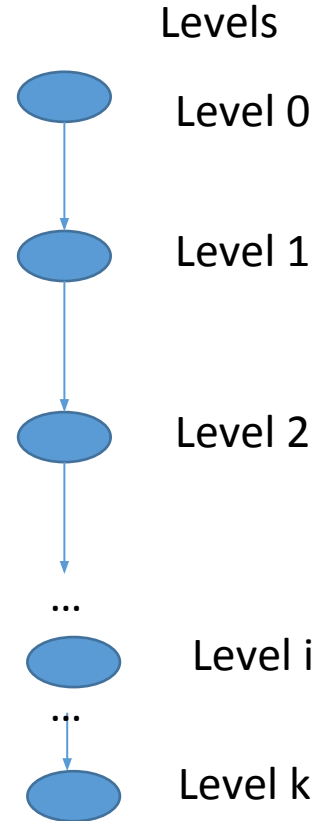


# Solving recurrence

Tree method

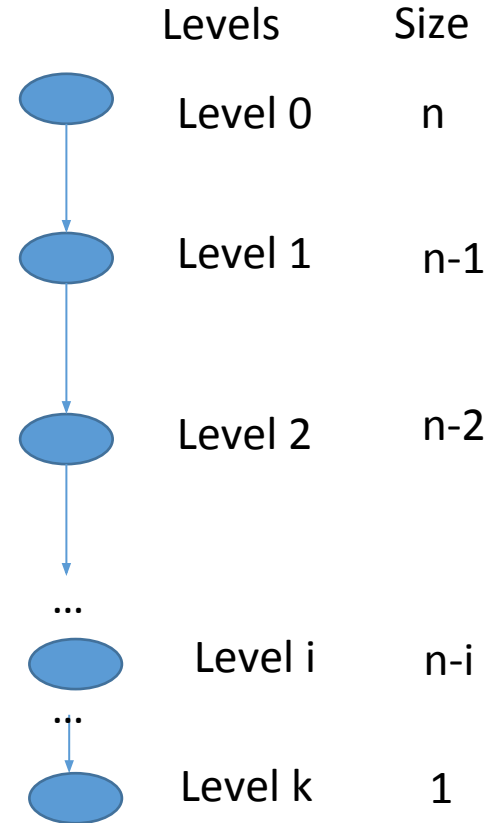
$$\begin{aligned} 1. \quad T(n) &= T(n-1) + d, \quad n > 1 \\ &= c, \quad n = 1 \end{aligned}$$

1.  $T(n) = T(n-1) + d, n > 1$   
 $= c, n = 1$



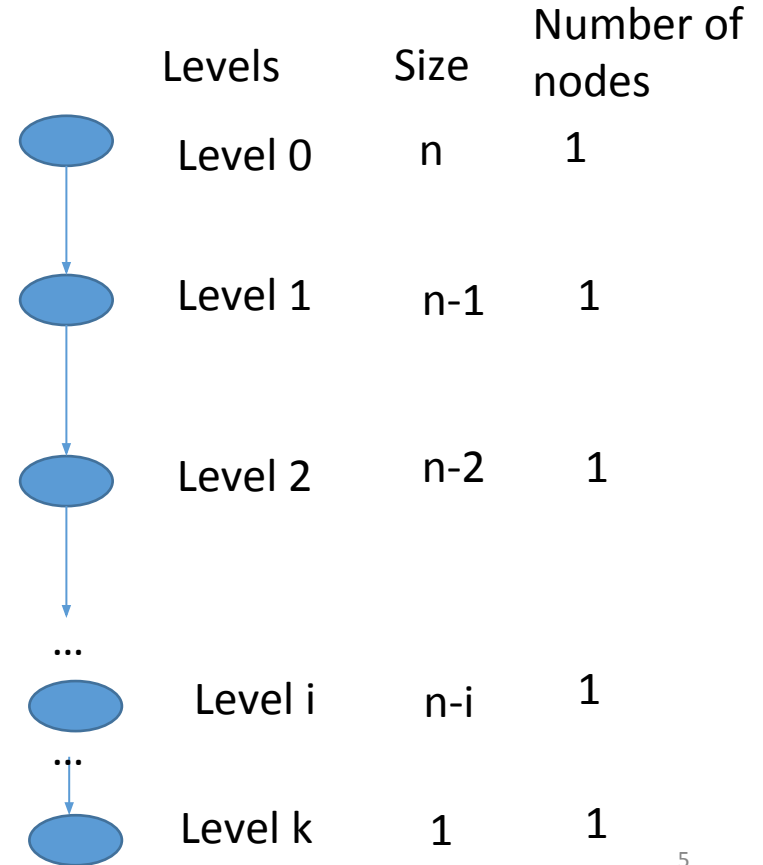
$$1. \quad T(n) = T(n-1) + d, \quad n > 1$$

$$= c, \quad n = 1$$



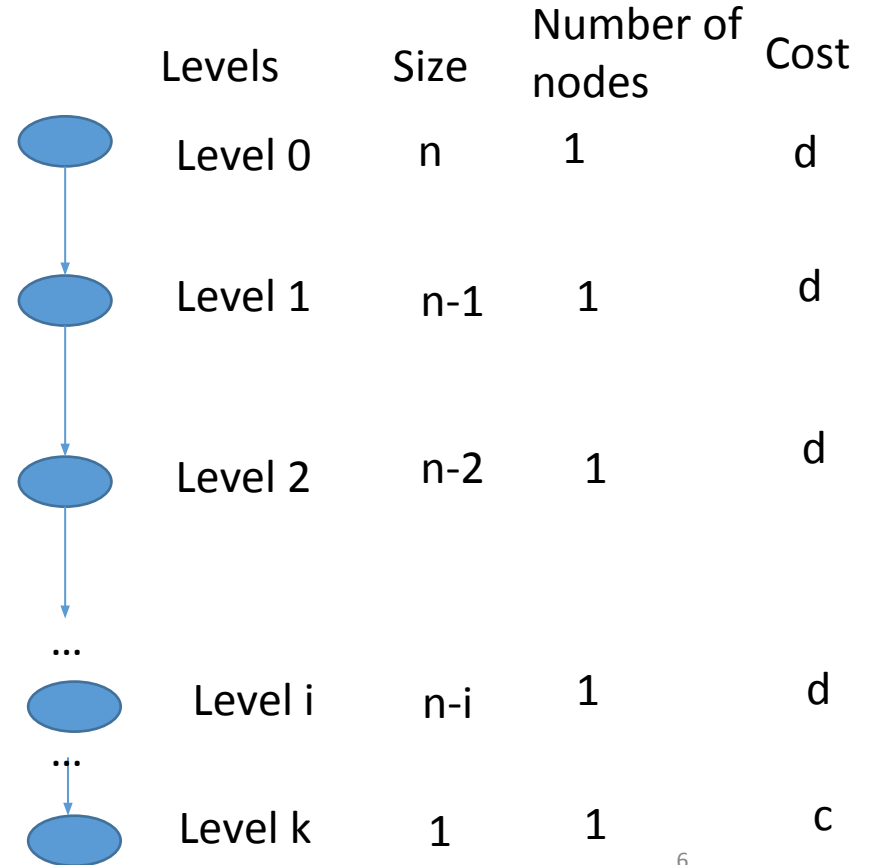
$$1. \quad T(n) = T(n-1) + d, \quad n > 1$$

$$= c, \quad n = 1$$



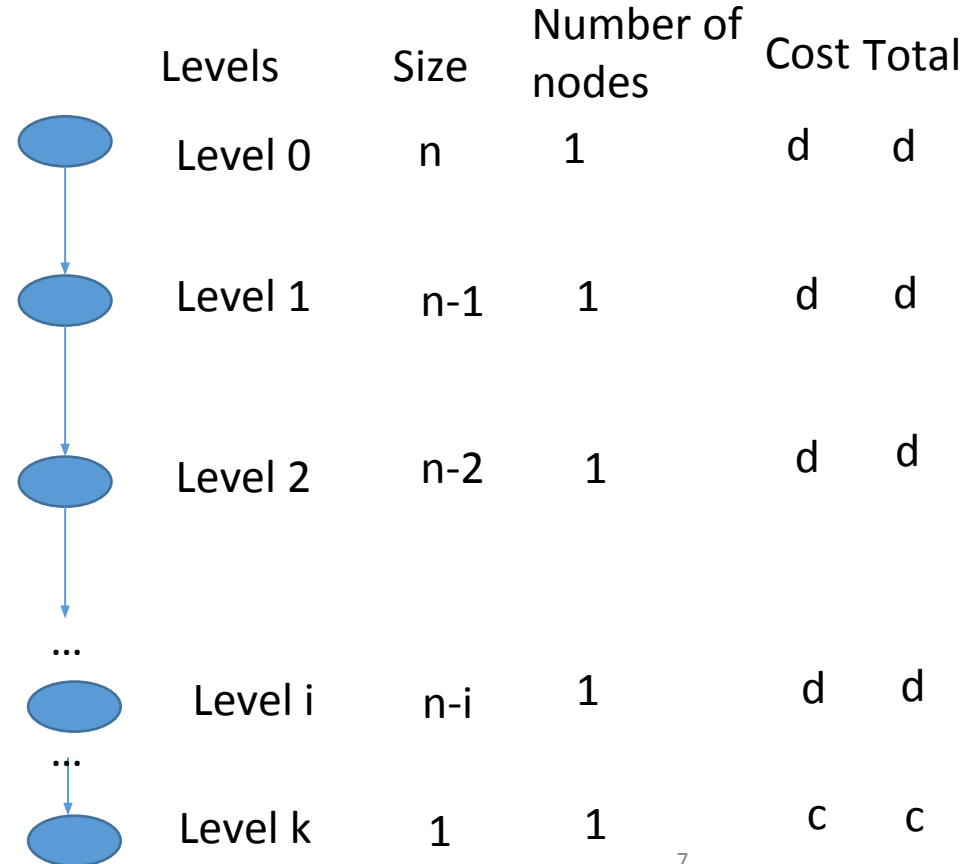
$$1. \quad T(n) = T(n-1) + d, \quad n > 1$$






$$= c, \quad n = 1$$








$$1. \quad T(n) = T(n-1) + d, \quad n > 1$$

$$= c, \quad n = 1$$



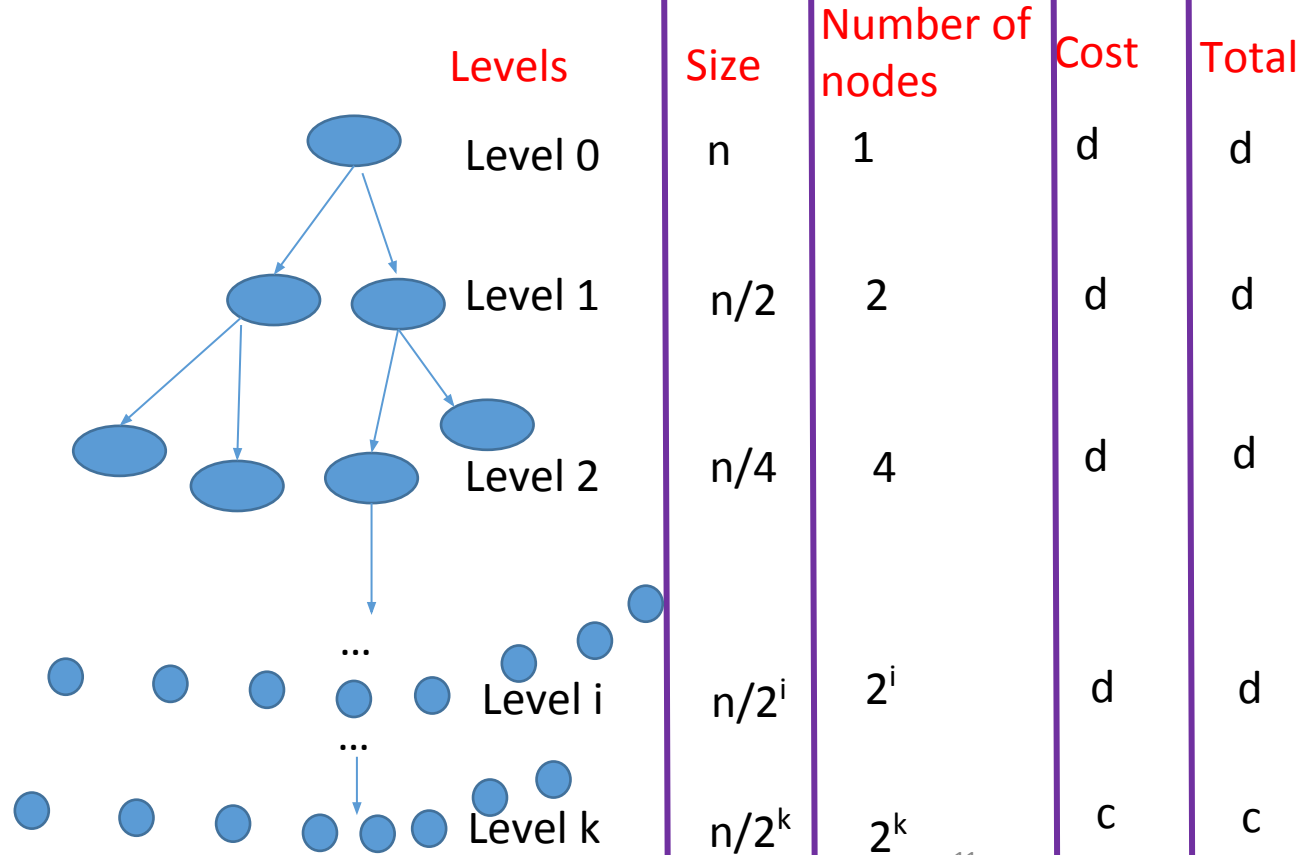
	Levels	Size	Number of nodes	Cost	Total
	Level 0	$n$	1	$d$	$d$
↓					
	Level 1	$n-1$	1	$d$	$d$
↓					
	Level 2	$n-2$	1	$d$	$d$
↓					
...					
	Level $i$	$n-i$	1	$d$	$d$
↓					
...					
	Level $k$	1	1	$c$	$c$

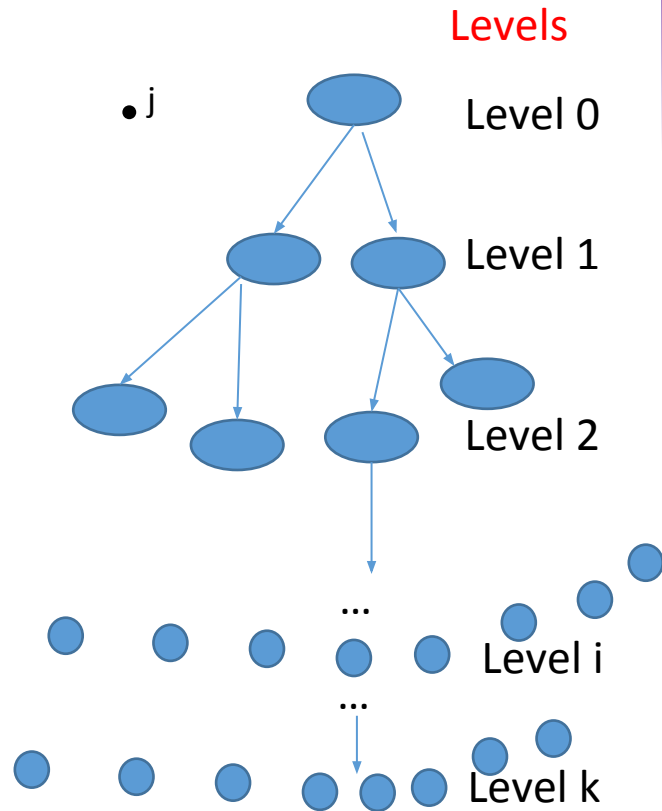


	Levels	Size	Number of nodes	Cost	Total
	Level 0	n	1	d	d
↓					
	Level 1	n-1	1	d	d
↓					
	Level 2	n-2	1	d	d
↓					
...					
	Level i	n-i	1	d	d
↓					
...					
	Level k	1	1	c	c

$$\begin{aligned}
 c + d(k) &= c + (n-1)d \\
 &= dn + c - d \\
 &= O(n)
 \end{aligned}$$

$$\begin{aligned} 2. \quad T(n) &= 2T(n/2) + d, \quad n > 1 \\ &= c, \quad n = 1 \end{aligned}$$

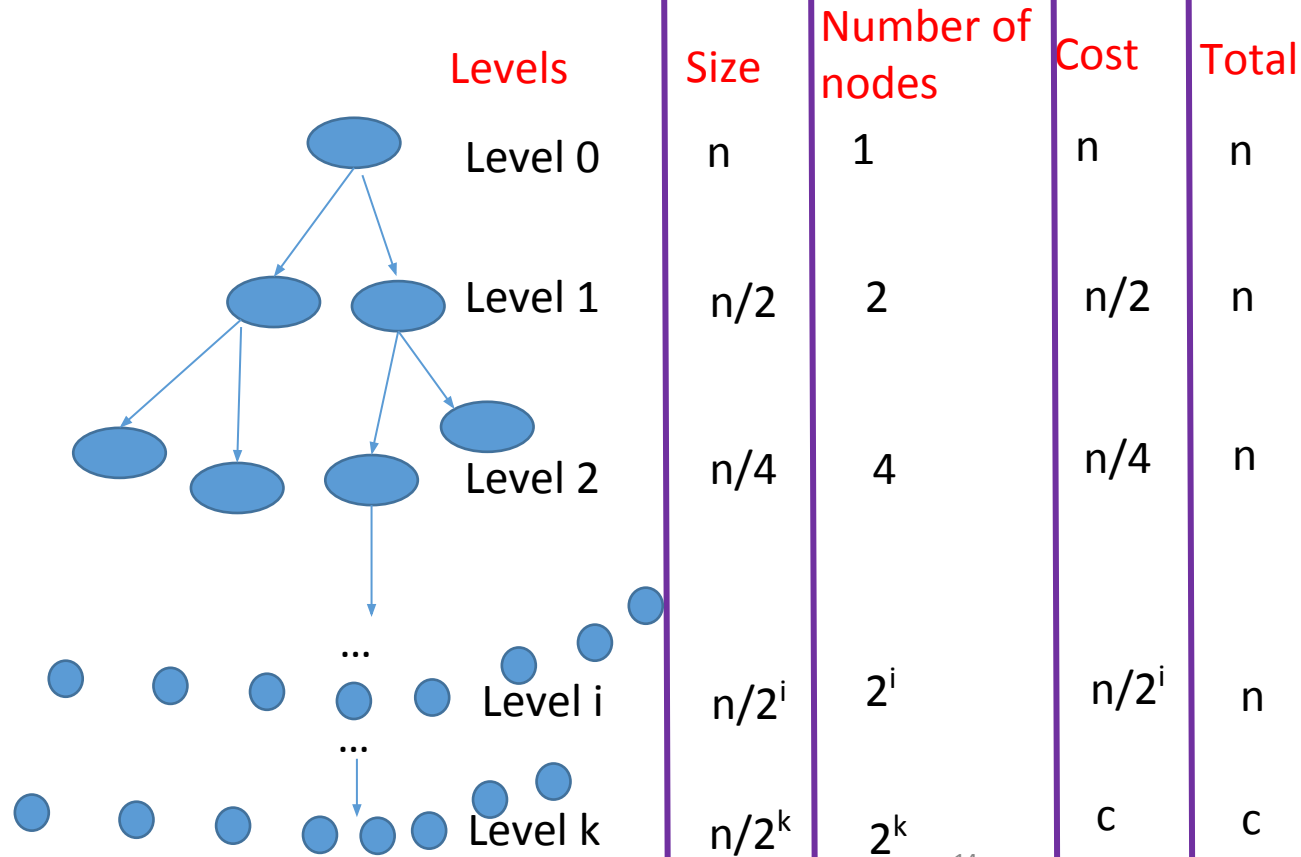


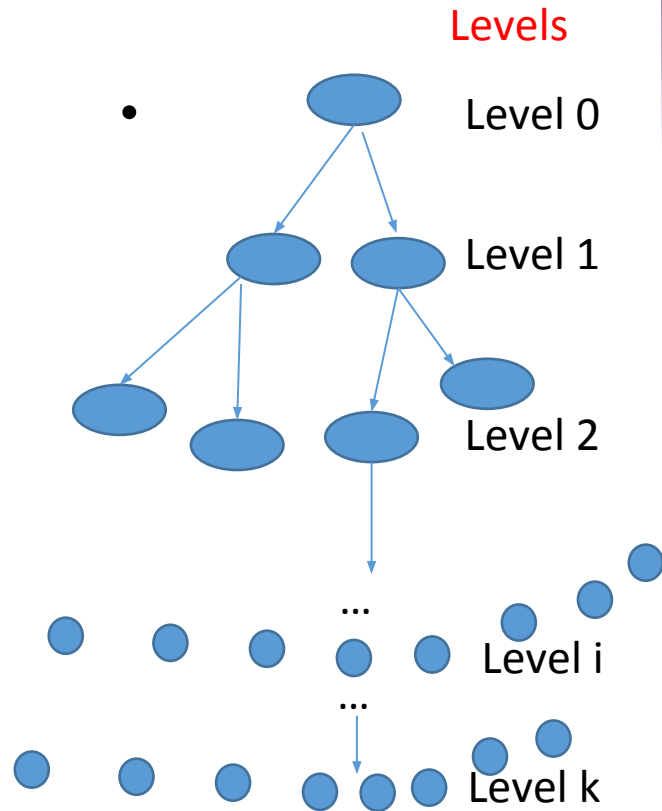


Size	Number of nodes	Cost	Total
n	1	d	d
n/2	2	d	d
n/4	4	d	d
...	...	...	...
n/2 <sup>i</sup>	2 <sup>i</sup>	d	d
...	...	...	...
n/2 <sup>k</sup>	2 <sup>k</sup>	c	c

$$\begin{aligned}
 & d + 2d + 4d + 8d + \dots + 2^{k-1}d + 2^k c \\
 &= d \sum_{j=0}^{k-1} 2^j + 2^k c \\
 &= c2^k + d(2^k - 1) \\
 &= 2^k(c + d) - d \\
 &= 2^{\log n}(c + d) - d \\
 &= n(c + d) - d \\
 &\leq n(c + d) = n * \text{Constant} = O(n)
 \end{aligned}$$

$$\begin{aligned} 3. \quad T(n) &= 2T(n/2) + n, \quad n > 1 \\ &= c, \quad n = 1 \end{aligned}$$





Levels

Level 0

Level 1

Level 2

Level i

Level k

Size

$n$

$n/2$

$n/4$

$n/2^i$

$n/2^k$

Number of  
nodes

1

2

4

$2^i$

$2^k$

Cost

$n$

$n/2$

$n/4$

$n/2^i$

$c$

Total

$n$

$n$


$n$


$n$

$c$

$$\begin{aligned}
 & n + n + n + n + \dots + n + c2^k \\
 &= \sum_{j=0}^{k-1} n + c2^k \\
 &= c2^k + n(k) \\
 &= 2^{\log n} c + n \log n \\
 &= nc + n \log n = O(n \log n)
 \end{aligned}$$

# Home Work

4.  $T(n) = T(n/3) + 5n, n > 1$   
 $= 6, n = 1$    $O(n)$

5.  $T(n) = 4T(n/3) + 6n^2, n > 1$   
 $= 11, n = 1$    $O(n^2)$