

5.1 Introduction

5.2 Binary Trees

5.3 Binary Trees Traversals

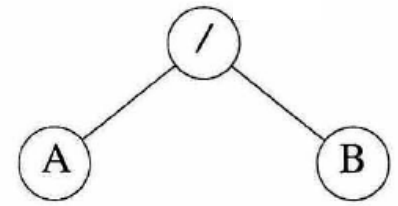
5.4 Additional Binary Tree Operations

5.5 Threaded Binary Trees

5.6 Heaps

5.7 Binary Search Trees

5.3 Binary Tree Traversal



- Tree traversal
 - Visiting each node in the tree exactly once ✱
- When traversing a binary tree
 - L, V, R : moving left, visiting the node, moving right
(= 왼쪽 서브트리 이동)
 - Six possible combinations of traversal
 - LVR, LRV, VLR, VRL, RVL, RLV
 - If we traverse left before right, only tree remains
 - LVR (inorder): ... 중위 순회
 - LRV (postorder): ... 후위 순회
 - VLR (preorder): ... 전위 순회

```
typedef struct node * treePointer;  
typedef struct node {  
    int data;  
    treePointer leftChild, rightChild;  
};
```

Expression:

$A / B * C * D + E$

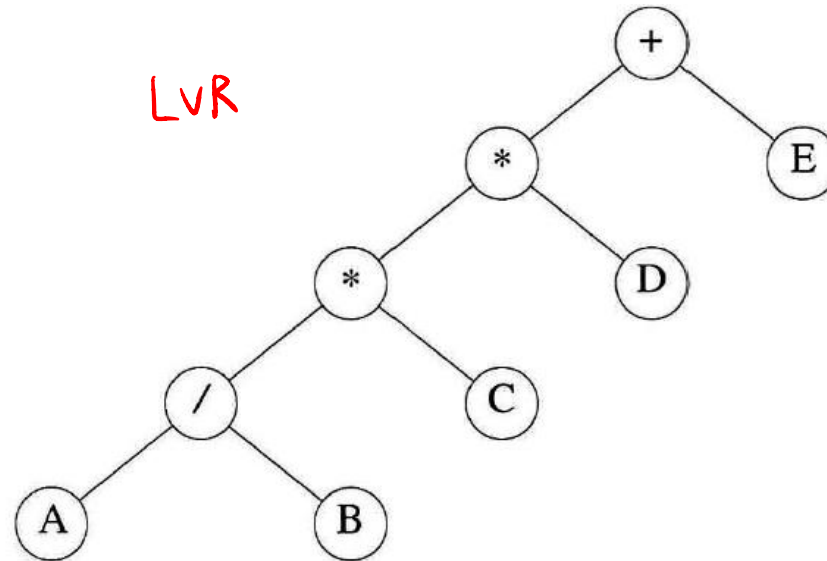
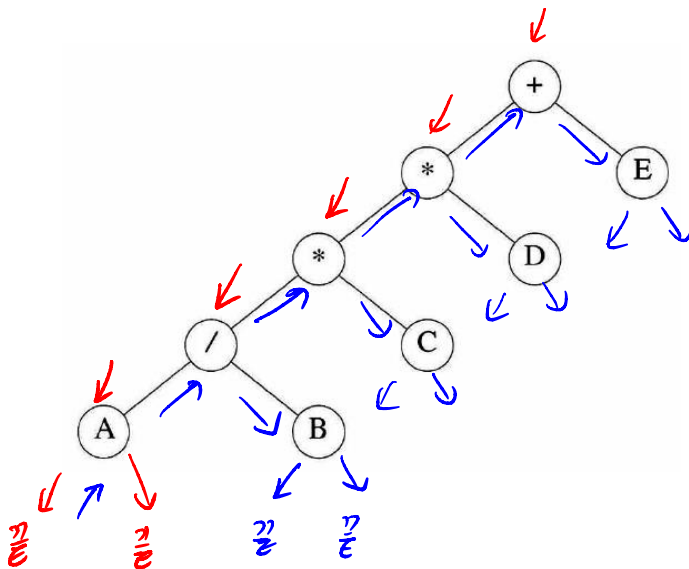


Figure 5.16: Binary tree with arithmetic expression

5.3.Inorder Traversal

- LVR :



→ A/B*C*D+E

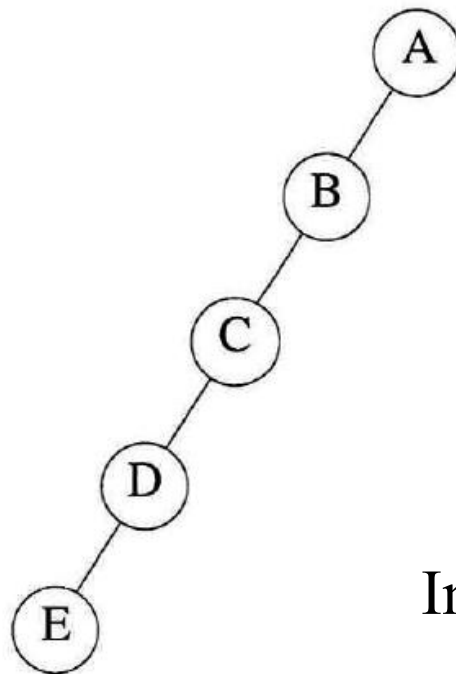
→ 2x9+1
 Empty Node
 9x9

→ A의 subtree(=NULL) 만났을 까지

```
void inorder(treePointer ptr)
{ /* inorder tree traversal */
  if (ptr) {
    inorder(ptr->leftChild);
    printf("%d", ptr->data);
    inorder(ptr->rightChild);
  }
}
```

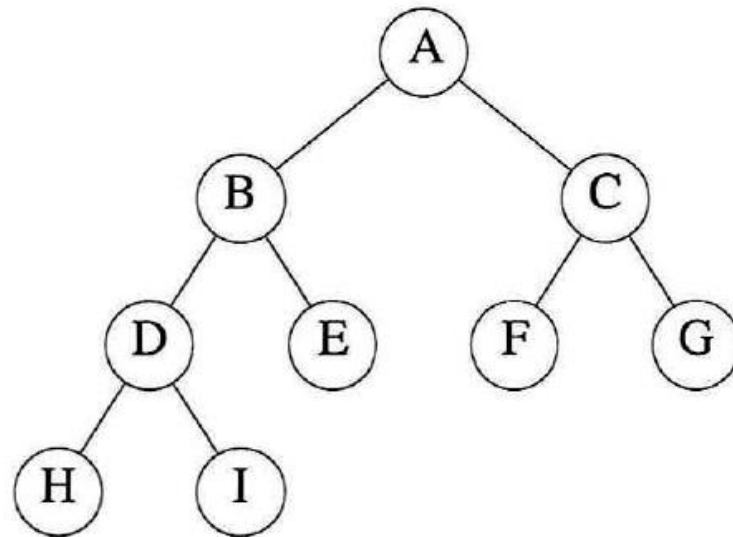
Program 5.1: Inorder traversal of a binary tree

→ # of inorder() : ...

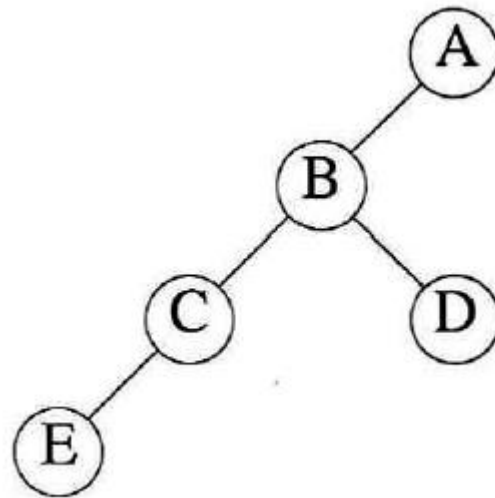


Inorder: ...

EDCBA

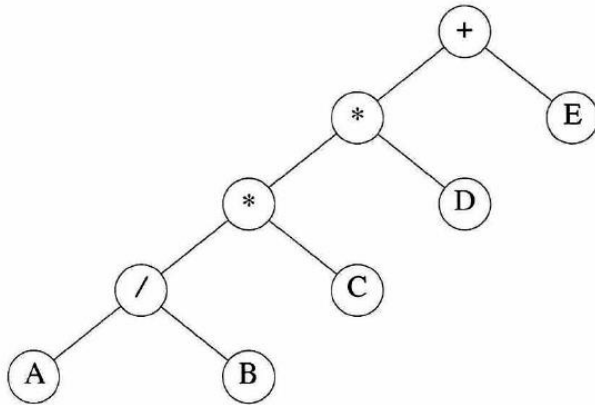


Inorder: ... HDIBEAFCA



Inorder: ... ECBDA

- LVR : $\rightarrow A/B * C * D + E$



```

void inorder(treePointer ptr)
{ /* inorder tree traversal */
    if (ptr) {
        inorder(ptr->leftChild);
        printf("%d", ptr->data);
        inorder(ptr->rightChild);
    }
}

```

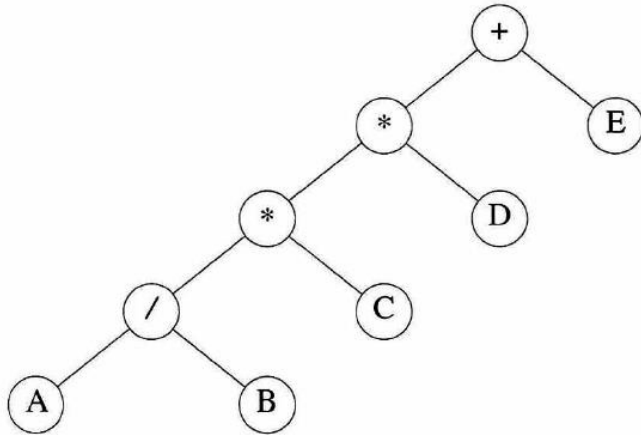
Program 5.1: Inorder traversal of a binary tree

Call of <i>inorder</i>	Value in root	Action	Call of <i>inorder</i>	Value in root	Action
1	+		11	C	
2	*		12	NULL	
3	*		11	C	printf
4	/		13	NULL	
5	A		2	*	printf
6	NULL		14	D	
5	A	printf	15	NULL	
7	NULL		14	D	printf
4	/	printf	16	NULL	
8	B		1	+	printf
9	NULL		17	E	
8	B	printf	18	NULL	
10	NULL		17	E	printf
3	*	printf	19	NULL	

5.3.2 Preorder Traversal

Root은 먼저 출력

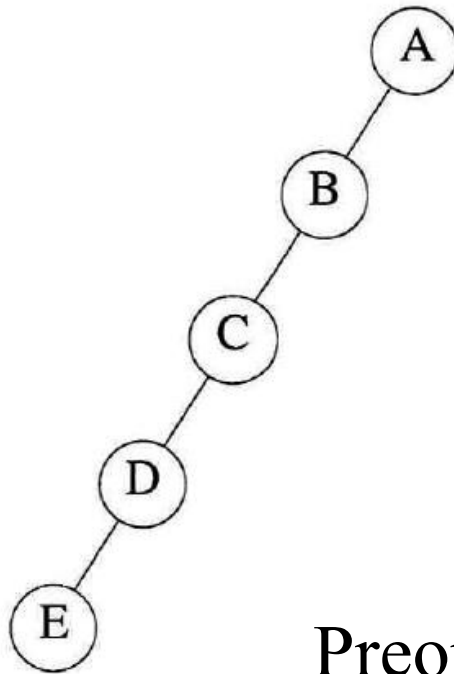
- VLR :



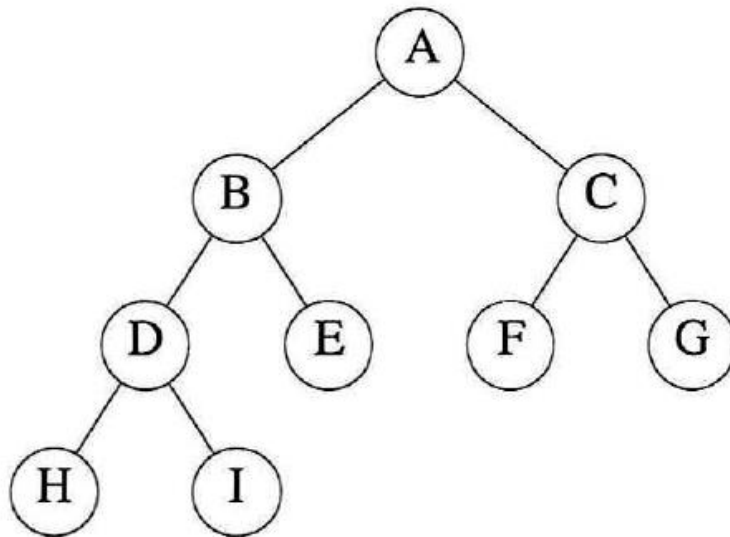
→ + * * / A B C D E

```
void preorder(treePointer ptr)
{ /* preorder tree traversal */
    if (ptr) {
        printf("%d", ptr->data);
        preorder(ptr->leftChild);
        preorder(ptr->rightChild);
    }
}
```

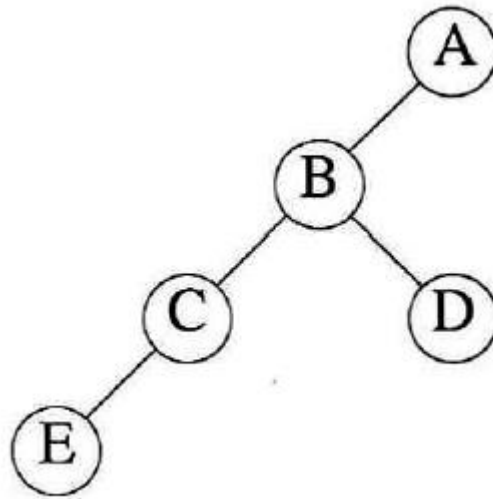
Program 5.2: Preorder traversal of a binary tree



Preorder (VLR): ... ~~AB~~ CDE



Preorder (VLR): ... ~~ABDHIE~~CFG

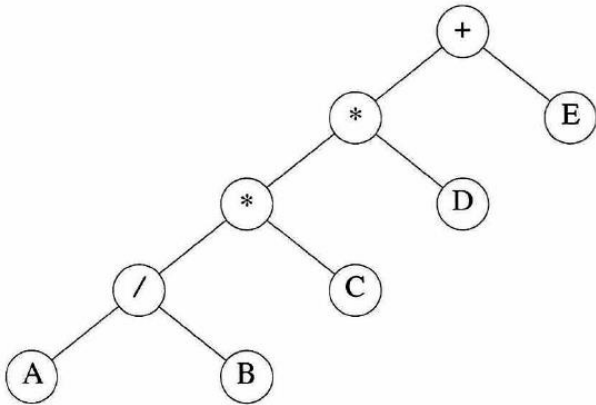


Preorder (VLR): ... ABCED

5.3.3 Postorder Traversal

계산식 증명

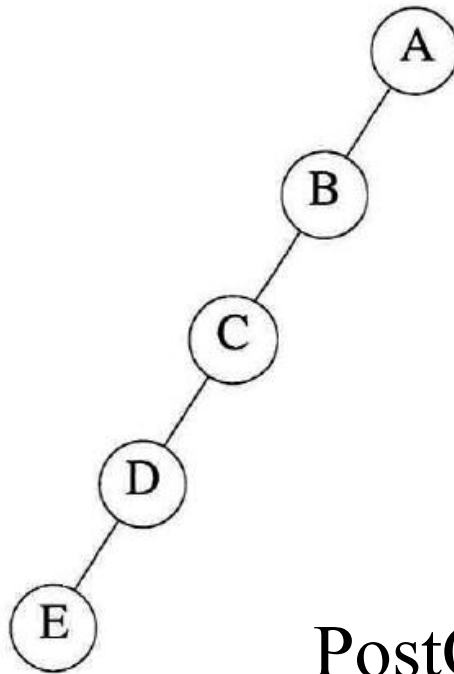
- LRV :



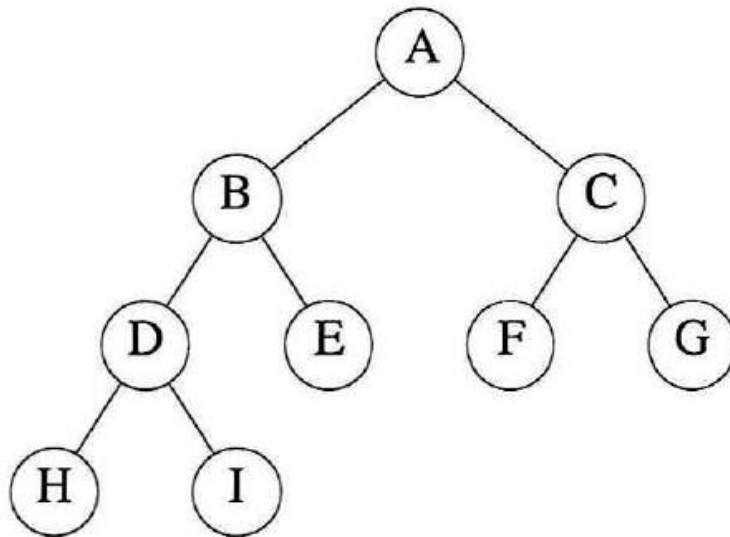
```
void postorder(treePointer ptr)
{ /* postorder tree traversal */
    if (ptr) {
        postorder(ptr->leftChild);
        postorder(ptr->rightChild);
        printf("%d", ptr->data);
    }
}
```

Program 5.3: Postorder traversal of a binary tree

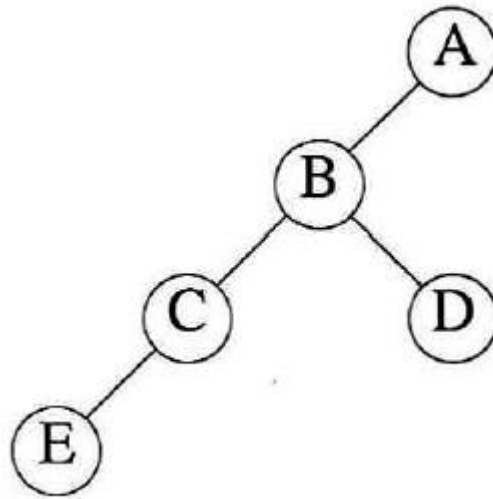
→ A B / C * D * E +



PostOrder (LRV): ... *EDCBA*



PostOrder (LRV): ... ~~HI~~DEBAFGC



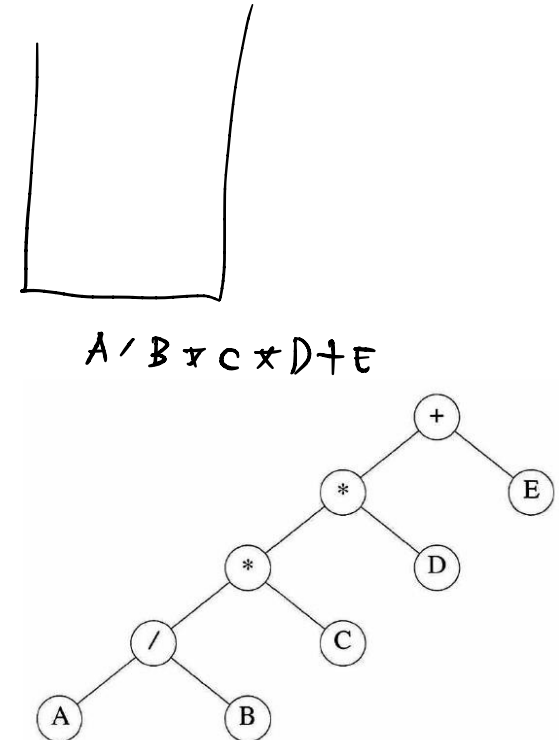
PostOrder (LRV): ... ECDBA

반복

5.3.4 Iterative Inorder Traversal

```
void iterInorder(treePointer node)
{
    int top = -1; /* initialize stack */
    treePointer stack[MAX_STACK_SIZE];
    for (;;) {
        for(; node; node = node->leftChild)
            push(node); /* add to stack */
        node = pop(); /* delete from stack */
        if (!node) break; /* empty stack */
        printf("%d", node->data);
        node = node->rightChild;
    }
}
```

Program 5.4: Iterative inorder traversal



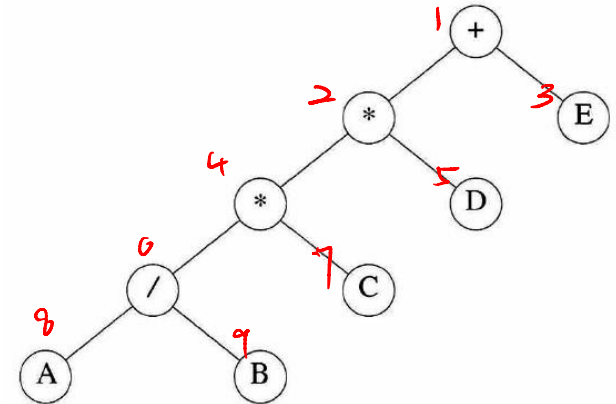
→ LVR: $A/B * C * D + E$

```

void iterInorder(treePointer node)
{
    int top = -1; /* initialize stack */
    treePointer stack[MAX-STACK-SIZE];
    for (;;) {
        for(; node; node = node->leftChild)
            push(node); /* add to stack */
        node = pop(); /* delete from stack */
        if (!node) break; /* empty stack */
        printf("%d", node->data);
        node = node->rightChild;
    }
}

```

Program 5.4: Iterative inorder traversal



→ LVR: A/B * C * D + E

Time complexity:

... $O(n)$

Space complexity:

... $O(n)$

↪ tree depth

이진 트리
높이 = depth

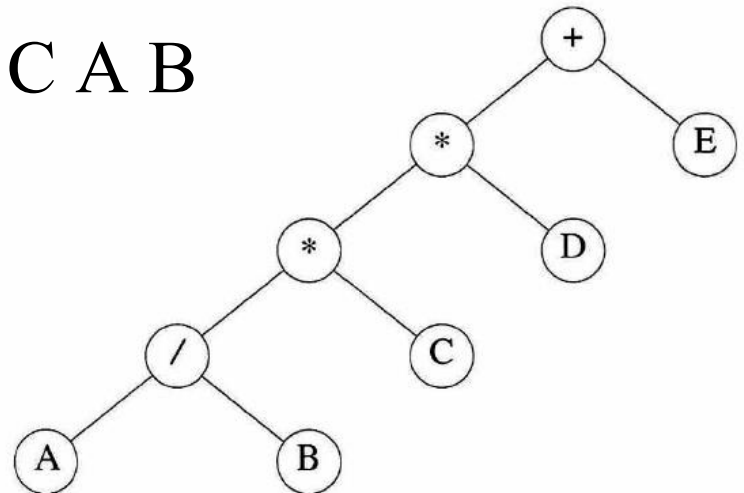
Queue로 구현

Pre, in, post order \rightarrow Stack, Recursion
Level-order \rightarrow Queue

5.3.5 Level-Order Traversal

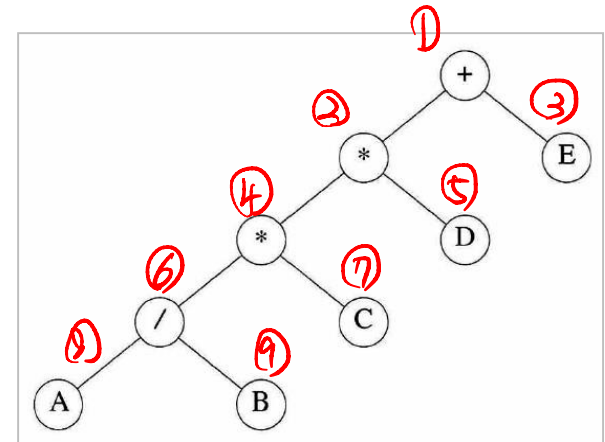
Level 별로 순차적으로

- Requires a queue
 - Visit the **root** first, the root's **left child**, followed by the root's **right child**
 - Continue, visiting the node at each new level from the leftmost node to the rightmost node
- LO Traversal: $+ * E * D / C A B$



Queue = FIFO

```
void levelOrder(treePointer ptr)
{ /* level order tree traversal */
    int front = rear = 0;
    treePointer queue[MAX_QUEUE_SIZE];
    if (!ptr) return; /* empty tree */
    addq(ptr);
    for (;;) {
        ptr = deleteq();
        if (ptr) {
            printf("%d", ptr->data);
            if (ptr->leftChild)
                addq(ptr->leftChild);
            if (ptr->rightChild)
                addq(ptr->rightChild);
        }
        else break;
    }
}
```



→ + * E * D / C A B

$+ * E * D$

/ C

Program 5.5: Level-order traversal of a binary tree

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5.4 Additional Binary Tree Operations

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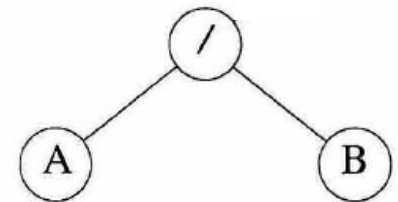
5.7 Binary Search Trees

5.4.1 Copying Binary Trees

- This function is a modified version of `postorder` (Program 5.3)

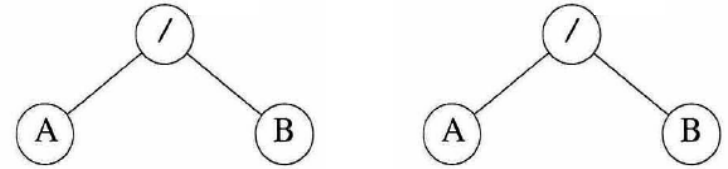
```
treePointer copy(treePointer original) {  
    /* this function returns a treePointer to an exact copy of the original tree */  
    treePointer temp;  
    if (original) {  
        MALLOC(temp, sizeof(*temp));  
        temp->leftChild = copy(original->leftChild);  
        temp->rightChild = copy(original->rightChild);  
        temp->data = original->data;  
        return temp;  
    }  
    return NULL;  
}
```

LRV 방식을 순회



Program 5.6: Copying a binary tree

5.4.2 Testing Equality



- Equivalent BTs
 - have the same structure and the same information in the corresponding nodes
 - Modification of preorder traversal

```
int equal(treePointer first, treePointer second)
{
    return ( (!first && !second) ||
             (first && second &&
              (first->data == second->data) &&
              equal(first->leftChild, second->leftChild) &&
              equal(first->rightChild, second->rightChild)
             ) );
}
```

Program 5.7: Testing for equality of binary trees

5.4.3 The Satisfiability Problem

- Expression

e.g) $(x_1 \wedge \neg x_2) \vee (\neg x_1 \wedge x_3) \vee \neg x_3$

- variables x_1, x_2, \dots, x_n : false/true
- operators: \wedge (and), \vee (or), \neg (not)

- Satisfiability problem:

- Asks if there is an assignment of values to the variables that causes the value of the expression to be true *기, 2, 3을 어떤 값으로만 true가 될까?*

- Solution?

- Let (x_1, \dots, x_n) take on all possible combinations 2^n of true and false values and check the formula for each combination;
- It will take exponential time

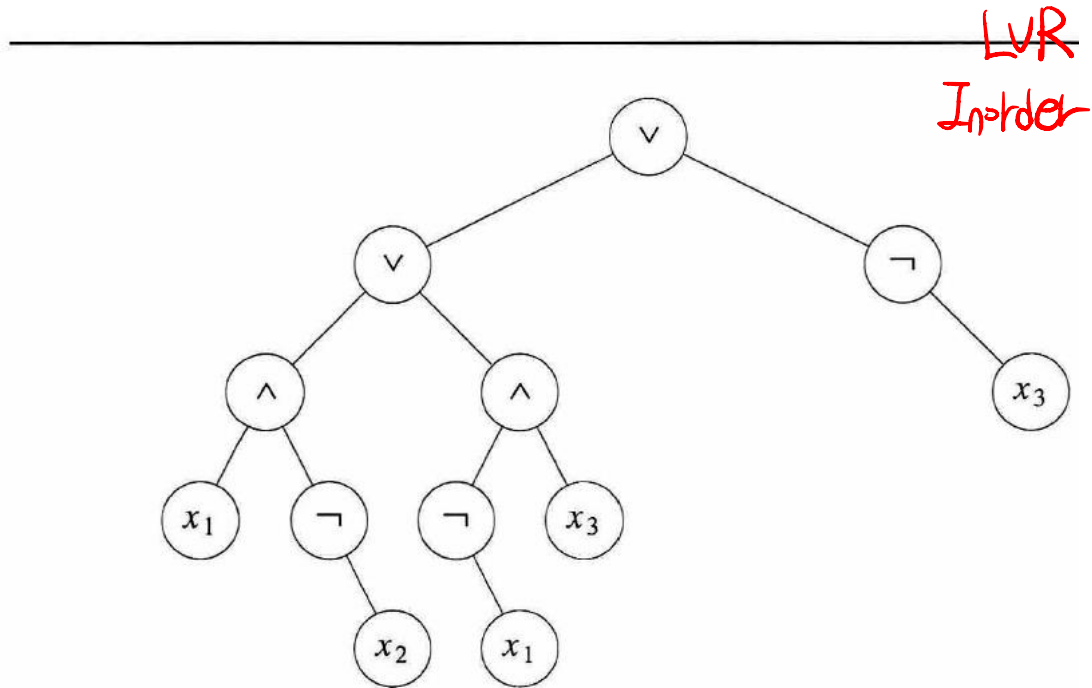


Figure 5.18: Propositional formula in a binary tree

$$(x_1 \wedge \neg x_2) \vee (\neg x_1 \wedge x_3) \vee \neg x_3$$

<i>leftChild</i>	<i>data</i>	<i>value</i>	<i>rightChild</i>
------------------	-------------	--------------	-------------------

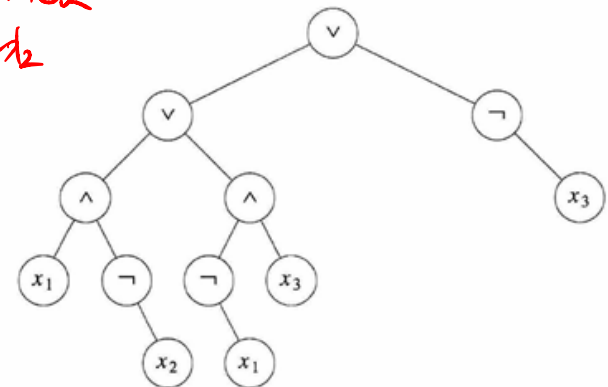
Figure 5.19: Node structure for the satisfiability problem

```

typedef enum {not, and, or, true, false} logical;
typedef struct node *treePointer;
typedef struct node {
    treePointer leftChild;
    logical      data;
    short int    value;
    treePointer rightChild;
} ;

```

← 중간단계의 계산값
예) ~1



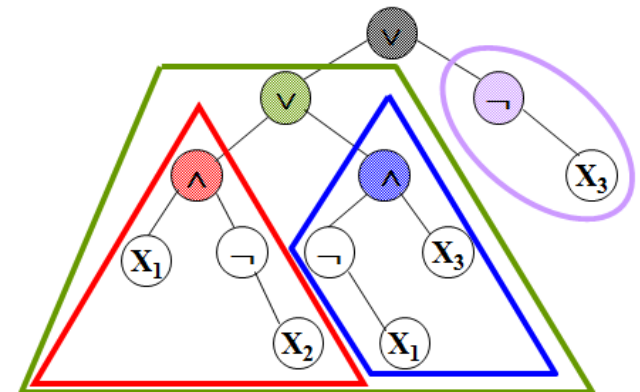
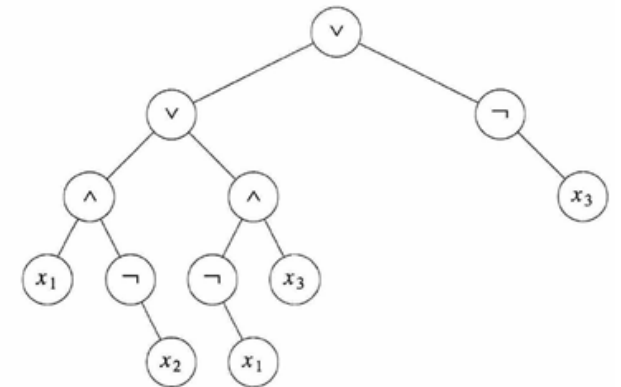
$$(x_1 \wedge \neg x_2) \vee (\neg x_1 \wedge x_3) \vee \neg x_3$$

```

for (all  $2^n$  possible combinations)
{
    generate the next combination;
    replace the variables by their values;
    evaluate root by traversing it in postorder;
    if (root->value) {
        printf(<combination>);
        return;
    }
}
printf("No satisfiable combination\n");

```

Program 5.8: First version of satisfiability algorithm



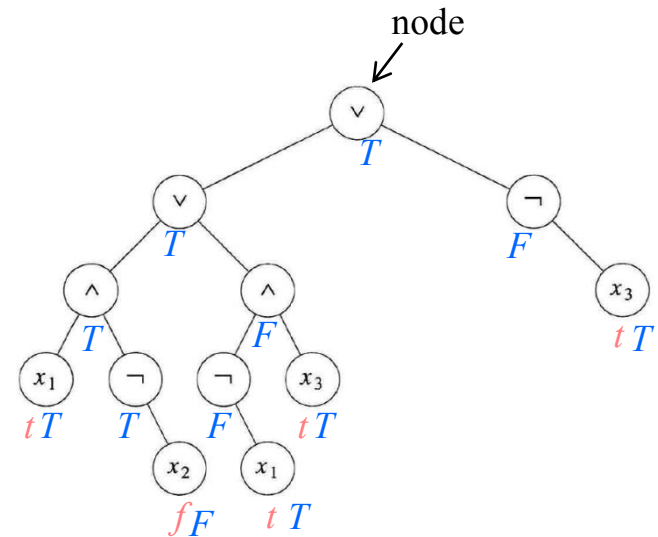
```

void postOrderEval(treePointer node)
{
    /* modified post order traversal to evaluate a
       propositional calculus tree */
    if (node) {
        postOrderEval(node->leftChild);
        postOrderEval(node->rightChild);
        switch(node->data) {
            case not: node->value =
                        !node->rightChild->value;
                        break;
            case and: node->value =
                        node->rightChild->value &&
                        node->leftChild->value;
                        break;
            case or:  node->value =
                        node->rightChild->value ||
                        node->leftChild->value;
                        break;
            case true: node->value = TRUE;
                       break;
            case false: node->value = FALSE;
                       break;
        }
    }
}

```

Handwritten notes:

- Red arrow pointing to `switch(node->data)` with text: **연산자** (operator)
- Red arrow pointing to `node->value` in the `not` case with text: **변수** (variable)
- Red arrow pointing to `node->value` in the `true` case with text: **변수** (variable)



ex) $(x_1, x_2, x_3) = (t, f, t)$

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