The following materials have been collected from the numerous sources such as Stanford CS106 and Harvard CS50 including my own and my students over the years of teaching and experiences of programming. Please help me to keep this tutorial up-to-date by reporting any issues or questions. Please send any comments or criticisms to idebtor@gmail.com. Your assistances and comments will be appreciated.

Profiling Lab – no coding^^

Table of Contents

Purpose of Assignment
Files provided
Step 1. Read profiling.cpp
Step 2. Build and run the profiling
Step 3. Compare the time complexity
Step 4. Compute the growth rate from the time measured
Step 5. Compute the time complexity for a million samples
Step 6. Plot the time complexity to compare
Submitting your solution
Files to submit
Due and Grade points

Purpose of Assignment

This project seeks to verify empirically the accuracy of those analysis's by measuring performance of each algorithm under specific conditions. Performance measurement or program profiling provides detailed empirical data on algorithm performance at different levels of granularity and measures.

"Program Profiling" measures, for example, the space (memory) or time complexity of a program, the usage of particular instructions, or the frequency and duration of function calls. Let us use the elapsed times printed by program execution even though it may not be as accurate as special profiling tools. With small input data size, all times will likely be 0.0000 because the clock interval is too large to measure the execution times. In that case, you should try to get sufficiently accurate results with various data sets and/or extra lines of code repetitions. Our focus on this assignment is to compare the time complexity of two sorting algorithms.

Files provided

profiling.pdf this file

profling.cpp do not changeprintLlst.cpp do not change

• sortx.exe, sortx use this program to measure the elapse time of sorting for 1 million samples

Step 1. Read profiling.cpp

Read the program **profiling.cpp** and be familiar with how it works.

There are two ways to start the program, profilingx.exe. Users may start it by the executable file. Then the program must prompt the user to enter "the number of the maximum sample numbers to sort". If the number entered is less than **STARTING_SAMPLES** (a magic number stored in sort.h), quit the program but with a proper message.

```
// sort.h
const int STARTING_SAMPLES = 1000;
```

Users may give the number of samples in the command line argument. We must check whether or not it is larger than **STARTING_SAMPLES**, exit the program if it is not with a proper message.

Step 2. Build and run the profiling

The quicksort really runs worst if the input data is already sorted. To test the worst-case quicksort, you must pass a sorted data. For other cases, you just pass the randomized data.

- How to compile: Using your own libsort.a as you made in the previous lab.
 \$ g++ profiling.cpp printList.cpp -I../../include -L../../lib -lsort -o profiling,
- How to run:

```
$ ./profiling.exe 30000  # PowerShell
$ profiling 30000  # cmd
```

- How to save the output into a file:
 - \$./profiling.exe 30000 > profiling.txt
 \$ profiling 30000 > profiling.txt
- How to increase the stack size

The worst-case quicksort may not finish completely since it requires a lot of stack memory. In this case, you must increase stack since it is only 1 megabyte by default. The following command increase the stack size to 16 megabytes.

By the way, these compiler option does **not** work in the Windows PowerShell nor Atom console, you must change PowerShell to **cmd windows** before run the command.

Also **mac users** may look it up more in googling since its syntax may be different a bit. You may try like this: -W1,-stack_size,0x10000000

```
$ cmd
$ g++ -Wl,--stack,16777216 profiling.cpp printList.cpp -o profiling -I../../include -
L../../lib -lsort
$ ./profiling
```

Step 3. Compare the time complexity

The code, profiling.cpp, are already invoking sorting functions as we need. Now we would like to compare the elapsed time of following cases:

- 1. insertionSort()
 - A. Best case O(n), Input data is already sorted
 - B. Average case $O(n^2)$, Input data is randomly ordered
 - C. Worst case $O(n^2)$, Input data is reversely ordered
- 2. quickSort()
 - A. Average case O(n log n), Input data is randomly ordered

Sample Run:

```
$ ./profiling 10000
The minimum number of entries is set to 500
Enter the number of max entries to sort: 10000
The maximum sample data size is 10000
        insertionSort(): already sorted - best case.
        Data will NOT be randomized before use.
                  repetitions
                                      sort(sec)
      1000
                        219992
                                       0.000005
      2000
                        121951
                                       0.000008
      3000
                         90212
                                       0.000011
      4000
                         70272
                                       0.000014
      5000
                                       0.000018
                         56832
      6000
                         48603
                                       0.000021
      7000
                         40910
                                       0.000024
      8000
                         35896
                                       0.000028
      9000
                                       0.000031
                         32081
     10000
                         28353
                                       0.000035
        insertionSort(): randomized - average case.
        Randomized Data will be used during sorting.
                  repetitions
                                      sort(sec)
      1000
                          1299
                                       0.000770
      2000
                           351
                                       0.002855
      3000
                           152
                                       0.006579
      4000
                            90
                                       0.011144
      5000
                            69
                                       0.014565
      6000
                            44
                                       0.023295
      7000
                            35
                                       0.028571
      8000
                            28
                                       0.036643
      9000
                            21
                                       0.047810
     10000
                            16
                                       0.063062
        insertionSort(): sorted reversed - worst case.
        Data will NOT be randomized before use.
         N
                  repetitions
                                      sort(sec)
      1000
                           551
                                       0.001815
      2000
                           202
                                       0.004960
      3000
                            92
                                       0.010967
      4000
                            46
                                       0.021761
      5000
                            33
                                       0.030485
      6000
                            23
                                       0.043783
      7000
                            18
                                       0.057889
      8000
                            13
                                       0.077769
      9000
                            11
                                       0.095909
     10000
                             9
                                       0.118778
         quickSort(): randomized - average case.
        Randomized Data will be used during sorting.
         N
                  repetitions
                                      sort(sec)
      1000
                          1054
                                       0.000949
      2000
                           586
                                       0.001706
      3000
                           361
                                       0.002773
      4000
                           242
                                       0.004145
      5000
                           210
                                       0.004781
      6000
                           172
                                       0.005814
      7000
                           146
                                       0.006897
      8000
                           120
                                       0.008350
                           107
      9000
                                       0.009346
     10000
                           101
                                       0.009941
```

Step 4. Compute the growth rate from the time measured

We safely assume that most algorithms approximately have the order of growth of the running time:

$T(N) \approx a N^b$

We want to compute the **constant** "a" and the **growth rate** "b" from data we got from profiling. Once we get a and b, we can compute T(N) where N is 1 billion or 1 million without running the profiling since it may take years or more.

To predict running times, multiply the last observed running time by 2^b and double N, continuing as long as desired. Let us compute b using the double ratio.

Since $T(N) \approx a N^b$, $T(2N) = a (2N)^b$, then

$$\frac{T(2N)}{T(N)} = \frac{a(2N)^b}{aN^b} = \frac{2^b(N)^b}{N^b} = 2^b$$

Log both sides

$$\log \frac{T(2N)}{T(N)} = \log 2^b$$

$$b = \log \frac{T(2N)}{T(N)}$$

In this example, let us choose N = 4000 or 2N = 8000, an average case of the insertion sort shown above. Recall that log we use here is **log base 2**.

$$b = log \frac{T(2N)}{T(N)} = log \frac{t2(8000)}{t1(4000)} = log \frac{0.036643}{0.011144} = 1.717$$

For b in the table below, you must show how you get your answer. You may use a calculator and compute up to two digits after the decimal separator. In my case, I have got 1.717 for the average case of InsertionSort. It will be close to 2.0 for the worst case of insertionSort, $1.2 \sim 1.5$ for the average case quickSort.

Now, we use this b = 1.717 to solve for α when N = 4000, T(N) = 0.011144 in the following:

$$T(N) = a N^{1.717}$$

$$0.011144 = a (4000)^{1.717}$$

$$a = \frac{0.011144}{(4000)^{1.717}}$$

$$a = 7.27 \times 10^{-9}$$

Therefore, we have the growth rate **b = 1.717**, **the constant** $a = 7.27 \times 10^{-9}$ for the insertion sort average case.

Step 5. Compute and measure the time complexity for a million samples

In this step, the question we want to answer is "How long will my program take, as function of the input size?" To help answer this question, we plot data with the problem size N on the x-axis and the running time T(N) on the y-axis. Let's suppose we have the running time, as a function of the input size,

$$T(N) = a N^b$$

where \mathbf{a} is a constant and \mathbf{b} is a growth rate. Even though we already computed a = 7.27 x 10⁻⁹ and b = 1.717 in the previous step, let's compute the constant \mathbf{a} again when N = 8000 as shown below,

$$T(8000) = 0.036643 = a 8000^{1.717}$$

$$a = \frac{0.036643}{8000^{1.717}} = 7.269 x 10^{-9}$$

$$T(N) = 7.27 x 10^{-9} N^{1.717}$$

In conclusion, we confirmed that two constant a's the same. Now you can estimate the elapsed time for one million samples or billion samples as well. With this a and b, we can compute an estimated time for T(N) for N = 1,000,000 or 1 billion samples without running the program, right?

Compute the estimated time for one million samples based your computation of the grow rate b and the constant a in your machine. In your report, show your exact steps how you compute the estimated time. Use the constant "a" that you computed with more samples. Use a proper time unit. For example, don't say 1,234.57 sec., but 20 min 35 sec.

Also fill the blanks with the elapsed times actually measured in your computer while running them with 1 million samples. **Use sortx.exe for this purpose.**

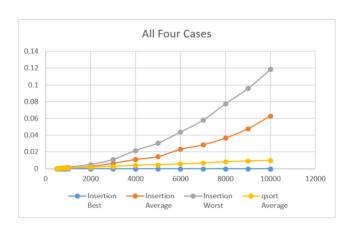
	$T(N) \approx a N^{b}$,		$T(N) \approx a N^{b}$,		$T(N) \approx a N^b$,	
	a =	b =	a =	b =	a =	b =
	insertionSort – Best		insertionSort – Average		insertionSort – Worst	
Ν	10,000	Time for Million	10,000	Time for Million	10,000	Time for Million
Time		Estimated:		Estimated:		Estimated:
Ν	20,000		20,000		20,000	
Time		Measured:		Measured:		Measured:

	$T(N) \approx a N^{b}$							
	a =	b =						
	Average qsort O(N log N): randomized input							
N	10,000	Time for Million						
Time		Estimated:						
Ν	20,000							
Time		Measured:						

Step 6. Plot the time complexity to compare

Plot the data sets that you got from the previous step to compare them graphically as shown below. You may use Excel Chart(분산형) to plot them. An output example combined data from InsertionSort and quickSort for plotting and report.

n	Insertion Best	Insertion Average	Insertion Worst	qsort Average
500	0.000003	0.000180	0.000341	0.000380
600	0.000003	0.000256	0.000446	0.000657
700	0.000004	0.000374	0.000667	0.000652
800	0.000004	0.000477	0.001075	0.000716
900	0.000004	0.000623	0.000990	0.001653
1000	0.000005	0.000770	0.001815	0.000949
2000	0.000008	0.002855	0.004960	0.001706
3000	0.000011	0.006579	0.010967	0.002773
4000	0.000014	0.011144	0.021761	0.004145
5000	0.000018	0.014565	0.030485	0.004781
6000	0.000021	0.023295	0.043783	0.005814
7000	0.000024	0.028571	0.057889	0.006897
8000	0.000028	0.036643	0.077769	0.008350
9000	0.000031	0.047810	0.095909	0.009346
10000	0.000035	0.063062	0.118778	0.009941



Submitting your solution

- On my honour, I pledge that I have neither received nor provided improper assistance in the completion of this assignment.
 - Signed: ______ Section: ____ Student Number: ____
- Make sure your code compiles and runs right before you submit it. Don't make "a tiny last-minute change" and assume your code still compiles. You will not receive sympathy for code that "almost" works.
- If you only manage to work out the Project problem partially before the deadline, you still need to turn it in. However, don't turn it in if it does not compile and run.
- Place your source files in the folder you and I are sharing.
- After submitting, if you realize one of your programs is flawed, you may fix it and submit again as
 long as it is before the deadline. You will have to resubmit any related files together, even if you
 only change one. You may submit as often as you like. Only the last version you submit before
 the deadline will be graded.

Files to submit

- ProfilingReport.docx or one in a readable format: your repost must include the followings:
 - 1. Screen capture of profiling exe output.
 - 2. Complete the performance analysis tables with your data estimated and measured. Show your computation steps for a, b and estimated time for all 4 cases.
 - 3. Include the excel chart and graph for comparing best/average/worst cases.
 - 4. Describe your observation of the time complexity of the insertion sort cases and quicksort case.

Due and Grade points

- Due: 11:55 pm
- 2 points