

On Compositional safety verification with Max-SMT

Fabian Böller with David Korzeniewski

RWTH Aachen

fabian.boeller@rwth-aachen.de

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Overview

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- 3 Algorithms
- 4 Example execution
- 5 Conclusion

Safety verification

Prove that an assertion is *always* true at a location

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Non-compositional safety verification

Safety verification where the whole program is analyzed in one step

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Non-compositional safety verification

Safety verification where the whole program is analyzed in one step

Compositional safety verification

Safety verification where program parts are analyzed semi-independently and composed

Scalability \leftrightarrow Loss in precision

Programs

Program

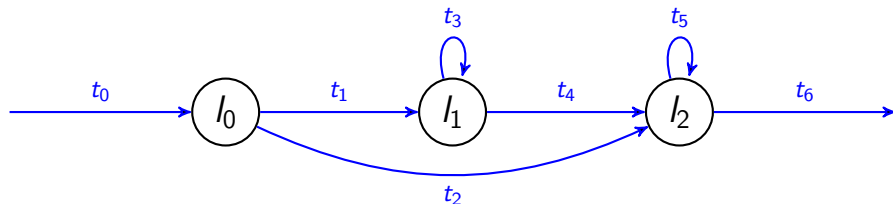
$$\mathcal{L} = \{\ell_0, \ell_1, \ell_2\}$$



Programs

Program

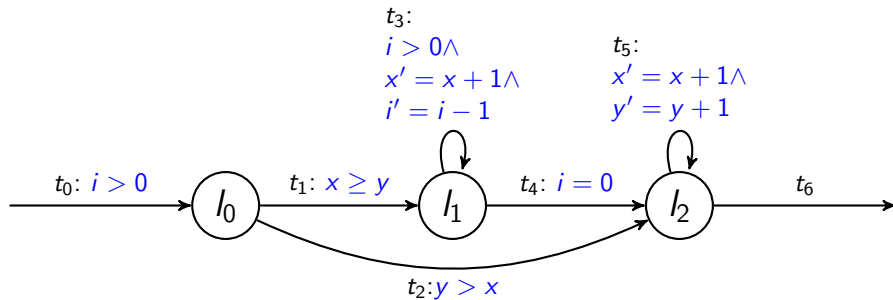
$\mathcal{L} = \{l_0, l_1, l_2\}$, $\mathcal{T} = \{t_i \mid i \in \{1, \dots, 6\}\}$



Programs

Program

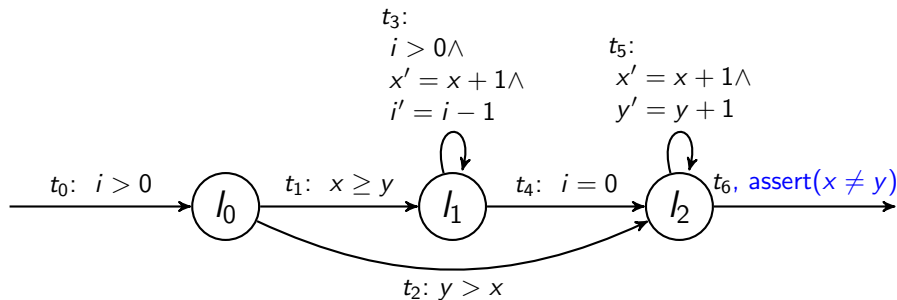
$\mathcal{L} = \{\ell_0, \ell_1, \ell_2\}$, $\mathcal{T} = \{t_i \mid i \in \{1, \dots, 6\}\}$, $\mathcal{V} = \{x, y, i\}$, $\mathcal{V}' = \{x', y', i'\}$



Programs

Program

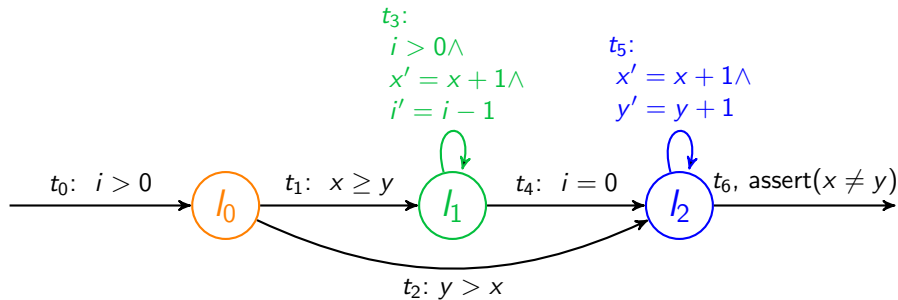
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Programs

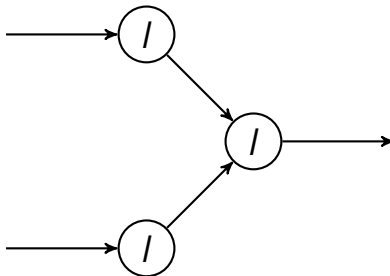
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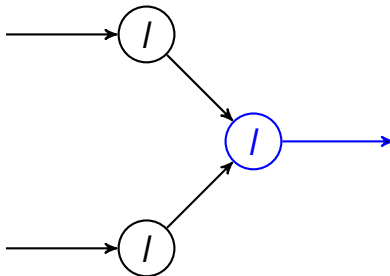
Idea

Prove that an assertion is satisfied by recursively checking all entry components



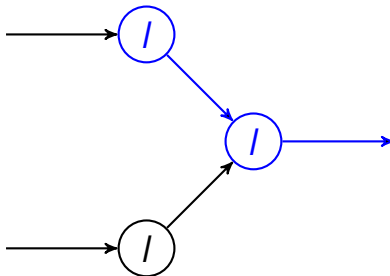
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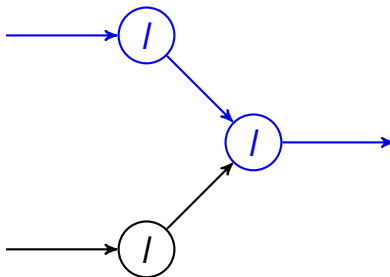
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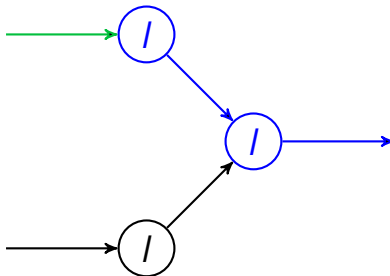
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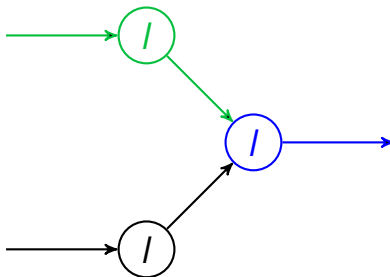
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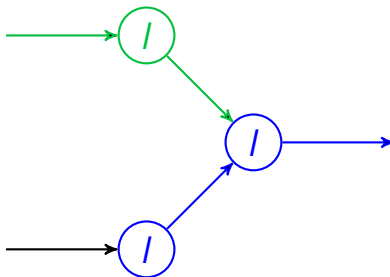
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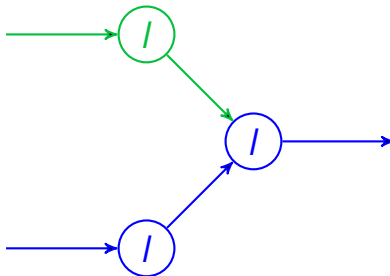
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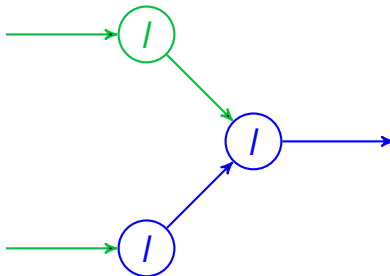
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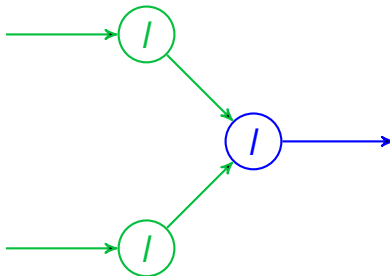
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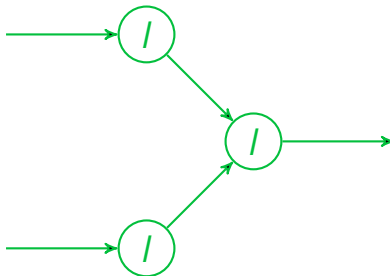
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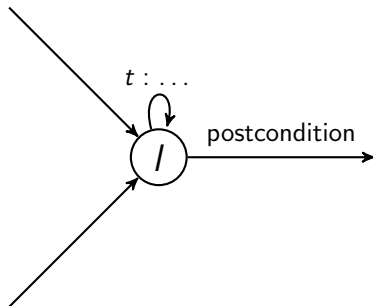
Idea

Prove that an assertion is satisfied by recursively checking all entry components



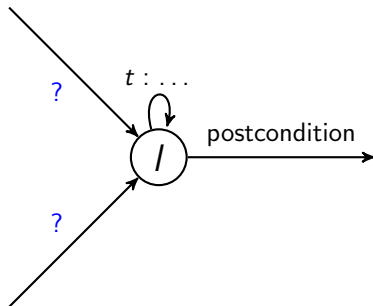
Idea

Find a precondition for a component such that all runs satisfying the precondition always imply the postcondition



Idea

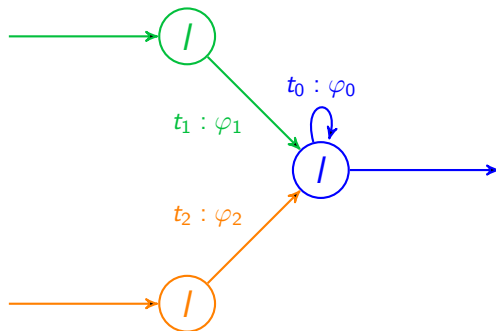
Find a precondition for a component such that all runs satisfying the precondition always imply the postcondition



Narrowing

Idea

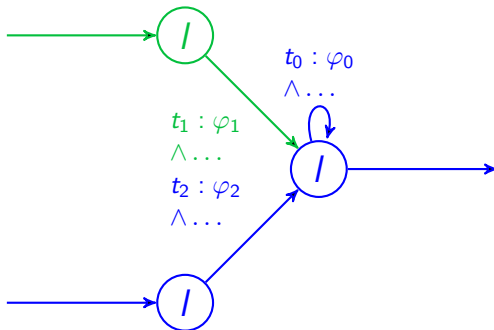
Manipulate the program such that new preconditions can be found



Narrowing

Idea

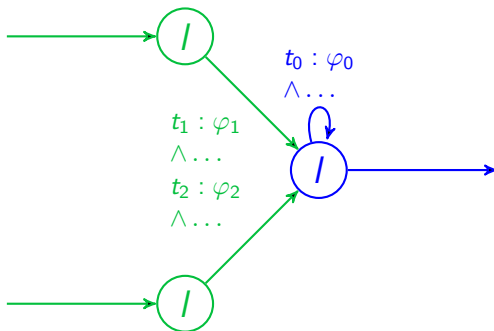
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Narrowing

Idea

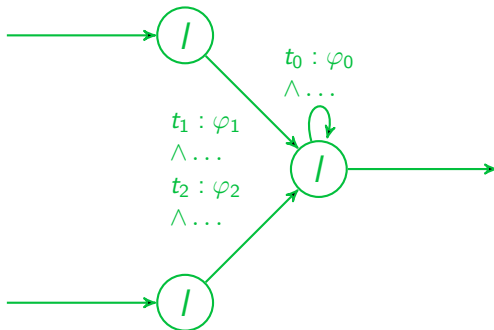
Manipulate the program such that new preconditions can be found



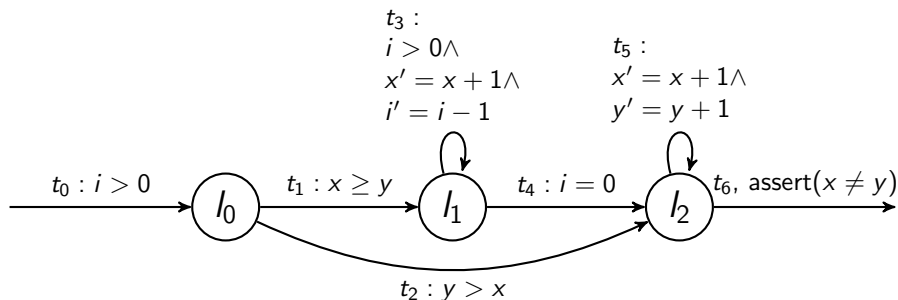
Narrowing

Idea

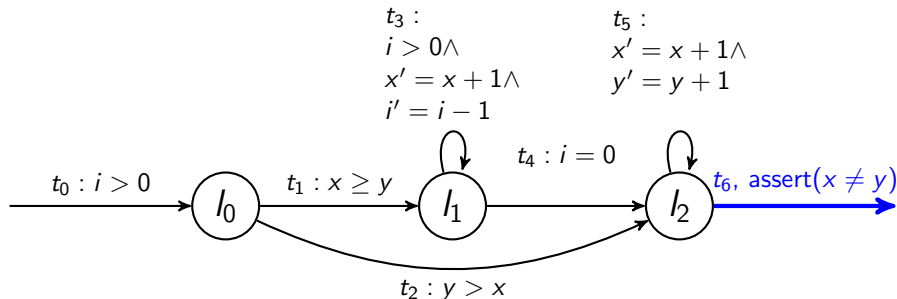
Manipulate the program such that new preconditions can be found



Example program



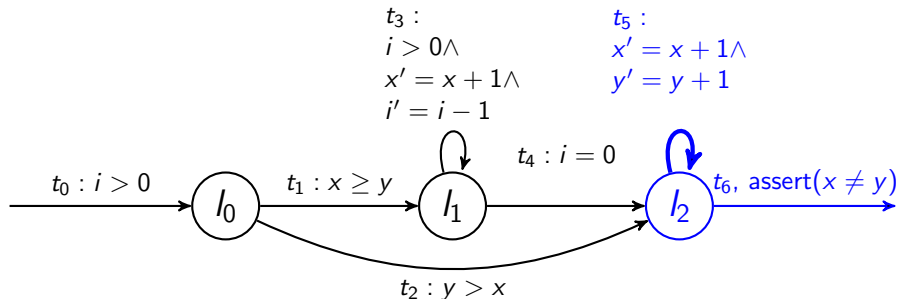
Example execution



Task

Prove that the program is safe for $x \neq y$ at t_6

Example execution



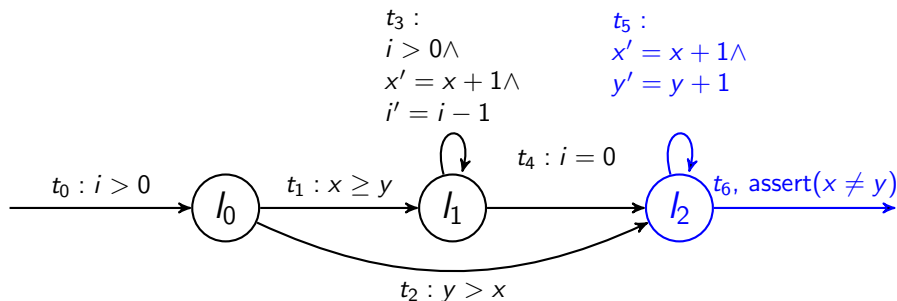
CheckSafe on $\{l_2\}$ for $x \neq y$

t_6 does not already imply $x \neq y$

t_6 is not an initial transition

Call CondSafe

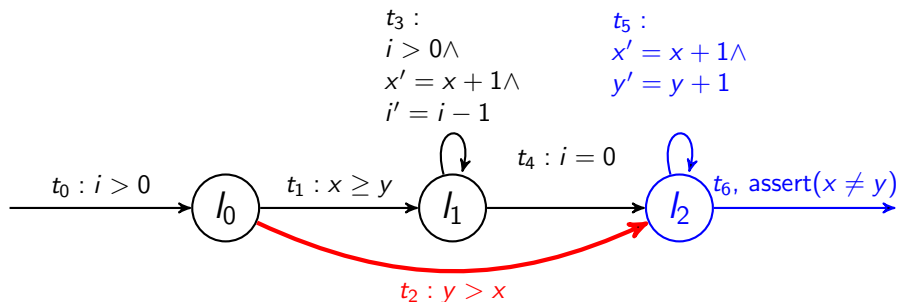
Example execution



Max-SMT on $\{\ell_2\}$ for $x \neq y$

$$I_{\ell_2,1}(\{x, y, i\}) \equiv i_{\ell_2,1} + i_{\ell_2,1,x} * x + i_{\ell_2,1,y} * y + i_{\ell_2,1,i} * i \leq 0$$

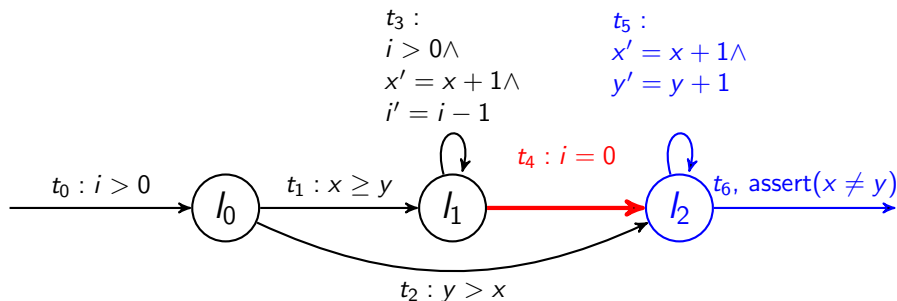
Example execution



Max-SMT on $\{l_2\}$ for $x \neq y$

$$\mathbb{I}_{t_2,1,1} \equiv y > x \wedge i' = i \wedge x' = x \wedge y' = y \Rightarrow l'_{l_2,1,1}$$

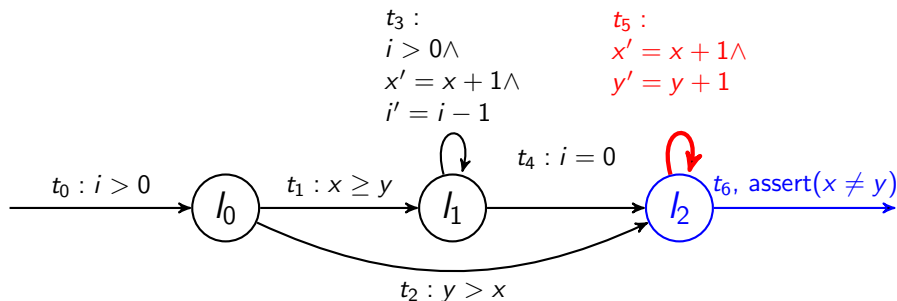
Example execution



Max-SMT on $\{\ell_2\}$ for $x \neq y$

$$\mathbb{I}_{t_4,1,1} \equiv i = 0 \wedge i' = i \wedge x' = x \wedge y' = y \Rightarrow I'_{\ell_2,1,1}$$

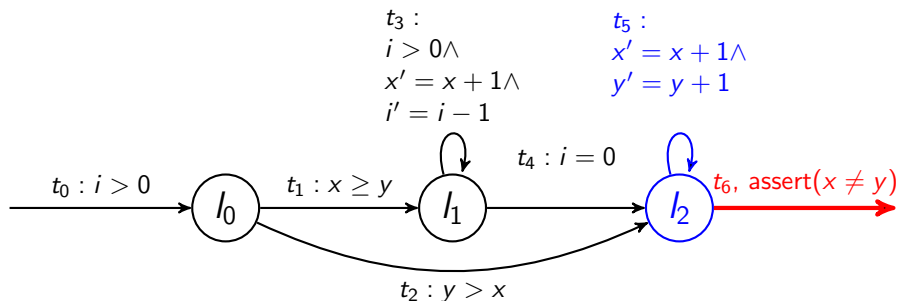
Example execution



Max-SMT on $\{\ell_2\}$ for $x \neq y$

$$\mathbb{C}_{t_5,1} \equiv l_{\ell_2,1} \wedge x' = x + 1 \wedge y' = y + 1 \wedge i' = i \Rightarrow l'_{\ell_2,1}$$

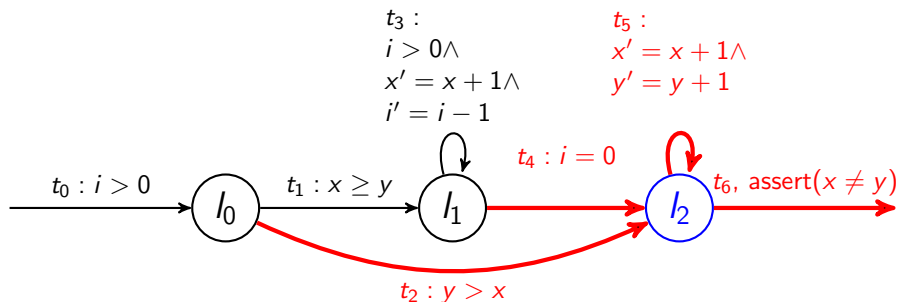
Example execution



Max-SMT on $\{\ell_2\}$ for $x \neq y$

$$S_1 \equiv l_{\ell_2,1} \wedge i' = i \wedge x' = x \wedge y' = y \Rightarrow x' \neq y'$$

Example execution

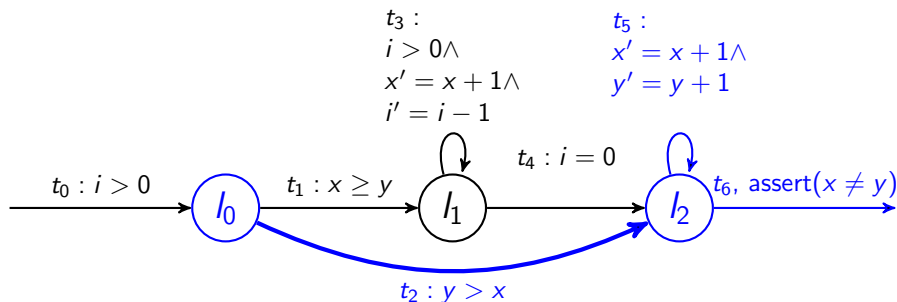


Max-SMT on $\{\ell_2\}$ for $x \neq y$

$\mathbb{F}_1 \equiv \mathbb{C}_{t_5,1} \wedge \mathbb{S}_1 \wedge (\mathbb{I}_{t_2,1,1} \vee \neg p_{t_2}) \wedge (\mathbb{I}_{t_4,1,1} \vee \neg p_{t_4}) \wedge [p_{t_2}, 1] \wedge [p_{t_4}, 1]$

Assume $x > y$ does satisfy \mathbb{F}_1

Example execution



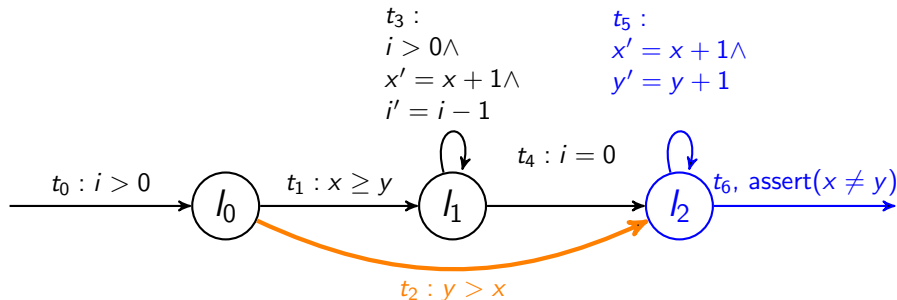
CheckSafe on $\{l_0\}$ for $x > y$

t_2 does not already imply $x > y$

t_2 is not an initial transition

Call CondSafe

Example execution

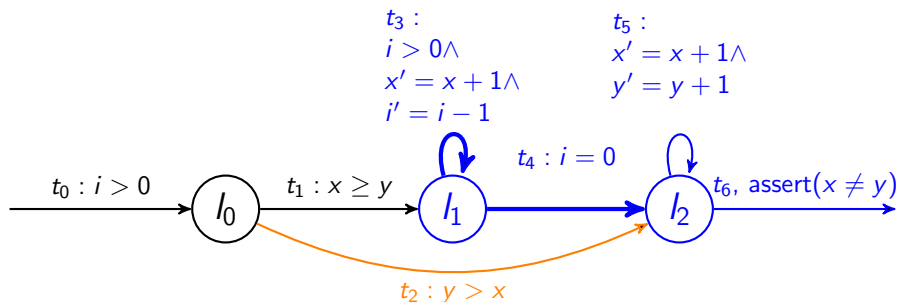


CheckSafe on $\{l_0\}$ for $x > y$

No precondition, since $y > x$ contradicts $x > y$

Path is maybe safe, but not for $x > y$

Example execution



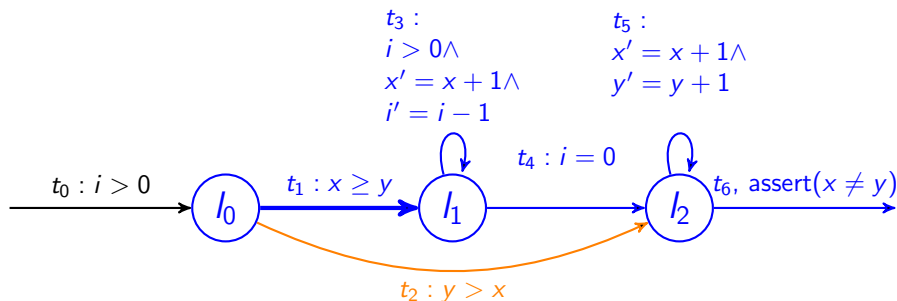
CheckSafe on $\{l_1\}$ for $x > y$

t_4 does not already imply $x > y$

t_4 is not an initial transition

Call CondSafe, get $i > 0 \wedge x \geq y$ as precondition

Example execution



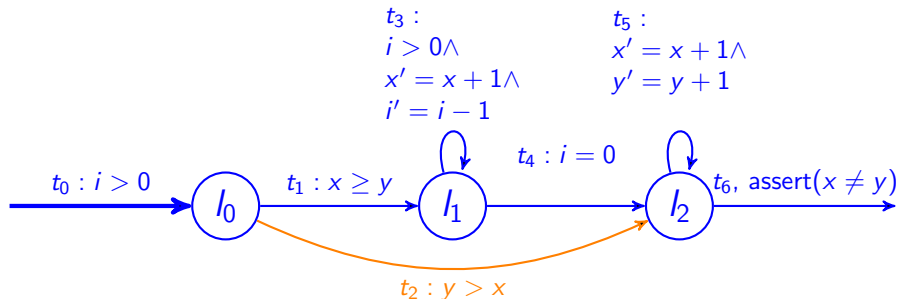
CheckSafe on $\{l_0\}$ for $i > 0$

t_1 does not already imply $i > 0$

t_1 is not an initial transition

Call CondSafe, get $i > 0$ as precondition

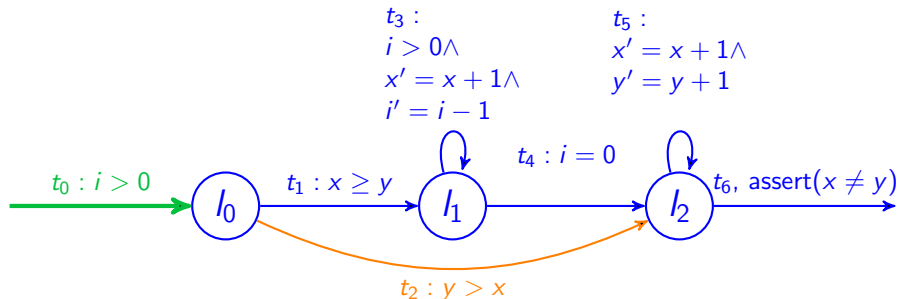
Example execution



CheckSafe on initial SCC for $i > 0$

t_0 does already imply $i > 0$

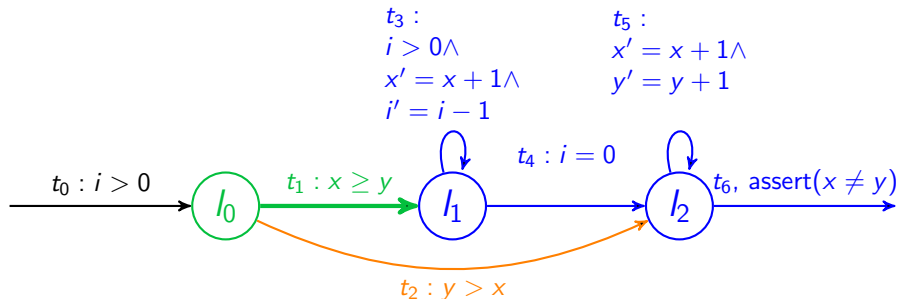
Example execution



CheckSafe on initial SCC for $i > 0$

Path is safe for $i > 0$

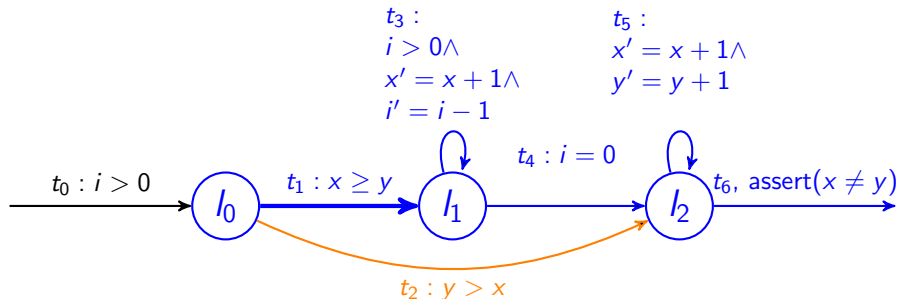
Example execution



CheckSafe on $\{\ell_0\}$ for $i > 0$

Path is safe for $i > 0$

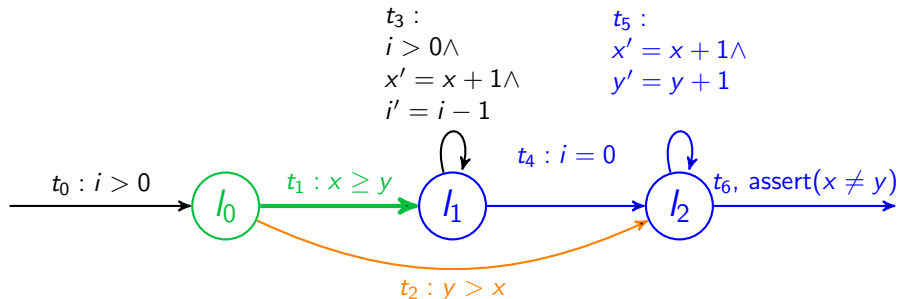
Example execution



CheckSafe on $\{\ell_0\}$ for $x \geq y$

t_1 does already imply $x \geq y$

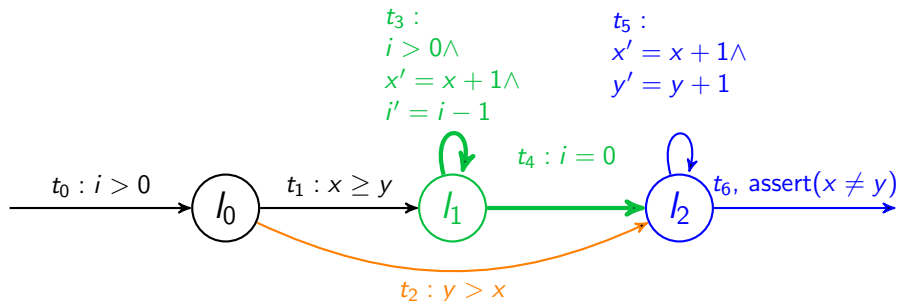
Example execution



CheckSafe on $\{l_0\}$ for $x \geq y$

Path is safe for $x \geq y$

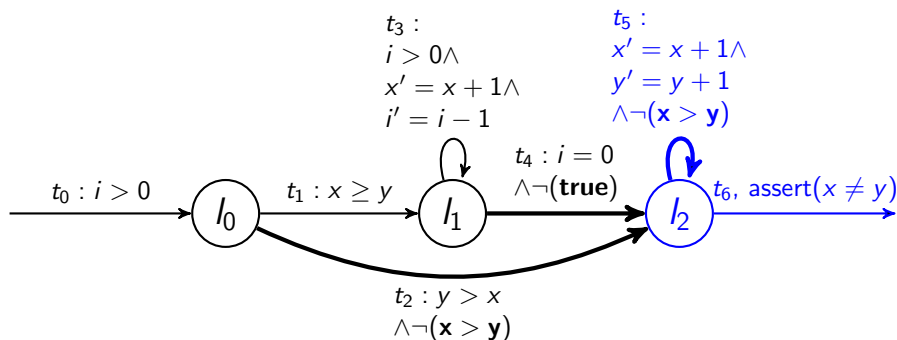
Example execution



CheckSafe on $\{l_1\}$ for $x > y$

Path is safe for $x > y$

Example execution

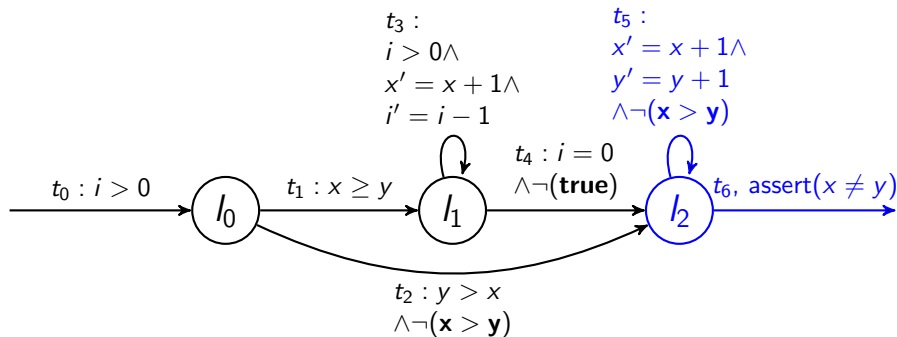


Narrow on $\{l_2\}$

Add $\neg(x > y)$ to t_2

Add $\neg(x > y)$ to t_5

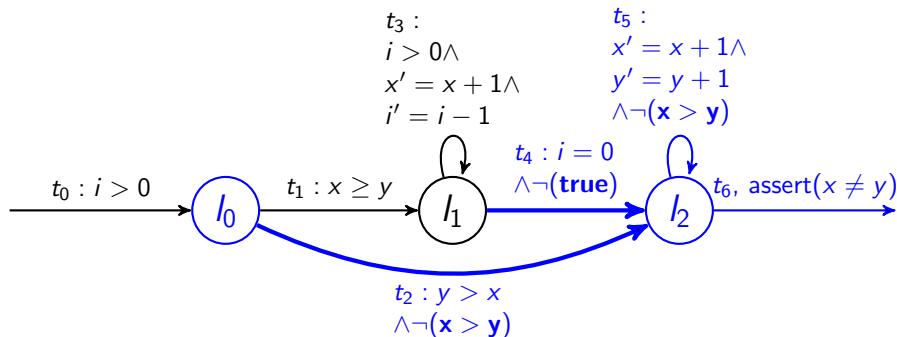
Example execution



CheckSafe on $\{\ell_2\}$ for $x \neq y$

Call CondSafe, get $y > x$ instead of $x > y$ as precondition

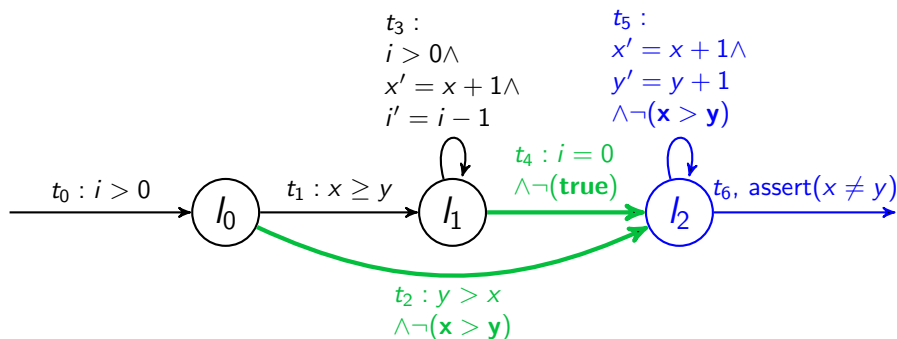
Example execution



CheckSafe on $\{l_0\}$ for $y > x$

t_2 does already imply $y > x$ t_4 does already imply $y > x$

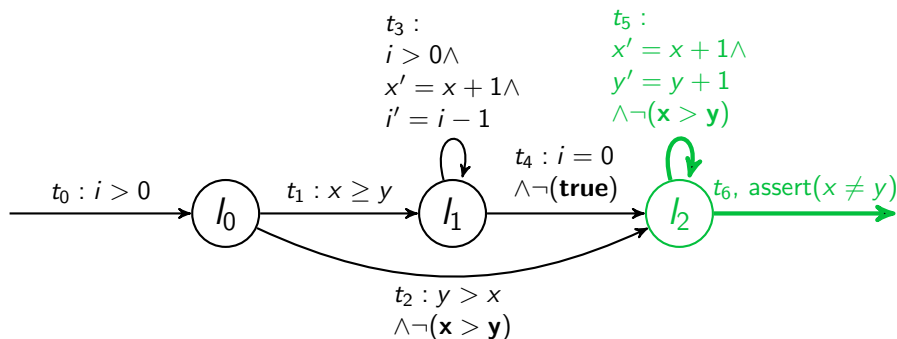
Example execution



CheckSafe on $\{\ell_0\}$ for $y > x$

Paths are safe for $y > x$

Example execution



CheckSafe on $\{l_0\}$ for $y > x$

Program is safe for $x \neq y$

Conclusion

We saw:

- 1 The exploration of multiple entry SCCs by CheckSafe
- 2 The effect of narrowing if a path to the SCC is proved safe and another could not be proved safe
- 3 The behavior of CheckSafe if a precondition with multiple conjunctions is found

We didn't saw:

- 1 SCCs with multiple transitions (effects consecution conditions in Max-SMT and entry locations in Narrowing)