

CS 106B, Lecture 9

Exhaustive Search and Backtracking

reading:

Programming Abstractions in C++, Chapters 8.2 - 8.3; 9

Exhaustive search

- **exhaustive search**: Exploring every possible combination from a set of choices or values.
 - often implemented recursively

Applications:

- producing all permutations of a set of values
 - enumerating all possible names, passwords, etc.
 - combinatorics and logic programming
- Often the search space consists of many ***decisions***, each of which has several available ***choices***.
 - Example: When enumerating all 5-letter strings, each of the 5 letters is a *decision*, and each of those decisions has 26 possible *choices*.

Exhaustive search

A general pseudo-code algorithm for exhaustive search:

Explore(*decisions*):

- if there are no more decisions to make: Stop.
- else, let's handle one decision ourselves, and the rest by recursion.
for each available choice *C* for this decision:
 - **Choose** *C*.
 - **Explore** the remaining decisions that could follow *C*.

Exercise: `printAllBinary`



`printBinary`

- Write a recursive function **`printAllBinary`** that accepts an integer number of digits and prints all binary numbers that have exactly that many digits, in ascending order, one per line.

– `printAllBinary(2);` `printAllBinary(3);`

00

01

10

11

000

001

010

011

100

101

110

111

- Use recursion; do not use any loops.
- How is this problem self-similar (recursive)?

Helper functions

- If the required function doesn't accept the parameters you need:
 - Write a **helper function** that accepts more parameters.
 - Extra params can represent current state, choices made, etc.

```
returnType functionName(params) {  
    ...  
    return helper(params, moreParams);  
}
```

```
returnType helper(params, moreParams) {  
    ...  
}  
  
(use moreParams to help solve the problem)
```

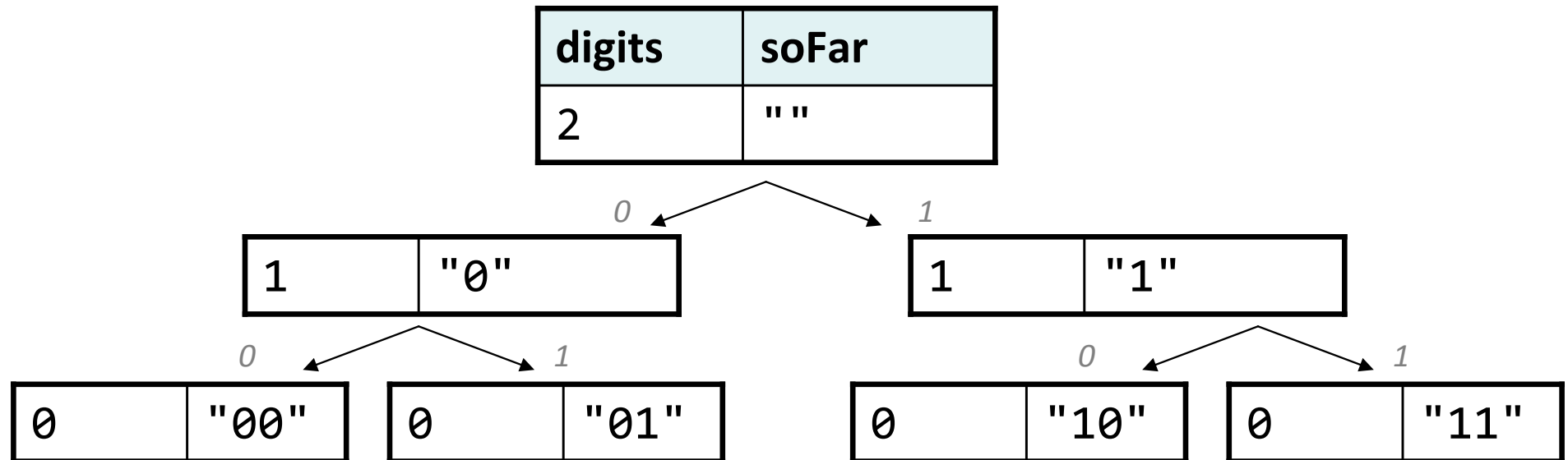
printAllBinary solution

```
void printAllBinary(int digits) {  
    printAllBinaryHelper(digits, "");  
}  
  
void printAllBinaryHelper(int digits, string soFar) {  
    if (digits == 0) {  
        cout << soFar << endl;  
    } else {  
        printAllBinaryHelper(digits - 1, soFar + "0");  
        printAllBinaryHelper(digits - 1, soFar + "1");  
    }  
}
```

- *Observation:* We write a 'helper' function to pass along the digits that were chosen so far by previous calls.

A tree of calls

- `printAllBinary(2);`



- This kind of diagram is called a *call tree* or *decision tree*.
- Think of each call as a choice or decision made by the algorithm:
 - Should I choose 0 as the next digit?
 - Should I choose 1 as the next digit?

The base case

```
void printAllBinaryHelper(int digits, string soFar) {  
    if (digits == 0) {  
        cout << soFar << endl;  
    } else {  
        printAllBinaryHelper(digits - 1, soFar + "0");  
        printAllBinaryHelper(digits - 1, soFar + "1");  
    }  
}
```

- The **base case** is where the code stops after doing its work.
 - pAB(3) -> pAB(2) -> pAB(1) -> pAB(0)
- Each call should keep track of the work it has done.
- Base case should print the result of the work done by prior calls.
 - Work is kept track of in some variable(s) - in this case, string soFar.

Exercise: `printDecimal`



`printDecimal`

- Write a recursive function **`printDecimal`** that accepts an integer number of digits and prints all base-10 numbers that have exactly that many digits, in ascending order, one per line.

– `printDecimal(2);`

00

01

02

..

98

99

`printDecimal(3);`

000

001

002

...

997

998

999

– Use recursion. *

printDecimal solution

```
void printDecimal(int digits) {  
    printDecimalHelper(digits, "");  
}  
  
void printDecimalHelper(int digits, string soFar) {  
    if (digits == 0) {  
        cout << soFar << endl;  
    } else {  
        for (int i = 0; i < 10; i++) {  
            printDecimalHelper(digits - 1, soFar +  
                                integerToString(i));  
        }  
    }  
}
```

- *Observation:* When the set of digit choices available is large, using a loop to enumerate them avoids redundancy. (*This is okay!*)

Backtracking

- **backtracking**: Finding solution(s) by trying partial solutions and then abandoning them if they are not suitable.
 - a "brute force" algorithmic technique (tries all paths)
 - often implemented recursively

Applications:

- producing all permutations of a set of values
- parsing languages
- games: anagrams, crosswords, word jumbles, 8 queens
- combinatorics and logic programming
- escaping from a maze

Backtracking algorithms

A general pseudo-code algorithm for backtracking problems:

Explore(*decisions*):

- if there are no more decisions to make: stop.
- else, let's handle one decision ourselves, and the rest by recursion.
for each available choice *C* for this decision:
 - **Choose** *C*.
 - **Explore** the remaining decisions that could follow *C*.
 - **Un-choose** *C*. (backtrack!)

Exercise: Dice roll sum



diceSum

- Write a function **diceSum** that accepts two integer parameters: a number of dice to roll, and a desired sum of all die values. Output all combinations of die values that add up to exactly that sum.

`diceSum(2, 7);`

{1, 6}
{2, 5}
{3, 4}
{4, 3}
{5, 2}
{6, 1}



`diceSum(3, 7);`

{1, 1, 5}
{1, 2, 4}
{1, 3, 3}
{1, 4, 2}
{1, 5, 1}
{2, 1, 4}
{2, 2, 3}
{2, 3, 2}
{2, 4, 1}
{3, 1, 3}
{3, 2, 2}
{3, 3, 1}
{4, 1, 2}
{4, 2, 1}
{5, 1, 1}

Easier: Dice rolls

- **Suggestion:** First just output all possible combinations of values that could appear on the dice.

`diceSum(2, 7);`

`diceSum(3, 7);`

{1, 1}	{3, 1}	{5, 1}
{1, 2}	{3, 2}	{5, 2}
{1, 3}	{3, 3}	{5, 3}
{1, 4}	{3, 4}	{5, 4}
{1, 5}	{3, 5}	{5, 5}
{1, 6}	{3, 6}	{5, 6}
{2, 1}	{4, 1}	{6, 1}
{2, 2}	{4, 2}	{6, 2}
{2, 3}	{4, 3}	{6, 3}
{2, 4}	{4, 4}	{6, 4}
{2, 5}	{4, 5}	{6, 5}
{2, 6}	{4, 6}	{6, 6}



{1, 1, 1}
{1, 1, 2}
{1, 1, 3}
{1, 1, 4}
{1, 1, 5}
{1, 1, 6}
{1, 2, 1}
{1, 2, 2}
...
{6, 6, 4}
{6, 6, 5}
{6, 6, 6}

- How is this problem recursive (self-similar)?

Examining the problem

- Generate all possible sequences of values:

- for (each possible first die value):

- for (each possible second die value):

- for (each possible third die value):

- ...

- print!

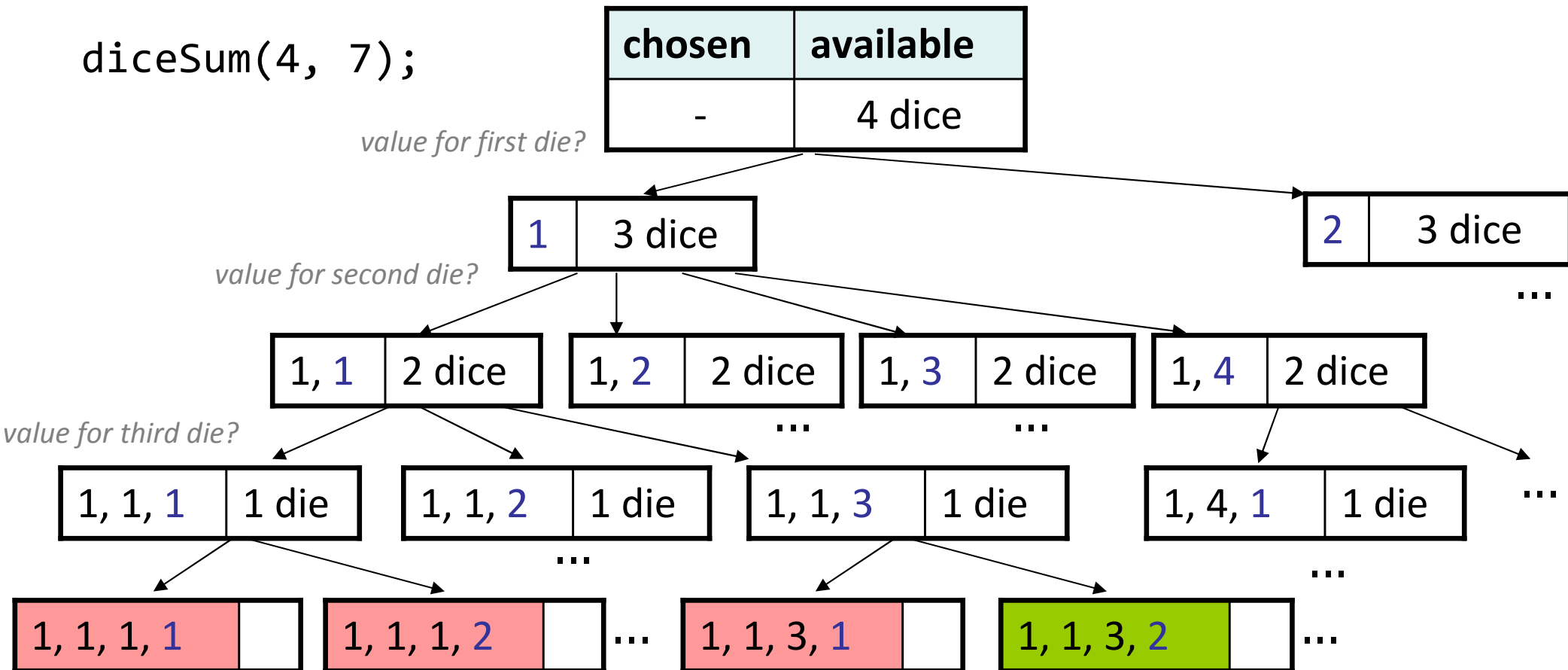


- This is called a **depth-first search**

- How many loops are needed?
 - How can we completely explore such a large search space?

A decision tree

diceSum(4, 7);



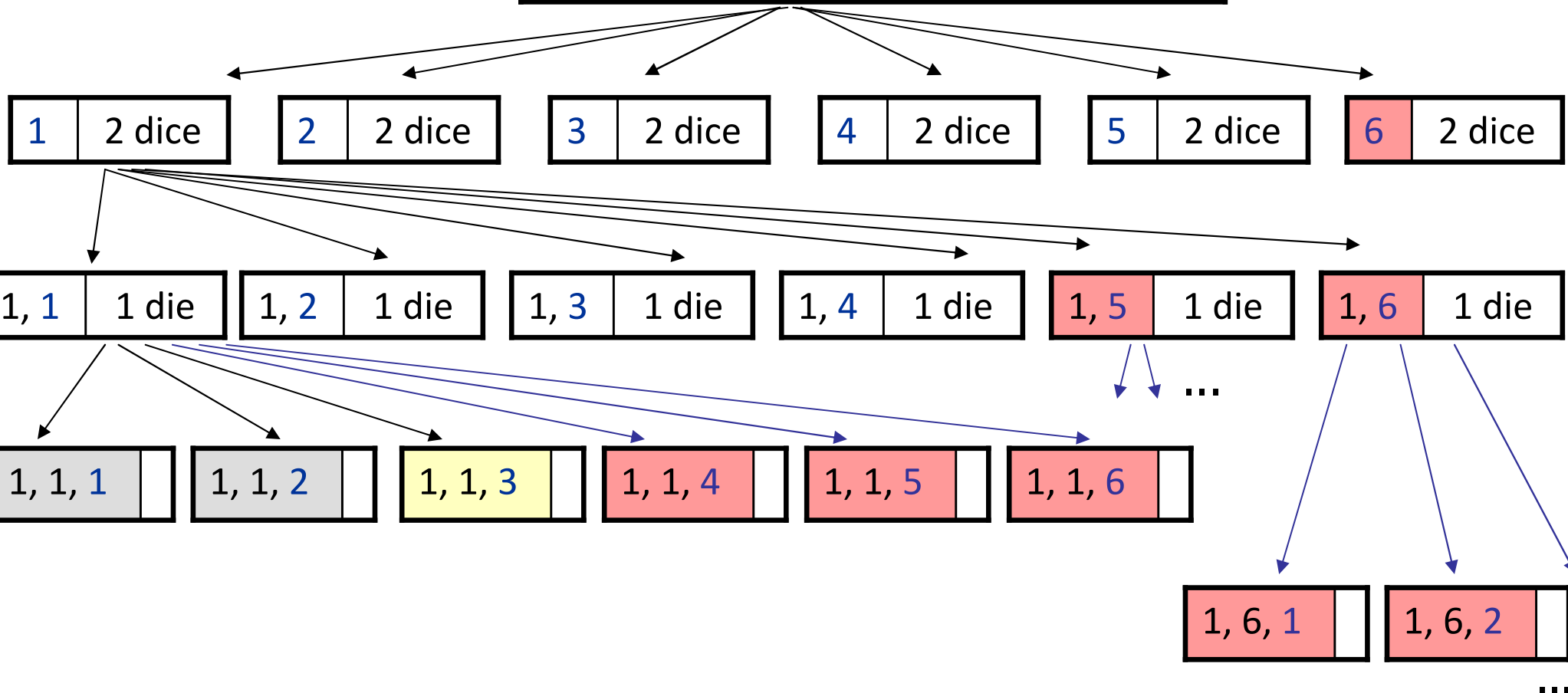
Initial solution

```
void diceSum(int dice, int desiredSum) {  
    Vector<int> chosen;  
    diceSumHelper(dice, desiredSum, chosen);  
}  
  
void diceSumHelper(int dice, int desiredSum, Vector<int>& chosen) {  
    if (dice == 0) {  
        if (sumAll(chosen) == desiredSum) {  
            cout << chosen << endl;           // base case  
        }  
    } else {  
        for (int i = 1; i <= 6; i++) {  
            chosen.add(i);                       // choose  
            diceSumHelper(dice - 1, desiredSum, chosen); // explore  
            chosen.remove(chosen.size() - 1);    // un-choose  
        }  
    }  
}  
  
int sumAll(const Vector<int>& v) { // adds the values in given vector  
    int sum = 0;  
    for (int k : v) { sum += k; }  
    return sum;  
}
```

Wasteful decision tree

diceSum(3, 5);

chosen	available	desired sum
-	3 dice	5



Optimizations

- We need not visit every branch of the decision tree.
 - Some branches are clearly not going to lead to success.
 - We can preemptively stop, or **prune**, these branches.
- Inefficiencies in our dice sum algorithm:
 - Sometimes the current sum is already **too high**.
 - (Even rolling 1 for all remaining dice would exceed the desired sum.)
 - Sometimes the current sum is already **too low**.
 - (Even rolling 6 for all remaining dice would exceed the desired sum.)
 - The code must **re-compute** the sum many times.
 - $(1+1+1 = \dots, 1+1+2 = \dots, 1+1+3 = \dots, 1+1+4 = \dots, \dots)$

diceSum solution

```
void diceSum(int dice, int desiredSum) {
    Vector<int> chosen;
    diceSumHelper(dice, 0, desiredSum, chosen);
}

void diceSumHelper(int dice, int sum, int desiredSum, Vector<int>& chosen) {
    if (dice == 0) {
        if (sum == desiredSum) {
            cout << chosen << endl;           // base case
        }
    } else if (sum + 1*dice <= desiredSum
               && sum + 6*dice >= desiredSum) {
        for (int i = 1; i <= 6; i++) {
            chosen.add(i);                      // choose
            diceSumHelper(dice - 1, sum + i, desiredSum, chosen); // explore
            chosen.remove(chosen.size() - 1);   // un-choose
        }
    }
}
```

For you to think about...

- How would you modify **diceSum** so that it prints only unique combinations of dice, ignoring order?
 - (e.g. don't print both {1, 1, 5} and {1, 5, 1})

`diceSum2(2, 7);`

{1, 6}
{2, 5}
{3, 4}
{4, 3}
{5, 2}
{6, 1}



`diceSum2(3, 7);`

{1, 1, 5}
{1, 2, 4}
{1, 3, 3}
{1, 4, 2}
{1, 5, 1}
{2, 1, 4}
{2, 2, 3}
{2, 3, 2}
{2, 4, 1}
{3, 1, 3}
{3, 2, 2}
{3, 3, 1}
{4, 1, 2}
{4, 2, 1}
{5, 1, 1}