

Homework Assignment: 3

Name: Jonathan Gaines

Date: 4/19/2017

1. Job Optimization

	Solution	Time Slot 1	Time Slot 2	profit
	1	Job 1	Job 3	55
	2	Job 3	Job 1	55
	3	Job 2	Job 1	65
	4	Job 2	Job 3	60
(a)	5	Job 4	Job 1	70
	6	Job 4	Job 3	65
	7	Job 1	N/A	30
	8	Job 2	N/A	35
	9	Job 3	N/A	25
	10	Job 4	N/A	40

(b) The optimal schedule has Job 4 in timeslot 1 and Job 1 in timeslot 2 for a profit of \$70.

(c) A high level greedy algorithm would choose the largest profit with a deadline of 1 or 2, then choose the largest profit with a deadline of 1. In this case, it would choose Job 4, then Job 1.

2. Dynamic Programming: Change Making

(a) The minimum number of coins needed to meet the amount is 3.

(b) Minimum coin combinations include {1, 2, 5} and {3, 3, 3}

(c)

n	0	1	2	3	4	5	6	7	8	9
f(n)	0	1	2	1	2	1	2	3	2	3

(d) Change-making(D[j], n):

```
f[0] = 0
for i = 1 to n do
    temp = ∞
    j = 1
    while j ≤ m and i ≥ D[j] do
        temp = min(f(i-D[j]), temp)
        j = j + 1
    f[i] = temp + 1
return f(n)
```

3. Dynamic Programming: Knapsack Problem

(a)

	0	1	2	3	4	5	6	7	8	9	10
0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	0	10	10	10	10	10	10
2	0	0	0	0	40	40	40	40	40	50	40
3	0	0	0	0	40	40	40	40	40	50	70
4	0	0	0	50	50	50	50	90	90	90	90

(b) The optimal subset has a value of \$90 and consists of items 2 and 4.

(c)

4. Greedy Algorithm

(a) Knapsack(vals[], weights[], w): n = len(vals[]) for j from 0 to w: m[0, j] = 0 for i from 1 to n: for j from 0 to w: if weights[i] ≤ j: m[i, j] = m[i-1, j] else: j = max(m[i-1, j], m[i-1, j - weights[i]] + vals[i]) return m[n, w]

(b) 10000110111

(c) BADFAD