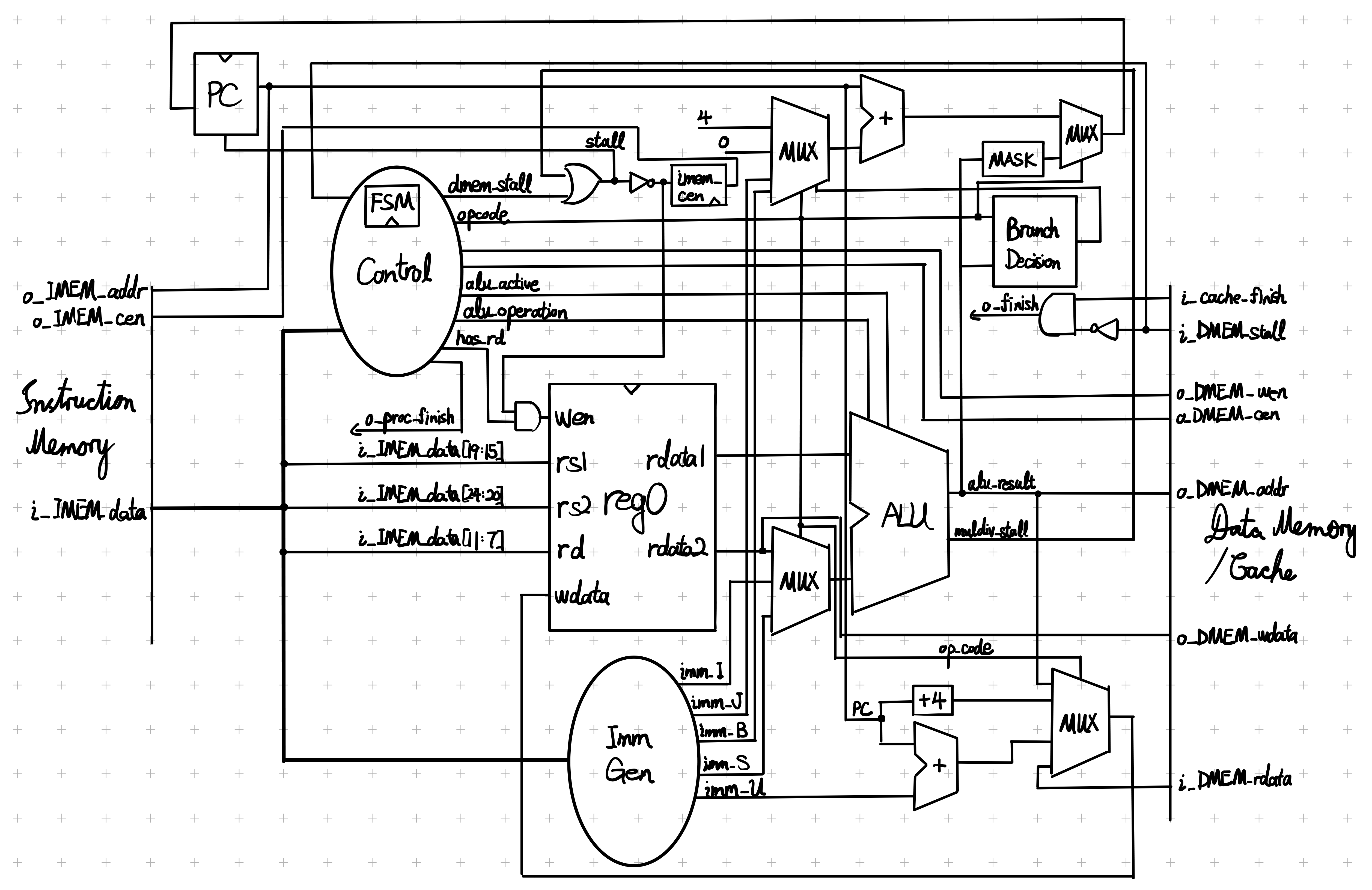
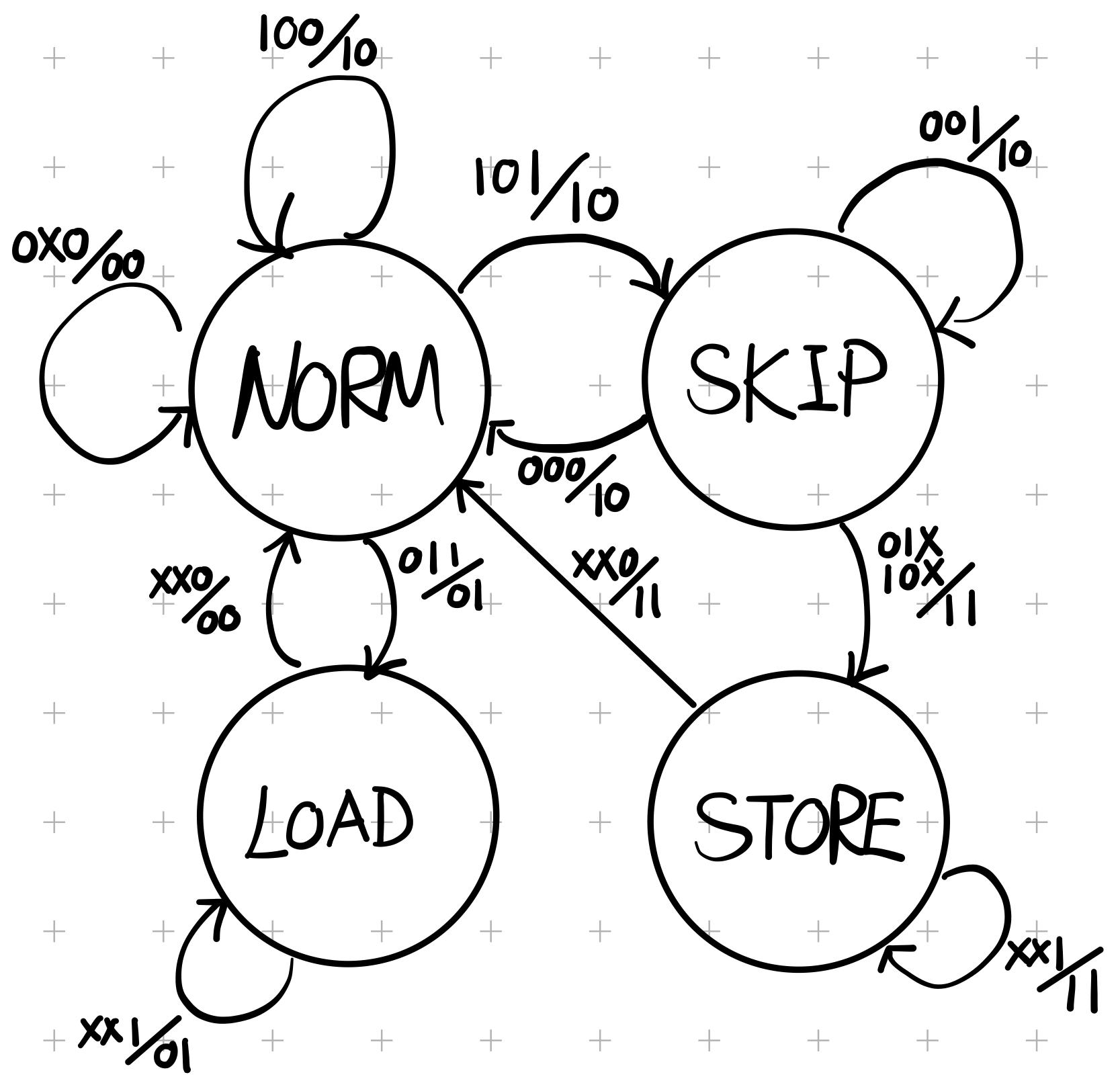
**Work description**

Block diagram:



FSM: arrow format:



For JAL, write PC+4 to R[rd], and PC increment by imm\_J.

For JALR, write PC+4 to R[rd] while ALU sums R[rs1] and imm\_I. Then the result has its LSB masked and stored to PC.

For AUIPC, write PC+imm\_U to R[rd].

For ECALL, set o\_proc\_finish, stop PC incrementation, and set o\_finish to !i\_DMEM\_stall && i\_cache\_finish.

For multicycle instruction, we stall PC, register file and instruction memory if needed, until the stalling signal turns low.

Since no write-back is needed when storing data to memory, we may skip waiting for !i\_DMEM\_stall until memory access is needed for the current instruction. This is implemented by the above Mealy machine.

It is observed that the data memory requires the input write enable signal to be kept high in order to produce an accurate stall signal, so o\_DMEM\_wen is also produced by the FSM.

**Cache design**

**Cache 1**

**Specification:**

2-way set associate cache

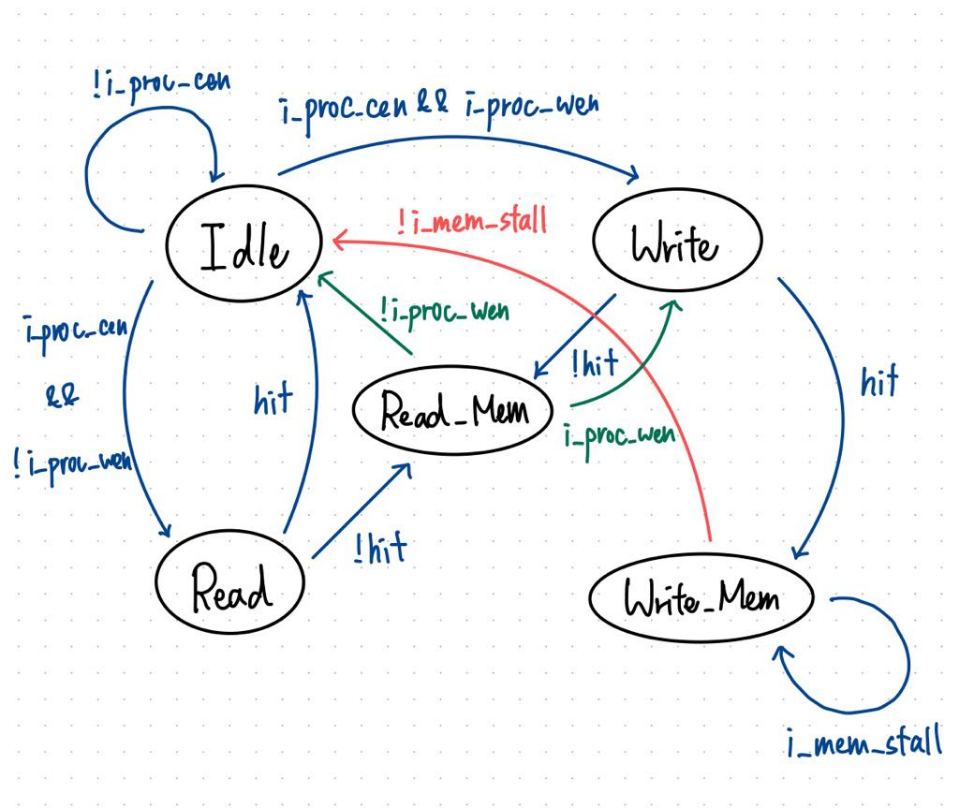
Number of Lines: 8

Line Size: 16 bytes

Write policy: Write through

Replace policy: Random

**Finite State Machine:**



**Cache 2**

**Specification**

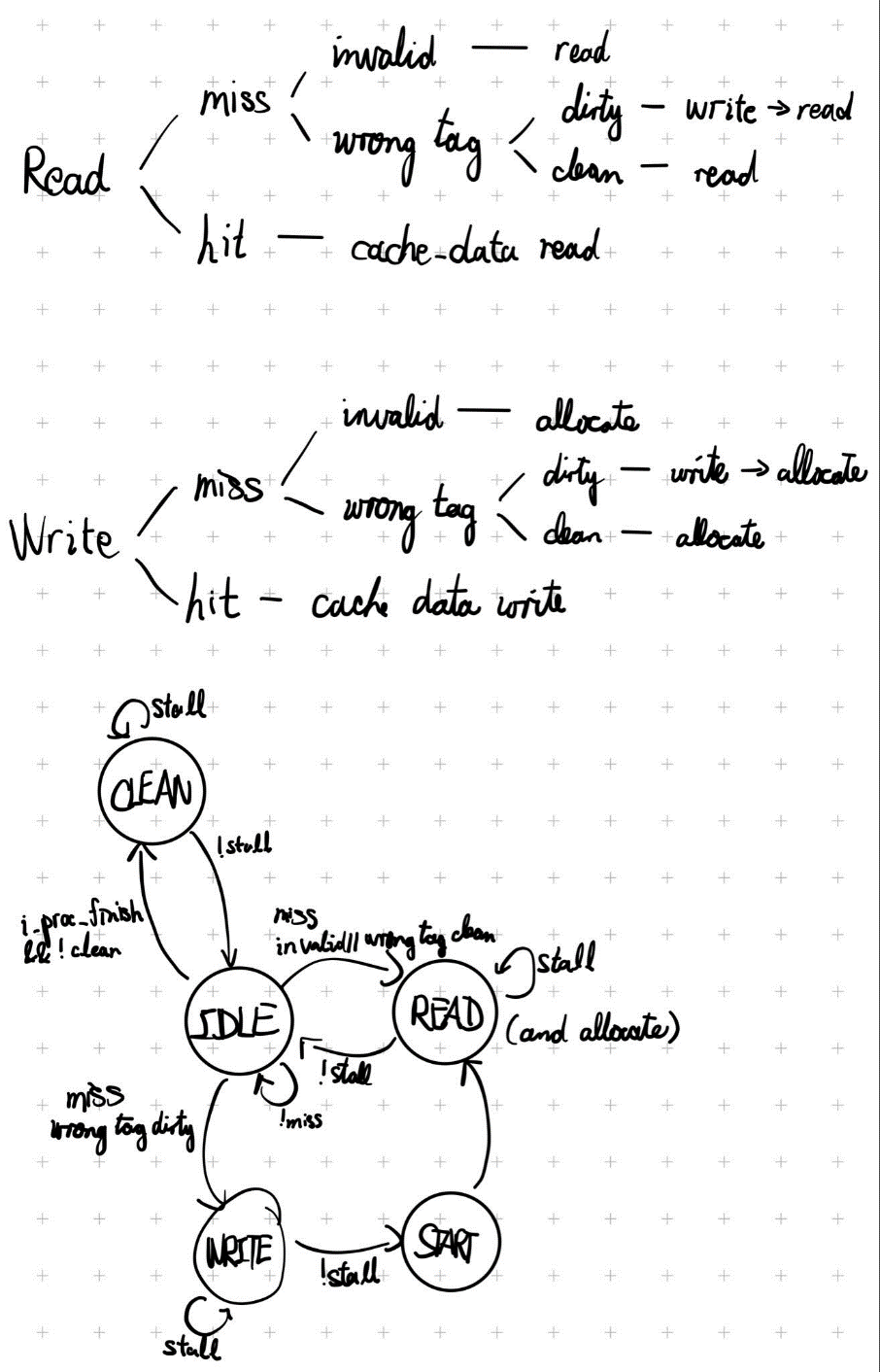
Direct mapped

Number of Blocks: 16

Line Size: 16 bytes

Write policy: Write back

**Finite State Machine:**



**Time performance improvement:**

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
| Instruction Set | I0 | I1 | I2 | I3 | Write Policy |
| No Cache | 77 | 290 | 202 | 1324 |  |
| Store skip | 74 | 261 | 183 | 1228 |  |
| 2-way associative | 63 | 266 | 198 | 742 | Through |
| Improvement | 1.22 | 1.09 | 1.02 | 1.78 |  |
| Store skip | 60 | 232 | 179 | 633 | Through |
| Improvement | 1.23 | 1.13 | 1.02 | 1.94 |  |
| Direct mapped | 55 | 180 | 142 | 426 | Back |
| Improvement | 1.40 | 1.61 | 1.42 | 3.11 |  |
| Store skip | 54 | 165 | 134 | 415 | Back |
| Improvement | 1.37 | 1.58 | 1.37 | 1.96 |  |

**Conclusion:**

Write policy influences more on clock cycle than Cache structure, we can implement a better cache by using write back in 2-way associate cache.

**Register Table**

**Chip:**

**一張含有 文字, 螢幕擷取畫面, 行 的圖片

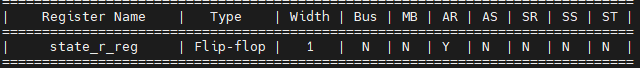
自動產生的描述**

**Reg\_file:**

**一張含有 螢幕擷取畫面, 文字, 字型, 行 的圖片

自動產生的描述**

**MULDIV\_unit**

****

**Cache:**

**一張含有 文字, 螢幕擷取畫面, 電路 的圖片

自動產生的描述**

**Work distribution table**

CPU: 吳子傑

2-way associative cache: 蔡宜豐

Direct mapped cache: 吳子傑

Report: 吳子傑 蔡宜豐