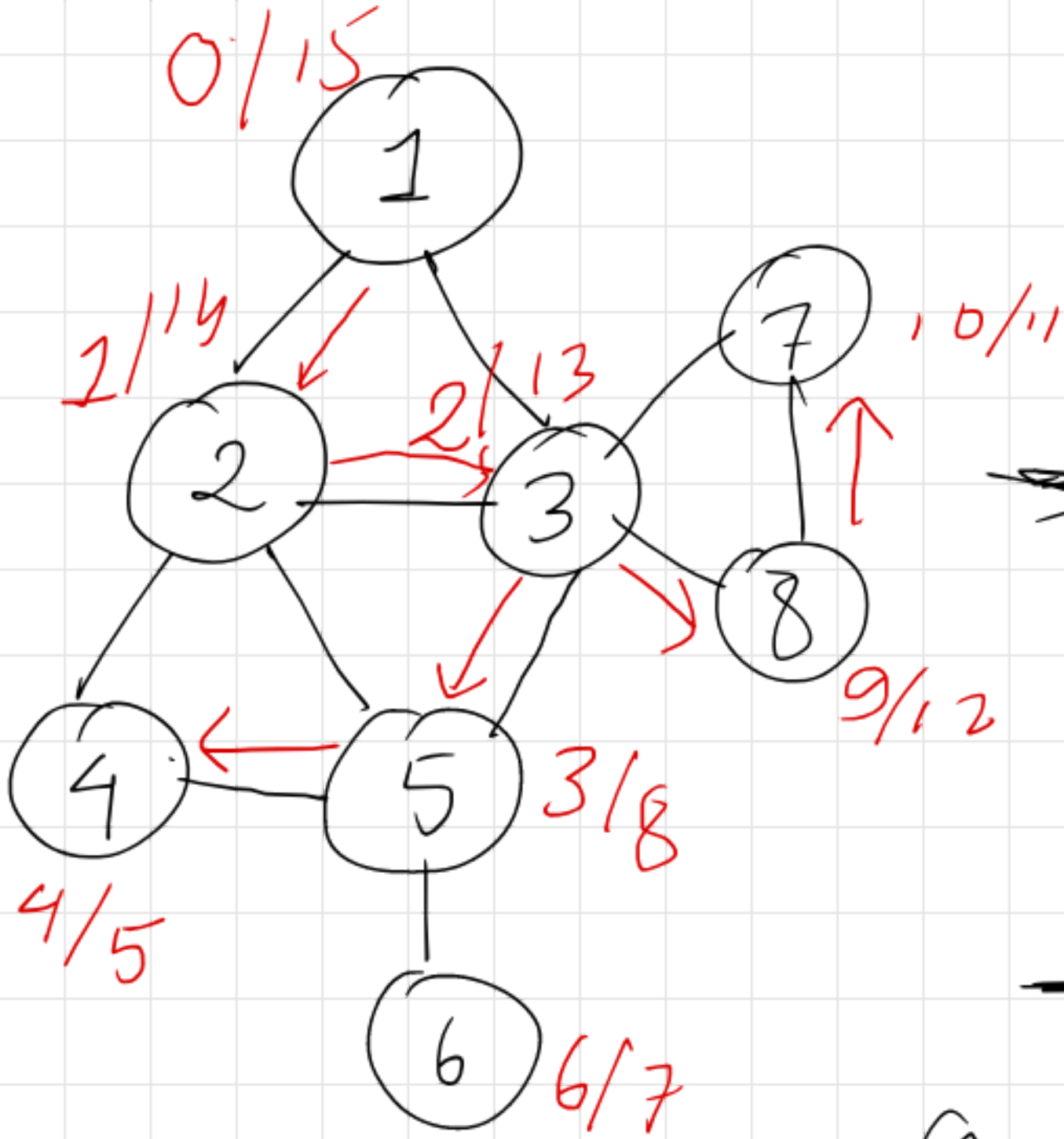


Depth First Search

Idea:

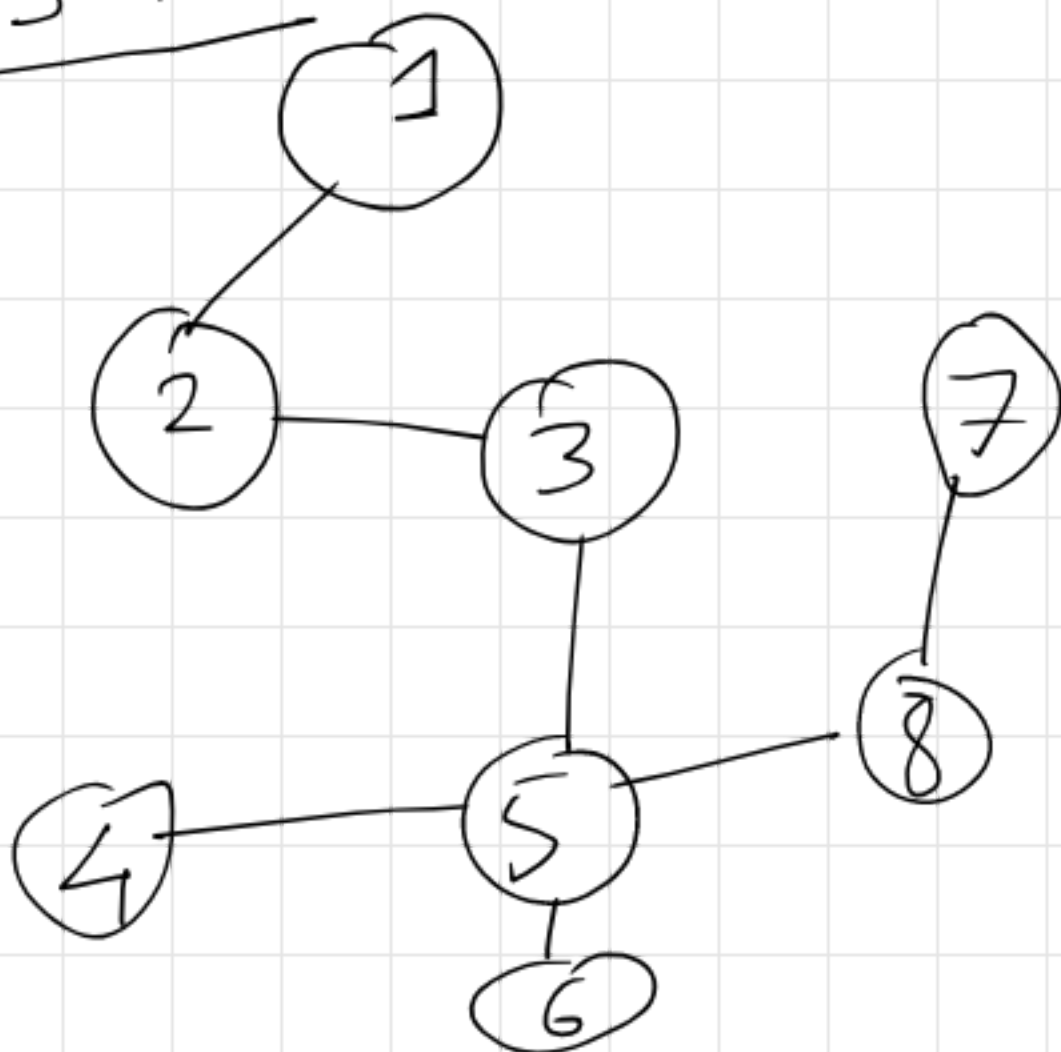


Timestamp
discovery
processing done

→ explore a node only if it is unexplored

→ backtrack from a node along the same edge that was taken to reach the node.

DFS Tree



$\frac{n-1 \text{ edges}}{\text{Connected acyclic}} \}$ Tree

0/15 (1)

1/14 (2)

2/13 (3)

3/8 (5)

4/5 (4)

9/12 (8)

10/11 (7)

6/7 (6)

Observation

1. discovery (ancestor)

< discovery (descendant)

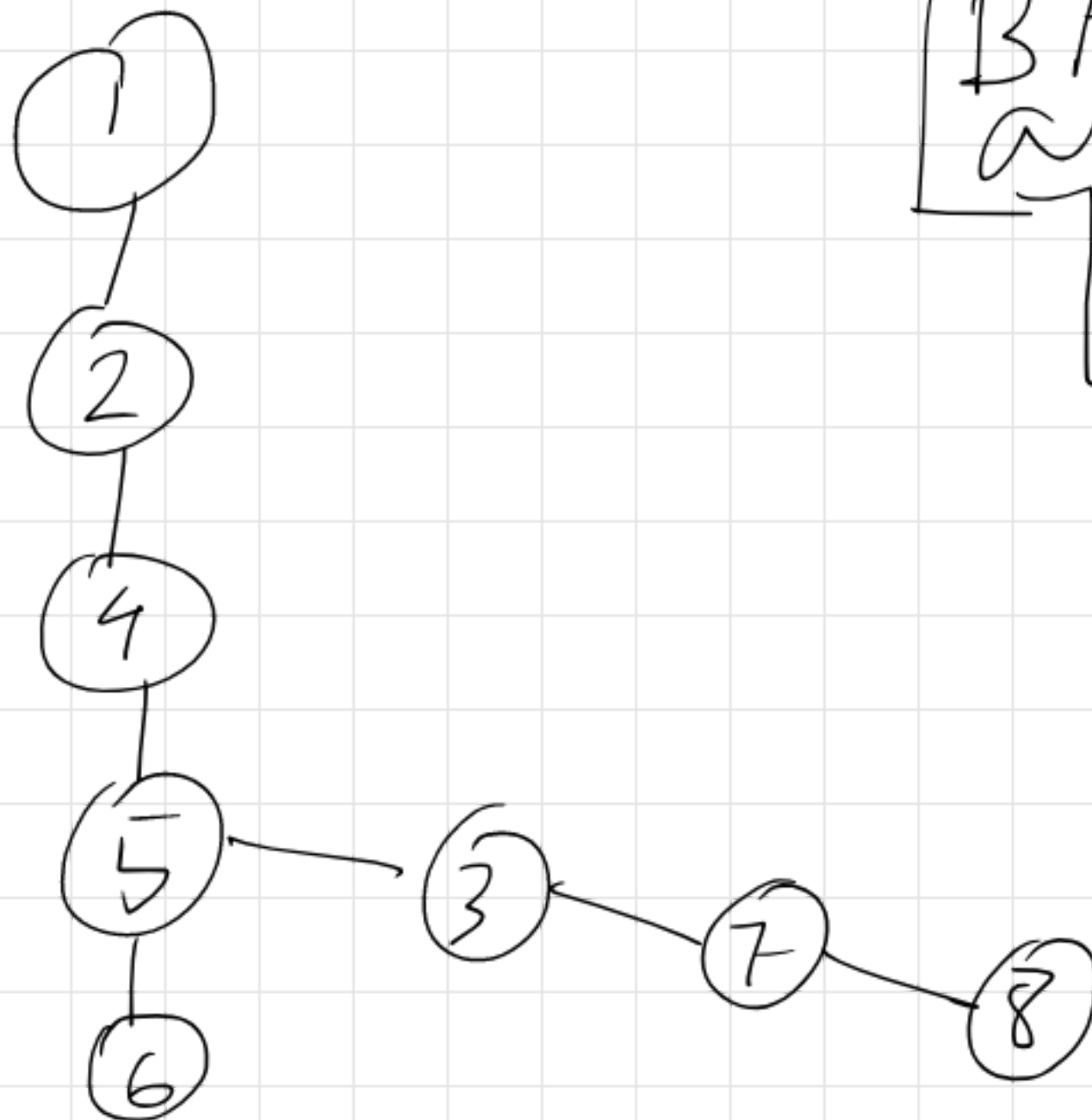
< finish (descendant)

< finish (ancestor)

Eg:

$\{ d(2) < d(6) < f(6) < f(2) \}$
 $\{ d(3) < d(5) < f(5) < f(3) \}$

Is DFS Tree unique?? \rightarrow NO

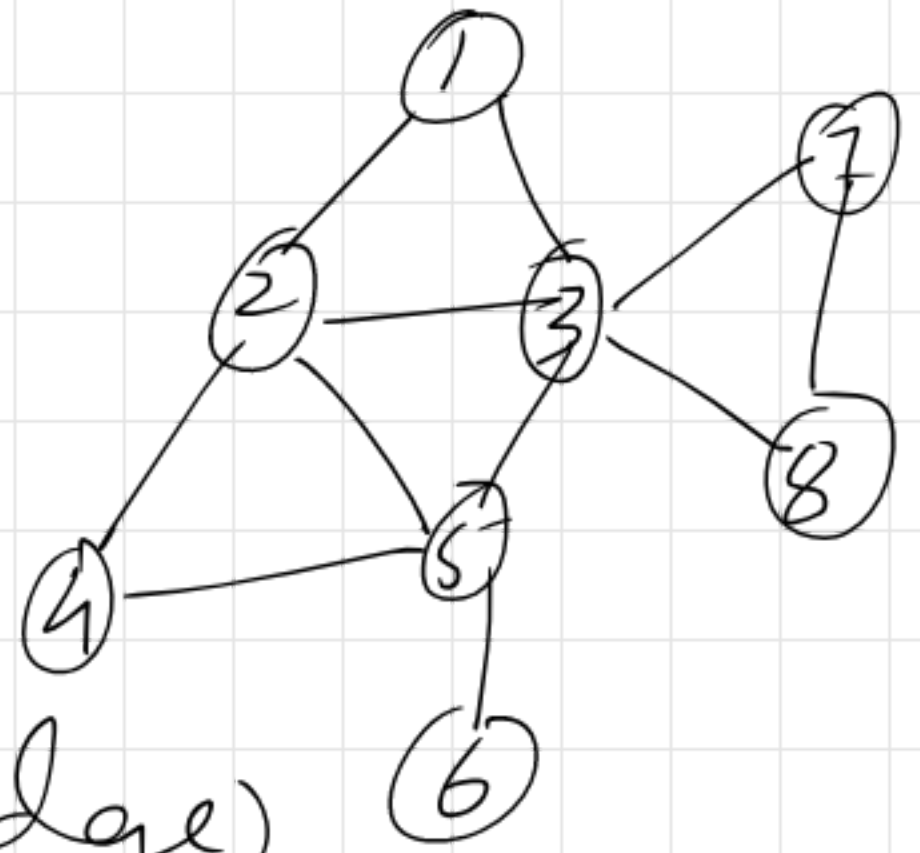
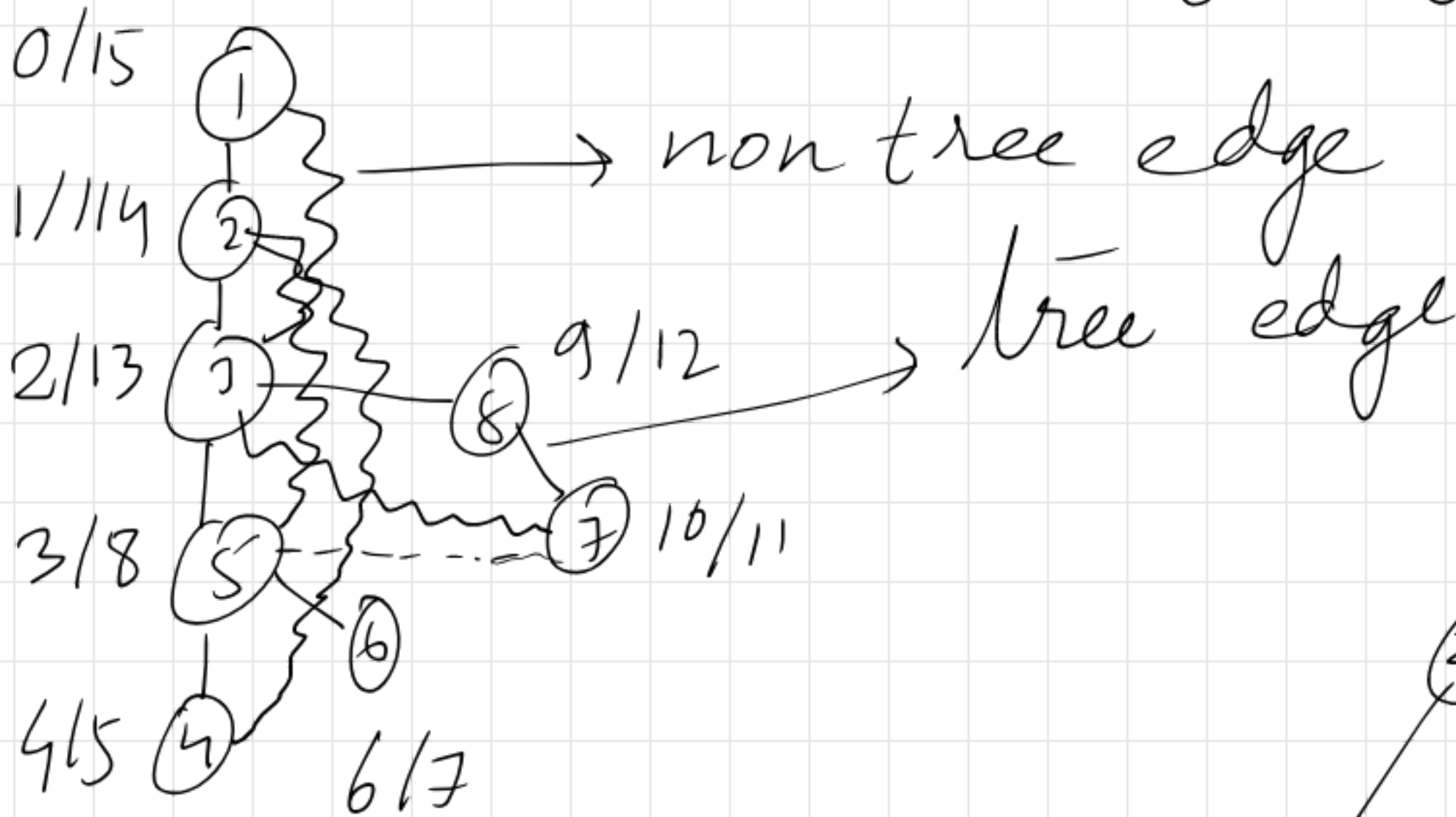


BFS tree is also not unique

DFS trees (narrow and deep)

BFS tree (short and broad)

$$\hookrightarrow (x, y) \in E \quad \begin{array}{l} x \in L_i \\ y \in L_j \end{array} \quad |i - j| \leq 1$$



~~~~~ (non tree edge)  
part of G

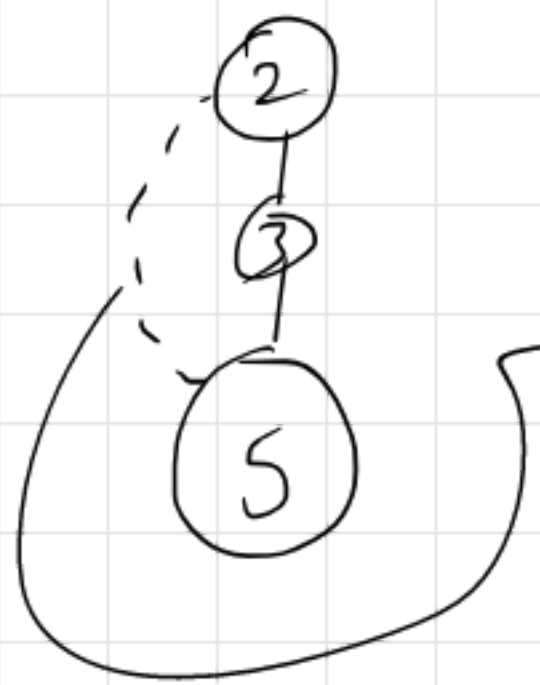
[Can we have  $(5, 7) \in E$  in  $G$ ?]  
Is it possible

$$\begin{array}{l} f(5) = 8 \\ s(7) = 10 \end{array} \quad \Bigg|$$

$(6, 7) \in E$  not possible

Non-tree edges connect only ancestors to descendant

non tree edges are called back edges



you realised this edge when you were exploring node (5) not (2)  
 thus the name  
 back edge

Claim:

Let  $T$  be a DFS tree. Let  $x$  and  $y$  be nodes in  $T$ . Let  $(x, y) \in E$ . Then one of  $x$  and  $y$  is an ancestor of the other.

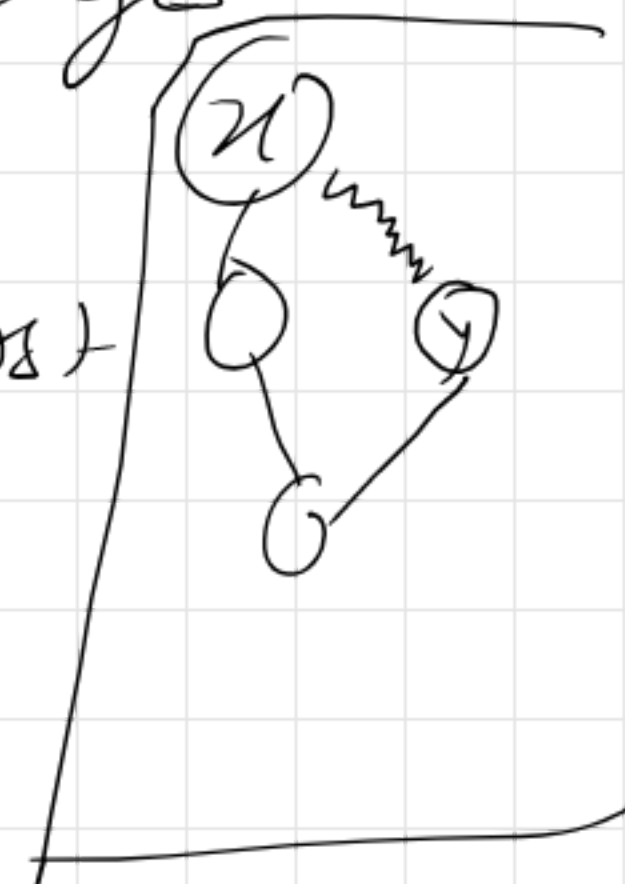
Pf:  $(x, y) \rightarrow$  tree edge (Claim) ✓

let  $(x, y) \rightarrow$  non tree edges

$(x, y) \notin T$

w.l.o.g algo explored  $x$  first

$d(x) < d(y) < f(x)$





# Implementing DFS

## Recursive

given  $G$  as  
adjacency list

$DFS(u)$

set  $explored[u] = \text{True}$

add  $u$  to  $R$ .

For each edge  $(u, v)$

if  $explored[v] = \text{false}$

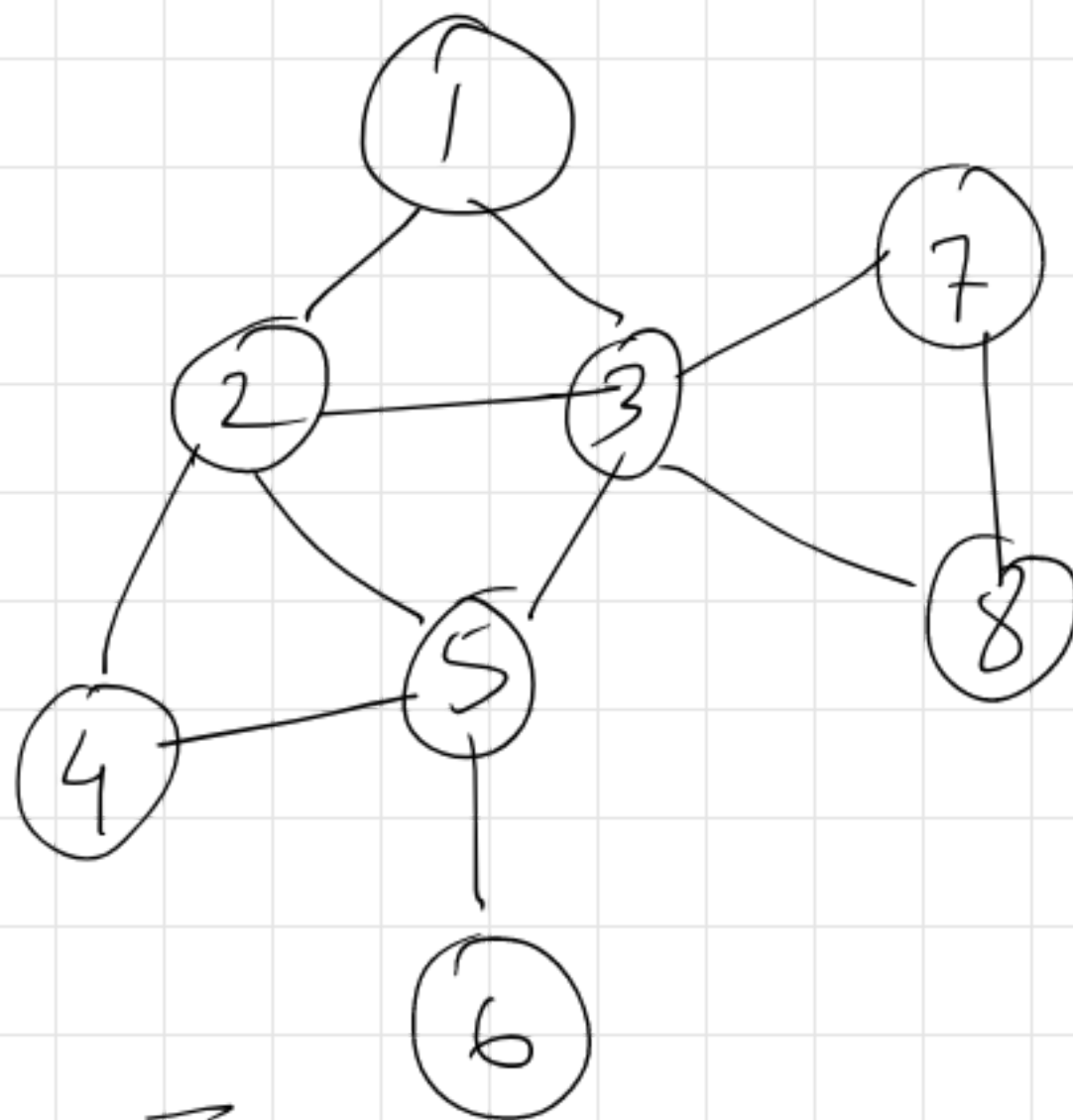
$DFS(v)$

endif

end for

// Check in Adj  
List( $u$ )

$DFS(1)$   
 $DFS(2)$   
 $DFS(3)$   
⋮



↑

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
| T | F | F | F | F | F | F | F |
|---|---|---|---|---|---|---|---|

# Non Recursive

DFS(s)

Initialized  $explored[v] = \text{false} \quad \forall v \in V$

Initialized  $S$  to be a stack

Push(s, S)

while (S  $\neq$  0)

$u \leftarrow \text{POP}(S) \rightarrow O(1)$

if  $explored[u] = \text{false}$

$explored[u] = \text{true}$

for each edge  $(u, v)$  incident to  $u$

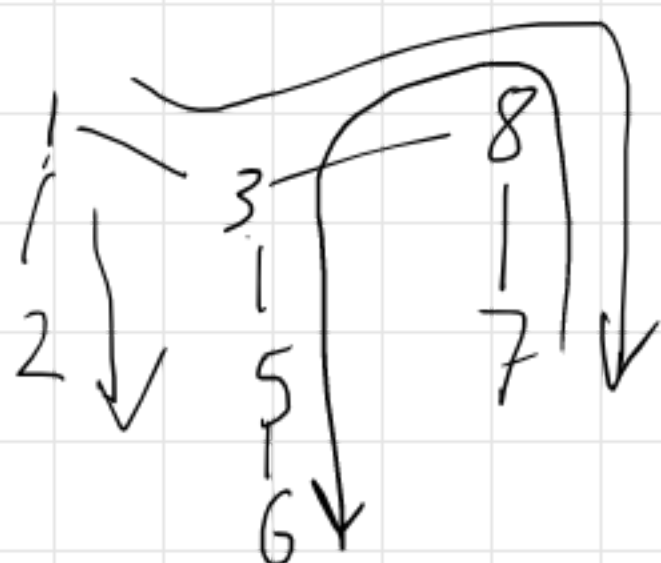
Push(s, v)  $\rightarrow O(1)$

end for

end if

end while

Want to create a DFS tree  
Push parent information  
in Stack



Time complexity  
#push  
 $= \sum \deg(u)$   
 $= O(n)$

in Adj list



