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SUBJECT:

DESIGN&ANALYSIS OF ALGORITHM

ASSIGNMENT NO.2

SUBMITTED TO: MAM SAIRA

Question No 1:

Solution:

Recurrence relations are mathematical expressions that define a sequence of values based on previous terms in the sequence. In simpler terms, they describe how to generate the next term in a sequence using one or more of the previous terms.

Fibonacci Sequence:

The Fibonacci sequence is perhaps the most famous example of a recurrence relation. It starts with two initial terms, 0 and 1, and each subsequent term is the sum of the two preceding terms. Mathematically, it can be defined as:

$F(n) = F(n-1) + F(n-2)$ for $n \geq 2$, with initial conditions $F(0) = 0$ and $F(1) = 1$.

Factorial Function:

The factorial function is defined recursively as the product of all positive integers up to a given integer n . Mathematically:

$n! = n \times (n-1)!$ for $n > 0$, with base case $0! = 1$.

Towers of Hanoi:

The Towers of Hanoi is a classic problem in computer science and mathematics. It involves moving a tower of disks from one peg to another, following certain rules, using a third peg as an auxiliary. The minimum number of moves required to solve the problem can be expressed recursively:

$T(n) = 2T(n-1) + 1$ for $n > 0$, with base case $T(0) = 0$.

Binary Search:

Binary search is a search algorithm that finds the position of a target value within a sorted array. It works by repeatedly dividing the search interval in half. The algorithm can be represented recursively:

$$BSearch(A, x, low, high) = \begin{cases} -1 & \text{if } low > high \\ BSearch(A, x, low, mid - 1) & \text{if } A[mid] > x \\ BSearch(A, x, mid + 1, high) & \text{if } A[mid] < x \\ mid & \text{if } A[mid] = x \end{cases}$$

Merge Sort:

Merge sort is a sorting algorithm that follows the divide-and-conquer strategy. It divides the input array into two halves, sorts each half independently, and then merges the sorted halves. The recurrence relation for merge sort can be expressed as:

$T(n) = 2T(n/2) + O(n)$, where $T(n)$ represents the time complexity of sorting an array of size n .

These examples demonstrate how recurrence relations can be used to define and analyze various mathematical problems and algorithms. They provide a powerful tool for understanding the behavior and complexity of algorithms and sequences.

Question No 2:

Solution:

Solving recurrence relations involves finding a closed-form expression for a sequence defined recursively. Several methods are commonly used:

1. **Substitution Method:** Guessing a form for the solution and then proving it correct by induction. This method is intuitive but may not always lead to the solution.
2. **Recurrence Tree Method:** Visualizing the recurrence as a tree, then analyzing the structure and summing up the costs. This method is useful for recurrences describing the cost of recursive algorithms.
3. **Master Theorem:** A powerful tool for solving recurrence relations that arise in the analysis of divide-and-conquer algorithms. It provides a straightforward way to determine the asymptotic behavior of solutions.

4. **Characteristic Equation Method:** Transforming the recurrence relation into a polynomial equation using generating functions or substitution. This method is particularly effective for linear homogeneous recurrence relations with constant coefficients.
5. **Generating Functions:** Representing sequences as power series and manipulating them algebraically. This method is versatile and can handle various types of recurrences, but it may involve complex algebraic manipulations.
6. **Matrix Method:** Representing the recurrence relation as a matrix equation and finding its eigenvalues and eigenvectors. This method is suitable for solving linear homogeneous recurrence relations with constant coefficients.
7. **Guess and Verify Method:** Proposing a solution based on observation or intuition, then proving its correctness by substitution into the recurrence relation.

Question No 3:

Solution:

Recurrence relations are equations that recursively define a sequence in terms of its previous terms. The general form of a linear recurrence relation of degree 2 is:

$$T(n) = a \cdot T(n-1) + b \cdot T(n-2) + c$$

Where:

- $T(n)$ is the value of the sequence at index n ,
- a , b and c are constants,
- $T(n-1)$ and $T(n-2)$ are the values of the sequence at indices $n-1$ and $n-2$ respectively.

The iterative method to solve such recurrence relations involves starting from the initial values of the sequence and iteratively calculating subsequent values until the desired value $T(n)$ is reached.

Here's how the iterative method works:

1. **Base Cases:** First, handle the base cases. These are the initial values of the sequence that are provided explicitly. For example, $T(0)$ and $T(1)$.
2. **Iteration:** Iterate over the indices starting from the index corresponding to the next value after the last initial value up to the desired value n . At each step, calculate the current value of the sequence using the recurrence relation based on the previous terms.
3. **Updating Previous Terms:** Update the values of the previous terms as you progress through the iteration. This ensures that the calculation for each step uses the correct values of the sequence.
4. **Return:** Finally, return the value of $T(n)$ calculated after the iteration.