

1. Job Optimization

| | Solution | Time Slot 1 | Time Slot 2 | profit |
|-----|----------|-------------|-------------|--------|
| | 1 | Job 1 | Job 3 | 55 |
| | 2 | Job 3 | Job 1 | 55 |
| | 3 | Job 2 | Job 1 | 65 |
| | 4 | Job 2 | Job 3 | 60 |
| (a) | 5 | Job 4 | Job 1 | 70 |
| | 6 | Job 4 | Job 3 | 65 |
| | 7 | Job 1 | N/A | 30 |
| | 8 | Job 2 | N/A | 35 |
| | 9 | Job 3 | N/A | 25 |
| | 10 | Job 4 | N/A | 40 |

- (b) The optimal schedule has Job 4 in timeslot 1 and Job 1 in timeslot 2 for a profit of \$70.
- (c) A high level greedy algorithm would choose the largest profit with a deadline of 1 or 2, then choose the largest profit with a deadline of 1. In this case, it would choose Job 4, then Job 1.

2. Dynamic Programming: Change Making

- (a) The minimum number of coins needed to meet the amount is 3.
- (b) Minimum coin combinations include {1, 2, 5} and {3, 3, 3}

(c)

| | | | | | | | | | | |
|--------|---|---|---|---|---|---|---|---|---|---|
| n | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| $f(n)$ | 0 | 1 | 2 | 1 | 2 | 1 | 2 | 3 | 2 | 3 |

- (d) Change-making($D[j]$, n):

```

    f[0] = 0
    for i = 1 to n do
        temp = ∞
        j = 1
        while j ≤ m and i ≥ D[j] do
            temp = min(f(i-D[j]), temp)
            j = j + 1
        f[i] = temp + 1
    return f(n)

```

3. Dynamic Programming: Knapsack Problem

- (a) 0 1 2 3 4 5 6 7 8 9 10
- (b)

(c)

4. Greedy Algorithm

(a)

(b)

(c)