

Bellman-Ford algorithm

The Bellman-Ford algorithm

- **Single-source** shortest-paths problem
- Edge weights may be **negative**.
- Given a **weighted, directed graph** $G = (V, E)$ with source s and weight function w
- The algorithm returns a **Boolean value** indicating whether or not there is a **negative-weight cycle** that is reachable from the source.
- If there is such a **cycle**, the algorithm indicates that **no solution exists**.
- If there is **no such cycle**, the algorithm produces the **shortest paths and their weights**.

Algorithm

BELLMAN-FORD(G, w, s)

```
1  INITIALIZE-SINGLE-SOURCE( $G, s$ )
2  for  $i = 1$  to  $|G.V| - 1$ 
3      for each edge  $(u, v) \in G.E$ 
4          RELAX( $u, v, w$ )
5  for each edge  $(u, v) \in G.E$ 
6      if  $v.d > u.d + w(u, v)$ 
7          return FALSE
8  return TRUE
```

initializing the d and π values of all vertices

$|V| - 1$ passes over the edges of the graph

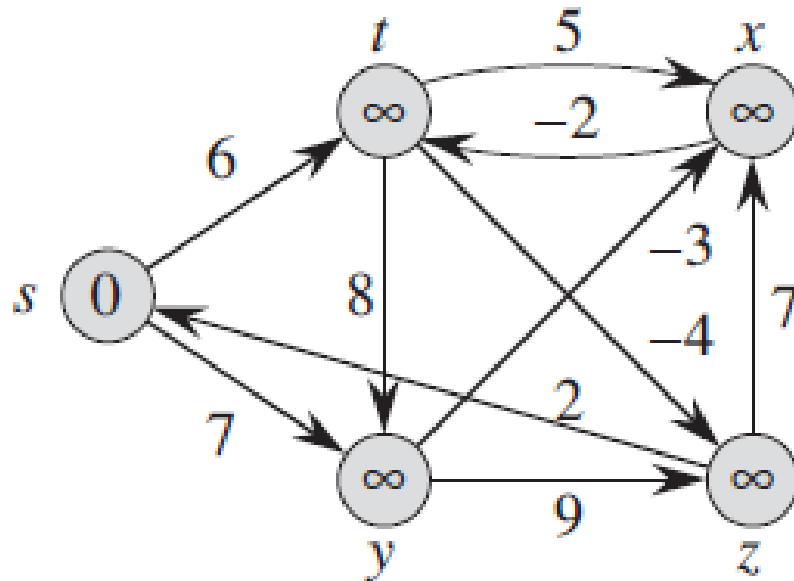
check for a negative-weight cycle

RELAX(u, v, w)

```
1  if  $v.d > u.d + w(u, v)$ 
2       $v.d = u.d + w(u, v)$ 
3       $v.\pi = u$ 
```

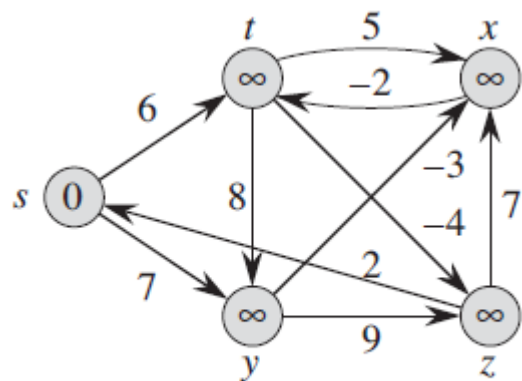
The Bellman-Ford algorithm runs in time $O(VE)$

Example

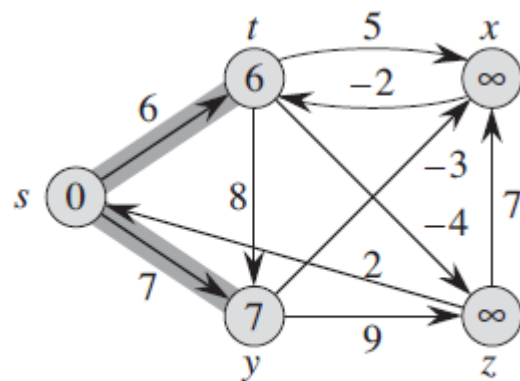


Each pass relaxes the edges in the order

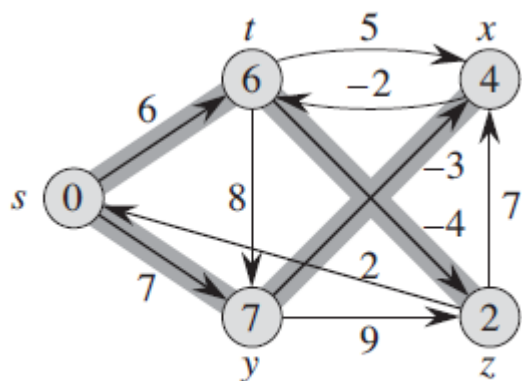
$(t, x), (t, y), (t, z), (x, t), (y, x), (y, z), (z, x), (z, s), (s, t), (s, y)$.



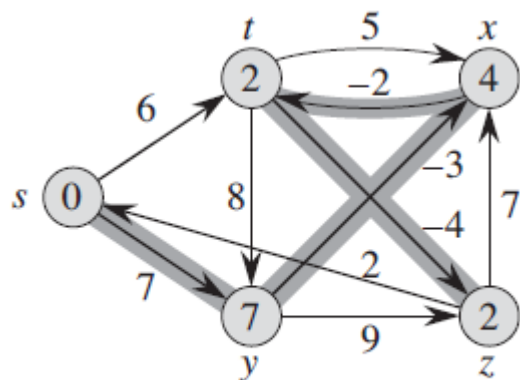
(a)



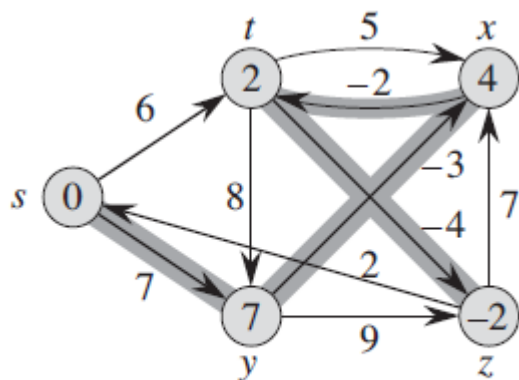
(b)



(c)



(d)



(e)

returns TRUE if graph G contains no negative-weight cycles

Homework

1. Find the shortest path tree from every node to node 1 for the graph of Fig.1 using the Bellman-ford

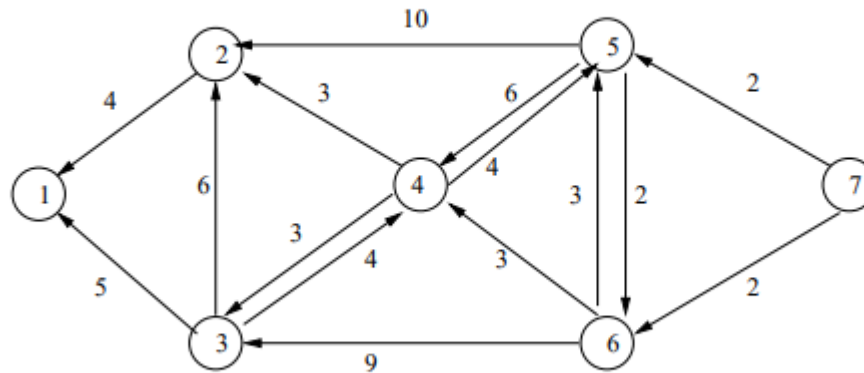


Figure 1: Graph for Problem 1

Homework

- Check negative cycle in the given graph

