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CIS7- FALL17

1. *Show that  $(p \rightarrow q)$  and  $(q' \rightarrow p')$  are logically equivalent without using truth tables or a "contrapositive" law (don't assume they are true)*

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|----|---|----------------------|
| a. | $(p' \vee q) \Leftrightarrow (q' \rightarrow p')$ | implication          |
|    | $(p' \vee q) \Leftrightarrow (q' \vee p')$        | implication          |
|    | $(p' \vee q) \Leftrightarrow (q \vee p')$         | double negation      |
|    | $(p' \vee q) \Leftrightarrow (p' \vee q)$         | communative property |

2. *Show that  $(p \rightarrow r) \wedge (q \rightarrow r) \Leftrightarrow (p \vee q) \rightarrow r$*

- |    |   |                      |
|----|---|----------------------|
| a. | $(p' \vee r) \wedge (q \rightarrow r) \Leftrightarrow (p \vee q) \rightarrow r$ | implication          |
|    | $(p' \vee r) \wedge (q' \vee r) \Leftrightarrow (p \vee q) \rightarrow r$       | implication          |
|    | $(p' \vee r) \wedge (q' \vee r) \Leftrightarrow (p \vee q)' \vee r$             | implication          |
|    | $(r \vee p') \wedge (r \vee q') \Leftrightarrow (p \vee q)' \vee r$             | communative property |
|    | $r \vee (p' \wedge q') \Leftrightarrow (p \vee q)' \vee r$                      | distribution         |
|    | $r \vee (p' \wedge q') \Leftrightarrow (p' \wedge q') \vee r$                   | negation             |
|    | $r \vee (p' \wedge q') \Leftrightarrow r \vee (p' \wedge q')$                   | communative property |

3. *Come up with 2 more questions like the ones above. Incorporate the logic laws we discussed in class. Your questions will be used in the next class' lab*

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|----|---|--------------|
| a. | $[(P \rightarrow Q) \wedge R] \Leftrightarrow [(P \vee R') \wedge (Q' \vee R')]'$ | negation     |
|    | $[(P' \vee Q) \wedge R] \Leftrightarrow (P' \wedge R) \vee (Q \wedge R)$          | implication  |
|    | $(P' \wedge R) \vee (Q \wedge R) \Leftrightarrow (P' \wedge R) \vee (Q \wedge R)$ | distribution |
| b. | $P \wedge (Q \rightarrow R) \Leftrightarrow (P \wedge Q') \vee (P \wedge R)$      |              |
|    | $P \wedge (Q' \vee R) \Leftrightarrow (P \wedge Q') \vee (P \wedge R)$            | implication  |
|    | $(P \wedge Q') \vee (P \wedge R) \Leftrightarrow (P \wedge Q') \vee (P \wedge R)$ | distribution |