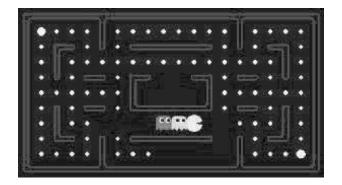
# Games II





#### Announcement

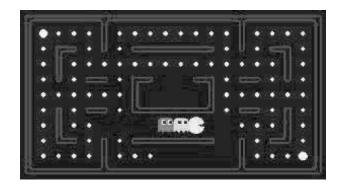
• Midterm is next week (Wednesday, 5/8, 6pm-8pm)

• Topics: all material up to and including today's lecture

• Logistics: look for detailed post on Ed



# Games: alpha-beta pruning recap

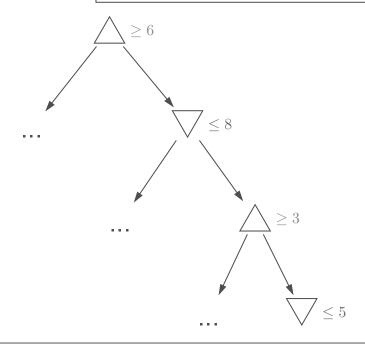


### Alpha-beta pruning



### Key idea: optimal path-

The optimal path is path that minimax policies take. Values of all nodes on path are the same.

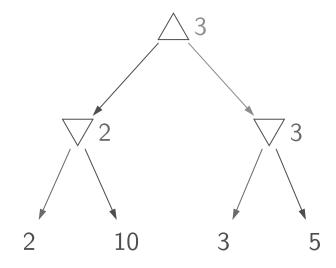


- $a_s$ : lower bound on value of max node s
- $b_s$ : upper bound on value of min node s
- Prune a node if its interval doesn't have non-trivial overlap with every ancestor (store  $\alpha_s = \max_{s' \leq s} a_{s'}$  and  $\beta_s = \min_{s' \prec s} b_{s'}$ )

# Move ordering

Pruning depends on order of actions.

Can't prune the 5 node:



[live solution: alpha-beta pruning]

### Move ordering

Which ordering to choose?

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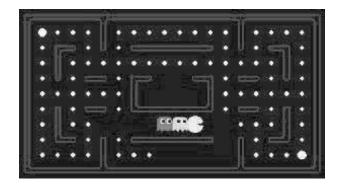
- Worst ordering:  $O(b^{2 \cdot d})$  time
- Best ordering:  $O((\sqrt{b-\frac{3}{4}}+\frac{1}{2})^{2\cdot d})\simeq O(b^{2\cdot 0.5d})$  time
- Random ordering:  $O(b^{2 \cdot 0.75d})$  time when b=2
- Random ordering:  $O((\frac{b-1+\sqrt{b^2+14b+1}}{4})^{2\cdot d})$  for general b

In practice, can use evaluation function Eval(s):

- Max nodes: order successors by decreasing Eval(s)
- Min nodes: order successors by increasing Eval(s)



# Games: TD-learning



#### **Evaluation function**

Old: hand-crafted



Example: chess-

Eval(s) = material + mobility + king-safety + center-controlmaterial =  $10^{100}(K - K') + 9(Q - Q') + 5(R - R') +$ 3(B - B' + N - N') + 1(P - P')mobility = 0.1(num-legal-moves - num-legal-moves'). . .

New: learn from data

$$\mathsf{Eval}(s) = V(s; \mathbf{w})$$

#### Model for evaluation functions

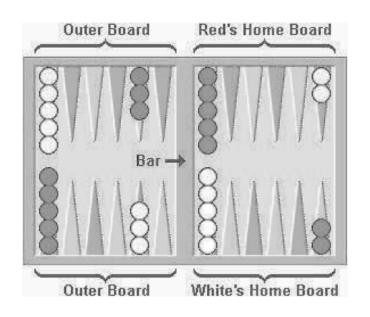
Linear:

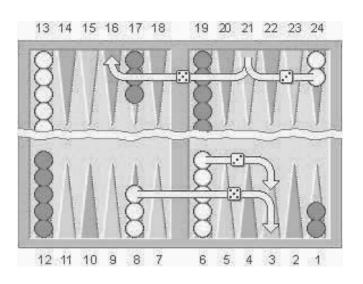
$$V(s; \mathbf{w}) = \mathbf{w} \cdot \phi(s)$$

Neural network:

$$V(s; \mathbf{w}, \mathbf{v}_{1:k}) = \sum_{j=1}^{k} w_j \sigma(\mathbf{v}_j \cdot \phi(s))$$

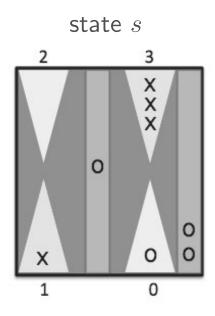
# Example: Backgammon





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## Features for Backgammon



```
Features \phi(s): [(\# \text{ o in column 0}) = 1] \colon 1[(\# \text{ o on bar})] \qquad \vdots \qquad 1[(\text{fraction o removed})] \qquad \vdots \qquad \frac{1}{2}[(\# \times \text{ in column 1}) = 1] \colon 1[(\# \times \text{ in column 3}) = 3] \colon 1[(\text{is it o's turn})] \qquad \vdots \qquad 1
```

### Generating data

Generate using policies based on current  $V(s; \mathbf{w})$ :

$$\pi_{\mathsf{agent}}(s; \mathbf{w}) = \arg \max_{a \in \mathsf{Actions}(s)} V(\mathsf{Succ}(s, a); \mathbf{w})$$

$$\pi_{\mathsf{opp}}(s; \mathbf{w}) = \arg\min_{a \in \mathsf{Actions}(s)} V(\mathsf{Succ}(s, a); \mathbf{w})$$

Note: don't need to randomize ( $\epsilon$ -greedy) because game is already stochastic (backgammon has dice) and there's function approximation

Generate episode:

$$s_0; a_1, r_1, s_1; a_2, r_2, s_2; a_3, r_3, s_3; \dots; a_n, r_n, s_n$$

## Learning algorithm

#### Episode:

$$s_0; a_1, r_1, s_1; a_2, r_2, s_2, a_3, r_3, s_3; \ldots, a_n, r_n, s_n$$

A small piece of experience:

Prediction:

$$V(s; \mathbf{w})$$

Target:

$$r + \gamma V(s'; \mathbf{w})$$

### General framework

Objective function:

$$\frac{1}{2}(\operatorname{prediction}(\mathbf{w}) - \operatorname{target})^2$$

Gradient:

$$(prediction(\mathbf{w}) - target)\nabla_{\mathbf{w}} prediction(\mathbf{w})$$

Update:

$$\mathbf{w} \leftarrow \mathbf{w} - \eta \underbrace{\left( \mathsf{prediction}(\mathbf{w}) - \mathsf{target} \right) \nabla_{\mathbf{w}} \mathsf{prediction}(\mathbf{w})}_{\mathsf{gradient}}$$

# Temporal difference (TD) learning



#### Algorithm: TD learning-

On each 
$$(s, a, r, s')$$
: 
$$\mathbf{w} \leftarrow \mathbf{w} - \eta \underbrace{[V(s; \mathbf{w})}_{\text{prediction}} - \underbrace{(r + \gamma V(s'; \mathbf{w}))}_{\text{target}}] \nabla_{\mathbf{w}} V(s; \mathbf{w})$$

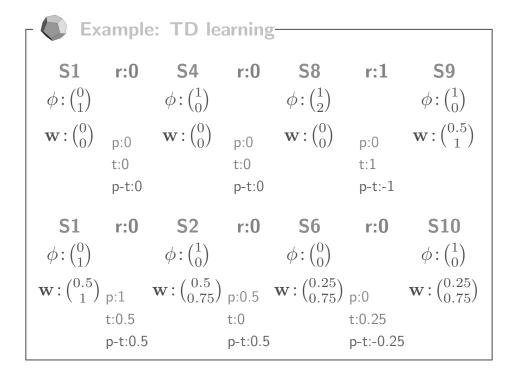
For linear functions:

$$V(s; \mathbf{w}) = \mathbf{w} \cdot \phi(s)$$

$$\nabla_{\mathbf{w}} V(s; \mathbf{w}) = \phi(s)$$

### Example of TD learning

Step size  $\eta = 0.5$ , discount  $\gamma = 1$ , reward is end utility



### Comparison



#### Algorithm: TD learning-

On each 
$$(s, a, r, s')$$
: 
$$\mathbf{w} \leftarrow \mathbf{w} - \eta [\underbrace{\hat{V}_{\pi}(s; \mathbf{w})}_{\text{prediction}} - \underbrace{(r + \gamma \hat{V}_{\pi}(s'; \mathbf{w}))}_{\text{target}}] \nabla_{\mathbf{w}} \hat{V}_{\pi}(s; \mathbf{w})$$



#### Algorithm: Q-learning-

On each 
$$(s, a, r, s')$$
: 
$$\mathbf{w} \leftarrow \mathbf{w} - \eta [\underbrace{\hat{Q}_{\text{opt}}(s, a; \mathbf{w})}_{\text{prediction}} - \underbrace{(r + \gamma \max_{a' \in \text{Actions}(s)} \hat{Q}_{\text{opt}}(s', a'; \mathbf{w}))}_{\text{target}}] \nabla_{\mathbf{w}} \hat{Q}_{\text{opt}}(s, a; \mathbf{w})$$

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### Comparison

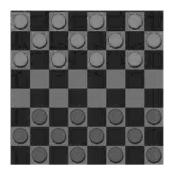
#### Q-learning:

- Operate on  $\hat{Q}_{\mathsf{opt}}(s, a; \mathbf{w})$
- Off-policy: value is based on estimate of optimal policy
- $\bullet$  To use, don't need to know MDP transitions  $T(s,a,s^\prime)$

#### TD learning:

- Operate on  $\hat{V}_{\pi}(s;\mathbf{w})$
- ullet On-policy: value is based on exploration policy (usually based on  $\hat{V}_{\pi}$ )
- $\bullet\,$  To use, need to know rules of the game  $\operatorname{Succ}(s,a)$

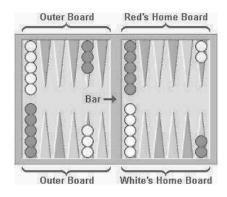
### Learning to play checkers



#### Arthur Samuel's checkers program [1959]:

- Learned by playing itself repeatedly (self-play)
- Smart features, linear evaluation function, use intermediate rewards
- Used alpha-beta pruning + search heuristics
- Reach human amateur level of play
- IBM 701: 9K of memory!

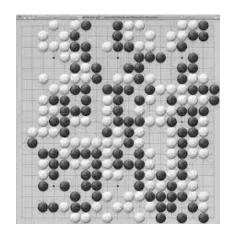
### Learning to play Backgammon



#### Gerald Tesauro's TD-Gammon [1992]:

- Learned weights by playing itself repeatedly (1 million times)
- Dumb features, neural network, no intermediate rewards
- Reached human expert level of play, provided new insights into opening

### Learning to play Go



#### AlphaGo Zero [2017]:

- Learned by self play (4.9 million games)
- Dumb features (stone positions), neural network, no intermediate rewards, Monte Carlo Tree Search
- Beat AlphaGo, which beat Le Sedol in 2016
- Provided new insights into the game



### Summary so far

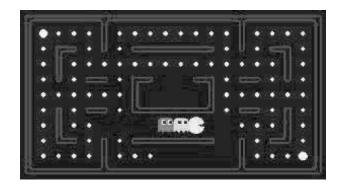
- Parametrize evaluation functions using features
- TD learning: learn an evaluation function

$$(prediction(w) - target)^2$$

Up next:



# Games: simultaneous games





### Question

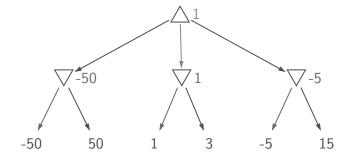
For a simultaneous two-player zero-sum game (like rock-paper-scissors), can you still be optimal if you reveal your strategy?

no

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#### Turn-based games:





#### Simultaneous games:



?

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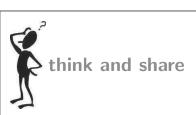
### Two-finger Morra



Example: two-finger Morra-

Players A and B each show 1 or 2 fingers. If both show 1, B gives A 2 dollars. If both show 2, B gives A 4 dollars. Otherwise, A gives B 3 dollars.

[play with a partner]



# Question

What was the outcome?

player A chose 1, player B chose 1
player A chose 1, player B chose 2
player A chose 2, player B chose 1
player A chose 2, player B chose 2



## Payoff matrix



#### Definition: single-move simultaneous game-

 $\mathsf{Players} = \{\mathsf{A},\mathsf{B}\}$ 

Actions: possible actions

V(a,b): A's utility if A chooses action a, B chooses b

(let V be payoff matrix)



Example: two-finger Morra payoff matrix

 $A \setminus B$  1 finger 2 fingers

1 finger 2 -3

2 fingers -3

# Strategies (policies)



Definition: pure strategy-

A pure strategy is a single action:  $a \in Actions$ 



Definition: mixed strategy-

A mixed strategy is a probability distribution  $0 \le \pi(a) \le 1$  for  $a \in \mathsf{Actions}$ 



Example: two-finger Morra strategies

Always 1:  $\pi = [1, 0]$ 

Always 2:  $\pi = [0, 1]$ 

Uniformly random:  $\pi = [\frac{1}{2}, \frac{1}{2}]$ 

#### Game evaluation



Definition: game evaluation

The **value** of the game if player A follows  $\pi_A$  and player B follows  $\pi_B$  is

$$V(\pi_A, \pi_B) = \sum_{a,b} \pi_A(a) \pi_B(b) V(a,b)$$



Example: two-finger Morra-

Player A always chooses 1:  $\pi_A = [1, 0]$ 

Player B picks randomly:  $\pi_{B} = [\frac{1}{2}, \frac{1}{2}]$ 

Value: 
$$-\frac{1}{2}$$

[whiteboard: matrix]

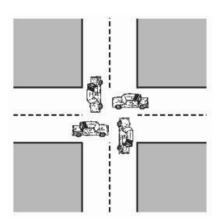
### How to optimize?

Game value:

$$V(\pi_A, \pi_B)$$

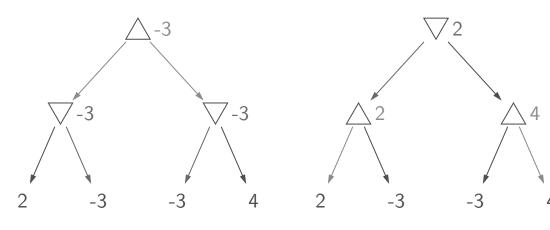
Challenge: player A wants to maximize, player B wants to minimize...

#### simultaneously



### Pure strategies: who goes first?

Player A goes first: Player B goes first:





Proposition: going second is no worse

$$\max_{a} \min_{b} V(a, b) \le \min_{b} \max_{a} V(a, b)$$

### Mixed strategies



Example: two-finger Morra-

Player A reveals:  $\pi_A = [\frac{1}{2}, \frac{1}{2}]$ 

Value  $V(\pi_A, \pi_B) = \pi_B(1)(-\frac{1}{2}) + \pi_B(2)(+\frac{1}{2})$ 

Optimal strategy for player B is  $\pi_B = [1,0]$  (pure!)



Proposition: second player can play pure strategy

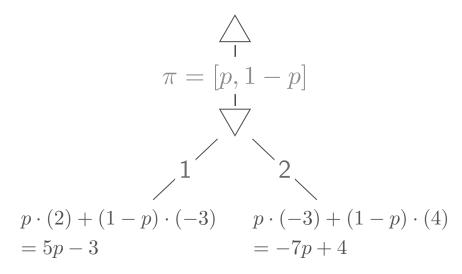
For any fixed mixed strategy  $\pi_A$ :

$$\min_{\pi_{\mathsf{B}}} V(\pi_{\mathsf{A}}, \pi_B)$$

can be attained by a pure strategy.

### Mixed strategies

Player A first reveals his/her mixed strategy

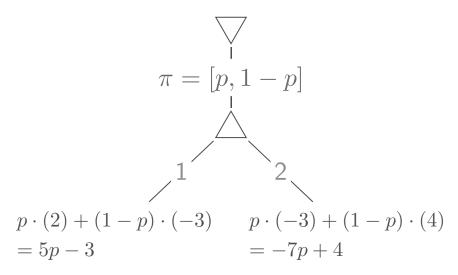


Minimax value of game:

$$\max_{0 \le p \le 1} \min\{5p - 3, -7p + 4\} = \boxed{-\frac{1}{12}} \text{ (with } p = \frac{7}{12}\text{)}$$

### Mixed strategies

Player B first reveals his/her mixed strategy



Minimax value of game:

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$$\min_{p \in [0,1]} \max\{5p - 3, -7p + 4\} = \boxed{-\frac{1}{12}} \text{ (with } p = \frac{7}{12}\text{)}$$

#### General theorem



#### Theorem: minimax theorem [von Neumann, 1928]

For every simultaneous two-player zero-sum game with a finite number of actions:

$$\max_{\pi_{\mathsf{A}}} \min_{\pi_{\mathsf{B}}} V(\pi_{\mathsf{A}}, \pi_{\mathsf{B}}) = \min_{\pi_{\mathsf{B}}} \max_{\pi_{\mathsf{A}}} V(\pi_{\mathsf{A}}, \pi_{\mathsf{B}}),$$

where  $\pi_A$ ,  $\pi_B$  range over **mixed strategies**.

Upshot: revealing your optimal mixed strategy doesn't hurt you!

Proof: linear programming duality

Algorithm: compute policies using linear programming



# Summary

• Challenge: deal with simultaneous min/max moves

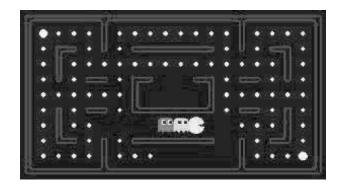
• Pure strategies: going second is better

• Mixed strategies: doesn't matter (von Neumann's minimax theorem)

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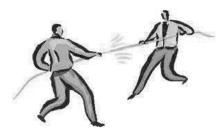


# Games: non-zero-sum games



# Utility functions

Competitive games: minimax (linear programming)



Collaborative games: pure maximization (plain search)



Real life: ?

#### Prisoner's dilemma



Example: Prisoner's dilemma-

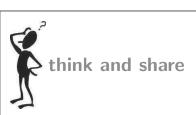
Prosecutor asks A and B individually if each will testify against the other.

If both testify, then both are sentenced to 5 years in jail.

If both refuse, then both are sentenced to 1 year in jail.

If only one testifies, then he/she gets out for free; the other gets a 10-year sentence.

[play with a partner]



# Question

What was the outcome?

player A testified, player B testified
player A refused, player B testified
player A testified, player B refused
player A refused, player B refused

#### Prisoner's dilemma



Example: payoff matrix-

 $\mathsf{B} \setminus \mathsf{A} \qquad \text{testify} \qquad \text{refuse}$ 

testify A = -5, B = -5 A = -10, B = 0

refuse A = 0, B = -10 A = -1, B = -1



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Definition: payoff matrix-

Let  $V_p(\pi_A, \pi_B)$  be the utility for player p.

### Nash equilibrium

Can't apply von Neumann's minimax theorem (not zero-sum), but get something weaker:



#### Definition: Nash equilibrium

A **Nash equilibrium** is  $(\pi_A^*, \pi_B^*)$  such that no player has an incentive to change his/her strategy:

$$V_{\mathsf{A}}(\pi_{\mathsf{A}}^*, \pi_B^*) \geq V_{\mathsf{A}}(\pi_{\mathsf{A}}, \pi_{\mathsf{B}}^*)$$
 for all  $\pi_A$ 

$$V_{\mathsf{B}}(\pi_{\mathsf{A}}^*, \pi_B^*) \geq V_{\mathsf{B}}(\pi_{\mathsf{A}}^*, \pi_{\mathsf{B}})$$
 for all  $\pi_B$ 



#### Theorem: Nash's existence theorem [1950]

In any finite-player game with finite number of actions, there exists at least one Nash equilibrium.

### Examples of Nash equilibria



**Example: Two-finger Morra-**

Nash equilibrium: A and B both play  $\pi = [\frac{7}{12}, \frac{5}{12}].$ 



Example: Collaborative two-finger Morra-

Two Nash equilibria:

- A and B both play 1 (value is 2).
- A and B both play 2 (value is 4).



Example: Prisoner's dilemma-

Nash equilibrium: A and B both testify.



### Summary so far

#### Simultaneous zero-sum games:

- von Neumann's minimax theorem
- Multiple minimax strategies, single game value

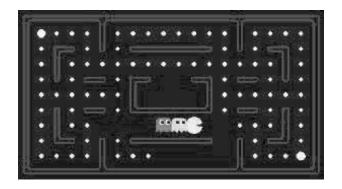
#### Simultaneous non-zero-sum games:

- Nash's existence theorem
- Multiple Nash equilibria, multiple game values

Huge literature in game theory / economics



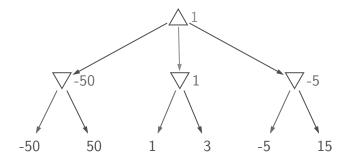
# Games: recap





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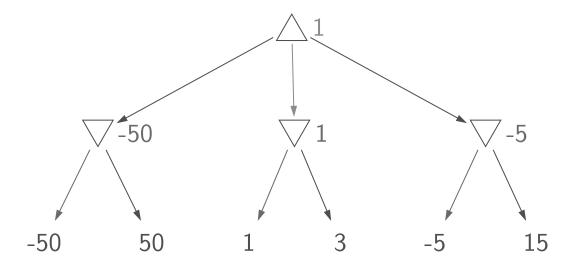
## Summary



- Game trees: model opponents, randomness
- Minimax: find optimal policy against an adversary
- Evaluation functions: domain-specific, approximate
- Alpha-beta pruning: domain-general, exact

### Review: minimax

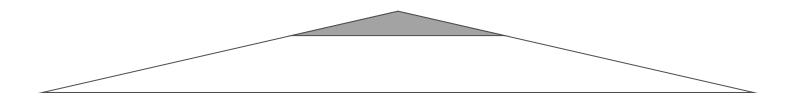
agent (max) versus opponent (min)



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### Review: depth-limited search



$$V_{\mathsf{minmax}}(s,d) = \begin{cases} \mathsf{Utility}(s) & \mathsf{IsEnd}(s) \\ \mathsf{Eval}(s) & d = 0 \\ \max_{a \in \mathsf{Actions}(s)} V_{\mathsf{minmax}}(\mathsf{Succ}(s,a),d) & \mathsf{Player}(s) = \mathsf{agent} \\ \min_{a \in \mathsf{Actions}(s)} V_{\mathsf{minmax}}(\mathsf{Succ}(s,a),d-1) & \mathsf{Player}(s) = \mathsf{opp} \end{cases}$$

Use: at state s, choose action resulting in  $V_{\rm minmax}(s,d_{\rm max})$ 



# Summary

Main challenge: not just one objective

• Minimax principle: guard against adversary in turn-based games

• Simultaneous non-zero-sum games: mixed strategies, Nash equilibria

• Strategy: search game tree + learned evaluation function



#### Chess

1997: IBM's Deep Blue defeated world champion Gary Kasparov

#### Fast computers:

- Alpha-beta search over 30 billion positions, depth 14
- Singular extensions up to depth 20

#### Domain knowledge:

- Evaluation function: 8000 features
- 4000 "opening book" moves, all endgames with 5 pieces
- 700,000 grandmaster games
- Null move heuristic: opponent gets to move twice



#### Checkers

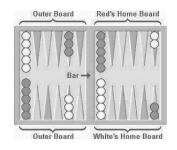
1990: Jonathan Schaeffer's Chinook defeated human champion; ran on standard PC

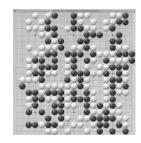
#### Closure:

- 2007: Checkers solved in the minimax sense (outcome is draw), but doesn't mean you can't win
- Alpha-beta search + 39 trillion endgame positions

# Backgammon and Go

Alpha-beta search isn't enough...





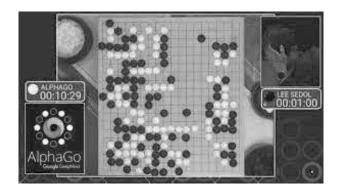
Challenge: large branching factor

• Backgammon: randomness from dice (can't prune!)

• Go: large board size (361 positions)

Solution: learning

### AlphaGo



- Supervised learning: on human games
- Reinforcement learning: on self-play games
- Evaluation function: convolutional neural network (value network)
- Policy: convolutional neural network (policy network)
- Monte Carlo Tree Search: search / lookahead

# Coordination games

Hanabi: players need to signal to each other and coordinate in a decentralized fashion to collaboratively win.

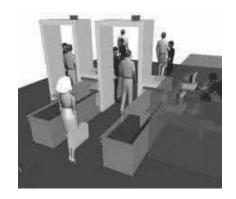


Hide-and-Seek: OpenAI has developed agents with emergent behaviors to play hide and seek.



# Other games

Security games: allocate limited resources to protect a valuable target. Used by TSA security, Coast Guard, protect wildlife against poachers, etc.

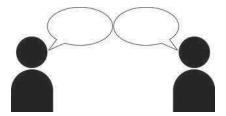


# Other games

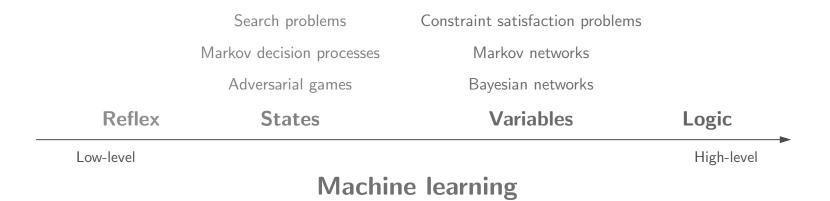
Resource allocation: users share a resource (e.g., network bandwidth); selfish interests leads to volunteer's dilemma



Language: people have speaking and listening strategies, mostly collaborative, applied to dialog systems



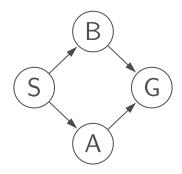
# Course plan



### State-based models

[Modeling]		
Framework	search problems	MDPs/games
Objective	minimum cost paths	maximum value policies
[Inference]		
Tree-based	backtracking	minimax/expectimax
<b>Graph-based</b>	DP, UCS, A*	value/policy iteration
[Learning]		
Methods	structured perceptron	Q-learning, TD learning

# State-based models: takeaway 1



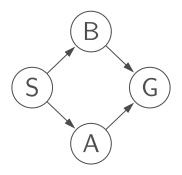


Key idea: specify locally, optimize globally

Modeling: specifies local interactions

Inference: find globally optimal solutions

# State-based models: takeaway 2





Key idea: state-

A **state** is a summary of all the past actions sufficient to choose future actions **optimally**.

Mindset: move through states (nodes) via actions (edges)

# Homework

due: next week

作业 4(optional)周5-图像识别 作业 5(optional)-周5-语音识别 作业 6(optional)-周5-文本识别