

STACKS AND QUEUES

3.1 Stacks

3.2 Stacks Using Dynamic Arrays

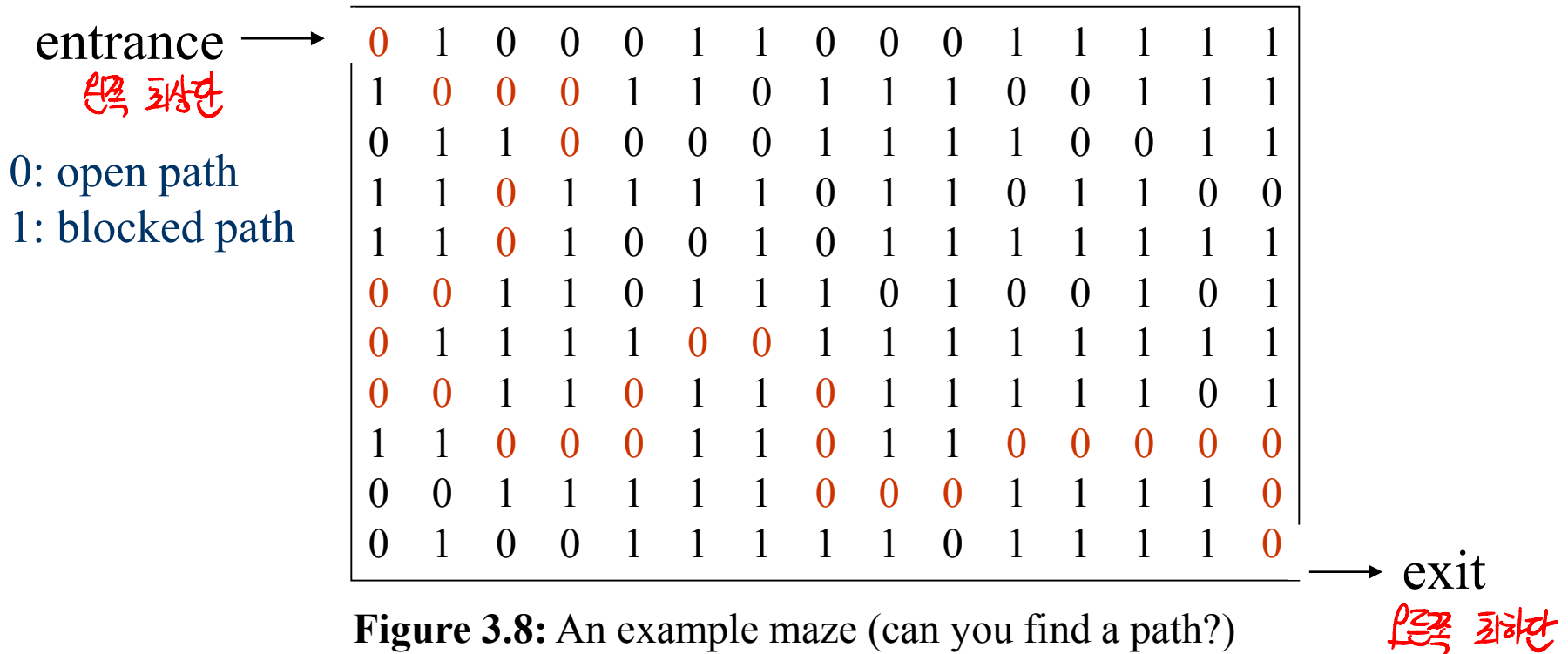
3.3 Queues

3.4 Circular Queues Using Dynamic Arrays

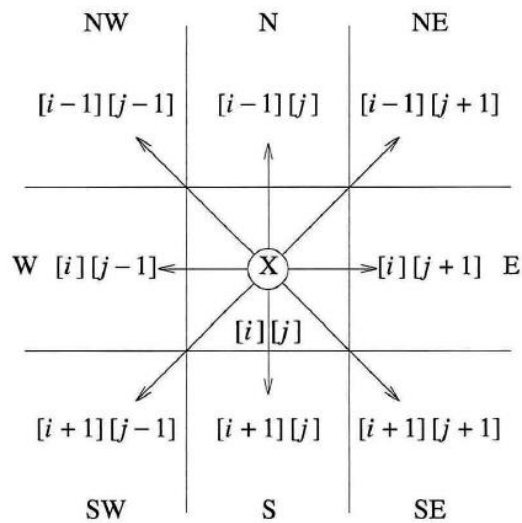
3.5 A Mazing Problem

3.6 Evaluation of Expressions

- Maze: Representation
 - Using a 2D array, $\text{maze}[m][p]$



- Current location X: **maze[i][j]**
 - Possible 8 moves



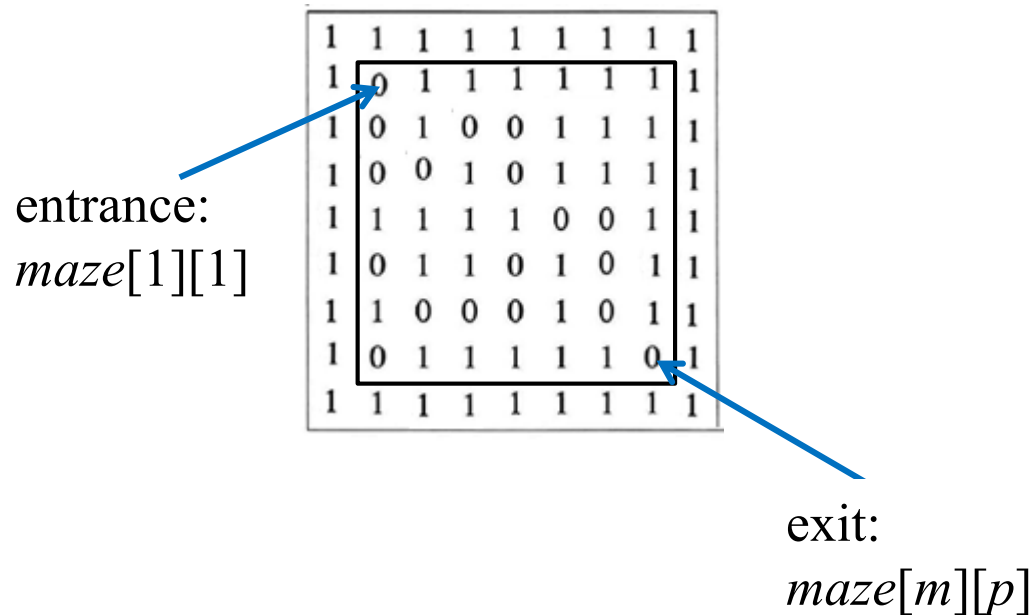
0	1	0	0	0	1	1	0	0	0	1	1	1	1	1
1	0	0	0	1	1	0	1	1	1	0	0	1	1	1
0	1	1	0	0	0	0	1	1	1	1	0	0	1	1
1	1	0	1	1	1	1	0	1	1	0	1	1	0	0
1	1	0	1	0	0	1	0	1	1	1	1	1	1	1
0	0	1	1	0	1	1	1	0	1	0	0	1	0	1
0	1	1	1	1	0	0	1	1	1	1	1	1	1	1
0	0	1	1	0	1	1	0	1	1	1	1	1	0	1
1	1	0	0	0	1	1	0	1	1	0	0	0	0	0
0	0	1	1	1	1	1	0	0	0	1	1	1	1	0
0	1	0	0	1	1	1	1	1	0	1	1	1	1	0

- But, not every position has eight neighbors



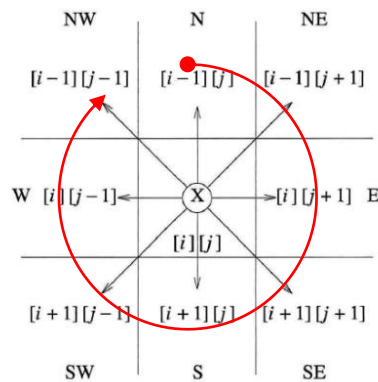
제한할 숫자 다름

- $m \times p$ maze $\rightarrow (m+2) \times (p+2)$ array
 - To avoid checking for the border conditions



- Possible directions: 1D array, *move*

```
typedef struct {
    short int vert;
    short int horiz;
} offsets;
offsets move[8]; /* array of moves for each direction */
```



Name	Dir	<i>move[dir].vert</i>	<i>move[dir].horiz</i>
N	0	-1	0
NE	1	-1	1
E	2	0	1
SE	3	1	1
S	4	1	0
SW	5	1	-1
W	6	0	-1
NW	7	-1	-1

- Current: **maze[row][col]** → Next: **maze[nextRow][nextCol]**
 - nextRow = row + move[dir].vert;
 - nextCol = col + move[dir].horiz;

- Records maze positions already checked: **2D array, mark**

– ex) when visiting a position, maze[row][col]

→ mark[row][col] = 1

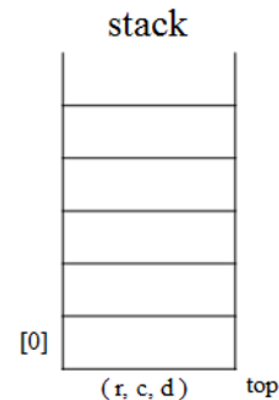
이미 방문했던 곳을 기록!

(= 길이 막혔을때 다시 돌아보기 위함)

	maze				mark			
entrance [1,1]	0	1	1	1	0	0	0	0
	0	1	1	0	0	0	0	0
	1	0	0	1	0	0	0	0
	1	1	0	1	0	0	0	0
	1	0	1	0	0	0	0	0
					exit [5,4]	0	0	0

- Keeps pass history: **Stack**

```
#define MAX_STACK_SIZE 100
typedef struct {
    short int row;
    short int col;
    short int dir;
} element;
element stack[MAX_STACK_SIZE];
```



```

initialize a stack to the maze's entrance coordinates and direction to north;
while (stack is not empty) {
    /* move to position at top of stack */
    <rol, col, dir> = pop from top of stack;
    while (there are more moves from current position) {
        <nextRow, nextCol> = coordinate of next move;
        dir = direction of move;
        if ((nextRow == EXIT_ROW) && (nextCol == EXIT_COL))
            success;
        if (maze[nextRow][nextCol] == 0) && mark[nextRow][nextCol] == 0) {
            /* legal move and haven't been there */
            mark[nextRow][nextCol] = 1;
            /* save current position and direction */
            push<row, col, dir> to the top of the stack;
            row = nextRow;
            col = nextCol;
            dir = north;
        }
    }
}
printf("No path found\n");

```

↙ 지근좌향 방향이 더 있을때까지

Program 3.11: Initial maze algorithm

```

void path(void) {
    int i, row, col, nextRow, nextCol, dir;
    int found=FALSE;
    element position;
    mark[1][1]=1; top=0;
    stack[0].row=1; stack[0].col=1;
    stack[0].dir=1;

```

```

    while (top>-1 && !found) {

```

```

        position = pop();

```

```

        row = position.row; col = position.col; dir = position.dir;

```

```

        while (dir<8 && !found) {

```

```

            /* move in direction dir */

```

```

            nextRow = row + move[dir].vert; nextCol = col + move[dir].horiz;

```

```

            if (nextRow==EXIT_ROW && nextCol==EXIT_COL)

```

```

                found = TRUE;

```

```

            else if ( !maze[nextRow][nextCol] && !mark[nextRow][nextCol]) {

```

```

                mark[nextRow][nextCol] = 1;

```

```

                ① position.row = row; position.col = col; position.dir = ++dir;

```

```

                push(position);

```

```

                ② row = nextRow; col = nextCol; dir = 0;

```

```

            }

```

```

            else ++dir;

```

```

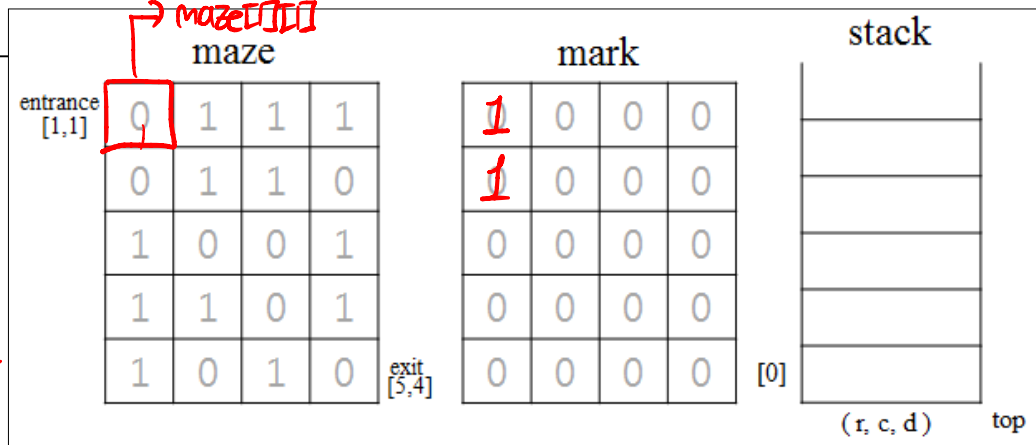
        }

```

```

    }

```



① 2.1.5가 아님
1.1.5인 이유!
이전 위치를 기록해,

② row와 col을 next로 바꿔줌

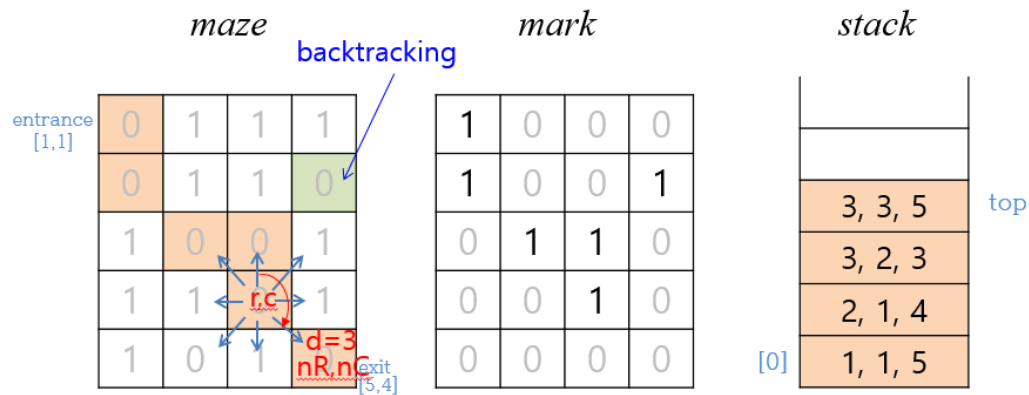
Program 3.12: Maze search function


```

if ( found ) {
    printf("The path is:\n");
    printf("row  col\n");
    for (i=0; i<=top; i++)
        printf("%2d%5d", stack[i].row, stack[i].col);
    printf("%2d%5d\n", row, col);
    printf("%2d%5d\n", EXIT_ROW, EXIT_COL);
}
else printf("The maze does not have a path\n");
}

```

Program 3.12: Maze search function



- Analysis of path

- Computing Time: $O(\dots) \Rightarrow O(M \times P)$

- Each position within the maze is visited no more than once

```

void path(void)
{
    /* output a path through the maze if such a path exists */
    int i, row, col, nextRow, nextCol, dir, found = FALSE;
    element position;
    mark[1][1] = 1; top = 0;
    stack[0].row = 1; stack[0].col = 1; stack[0].dir = 1;
    while (top > -1 && !found) {
        position = pop();
        row = position.row; col = position.col;
        dir = position.dir;
        while (dir < 8 && !found) {
            /* move in direction dir */
            nextRow = row + move[dir].vert;
            nextCol = col + move[dir].horiz;
            if (nextRow == EXIT-ROW && nextCol == EXIT-COL)
                found = TRUE;
            else if (!maze[nextRow][nextCol] &&
                !mark[nextRow][nextCol]) {
                mark[nextRow][nextCol] = 1;
                position.row = row; position.col = col;
                position.dir = ++dir;
                push(position);
                row = nextRow; col = nextCol; dir = 0;
            }
            else ++dir;
        }
    }
    if (found) {
        printf("The path is:\n");
        printf("row col\n");
        for (i = 0; i <= top; i++)
            printf("%2d%5d", stack[i].row, stack[i].col);
        printf("%2d%5d\n", row, col);
        printf("%2d%5d\n", EXIT-ROW, EXIT-COL);
    }
    else printf("The maze does not have a path\n");
}

```

Program 3.12: Maze search function

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3.6 Evaluation of Expressions

3.6.1 Expressions

- Expression statements
 - e.g)
 - $((rear+1==front) || ((rear==MAX_QUEUE_SIZE-1) \&\&!front))$
 - $x = a / b - c + d * e - a * c$
 - Contain operators, operands, and parentheses
- Understanding the meaning
 - Figure out the *order* in which the operations are performed
 - e.g. If $a = 4, b = c = 2, d = e = 3$
 - $x = ((a/b)-c)+(d*e)-(a*c) = ((4/2)-2)+(3*3)-(4*2) = 1 \quad (O)$
 - $x = (a/(b-c+d))*(e-a)*c = (4/(2-2+3))*(3-4)*2 = -2.66666... \quad (X)$

Token	Operator	Precedence ¹	Associativity
() [] → .	function call array element struct or union member	17	left-to-right
-- ++	decrement, increment ²	16	left-to-right
-- ++ ! ~ - + & * sizeof	decrement, increment ³ logical not one's complement unary minus or plus address or indirection size (in bytes)	15	right-to-left
(type)	type cast	14	right-to-left
* / %	multiplicative	13	left-to-right
+ -	binary add or subtract	12	left-to-right
<< >>	shift	11	left-to-right
> >= < <=	relational	10	left-to-right
== !=	equality	9	left-to-right
&	bitwise and	8	left-to-right
^	bitwise exclusive or	7	left-to-right
	bitwise or	6	left-to-right
&&	logical and	5	left-to-right
	logical or	4	left-to-right
?:	conditional	3	right-to-left
= += -= /= *= %= <<= >>= &= ^= =	assignment	2	right-to-left
,	comma	1	left-to-right

1. The precedence column is taken from Harbison and Steele.

2. Postfix form

3. Prefix form

Figure 3.12: Precedence hierarchy for C

3.6.2 Evaluating Postfix Expressions

- Infix notation 중위표기법
 - The standard way of writing expressions
 - Placed a binary operator in-between its two operands
 - Not used by compilers to evaluate expressions
- 후위 표기법 Postfix notation → 컴퓨터는 이렇게 인식함
 - Parenthesis-free notation
 - Each operator appears after its operands

Infix	Postfix
$2+3*4$	$2\ 3\ 4\ *+$
$a*b+5$	$ab\ *5+$
$(1+2)*7$	$1\ 2+7*$
$a*b/c$	$ab\ *c/$
$((a/(b-c+d))*(e-a))*c$	$abc\ -d\ +/ea\ -*c*$
$a/b-c+d*e-a*c$	$ab\ /c-de*+ac*-$

f
Figure 3.13: Infix and postfix notation

- Evaluating postfix expressions: p.129

- Input string: $6\ 2/3-4\ 2*+$

Postfix \rightarrow infix

To evaluate an expression :

- 1) make a single left-to-right scan of it
- 2) place the operands on a stack until we find an operator
- 3) remove, from the stack, the correct number of operands for the operator
- 4) perform the operation, and place the result back on the stack
- 5) continue in this fashion until we reach the end of the expression.
- 6) We then remove the answer from the top of the stack

Token	Stack			Top
	[0]	[1]	[2]	
6	6			0
2	6	2		1
/	6/2			0
3	6/2	3		1
-	6/2-3			0
4	6/2-3	4		1
2	6/2-3	4	2	2
*	6/2-3	4*2		1
+	6/2-3+4*2			0

Figure 3.14: Postfix evaluation

- Representation : stack, expression
 - Assumptions
 - Operators: +, -, *, /, %
 - Operands: single digit integer

```
#define MAX_STACK_SIZE 100
#define MAX_EXPR_SIZE 100

int stack[MAX_STACK_SIZE]; /* global stack */
char expr[MAX_EXPR_SIZE]; /* global input string
                           (a postfix expression) */
```

expr: 6 2/3-4 2*+

↑
52 61 442 34 6


```
token = getToken( &symbol, &n );
```

expr: 6 2/3-4 2*+

```
precedence getToken(char *symbol, int *n)
{
    *symbol = expr[(*n)++];
    switch ( *symbol ) {
        case '(': return lparen;
        case ')': return rparen;
        case '+': return plus;
        case '-': return minus;
        case '/': return divide;
        case '*': return times;
        case '%': return mod;
        case '\0': return eos;
        default : return operand; /* no error checking,
                                   default is operand */
    }
}
```

```
typedef enum {lparen, rparen, plus,
              minus, times, divide,
              mod, eos, operand
} precedence;
```

Program 3.14: Function to get a token from the input string

```

int eval(void) {
    precedence token;
    char symbol;
    int op1, op2;
    int n = 0; /* counter for the expression string */
    int top = -1;
    token = getToken(&symbol, &n);
    while ( token != eos ) {
        if (token == operand) → token이 피연산자이면.. (숫자들)
            push( symbol - '0' ); /* convert: char → integer */
        else {
            op2 = pop();    op1 = pop();
            switch(token) {
                case plus:  push(op1+op2); break;
                case minus: push(op1-op2); break;
                case times:  push(op1*op2); break;
                case divide: push(op1/op2); break;
                case mod:    push(op1%op2);
            }
        }
        token = getToken(&symbol, &n);
    }
    return pop(); /* return result */
}

```

```

typedef enum {lparan, rparan, plus,
              minus, times, divide,
              mod, eos, operand
} precedence;

```

expr: 6 2/3-4 2*+

Input이 postfix
(후위표기법)

→ 아스키코드!

→ 스택에서 먼저꺼낸 피연산자가 2번째 피연산자와 뺄셈

Program 3.13: Function to evaluate a postfix expression

3.6.3 Infix to Postfix

- Algorithm

(1) Fully parenthesize the expression

(2) Move all binary operators so that they replace their corresponding right parentheses

(3) Delete all parentheses *괄호를 모두 없애 줌*

e.g) $a/b - c + d * e - a * c$

(1) $(((((a/b) - c) + (d * e)) - (a * c)))$

(2) $(((((a\ b / c - (d\ e * + (a\ c * -$

(3) $a\ b / c - d\ e * + a\ c * -$

$a * (b + c) * d$

$\rightarrow a\ b\ c + * d *$

- Note

- The order of operands is the same in infix and postfix;

$a/b-c+d*e-a*c$
 $\rightarrow ab/c-de*+ac*-$

- We can form the postfix equivalent the infix expression by
scanning left-to-right

“우선순위가 높은 연산자는 우선순위가 낮은 연산자 위에 올라서서 먼저 자리를 잡지 못하게 한다.”

- Ex 3.3 [Simple expression]: $a + b * c$ (=우선순위가 높으면 stack!)
 - Operands: Passed to the output immediately
 - Operators: Stacked if $ICP > ISP$,
Unstacked if not
(ICP : incoming precedence, ISP : in-stack precedence)
 - Unstacking occurs only when we reach *eos*

Token	Stack			Top	Output
	[0]	[1]	[2]		
<i>a</i>				-1	<i>a</i>
+	+			0	<i>a</i>
<i>b</i>	+			0	<i>ab</i>
*	+	*		1	<i>ab</i>
<i>c</i>	+	*		1	<i>abc</i>
<i>eos</i>				-1	<i>abc*+</i>

Figure 3.15: Translation of $a + b * c$ to postfix

- $a/b - c + d * e - a * c \rightarrow \dots$
 $ab/c - de * + ac * -$

괄호가 있는 경우

- Ex 3.4 [Parenthesized expression]:

$a * (b + c) * d$

Token	Stack			Top	Output
	[0]	[1]	[2]		
<i>a</i>				-1	<i>a</i>
*	*			0	<i>a</i>
(*	(1	<i>a</i>
<i>b</i>	*	(1	<i>ab</i>
+	*	(+	2	<i>ab</i>
<i>c</i>	*	(+	2	<i>abc</i>
)	*			0	<i>abc +</i>
*	*			0	<i>abc +*</i>
<i>d</i>	*			0	<i>abc +*d</i>
<i>eos</i>				0	<i>abc +*d*</i>

Figure 3.16: Translation of $a*(b+c)*d$ to postfix

left parenthesis:

It behaves like a **low-precedence** operator when it is on the stack;

We stack operators until we reach the right parenthesis;

$e/(f+a*d)+c$

$\rightarrow \dots$

/
stack e f a d * + / c +

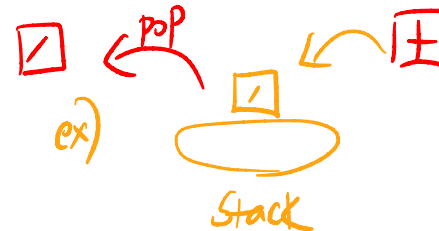
- Implementation

- Uses two types of precedence
 - in-stack precedence (isp), incoming precedence (icp)

```
/* isp and icp arrays -- index is value of precedence
lparen, rparen, plus, minus, times, divide, mod, eos * /
int isp[] = { 0,19,12,12,13,13,13,0};
int icp[] = {20,19,12,12,13,13,13,0};
```

- ‘(‘ has low isp, and high icp → 들어올때는 다 높게버림
- A operator is **removed** from the stack
only if ICP <= ISP

stack에 있는 operator의 우선위가 높을때, 뒤에서부터




```

void postfix(void) {
    char symbol; precedence token;
    int n = 0;
    int top = 0; /* place eos on stack */
    stack[0] = eos;
    for ( token=getToken(&symbol,&n); token!=eos; token=getToken(&symbol,&n) ) {
        if ( token == operand )
            printf("%c", symbol);
        else if ( token == rparen ) {
            /* unstack tokens until left parenthesis */
            while ( stack[top] != lparen )
                printToken( pop() );
            pop(); /* discard the left parenthesis */
        }
        else {
            /* remove and print symbols whose isp is greater
               than or equal to the current token's icp */
            while ( isp[ stack[top] ] >= icp[ token ] )
                printToken( pop() );
            push( token );
        }
    }
    while ( (token=pop()) != eos ) printToken( token );
    printf("\n");
}

```

expr: e/(f+a*d)+c

→ 피연산자이면 stack에 넣기

→) 이면 (이 세 연산자를 pop()해줌

→ 일반적인 경우, stack에 있는 연산자가 우선도가 높으면 pop()해버림.

→ 남아있는 연산자들을 모두 pop()하고 print

```
typedef enum { lparen, rparen, plus, minus, times, divide, mod, eos, operand} precedence;  
int stack[MAX_STACK_SIZE]; /* global stack */  
char expr[MAX_EXPR_SIZE]; /* global input string */
```

```
precedence getToken(char *symbol, int *n)  
{  
    *symbol = expr[(*n)++];  
  
    switch (*symbol) {  
        case '(': return lparen;  
        case ')': return rparen;  
        case '+': return plus;  
        case '-': return minus;  
        case '/': return divide;  
        case '*': return times;  
        case '%': return mod;  
        case '\0': return eos;  
        default : return operand;  
    }  
}
```

Program 3.14: Function to get a token from the input string

- Analysis of postfix

- n : the number of tokens in the expression
- Time complexity: ...

expr: $e/(f+a*d)+c$

```
void postfix(void)
{
    /* output the postfix of the expression. The expression
       string, the stack, and top are global */
    char symbol;
    precedence token;
    int n = 0;
    int top = 0;    /* place eos on stack */
    stack[0] = eos;
    for (token = getToken(&symbol, &n); token != eos;
         token = getToken(&symbol, &n))
    {
        if (token == operand)
            printf("%c", symbol);
        else if (token == rparen) {
            /* unstack tokens until left parenthesis */
            while (stack[top] != lparen)
                printToken(pop());
            pop(); /* discard the left parenthesis */
        }
        else {
            /* remove and print symbols whose isp is greater
               than or equal to the current token's icp */
            while (isp[stack[top]] >= icp[token])
                printToken(pop());
            push(token);
        }
    }
    while ( (token = pop()) != eos)
        printToken(token);
    printf("\n");
}
```

시작 1
 $\theta(n)$
↗

1번 돌까?

\exists lower bound and \exists upper bound

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