convex IECP $\min_{\vec{x} \in \mathbb{R}^n} f(\vec{x})$ s.t. $A\vec{x} \leq \vec{b}$

1. Pick $\vec{x}^{(0)}$, and width tolerance $\epsilon > 0$. Set $\ell_0 = -\infty, h_0 = f(\vec{x}^{(0)})$.
2. while $h_k - \ell_k > \epsilon$: solve $\hat{\vec{x}}^{(k)} = \text{argmin } z(\vec{x}) = f(\vec{x}^{(k)}) + \nabla f(\vec{x}^{(k)})^T (\vec{x} - \vec{x}^{(k)})$ s.t. $A\vec{x} \leq \vec{b}$ constrained linear subproblem.
3. $\ell_{k+1} = \max(\ell_k, z(\hat{\vec{x}}^{(k)})) = \max\{\ell_0 \dots \ell_k, z(\hat{\vec{x}}^{(k)})\}$ nondecreasing.
4. Optional line search: feasible direction $\vec{d}^{(k)} := \hat{\vec{x}}^{(k)} - \vec{x}^{(k)} \Rightarrow t_k = \text{argmin}_{t \in [0, 1]} \{f(\vec{x}^{(k)} + t \vec{d}^{(k)})\}$
5. $\vec{x}^{(k+1)} = \vec{x}^{(k)} + t_k \vec{d}^{(k)}; h_{k+1} = f(\vec{x}^{(k+1)})$ (∴ h_i may fluctuate)
6. **return** $\vec{x}^{(f)}$. Corollaries: 1. f convex \Rightarrow Taylor's approximate $z(\vec{x}) \leq f(\vec{x})$ | 2. $\vec{x}^{(k)}$ minimises $f \Leftrightarrow \vec{x}^{(k)}$ minimises $z_k(\vec{x})$

Quadratic Penalty Method for (not-necessarily linear/convex) ECP $\min_{\vec{x} \in \mathbb{R}^n} f(\vec{x})$ s.t. $g(\vec{x}) = \vec{0}$

1. Pick $\vec{x}^{(0)}$, and constraints tolerance $\epsilon > 0$. Conventionally set initial penalty $\mu_0 = 1$.
2. Crucially formulate $\min_{\vec{x} \in \mathbb{R}^n} Q(\vec{x}; \mu_k) := f(\vec{x}) + \frac{1}{2\mu_k} \sum g_i^2(\vec{x})$ unconstrained subproblem.
3. while $\|g(\vec{x}^{(k)})\| \geq \epsilon$: solve $\vec{x}^{(k+1)} = \text{argmin} Q(\vec{x}; \mu_k)$ unconstrained subproblem, oft-via same ϵ and Newton's Method.
4. $\mu_{k+1} = \rho \mu_k$ for some $\rho \in (0, 1)$ i.e, penalise positive constraint values more.
5. **return** $\vec{x}^{(f)}$.

Corollaries: 1. $(\forall k \in \mathbb{N}_0)(\vec{x}^{(k+1)})$ exactly $\min Q(\vec{x}; \mu_k)$ at $\epsilon_k = 0) \wedge \mu_k \rightarrow 0^+ \Rightarrow \{\vec{x}^{(k)}\} \rightarrow \vec{x}^*$ exactly $\min f$.
2. If $(\forall k \in \mathbb{N}_0)(\vec{x}^{(k+1)})$ approx $\min Q(\vec{x}; \mu_k)$ s.t. $\|\nabla_{\vec{x}} Q(\vec{x}^{(k+1)}; \mu_k)\| \leq \epsilon_k \rightarrow 0^+) \wedge \mu_k \rightarrow 0^+$, then $\{\vec{x}^{(k)}\} \rightarrow \vec{x}^*$ regular feasible $\Leftrightarrow \vec{x}^*$ KKT with $\lambda_i^* = \lim_{k \rightarrow \infty} \frac{g_i(\vec{x}^{(k+1)})}{\mu_k}$. Note \vec{x}^* may not be regular/feasible.

Augmented Lagrangian Method for (not-necessarily linear/convex) ECP $\min_{\vec{x} \in \mathbb{R}^n} f(\vec{x})$ s.t. $g(\vec{x}) = \vec{0}$

1. Pick $\vec{x}^{(0)}$, and constraints tolerance $\epsilon > 0$. Conventionally set initial penalty $\mu_0 = 1$.
2. Crucially formulate $\min_{\vec{x} \in \mathbb{R}^n} L_A(\vec{x}; \vec{\lambda}^{(k)}, \mu_k) := f(\vec{x}) + \sum \lambda_i^{(k)} g_i(\vec{x}) + \frac{1}{2\mu_k} \sum g_i^2(\vec{x})$ unconstrained subproblem.
3. while $\|g(\vec{x}^{(k)})\| \geq \epsilon$: solve $\vec{x}^{(k+1)} = \text{argmin} L_A(\vec{x}; \vec{\lambda}^{(k)}, \mu_k)$ unconstrained subproblem, oft same ϵ and Newton Method.
4. $\forall i(\lambda_i^{(k+1)} = \lambda_i^{(k)} + \frac{g_i(\vec{x}^{(k+1)})}{\mu_k}), \mu_{k+1} = \rho \mu_k$ for some $\rho \in (0, 1)$.
5. **return** $\vec{x}^{(f)}$. Step 2 can check $\|\nabla_{\vec{x}} L_A(\vec{x}^{(k)}; \vec{\lambda}^{(k)}, \mu_k)\| < \epsilon$ instead.

Corollary: \forall KKT \vec{x}^* with $\vec{\lambda}^*$ to non-augmented $L(\vec{x}, \vec{\lambda})$, if \vec{x}^* regular local min of $f \wedge H_L(\vec{x}^*, \vec{\lambda}^*)$ p.d. on $C(\vec{x}^*, \vec{\lambda}^*)$, then

1. $(\exists \gamma > 0)(\forall \mu \in (0, \gamma])(\vec{x}^*$ strict local minimiser of $L_A)$ (∴ $\mu_k \rightarrow 0^+$ not needed, allowing better κ)
2. $(\exists \delta, \epsilon, M > 0)(\mu_k \|\vec{\lambda}^{(k)} - \vec{\lambda}^*\| \leq \delta \wedge \mu_k \in (0, \gamma) \Rightarrow \exists!$ global minimiser $\vec{x}^{(k+1)}$ to $\min_{\|\vec{x} - \vec{x}^*\| \leq \epsilon} L_A(\vec{x}; \vec{\lambda}^{(k)}, \mu_k) \wedge \|\vec{x}^{(k+1)} - \vec{x}^*\| \leq M \mu_k \|\vec{\lambda}^{(k)} - \vec{\lambda}^*\| \wedge \vec{\lambda}^{(k)} \rightarrow \vec{\lambda}^*$ Q-linearly $\wedge \vec{x}^{(k+1)} \rightarrow \vec{x}^*$ R-linearly.

Barrier Function Method

1. Pick $\vec{x}^{(0)}$, and constraints tolerance $\epsilon > 0$. Conventionally set initial penalty $\mu_0 = 1$.
2. Crucially formulate $\min_{\vec{x} \in \mathbb{R}^n} P(\vec{x}; \mu_k) := f(\vec{x}) + \mu_k B(\vec{x})$, barrier $B(\vec{x}) := \sum \phi(-g_i(\vec{x}))$ s.t. $\phi : \mathbb{R}^+ \rightarrow \mathbb{R}^+ \wedge \phi' < 0 \wedge \lim_{g_i \rightarrow 0^-} \phi(-g_i(\vec{x})) \rightarrow \infty$.
3. while $\|\nabla_{\vec{x}} P(\vec{x}^{(k)}; \mu_k)\| > \epsilon$: solve $\vec{x}^{(k+1)} = \text{argmin} P(\vec{x}; \mu_k)$ unconstrained subproblem, oft same ϵ , Newton Method.
4. $\mu_{k+1} = \rho \mu_k$ for some $\rho \in (0, 1)$, or dynamically set $\{\epsilon_k\}_{k \in \mathbb{N}^+}$.
5. **return** $\vec{x}^{(f)}$.

Corollaries: 1. If $(\exists \text{ global min } \vec{x}^* \text{ of } f)(\forall r > 0)(N_r(\vec{x}^*) \cap X \neq \emptyset \text{ i.e, } \vec{x}^* \text{ limit point of a convergent subseq}\{x_\mu^*\})$, then $\lim_{\mu \rightarrow 0^+} \inf_{\vec{x} \in X} \{P(\vec{x}; \mu)\} = f(\vec{x}^*) \wedge \lim_{\mu \rightarrow 0^+} \mu B(x_\mu^*) = 0 \wedge$ any convergent subseq to a global min of f
2. $(\forall \text{convergent subseq s.t. } \{\vec{x}^{(k)}\} \rightarrow \vec{x}^*)(\vec{x}^*$ regular $\Rightarrow \vec{x}^*$ KKT with $\lambda_i^* = \lim_{k \rightarrow \infty} \mu_k \nabla \phi(-g_i(\vec{x}^{(k)}))$)
(2): Barrier function, if convergent/terminating, outputs approx-KKT-minimisers practically.

LEUNG HUAN [Student ID: A0269854J]
MA3236 Non-Linear Programming