# CSCE 222 [505] Discrete Structures for Computing Fall 2015 – Philip C. Ritchey

### Problem Set 4

Due dates: Electronic submission of LATEX and PDF files of this homework is due on 30 September 2015 (Wednesday) before 11:30 a.m. on eCampus (http://ecampus.tamu.edu).

Name	Problems
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9-30-15	
UIN: 124002210	

Resources. Discrete Mathematics and its Applications by Rosen Chapter 5 http://math.stackexchange.com/questions/241060/structural-inductions http://highered.mheducation.com/sites/0073383090/student-view0/index.html https://www.youtube.com/watch?v=dMn5w4-ztSw https://www.youtube.com/watch?v=IFqna5F0kW8 https://www.cs.cmu.edu/adamchik/21-127/lectures/induction-1-print.pdf http://www.inf.ed.ac.uk/teaching/courses/dmmr/slides/13-14/Ch5.pdf http://math.stackexchange.com/questions/855680/discrete-math-induction-problem http://math.stackexchange.com/questions/517440/whats-the-difference-betweensimple-induction-and-strong-induction

**Problem 1.** (10 points) Section 5.1, Exercise 4

#### Solution. :

p(n) = 
$$1^3 + 2^3 + n^3 = (\frac{n(n+1)}{2})^2$$
  
A)  $p(1) = 1^3 = (\frac{1(1+1)}{2})^2$   
B) Basis Step: $1^3 = (\frac{1(1+1)}{2})^2$ 

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$$p(1) = 1^3 = (\frac{1(1+1)}{2})^2$$

B) Basis Step: 
$$1^3 = (\frac{1(1+1)}{2})^2$$

1=1 so p(1) is True

C) IH: Assume 
$$p(k) = 1^2 + ... + k^3 = (\frac{k(k+1)}{2})^2$$

D) Prove: 
$$p(k) \rightarrow p(k+1)$$

D) Prove: 
$$p(k) \to p(k+1)$$
  
Show:  $p(k) = 1^3 + \dots + k^3 + (k+1)^3 = (\frac{k+1(k+2)}{2})^2$   
E)  $p(k) = k^3 = (\frac{k(k+1)}{2})^2$   
so:  $k^3 + (k+1)^3 = (\frac{k(k+1)}{2})^2 + (k+1)^3$  by IH  
 $= (\frac{k(k+1)}{2})(\frac{k(k+1)}{2}) + (k+1)^3$   
 $= (\frac{(k^2+1)(k^2+1)}{4}) + (k+1)^3$   
 $= 1/4k^2(k+1)^2 + (k+1)^2 + (k+1)$   
 $= 1/4(k+1)^2(k^2+4k+4)$ 

E) 
$$p(k) = k^3 = (\frac{k(k+1)}{2})^2$$

so: 
$$k^3 + (k+1)^3 = (\frac{k(k+1)}{2})^2 + (k+1)^3$$
 by IH

$$= (\frac{k(k+1)}{2})(\frac{k(k+1)}{2}) + (k+1)^3$$

$$=\left(\frac{(k^2+1)(k^2+1)}{4}\right)+(k+1)^3$$

$$= \frac{4}{1/4k^2(k+1)^2 + (k+1)^2 + (k+1)}$$

$$= 1/4(k+1)^2(k^2+4k+4)$$

$$= 1/4(k+1)^2(k+2)^2$$

$$= \frac{1}{4(k+1)^2(k+2)^2}$$

 $= \left(\frac{k+1(k+2)}{2}\right)^2$ 

Therefore p(k+1) is true by PMI.

F) The basis step is true and our inductive step is also true, so by PMI the formula is true whenever n is a positive number.

# **Problem 2.** (10 points) Supplementary Exercise 10

$$p(n) = n^3 + (n+1)^3 + (n+2)^3 divisible by 9$$
  
Basis:  $p(0) = 9$  divisible by 9? YES so  $p(0)$  is true Inductive:  $p(k) \to p(k+1)$   
Assume  $p(k) = k^3 + (k+1)^3 + (k+2)^3$  divisible by 9  
Show  $p(k+1)$  is true:  $p(k+1) - p(k)$  using IH  
 $= (k+1)^3 + (k+2)^3 + (k+3)^3 - k^3 - (k+1)^3 - (k+2)^3$   
 $= (k+3)^3 - k^3$   
 $= (k^2 + 6k + 9)(k+3) - k^2$  by expanding  $(k+3)^3$   
 $= 9k^2 + 27k + 27$   
 $= 9(k^2 + 3k + 3)$   
so we can see it is divisible by 9.  
therefore  $p(k)$  and  $p(k+1)$  are divisible by 9.

# Problem 3. (10 points) Supplementary Exercise 20

#### Solution. :

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first few numbers of fibonacci:
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(0,1,1,2,3,5,8,13,21)

Basis: f(0) = 0

0 is divisible by of 3. means position 0 of the sequence which is the first element in the sequence. So:

f(4) = 3 is a divisible of 3 so basis step is completed.

Inductive: assume f(k) is divisible by 3

show f(k+4) is divisible by 3

f(k+4) = f(k+3) + f(k+2)

= f(k+2) + f(k+1) + f(k+1) + f(k)

= f(k+1) + f(k) + f(k+1) + f(k+1) + f(k)

= 2f(k) + 3f(k+1)

by IH f(k) is divisible by 3 and we know f(k+1) is some integer times 3 which is divisible by 3

so f(k+4) is divisible by 3 by PMI

# Problem 4. (10 points) Section 5.2, Exercise 4

- A)Basis step:
- p(18) true b/c 1 4 cent stamp and 2 7 cent stamps
- p(19) true b/c 3 4 cent stamp and 1 7 cent stamp
- p(20) true b/c 5 4 cent stamps
- p(21) true b/c 3 7 cent stamps
- B) Inductive: Assume p(i) for  $18 \le i \le k$
- C) Show (k+1) with 4 cent and 7 cent stamps
- D) (k+1)cent = (k-3)cent and 4 cents.. k-3+4=k+1 so holds true
- E) True because both the basis and inductive steps were completed

**Problem 5.** (10 points) Section 5.2, Exercise 32

# Solution. :

The flaw is in the induction step from the 4 cent case, we can form it using only  $1\ 4$  cent stamp so there is no need for  $2\ 4$  cent stamps.

# **Problem 6.** (10 points) Supplementary Exercise 52

- A) No, not well ordered because it has infinite integers.
- B) Yes, it is well-ordered because least would be -99.
- C) No, not well ordered because it has positive infinite rationals with no least element.
- D) Yes, it is well ordered because least element would be 1/100.

# **Problem 7.** (10 points) Section 5.3, Exercise 13

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\begin{split} p(n) &= f_1 + f_3 + \ldots + f_{2n-1} = f_{2n} \\ \text{Basis step} : \ & \text{p}(1) = f_1 = 1 = f_2 \\ \text{Inductive step:} \ & \text{p}(k) \to \text{p}(k+1) \\ \text{Assume} \ & p(k) : f_1 + f_3 + \ldots + f_{2k-1} = f_{2k} \\ \text{Show p}(k+1) : \ & f_1 + f_3 + \ldots + f_{2k-1} + f_{2k+1} = f_{2k} + f_{2k+1} \\ \text{by IH the right hand side} &= f_{2k} + f_{2k+1} \\ &= f_{2k+2} \\ &= f_{2(k+1)} \\ \text{so:} \\ & \text{p}(k+1) \text{ is true by PMI} \end{split}
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# **Problem 8.** (10 points) Section 5.3, Exercise 32

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\begin{aligned} & \textbf{Solution.}: \\ & \Sigma = (0,1) \\ & x \in \Sigma \text{ , } t \in \Sigma^* \\ & \text{ones}(tx) = \text{ones}(t) + x \\ & B) \text{ Base step: } t = x \\ & \text{ones}(sx) = \text{ones}(s) + 0 \\ & = \text{ones}(s) + \text{ones}(x) \text{ b/c ones}(x) = 0 \\ & \text{Assume ones}(st) = \text{ones}(s) + \text{ones}(t) \\ & \text{Recursive step: } y \in \Sigma \text{ s,t} \in \Sigma^* \\ & \text{ones}(s+y) = \text{ones}(st) + y \\ & = \text{ones}(s) + \text{ones}(t) + y \text{ b/c from above we get ones}(s) + \text{ones}(ty) \end{aligned}
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There fore by structural induction it is proved, because of our assumption we used in recursive step.

# Problem 9. (10 points) Section 5.3, Exercise 44

# Solution. :

Basis step: since no internal vertices i(T) = 0 so l(T) = 1= 1 + i(T)l(T) = 1 + i(T) is True Assume: l(s) = 1 + i(S)Recursive step: If it is true for  $T_1$  and  $T_2$  then it is true for  $T_1$  know:  $i(T) = i(T_1) + i(T_2) + 1$  $l(T) = l(T_1) + l(T_2)$  $= i(T_1) + 1 + i(T_2) + 1$  from assumption = i(T) + 1 using  $i(T) = i(T_1) + i(T_2) + 1$ l(T) = i(T) + 1 is true by structural induction Problem 10. (10 points) Section 5.3, Exercise 64

#### Solution. :

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f(n) = n/2 \dots it is basically f divided by 2 divide by 2^k to iterate through function so: f^k(n) = n/2^k f^k(n) \le 1 n/2^k \le 1 n \le 2^k log_2(n) \le k so: f_1^k(n) = log_2(n)
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#### Checklist:

- 1. Did you type your full name and UIN and those of any collaborators?
- 2. Did you abide by the Aggie Honor Code?
- 3. Did you solve all problems and start a new page for each?
- 4. Did you submit
  - (a) your LATEX source file?
  - (b) your PDF file?