

UNIVERSITY OF COLOMBO, SRI LANKA



UNIVERSITY OF COLOMBO SCHOOL OF COMPUTING

DEGREE OF BACHELOR OF INFORMATION TECHNOLOGY (EXTERNAL)

Academic Year 2008/2009 –3rd Year Examination – Semester 6

IT6403 - Database Systems II

Structured Question Paper

13th September, 2009 (TWO HOURS)

To be completed by the candidate									
BIT Examination	Index	No:							

Important Instructions:

- The duration of the paper is 2 (two) hours.
- The medium of instruction and questions is English.
- This paper has 4 questions and 15 pages.
- Answer all questions (25 marks each).
- Write your answers in English using the space provided in this question paper.
- Do not tear off any part of this answer book.
- Under no circumstances may this book, used or unused, be removed from the Examination Hall by a candidate.
- Note that questions appear on both sides of the paper. If a page is not printed, please inform the supervisor immediately.
- Non-programmable Calculators may be used.

Questions Answered

Indicate by a cross (x), (e.g) (x)) the numbers of the **four** questions answered.

	Ques				
To be completed by the candidate by marking a cross (x).	1	2	3	4	
To be completed by the examiners:					

Index No:																					
-----------	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--

1) (a) Although Basic Timestamp Ordering algorithm causes deadlocks, it ensures that the schedules are both conflict serializable and recoverable. Discuss whether one can agree/disagree with this statement. Justify your answer.

(04 marks)

ANSWER IN THIS BOX
Basic Timestamp Ordering (TO) algorithm does not cause deadlocks
since T waits for T' only if TS(T) > TS(T').
The schedules produced by basic TO are guaranteed to be conflict serializable
since whenever the basic TO algorithm detects two conflicting operations
which occur in the incorrect order, it rejects the later of the two operations
by aborting the transaction that issued it.
The schedules produced by basic TO are not guaranteed to be recoverable and
an additional protocol (i.e. strict TO) must be enforced to ensure that the
schedules are recoverable.

(b) The following schedule consists of four transactions T_1 , T_2 , T_3 and T_4 where r_i , w_i and c_i means the read, write and commit operations of the transaction T_i respectively.

$$r_1(A)$$
, $w_1(A)$, $r_2(A)$, $r_3(B)$, $w_3(B)$, $w_2(A)$, $r_4(B)$, $w_1(B)$, c_1 , c_2 , c_4 , c_3 .

(i) Produce the precedence graph and determine whether the given schedule is conflict serializable. If so give the corresponding serial schedule.

ANSWER IN THIS BOX

T1

T2

T3

Contd.

- Indo-	lo:
The precedence graph does not have a cycle and	
hence the schedule is conflict serializable.	
The equivalent serial schedule is T ₃ , T ₄ , T ₁ ,T ₂ .	
State giving reasons whether the schedule given above is view serializable.	(02 mayle
ANSWER IN THIS BOX	(02 marks
The given schedule is conflict serializable and hence it is view se	rializable.
The given contend to commot containable and notice it to view sol	
This is due to the reason that all conflict serializable schedules ar	е
view serializable.	
State giving reasons whether the schedule given above is recoverable. I recoverable, give a modified schedule that will be recoverable.	
State giving reasons whether the schedule given above is recoverable. I recoverable, give a modified schedule that will be recoverable.	
) State giving reasons whether the schedule given above is recoverable. I	
State giving reasons whether the schedule given above is recoverable. I recoverable, give a modified schedule that will be recoverable.	(03 marks
State giving reasons whether the schedule given above is recoverable. I recoverable, give a modified schedule that will be recoverable. ANSWER IN THIS BOX The schedule is not recoverable.	(03 marks
State giving reasons whether the schedule given above is recoverable. I recoverable, give a modified schedule that will be recoverable. ANSWER IN THIS BOX The schedule is not recoverable.	(03 marks
State giving reasons whether the schedule given above is recoverable. I recoverable, give a modified schedule that will be recoverable. ANSWER IN THIS BOX The schedule is not recoverable. A modified schedule that is recoverable is r ₁ (A), w ₁ (A), r ₂ (A).	(03 marks

T 1	3 T	
Index	No:	

The given schedule is conflict serializable and it is equivalent to the serial schedule. T ₃ , T ₁ , T ₂ . However, this schedule cannot be produced through 2PL since locks acquired to cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. onsider two Transactions T ₁ and T ₂ which are received by a system at time 1 and 2, respect hat will happen if the following schedule is executed according to the timestamp ordering prote $W_1(X)$, $V_2(X)$, $V_2(Y)$, $V_3(Y)$ (04 m ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad V_3(X) \qquad V_4(X) \qquad V_5(X) \qquad V_7(X) \qquad V_7$	The given schedule is conflict serializable and it is equivalent to the serial schedule. T ₃ , T ₁ , T ₂ . However, this schedule cannot be produced through 2PL since locks acquired cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. onsider two Transactions T ₁ and T ₂ which are received by a system at time 1 and 2, respectant will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein the following schedule is executed according to the timestamp ordering protein the following schedule is executed according to the timestamp ordering protein the following schedule is executed according to the timestamp ordering protein the following schedule is executed according to the timestamp ordering protein the following schedule is executed by 2PL.		$r_1(X), w_2(X), r_3(Y), w_1(Y)$	(04 m)
However, this schedule cannot be produced through 2PL since locks acquired to cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. In the serial schedules can be executed by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $Y_2(X)$, $Y_2(Y)$, $Y_3(Y)$, $Y_4(Y)$. ANSWER IN THIS BOX To the schedule cannot be produced through 2PL since locks acquired to cannot be released as there is no deadlock situation for the given schedule.	However, this schedule cannot be produced through 2PL since locks acquired cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respectant will happen if the following schedule is executed according to the timestamp ordering protein $W_1(X)$, $V_2(X)$, $V_2(Y)$, $V_3(Y)$, $V_3(Y$	ANSWER IN	THIS BOX	(04 m
However, this schedule cannot be produced through 2PL since locks acquired to cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. In the serial schedules can be executed by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the timestamp ordering protocomic with the serial schedules according to the s	However, this schedule cannot be produced through 2PL since locks acquired cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. In the serial schedules can be executed by a system at time 1 and 2, respectant will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed according to the timestamp ordering protein will happen if the following schedule is executed by a system at time 1 and 2, respectively.			
However, this schedule cannot be produced through 2PL since locks acquired I cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. In the serial serial schedules can be executed by a system at time 1 and 2, respect not will happen if the following schedule is executed according to the timestamp ordering protection $w_1(X)$, $v_2(X)$, $v_2(Y)$, $w_1(Y)$ (04 minus ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad w_1(X)$ W-time(X) = 1	However, this schedule cannot be produced through 2PL since locks acquired cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. Onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect nat will happen if the following schedule is executed according to the timestamp ordering prote $W_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad W_1(X)$ W-time(X) = 1 $r_2(X)$	The given sch	redule is conflict serializable and it is equiv	alent to the serial sched
cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. Onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect nat will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 m ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad W_1(X)$ W-time(X) = 1	cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respectant will happen if the following schedule is executed according to the timestamp ordering prote $W_1(X)$, $T_2(X)$, $T_2(Y)$, $W_1(Y)$ ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad T_1 \qquad T_2 \qquad T_2$ $W_1(X)$ W-time(X) = 1 $T_2(X)$	$T_3, T_1, T_2.$		
cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect nat will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 m ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad W_1(X)$ W-time(X) = 1	cannot be released as there is no deadlock situation for the given schedule. Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. Desider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect nat will happen if the following schedule is executed according to the timestamp ordering prote $W_1(X)$, $V_2(X)$, $V_2(Y)$, $V_1(Y)$ (04 manually			
Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. Simplifying two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $r_2(X)$, $r_2(Y)$, $W_1(Y)$ ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad W_1(X)$ W-time(X) = 1	Although 2PL ensures conflict serializability it does not mean that all conflict serializable schedules can be executed by 2PL. Sensider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect nat will happen if the following schedule is executed according to the timestamp ordering prote $W_1(X)$, $V_2(X)$, $V_2(Y)$, $V_1(Y)$ (04 manual energy and $V_1(X)$). To the following schedule $V_1(X)$, $V_2(Y)$, $V_1(Y)$ (14 manual energy and $V_1(X)$). We time $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the following schedule $V_1(X)$ and $V_2(X)$ are the following schedule $V_1(X)$ are the fo	However, this	schedule cannot be produced through 2PI	L since locks acquired b
conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 m ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad \qquad T_1 \qquad T_2$ $w_1(X)$ W-time(X) = 1	conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respectively happen if the following schedule is executed according to the timestamp ordering protection $w_1(X)$, $v_2(X)$, $v_2(Y)$, $v_1(Y)$ ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad v_1(X)$ W-time(X) = 1 $r_2(X)$	cannot be rele	eased as there is no deadlock situation for	the given schedule.
conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $Y_2(X)$, $Y_2(Y)$, Y_2	conflict serializable schedules can be executed by 2PL. consider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect hat will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad \qquad T_1 \qquad \qquad T_2 \qquad \qquad T_2 \qquad \qquad T_1 \qquad \qquad T_2 \qquad T_2 \qquad \qquad T_1 \qquad \qquad T_2 \qquad T_2 \qquad T_3 \qquad T_4 \qquad T_5 \qquad T_5 \qquad T_6 \qquad T_6 \qquad T_7 \qquad T_8 \qquad T_9 \qquad T$			
conflict serializable schedules can be executed by 2PL. onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering proto $W_1(X)$, $Y_2(X)$, $Y_2(Y)$, Y_2	conflict serializable schedules can be executed by 2PL. consider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect hat will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad \qquad T_1 \qquad \qquad T_2 \qquad \qquad T_2 \qquad \qquad T_1 \qquad \qquad T_2 \qquad T_2 \qquad \qquad T_1 \qquad \qquad T_2 \qquad T_2 \qquad T_3 \qquad T_4 \qquad T_5 \qquad T_5 \qquad T_6 \qquad T_6 \qquad T_7 \qquad T_8 \qquad T_9 \qquad T$	Although 2DI	oncurse conflict carializability it does not a	man that all
consider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respect that will happen if the following schedule is executed according to the timestamp ordering proto $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 m ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad w_1(X)$ W-time(X) = 1	onsider two Transactions T_1 and T_2 which are received by a system at time 1 and 2, respectively happen if the following schedule is executed according to the timestamp ordering prote $W_1(X)$, $r_2(X)$, $r_2(Y)$, $W_1(Y)$ (04 n ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad W_1(X)$ W-time(X) = 1 $r_2(X)$	Altilougii 2FL	ensures commict semanzability it does not i	mean mat an
that will happen if the following schedule is executed according to the timestamp ordering proto $w_1(X), r_2(X), r_2(Y), w_1(Y)$ (04m) $ANSWER IN THIS BOX$ $T_1 \qquad T_2$ $w_1(X)$ $W-time(X) = 1$	that will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$ and $w_1(X)$ (15 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$ and $w_2(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$, $w_4(X$	conflict serial	izable schedules can be executed by 2PL.	
that will happen if the following schedule is executed according to the timestamp ordering proto $w_1(X), r_2(X), r_2(Y), w_1(Y)$ (04m) $ANSWER IN THIS BOX$ $T_1 \qquad T_2$ $w_1(X)$ $W-time(X) = 1$	that will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$ and $w_1(X)$ (15 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$ and $w_2(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$, $w_4(X$			
that will happen if the following schedule is executed according to the timestamp ordering proto $w_1(X), r_2(X), r_2(Y), w_1(Y)$ (04m) $ANSWER IN THIS BOX$ $T_1 \qquad T_2$ $w_1(X)$ $W-time(X) = 1$	that will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$ and $w_1(X)$ (15 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$ and $w_2(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$, $w_4(X$			
hat will happen if the following schedule is executed according to the timestamp ordering proto $w_1(X), r_2(X), r_2(Y), w_1(Y)$ (04m) $ANSWER IN THIS BOX$ $T_1 \qquad T_2$ $w_1(X)$ $W-time(X) = 1$	hat will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 n ANSWER IN THIS BOX $T_1 \qquad T_2 \qquad w_1(X)$ W-time(X) = 1 $r_2(X)$			
that will happen if the following schedule is executed according to the timestamp ordering proto $w_1(X), r_2(X), r_2(Y), w_1(Y)$ (04m) $ANSWER IN THIS BOX$ $T_1 \qquad T_2$ $w_1(X)$ $W-time(X) = 1$	that will happen if the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ (04 manual example of the following schedule is executed according to the timestamp ordering prote $w_1(X)$ and $w_1(X)$ (15 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$ and $w_2(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$ (16 manual example of the following schedule is executed according to the timestamp ordering protein $w_1(X)$, $w_2(X)$, $w_3(X)$, $w_4(X)$, $w_4(X$			
$w_1(X), r_2(X), r_2(Y), w_1(Y) \end{tabular} \begin{tabular}{c} (04 m) \\ \hline ANSWER IN THIS BOX \end{tabular}$	$w_1(X), r_2(X), r_2(Y), w_1(Y) \end{tabular} $			
ANSWER IN THIS BOX $T_1 \qquad T_2 \\ w_1(X)$ $W-time(X) = 1$	ANSWER IN THIS BOX $T_1 \qquad T_2 \\ w_1(X) \\ W\text{-time}(X) = 1$ $r_2(X)$		· -	
ANSWER IN THIS BOX $T_1 \qquad T_2 \\ w_1(X)$ $W-time(X) = 1$	ANSWER IN THIS BOX $T_1 \qquad T_2 \\ w_1(X) \\ W\text{-time}(X) = 1$ $r_2(X)$		f the following schedule is executed according to the	
w ₁ (X) W-time(X) = 1	$w_1(X)$ $W-time(X) = 1$ $r_2(X)$		f the following schedule is executed according to the	ne timestamp ordering proto
w ₁ (X) W-time(X) = 1	$w_1(X)$ $W-time(X) = 1$ $r_2(X)$	nat will happen if	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$	ne timestamp ordering proto
w ₁ (X) W-time(X) = 1	$w_1(X)$ $W-time(X) = 1$ $r_2(X)$	nat will happen if	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$	ne timestamp ordering proto
w ₁ (X) W-time(X) = 1	$w_1(X)$ $W-time(X) = 1$ $r_2(X)$	nat will happen if	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$	
w ₁ (X) W-time(X) = 1	$w_1(X)$ $W-time(X) = 1$ $r_2(X)$	nat will happen if	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$	ne timestamp ordering proto
w ₁ (X) W-time(X) = 1	$w_1(X)$ $W-time(X) = 1$ $r_2(X)$	nat will happen if	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$	ne timestamp ordering proto
W-time(X) = 1	W-time(X) = 1 $r_2(X)$	nat will happen if	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$	ne timestamp ordering proto
	r ₂ (X)	ANSWER IN	f the following schedule is executed according to the w ₁ (X), r ₂ (X), r ₂ (Y), w ₁ (Y) THIS BOX	ne timestamp ordering proto
	r ₂ (X)	ANSWER IN	f the following schedule is executed according to the $w_1(X)$, $r_2(X)$, $r_2(Y)$, $w_1(Y)$ THIS BOX T ₂	ne timestamp ordering proto
		ANSWER IN T ₁ w ₁ (X)	f the following schedule is executed according to the w ₁ (X), r ₂ (X), r ₂ (Y), w ₁ (Y) THIS BOX T ₂	ne timestamp ordering proto

	Index	No:
	r ₂ (Y)	
	R-time(Y) = 2	
	w ₁ (Y)	
	rejected	
T1 is re	olled back, which causes T ₂ to be rolled back too,	
' o		
becaus	se it had read the value of X written by T ₁ .	
onsider t	he following four recovery concepts.	
onsider t	he following four recovery concepts. NO-UNDO / REDO	
(i) (ii)	NO-UNDO / REDO UNDO / NO-REDO	
(i) (ii)	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO	
(i) (ii) (iii) (iv)	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO	
(i) (ii) (iii) (iv)	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts.	(03 ma
(i) (ii) (iii) (iv)	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO	(03 ma
(i) (ii) (iii) (iv) Name the	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts.	(03 ma
(i) (ii) (iii) (iv) Name the	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts.	(03 ma
(i) (iii) (iv) Name the	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts. VER IN THIS BOX Deferred Update	(03 ma
(i) (iii) (iv) Name the	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts.	(03 ma
(i) (iii) (iv) Name the	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts. VER IN THIS BOX Deferred Update	(03 ma
(i) (iii) (iv) Name the ANSW (i)	NO-UNDO / REDO UNDO / NO-REDO UNDO / REDO NO-UNDO / NO-REDO e recovery technique for each of the above recovery concepts. VER IN THIS BOX Deferred Update	(03 ma

(iv) Shadow Paging OR

Deferred Update

	Index No:
(a)	To create an index on a relation, a choice between hash index and B+-tree index has to be made. For each of the two indexing methods describe when it would be more suitable then the other.
	(03 marks)
	ANSWER IN THIS BOX
	B+-tree index - If it is likely that ranged queries are going to be performed often,
	then we should use a B+-tree on the index for the relation
	since hash indices cannot perform range queries.
	Hash index - If it is more likely that we are only going to perform equality queries,
	for example the case of social security numbers,
	then hash indices are the best choice
	since they allow for the faster retrieval than B+-trees by 2-3 I/Os per request.
(b)	 (i) Search for records based on a range of field values and the data would not be modified often. (ii) Search for records based on a range of field values and the data would be modified often. (iii) Perform data insertions and deletions often and the order of records does not matter with respect to data retrieval. (iv) Search for a particular record based on a given field value.
	For each of the above operations, which file organization/index is most suited?
	ANSWER IN THIS BOX (04 marks)
	(i) Sorted file organization on the given field
	(ii) B+ tree index on the given field
	(iii) Heap file.
	(iv) Hash indexed on the given field

Index 1	No:				
---------	-----	--	--	--	--

			schema.

Employee (Eid, Ename, Salary, Age, Did) Department (Did, Dname, Floor, Mgrid) Finance (Did, Budget, Expenses)

Assume that each Employee record is 40 bytes long, each Department record is 25 bytes long, and each Finance record is 20 bytes long on average. There are 50,000 tuples in Employee and 5,000 tuples in Department. The file system supports 2,000 byte pages. Each employee is working for one department. Assume uniform distribution of values with respect to any given set of tuples and 20% of employees are earning more than Rs.50,000 while 10% of all employees are over 45 years of age. The company owns five floors in the building and due to uniform distribution, given a set of tuples with Did (department identity), 20% of the tuples of the given set would belong to each floor.

For each of the following queries, which file organization/index choices would you choose to speed them up?

(i) SELECT Ename, Age, Salary FROM Employee;

(03 marks)

ANSWER IN THIS BOX	
An unclustered hash index on Ename, Age, Salary attributes of Employee	
which is an index only plan	
OR	
Index will not be created since query requires to access all the Employee records and	<u> </u>
the records are accessed using a file scan.	

(ii) SELECT Did FROM Department WHERE Floor = 5 AND MgridBudget < 1500000;

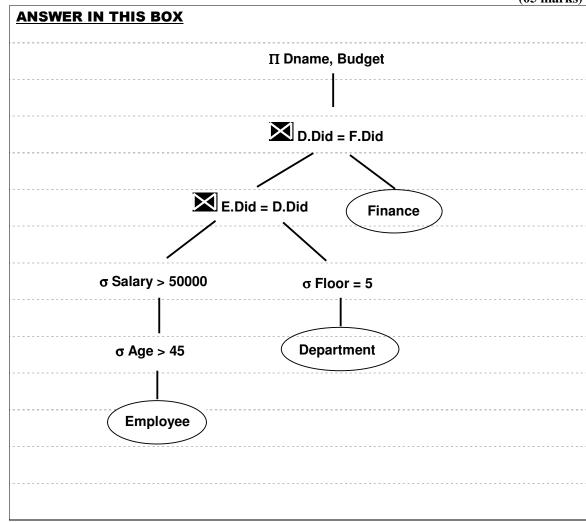
ANSWER IN THIS BOX	(03 mark
A clustered dense B+ tree index on <floor, mgrid=""> attributes</floor,>	of Department,
then the records would be ordered on these attributes assisti	ng the given query.
When executing this query, the first record with floor = 5 mus	st be retrieved, and
then the other records with Mgrid < 150 can be found in the o	order of Mgrid.
This is the best plan for this query and is not an index-only pl	lan.

(iii) Consider the following query executed without using any indexes;

```
SELECT D.Dname, F.Budget
FROM Employee E, Department D, Finance F
WHERE E.Did=D.Did AND D.Did=F.Did AND D.Floor = 5
AND E.Salary > 50000 AND E.Age > 45;
```

Draw the optimized query tree for the above query.

(05 marks)



(iv) Suggest indices to produce the best plan for the execution of the query in (c)(iii). Explain the query plan illustrating where necessary the number of tuples processed and how the indices would be used in the query plan to minimise the processing cost.

(07 marks)

ANSWER IN THIS BOX

Indices

- Clustered B+ tree index on age of Employee
- Hash/B+ index on Did of Department
- Hash/B+ index on Did of Finance

Contd.

]	ĺ	1	d	e	X	(ľ	V	C):																					
-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	

First, retrieve the tuples from Employee with Age > 45
using the B-tree index on Age and based on the estimation 5,000 (10% of 50,000)
such tuples will be found.
The index on Age is more selective than an index on Salary which would have had to
select 10,000 tuples (20% of 50,000).
Out of these retrieved 5,000 tuples, those who earn more than Rs. 50,000
would be selected on the fly. There are 1,000 (20% of 5,000) such tuples.
Pipeline these 1,000 tuples one at a time to Department and using index on Did
it is possible to find at most one matching tuple from the Department tuples
(index nested loop). Consequently, there will be 1,000 tuples.
Out of the retrieved 1,000 tuples, the tuples corresponding to 5th floor
would be selected on the fly and there will be 200 (20% of 1,000) such tuples.
Pipeline the estimated 200 tuples one at a time to Finance and use the index on Did
to retrieve at most one Finance tuple for each of the 200 tuples.

Index No:	
 1	

3)	(a) (i) In data mining one of the techniques of data analysing and decision making is through Market
	Basket Analysis? Briefly explain what it is.

(02 marks)

ANSWER I	N THIS BOX	
Market Bas	ket Analysis is a modelling technique based upon the theory that	
if you buy a	certain group of items,	
you are mo	re (or less) likely to buy another group of items.	

(ii) What is an association rule? What is its role in data mining?

(02 marks)

ANSWER IN THIS BOX

An association rule is defined as a statement of the form

{X1,X2,...,Xn} -> {Y1,Y2,...,Yn}, which means that Y1,Y2,...,Yn is present in the

transaction if X1,X2,...,Xn are all in the transaction. It helps to establish relationships

among data sets. For instance the confidence level of a data item against a data set.

(b) A market basket analysis is to be performed to find the relationships between a set of items = {milk, tea, coffee, sugar, juice}. After inspecting eight baskets for these items the following were found.

```
B1 = {milk, tea, sugar} B2 = {milk, coffee, sugar}
B3 = {milk, coffee, juice}
                             B4 = {milk, sugar, juice}
B5 = \{tea, sugar\}
                             B6 = {tea, sugar, juice}
B7 = \{tea, juice\}
                              B8 = {coffee, sugar}
```

(i) Define support in terms of a pair of items A and B. Identify how many of the same pair should be present to achieve a support threshold of 35%.

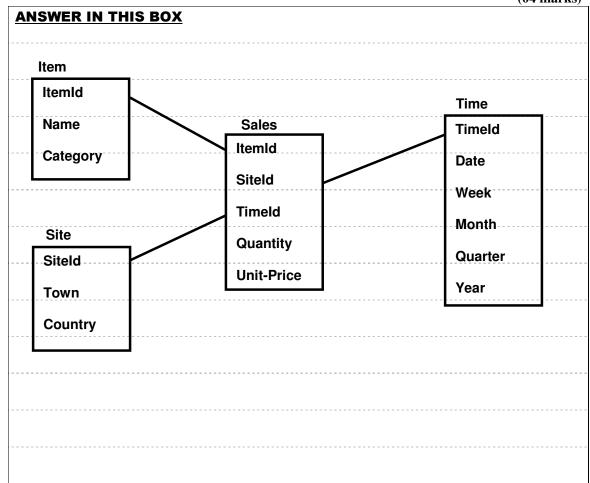
ANSWER IN THIS BOX	(03 marks)
The support of A=>B is the percentage of the transactions which contains	
both A and B.	
A pair must appear in 3 of the 8 baskets to have a support of 35%.	

Index	No:											

ANSWED IN THIS BOY	(03 mai
ANSWER IN THIS BOX	
There are ten possible pairs, each with the following occurrences.	
{milk, tea} 1, {milk, coffee} 2, {milk, sugar} 3, {milk, juice} 2,	
{tea, coffee} 0, {tea, sugar} 3, {tea, juice} 2,	
{coffee, sugar} 2, {coffee, juice} 1,	
{sugar, juice} 2.	
{milk, sugar} and {tea, sugar} are the pairs which appear 3 times.	
(iii) What is the confidence of tea, given milk and sugar (i.e. association rule {milk	x, sugar} -> tea)?
ANSWER IN THIS BOX	(02 1114)
Thus the confidence of tea given milk and sugar is = 33%.	
Thus the confidence of tea given milk and sugar is = 33%. Consider the following multi dimensional data model of a data warehouse on sales Item(ItemId, Name, Category) Site(SiteId, Town, Country) Sales(ItemId, SiteId, TimeId, Quantity, Unit-Pri Time(TimeId, Date, Week, Month, Quarter, Year)	ce)
Thus the confidence of tea given milk and sugar is = 33%. Consider the following multi dimensional data model of a data warehouse on sales Item(ItemId, Name, Category) Site(SiteId, Town, Country) Sales(ItemId, SiteId, TimeId, Quantity, Unit-Pri Time(TimeId, Date, Week, Month, Quarter, Year) (i) What would be the Fact and Dimension tables for the above data warehouse?	ce)
Thus the confidence of tea given milk and sugar is = 33%. Consider the following multi dimensional data model of a data warehouse on sales Item(ItemId, Name, Category) Site(SiteId, Town, Country) Sales(ItemId, SiteId, TimeId, Quantity, Unit-Pritime(TimeId, Date, Week, Month, Quarter, Year) (i) What would be the Fact and Dimension tables for the above data warehouse? ANSWER IN THIS BOX Fact Table: Sales.	ce) (02 mai
Site(<u>SiteId</u> , Town, Country) Sales(<u>ItemId</u> , <u>SiteId</u> , <u>TimeId</u> , Quantity, Unit-Pri Time(<u>TimeId</u> , Date, Week, Month, Quarter, Year) (i) What would be the Fact and Dimension tables for the above data warehouse? ANSWER IN THIS BOX Fact Table: Sales.	се) (02 ma r
Thus the confidence of tea given milk and sugar is = 33%. Consider the following multi dimensional data model of a data warehouse on sales Item(ItemId, Name, Category) Site(SiteId, Town, Country) Sales(ItemId, SiteId, TimeId, Quantity, Unit-Pritime(TimeId, Date, Week, Month, Quarter, Year) (i) What would be the Fact and Dimension tables for the above data warehouse? ANSWER IN THIS BOX Fact Table: Sales.	ce) (02 ma

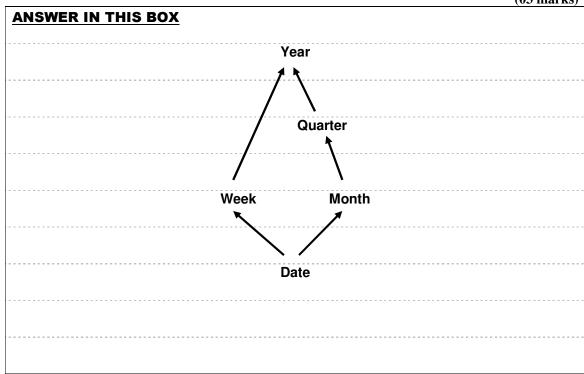
(ii) Draw a star schema for the data warehouse in (c)?

(04 marks)



(iii) Draw a dimension/concept hierarchy for Time (of the above given data warehouse).

(03 marks)



Index	Ma.								
inaex	INO:			 			 		

	(04 mark
ANSWER I	IN THIS BOX
OF! FOT To	Voor CUM/Organiih-#Unit Duigo)
SELECT 10	own, Year, SUM(Quantity*Unit-Price)
FROM Sale	s s, Site si, Time t
WHERE s.S	SiteId=si.SiteId and s.TimeId=t.TimeId
GROUP BY	Town, Year
	rtment maintains the Employee table to pay salaries of their employees. HR department oyees to various projects and maintains Project and Assignment tables. The three tables elow.
	ee(EmpNo, Name, Title, Salary) t(ProjId, Pname, Budget, Location)
	ment(ProjId, EmpNo, Responsibility, Duration)
majority	the projects are based in "Colombo" and hence Location = "Colombo" appears in the of queries. Write down the fragments to perform primary fragmentation on Project give e predicate.
	(03 mark
ANSWER	IN THIS BOX
ANSWER I	IN THIS BOX
	IN THIS BOX
Proj1 = σ_{Lo}	ocation="Colombo" Project
Proj1 = σ_{Lo}	
Proj1 = σ_{Lo}	ocation="Colombo" Project
Proj1 = σ_{Lo} Proj2 = σ_{Lo} (ii) Write do	ocation="Colombo" Project ocation<>"Colombo" Project own the fragments if we perform derived horizontal fragmentation on Assignment base agments of Project in (i).
Proj1 = σ _{Lo} Proj2 = σ _{Lo} (ii) Write do on the fra	ocation="Colombo" Project ocation<>"Colombo" Project own the fragments if we perform derived horizontal fragmentation on Assignment base agments of Project in (i).
Proj1 = σ _{Lo} Proj2 = σ _{Lo} (ii) Write do on the fra	ocation="Colombo" Project over the fragments if we perform derived horizontal fragmentation on Assignment base agments of Project in (i). (03 mark IN THIS BOX
Proj1 = σ _{Lo} Proj2 = σ _{Lo} (ii) Write do on the fra	ocation="Colombo" Project ocation<>"Colombo" Project own the fragments if we perform derived horizontal fragmentation on Assignment base agments of Project in (i).
Proj1 = σ _{Lo} Proj2 = σ _{Lo} (ii) Write do on the fra ANSWER I	ocation="Colombo" Project over the fragments if we perform derived horizontal fragmentation on Assignment base agments of Project in (i). (03 mark IN THIS BOX

4)

Index	Ma.										
inaex	INO.										

ANSWER IN THIS BOX	(03 marks)
ANOWER IN THIS BOX	
SELECT Pname, Budget	
FROM Proj1	
UNION	
SELECT Pname, Budget	
FROM Proj2	
(iv) Write a query to identify responsibilities with respect to projects outside "	
ANSWER IN THIS BOX	(03 marks
OR SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a	
SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a	s write a query or quer ocated in "Colombo".
SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a WHERE p.ProjId=a.ProjId (v) Assume that there are two databases called HR and Finance where HI Assignments while Finance manages Employee data. (I) If these two databases are to be managed as distributed databases	s write a query or quer
SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a WHERE p.ProjId=a.ProjId (v) Assume that there are two databases called HR and Finance where HI Assignments while Finance manages Employee data. (I) If these two databases are to be managed as distributed databases plan to retrieve name with salary of Employees working in projects to ANSWER IN THIS BOX	s write a query or querocated in "Colombo".
SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a WHERE p.ProjId=a.ProjId (v) Assume that there are two databases called HR and Finance where HI Assignments while Finance manages Employee data. (I) If these two databases are to be managed as distributed databases plan to retrieve name with salary of Employees working in projects to	s write a query or querocated in "Colombo".
SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a WHERE p.ProjId=a.ProjId (v) Assume that there are two databases called HR and Finance where HI Assignments while Finance manages Employee data. (I) If these two databases are to be managed as distributed databases plan to retrieve name with salary of Employees working in projects to ANSWER IN THIS BOX	s write a query or querocated in "Colombo".
SELECT DISTINCT Responsibility FROM Proj2 p, Assignment a WHERE p.ProjId=a.ProjId (v) Assume that there are two databases called HR and Finance where HI Assignments while Finance manages Employee data. (I) If these two databases are to be managed as distributed databases plan to retrieve name with salary of Employees working in projects to ANSWER IN THIS BOX SELECT Name, Salary FROM Finance.Employee, HR.Asg1	s write a query or querocated in "Colombo".

Index	No.									

	(II) If these two databases are to be managed as multi databases write a query or query plan to retrieve name with salary of Employees working in projects located in "Colombo".
	ANSWER IN THIS BOX
	SELECT Name, Salary FROM Finance.Employee
	WHERE Finance.Employee.EmpNo IN
	(SELECT HR.Asg1.EmpNo FROM HR.Asg1)
	Data from HR Asg1 would be exported and then joined with the Employee table.
b)	What is Embedded SQL?
	ANSWER IN THIS BOX
	Embedded SQL is a method of combining the computing power of a
	
	programming language and the database manipulation capabilities of SQL.
c)	Describe Object Definition Language (ODL) and Object Query Language (OQL) used in Object-Oriented Database Management Systems. For each of these two object languages, what is the corresponding language provided by relational database management system called? (03 marks)
	ANSWER IN THIS BOX
	Object Definition Language (ODI) is used for defining interference to object types
	Object Definition Language (ODL) is used for defining interfaces to object types.
	This is similar to the Data Definition Language (DDL) of Relational DBMS.
	OQL is used for the retrieval and manipulation of objects
	This is similar to the Data Manipulation Language (DML) of Relational DBMS.
}	
