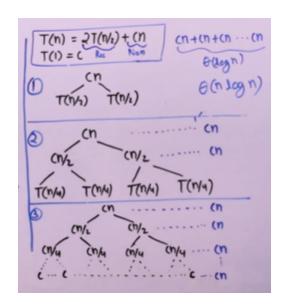
```
Asymptotic Analysis – measures the order of growth of an algorithm in terms of the input size
```

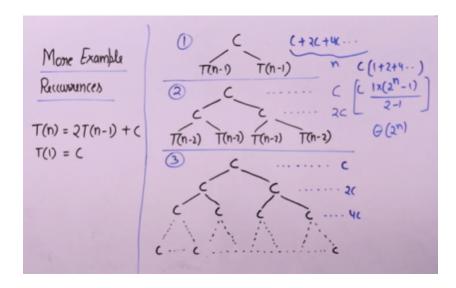
```
1 < \log \log n < \log n < n^{1/3} < n^{1/2} < n < n^2 < n^3 < n^4 < 2^n < n^n
Big O Notation – upper/exact bound
Theta Notation – exact bound
Omega Notation – lower/exact bound
for (int i = 0; i < n; i = i + c) {
  // θ (1) work
}
\theta (n)
for (int i = 0; i < n; i = i - c) {
  // \theta (1) work
}
\theta (n)
for (int i = 0; i < n; i = i * c) {
  // \theta (1) work
\theta (log n)
for (int i = 0; i < n; i = i / c) {
  // \theta (1) work
}
\theta (log n)
for (int i = 0; i < n; i = pow(i, c)) {
  // \theta (1) work
\theta (log log n)
add subsequent loop values
multiply nested loop values
```

```
void fun (int n) {
   if (n \le 0)
     return;
  // θ (1) work
  fun (n / 2);
   fun (n / 2);
}
T(n) = 2T(n/2) + \theta(1)
T(0) = \theta(1)
void fun (int n) {
   if (n \le 0)
     return;
  for (int i = 0; i < n; i = i++) {
     // \theta (1) work
  }
  fun (n / 2);
  fun (n / 3);
T(n) = T(n/2) + T(n/3) + \theta(n)
T(0) = \theta(1)
void fun (int n) {
   if (n \le 1)
     return;
  // \theta (1) work
  fun (n - 1);
T(n) = T(n-1) + \theta(1)
T(1) = \theta(1)
```

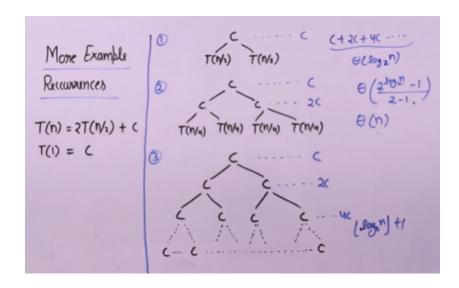


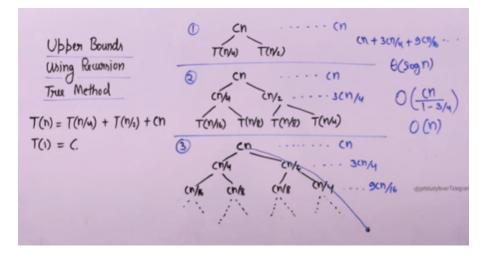
$$T (n) = 2T (n / 2) + c_1 n$$
  
 $T (1) = c_2$ 

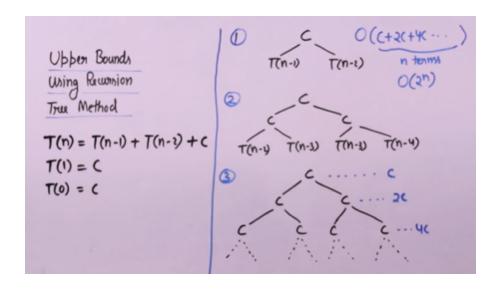
θ (nlog n)



Mone Example	0	C (+(+···(
Recommended $T(n) = T(n/2) + C$ $T(1) = C$	averTelegram	C C & (log 2h)  C C  I  T(n/4)
	3	¿ c





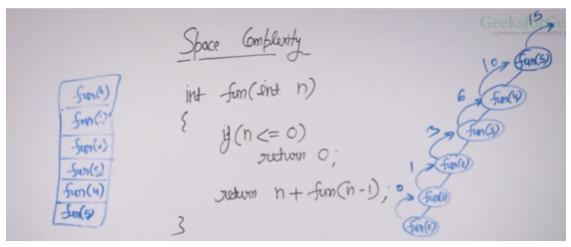


## Space Complexity -

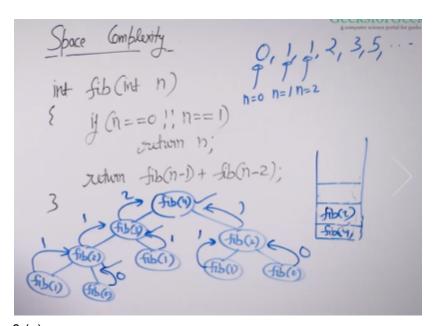
```
int get_sum1 (int n) {
  return n * (n + 1) / 2;
}
θ (1)
int get_sum2 (int n) {
  int sum = 0;
  for (int i = 1; i \le n; i++)
     sum = sum + i;
  return sum;
}
\theta (1)
int arr_sum (int a [], int n) {
  int sum = 0;
  for (int i = 0; i < n; i++)
     sum += a [i];
  return sum;
θ (n)
```

## Auxiliary Space - measures the order of growth of extra space

```
int arr_sum (int a [], int n) {
  int sum = 0;
  for (int i = 0; i < n; i++)
     sum += a [i];
  return sum;
}
θ (1)</pre>
```



θ (n)



θ (n)