

CSCE 222 [505] Discrete Structures for Computing
Fall 2015 – Philip C. Ritchey

Problem Set 3

Due dates: Electronic submission of L^AT_EX and PDF files of this homework is due on **23 September 2015 (Wednesday) before 11:30 a.m.** on eCampus (<http://ecampus.tamu.edu>).

Name	Problems
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9-22-15	

Resources. Discrete Structures and its Applications - Rosen Chapter 3.

<https://rob-bell.net/2009/06/a-beginners-guide-to-big-o-notation/>

<http://www.trans4mind.com/personal-development/mathematics/series/sumNaturalSquares.htm>

<http://www.ee.ryerson.ca/courses/coe428/sorting/insertionsort.html>

<http://codeyarns.com/2012/06/29/inserting-an-algorithm-in-latex/>

Problem 1. (10 points) Section 3.1, Exercise 4

Solution. :

Algorithm 1 Find largest difference in a list of integers

```
1: procedure DIFF( $x_1, \dots, x_n : integers$ )
2:    $large := x_1 - x_2$ 
3:   for  $i := 2$  to  $n-1$  do
4:     if  $large < x_{i+1} - x_i$  then  $large := x_{i+1} - x_i$ 
5:     end if
6:   end for
7:   Return  $large$ 
8: end procedure
```

Problem 2. (10 points) Supplementary Exercise 2

Solution. :

A)

Algorithm 2 Find first and second largest element in list of integers

```
1: procedure FIND TWO( $x_1, \dots, x_n : \text{integers}$ )
2:    $first := x_1$ 
3:   for  $i := 2$  to  $n$  do
4:     if  $first < x_i$  then  $first := x_i$ 
5:   end if
6: end for
7:    $second := x_1$ 
8:   for  $i := 2$  to  $n$  do
9:     if  $second < x_i$  and  $second \neq first$  then  $second := x_i$ 
10:  end if
11: end for
12:  Return  $first$ 
13:  Return  $second$ 
14: end procedure
```

B)

$O(n)$ for both for loops so number of comparisons is $O(n)$.

Problem 3. (10 points) Section 3.1, Exercise 38

Solution. :

Insertion Sort

Original list: 6, 2, 3, 1, 5, 4

Checks: $6 > 2$

After first check: 2, 6, 3, 1, 5, 4

Checks: $3 > 2, 3 > 6$

After second check: 2, 3, 6, 1, 5, 4

Checks: $1 > 2, 1 < 3, 1 < 6$

After third check: 1, 2, 3, 6, 5, 4

Checks: $5 > 1, 5 > 2, 5 > 3, 5 < 6$

After fourth check: 1, 2, 3, 5, 6, 4

Checks: $4 > 1, 4 > 2, 4 > 3, 4 < 5, 4 < 6$

After fifth check: 1, 2, 3, 4, 5, 6

List is now sorted after the fifth check.

Problem 4. (10 points) Section 3.1, Exercise 56

Solution. :

We can use proof by counterexample:

Our example could be : 20 cents.

We first start with the normal way of counting change.

20 cents is two dimes which is two coins.

Next, we use the 12 cent method to see how many coins it takes to satisfy 20 cents.

Using the 12 cent method, we get 1 12 cent, 1 nickel, and 3 pennies which is 5 coins.

Therefore, it does not produce change using the fewest coins possible by our proof by counterexample.

Problem 5. (10 points) Supplementary Exercise 18

Solution. :

$$\sum_{j=1}^n j(j+1)$$

$$= \sum_{j=1}^n j^2 + \sum_{j=1}^n j \text{ using: } \sum_{j=1}^n j = \frac{n(n+1)}{2}$$

$$\text{and } \sum_{j=1}^n j^2 = \frac{n(n+1)(2n+1)}{6}$$

$$= \frac{n(n+1)(2n+1)}{6} + \frac{n(n+1)}{2}$$

$$= \frac{n(n+1)}{2} \left(\frac{2n+1}{3} + 1 \right) \text{ by factoring}$$

$$= \frac{n(n+1)(n+2)}{3}$$

$$\text{So: } f(n) = \frac{n(n+1)(n+2)}{3} \leq 2n^3$$

We know $c = 2$ and $k = 1$

Therefore the Big O notation would be $O(n^3)$

Problem 6. (10 points) Supplementary Exercise 26

Solution. :

For this problem we need to know that n grows faster than $\log(n)$ and that $\log(\log(n)) < \log(n)$. We can see this by graphing these functions on a graphing calculator.

Looking at the growth rate of the functions given to order. We only care about the exponents so:

$n^{\log(n)}, n \log(n) \log \log(n), n(\log(n))^{3/2}, n^{4/3} \log(n)^2, n^{3/2}, 2^{100n}, 2^{n^2}, 2^{2n}, 2^{n!}$

Problem 7. (10 points) Section 3.2, Exercise 30

Solution. :

A) $3x+7 = O(x)$ b/c polynomial degree is 1

$$x \leq 3x$$

$$x \leq 3x+7$$

$$c = 3, k=1$$

$$O(3x+7)$$

$3x+7$ and x are same order

B) $2x^2 + x - 7 = O(x^2)$ b/c polynomial degree is 2

$$x^2 < 2x^2$$

$$x^2 < 2x^2 + x - 7$$

$$c=2, k=7$$

$$O(2x^2 + x - 7)$$

$2x^2 + x - 7$ and x^2 are same order

C) $[x + \frac{1}{2}] = O(x)$

$$x \leq [x + \frac{1}{2}]$$

$$c=1, k=1$$

$$O([x + \frac{1}{2}])$$

$[x + \frac{1}{2}]$ and x are same order

D) $\log(x^2 + 1) \leq \log(2x^2)$

$$\log(2x^2) = 2\log(x)$$

$$c = 2, k=1$$

$$= O(\log(x))$$

$$\log(x) \leq \log(x^2 + 1)$$

$$c=1, k=1$$

$$O(\log(x^2 + 1))$$

$\log(x^2 + 1)$ and $\log(x)$ are same order

E) $\log_{10}(x) \leq \log(x)$

$$O(\log(x))$$

$$\log(x) \leq \log_{10}(x)$$

$$O(\log_{10}(x))$$

$\log_{10}(x)$ and $\log(x)$ are same order

Problem 8. (10 points) Supplementary Exercise 30

Solution. :

A)

Algorithm 3 sorts integers and checks for each pair of terms whether their difference is in the sequence

```
1: procedure CHECK( $x_1 \dots x_n : integers$ )
2:   Sort ( $x_1 \dots x_n$ ) from high to low
3:   for  $i := 1$  to  $n-1$  do  $sub := x_{i+1} - x_i$ 
4:     for  $k := i+1$  to  $n$  do
5:       if  $x_k := sub$  then Return true
6:     end if
7:   end for
8: end for
9: end procedure
```

B) $O(n^2)$ because nested for loops, and this algorithm is better than the brute force since it more efficient.

Problem 9. (10 points) Section 3.3, Exercise 16

Solution. :

86400 seconds in one day which is approximately 10^5

so: $10^5/10^{-11} = 10^{16}$

A) $\log(n) \leq 10^{16}$

$$n \leq 2^{10^{16}}$$

B) $1000n \leq 10^{16}$

$$n \leq 10^{13}$$

$$10^{13}$$

C) $n^2 \leq 10^{16}$

$$n \leq 10^8$$

$$10^8$$

D) $1000n^2 \leq 10^{16}$

$$n^2 \leq 10^{13}$$

$$n \leq 10^{13/2}$$

$$10^{13/2}$$

E) $n^3 \leq 10^{16}$

$$10^{16/3}$$

F) $2^n \leq 10^{16}$

$$n \leq \log 10^{16}$$

$$\log 10^{16}$$

G) $2^{2n} \leq 10^{16}$

$$n \leq \frac{\log 10^{16}}{2}$$

$$\frac{\log 10^{16}}{2}$$

H) $2^{2^n} \leq 10^{16}$

$$2^n \leq \log 10^{16}$$

$$\log(\log 10^{16})$$

Problem 10. (10 points) Section 3.3, Exercise 46

Solution. :

A)

To use a brute-force algorithm we need to find a match for the first character of target then check each character afterwards. Using a while and a for loop to check for matching character of target and to check all the characters afterwards can help.

B)

Algorithm 4 Finding match for first character and checking successive characters for a match

```
1: procedure MATCH(x: target , y: text)
2:   i := 1
3:   while  $x_1 = y_i$  AND  $i \leq n$  do TorF := 1
4:     for k := 1 to m do
5:       if  $x_k \neq y_{i+k}$  then TorF := 0
6:     end if
7:   end for i := i+1
8:   end while
9: end procedure
```

C) Since outer loop runs n times and inner loop runs m times the big O is $O(nm)$.

Aggie Honor Statement: On my honor, as an Aggie, I have neither given nor received any unauthorized aid on any portion of the academic work included in this assignment. Furthermore, I have disclosed all resources (people, books, web sites, etc.) that have been used to prepare this homework.

Checklist:

1. Did you type your full name and UIN and those of any collaborators?
2. Did you abide by the Aggie Honor Code?
3. Did you solve all problems and start a new page for each?
4. Did you submit
 - (a) your \LaTeX source file?
 - (b) your PDF file?