Assignment 2 Problem 1

A *stable* sorting algorithm is one in which the relative order of all identical elements (or keys) is the same in the output as it was in the input. Prove, using a well-chosen example, that heapsort is not a stable sorting algorithm.

solution:

Let a = 0, b = 1, c = 1, d = 2 e = 3 f = 4 and g = 5

Let the original array A to be $A[7] = \{g, f, e, d, c, b, a\}$

Suppose the heapsort is a stable soring, then the output should be $A[7] = \{a, c, b, d, e, f, g\}$ Assuming the HeapSort we used have the following peseudocode

$$HeapSort(A, n)$$
 (1)

$$n = A.size (2)$$

$$for \quad i = parent(last(n)) \quad down \quad to \quad 0 \quad do$$
 (3)

$$fix - down(A, n, i) \tag{4}$$

$$while \quad n > 1 \tag{5}$$

$$swap items at A[root()] and A[last(n)]$$
 (6)

$$decrease n$$
 (7)

$$fix - down(A, n, root()) \tag{8}$$

(9)

$$fix - down(A, n, k) \tag{10}$$

while
$$k$$
 is not a leaf do (11)

$$j = left - child(k) \tag{12}$$

$$if(j! = last(n) \quad and \quad A[j+1] > A[j]) \tag{13}$$

$$j = j + 1 \tag{14}$$

$$if(A[k] \ge A[j]) \quad break$$
 (15)

$$swap(A[j], A[k]) \tag{16}$$

$$k = j (17)$$

(18)

After doing the heapify part of the code (ie. line 2 to line 4 in the previous peseudocode), the entries in array A is now $\{g, f, e, d, c, b, a\}$ which does not change since it is already a maximum heap.

Then, we repeatedly find maximum (ie. line 5 to line 8) after the 1st loop, we get $A[7] = \{f, d, e, a, c, b, g\}$

after the 2nd loop, we get $A[7] = \{e, d, b, a, c, f, g\}$

after the 3rd loop, we get $A[7] = \{d, c, b, a, e, f, g\}$

after the 4th loop, we get $A[7] = \{c, a, b, d, e, f, g\}$

after the 5th loop, we get $A[7] = \{b, a, c, d, e, f, g\}$ after the 6th loop, we get $A[7] = \{a, b, c, d, e, f, g\}$ which is also the final output Notice that this is not identical to our expected output, therefore, we can conclude that the heapsort is not a stable sorting algorithm. QED.