## **Homework 4**

## **Algorithm Design**

- 6-1
  - (a) 5—3—3—4—3 通过算法得到最大独立集为 9,但实际的最大独立集为11
  - (b) 10—1—1—9 通过算法得到最大的独立集为 11, 但实际的最大独立集为19

(c)

```
int findMax(int num[]){
   int preMax = 0;
   int currMax = 0;
   for(int x : num){
      int temp = currMax;
      currMax = Max( prevMax + x, currMax );
      prevMax = temp;
   }
   return currMax;
}
```

```
f(k) = max(f(k-2)) + A_k, f(k-1)
```

复杂度分析:

时间复杂度: O(n)

空间复杂度: O(1)

• 6-2

(a)

|   | Week 1 | Week 2 | Week 3 |
|---|--------|--------|--------|
| l | 2      | 2      | 2      |
| h | 1      | 5      | 10     |

通过算法得到的最大利益为7,但实际的最大利益为12

(b)

```
int findMax(int L[], int H[]){
    L(1) = L[1];
    H(1) = 0;
    loop:
        L(i) = L[i] + max(L(i-1), H(i-1));
        H(i) = H[i] + max(L(i-2), H(i-2));

return max(L(i), H(i));
}
```

```
int findMin(int s[]){
    OPT(1) = s[1];
    while(i < num){
        OPT(i) = min(rs[i]+OPT(i-1), 4c+OPT(i-4));
        i++;
    }
    return OPT(num);
}</pre>
```

## • 6-13

we build a graph with the nodes means each stock, and the directed edge( i, j ) for each pair of stocks. we put a cost of  $-log_{r_{ij}}$  on edge (i, j).

A trading cycle C in G is an opportunity cycle 
$$\inf\prod_{(i,j)\in C}r_{ij}>1$$
 or  $\inf\sum_{(i,j)\in C}log_{ij}>0$ 

A trading cycle C in G is an opportunity cycle iff it is a negative cycle. Thus we can use the polynomial-time algorithm for negative-sycle detection to determine whether an opportunity cycle exists.