

Assignment 5 Problem 4

- Construct the last occurrence function L for pattern $P = \text{adobodoa}$ where $\Sigma = \{a, b, c, d, o, t\}$.
- Trace the search for P in $T = \text{dotadotadotdotadobodoaodot}$ using the Boyer-Moore algorithm. Indicate in a table such as Table 2 which characters of P were compared with which characters of T . Follow the example from the lecture video in Module 9. Place each character of P in the column of the compared-to character of T . Put brackets around the character if they are known to match from the previous step (similar to the examples in the slides). Use a new row when sliding the pattern. You may not need all rows in the table. Add more rows to the table if you need more.
- For any $m \geq 1$ and any $n \geq m$, give a pattern P and a text T such that the Boyer-Moore algorithm looks at exactly $\Theta(n/m)$ characters. Justify your answer.

Solution:

- a) $L(a) = 7$
 $L(b) = 3$
 $L(c) = -1$
 $L(d) = 5$
 $L(o) = 6$
 $L(t) = -1$
- b)

[illegible]

Table 2: Table for Boyer-Moore problem.

- c) For any pattern P of length m, and text T of length n. If all characters on position p in T where ($p \equiv m - 1 \text{ mod } m$) are the characters that does not occur in P (i.e. $L(T[p]) = -1$). Then the Boyer-Moore algorithm looks at exactly $\Theta(n/m)$ characters because in this case, Boyer-Moore algorithm will only look at the characters that are in position p, and will find that it does not occur, then it will shift i into $i + m$. Since each time it shifts m position, totally, it will only check $\Theta(n/m)$ times.