

# NJIT



New Jersey's Science &  
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***THE EDGE IN KNOWLEDGE***

# **CS 280**

# **Programming Language**

# **Concepts**

## **Types**

# Introduction

- A *data type* defines a collection of data objects and a set of predefined operations on those objects
- A *descriptor* is the collection of the attributes of a variable
- An *object* represents an instance of a user-defined (abstract data) type
- Design issue for all data types: What operations are defined and how are they specified?

# Primitive Data Types

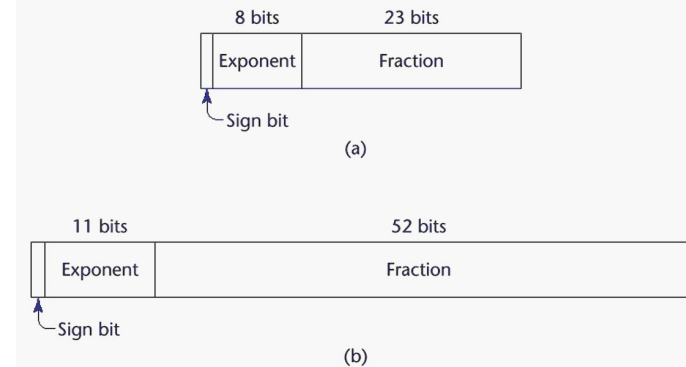
- Almost all programming languages provide a set of *primitive data types*
- Primitive data types: Those not defined in terms of other data types
- Some primitive data types are merely reflections of the hardware
- Others require only a little non-hardware support for their implementation

# Primitive Data Types: Integer

- Almost always an exact reflection of the hardware so the mapping is trivial
- There may be as many as eight different integer types in a language
- Java's signed integer sizes: **byte**, **short**, **int**, **long**

# Primitive Data Types: Floating Point

- Model real numbers, but only as approximations
- Languages for scientific use support at least two floating-point types (e.g., **float** and **double**; sometimes more)
- Usually exactly like the hardware, but not always
- IEEE Floating-Point Standard 754



# Primitive Data Types: Complex

- Some languages support a complex type, e.g., C99, Fortran, and Python
- Each value consists of two floats, the real part and the imaginary part
- Literal form (in Python):  
 $(7 + 3j)$ , where 7 is the real part and 3 is the imaginary part

# Primitive Data Types: Decimal

- For business applications (money)
  - Essential to COBOL
  - C# offers a decimal data type
- Store a fixed number of decimal digits, in coded form (BCD)
- *Advantage*: accuracy
- *Disadvantages*: limited range, wastes memory



# Primitive Data Types: Boolean

- Simplest of all
- Range of values: two elements, one for “true” and one for “false”
- Could be implemented as bits, but often as bytes
  - Advantage: readability

# Primitive Data Types: Character

- Stored as numeric encodings
- Most commonly used coding: ASCII
- An alternative, 16-bit coding: Unicode (UCS-2)
  - Includes characters from most natural languages
  - Originally used in Java
  - C# and JavaScript also support Unicode
- 32-bit Unicode (UCS-4)
  - Supported by Fortran, starting with 2003

# Character String Types

- Values are sequences of characters
- Design issues:
  - Is it a primitive type or just a special kind of array?
  - Should the length of strings be static or dynamic?

# Character String Types: Operations

- Typical operations:
  - Assignment and copying
  - Comparison (=, >, etc.)
  - Catenation
  - Substring reference
  - Pattern matching

# Character String Type in Certain Languages

- C and C++
  - Not primitive
  - Use `char` arrays and a library of functions that provide operations
  - A `string` class is provided in C++
- SNOBOL4 (a string manipulation language)
  - Primitive
  - Many operations, including elaborate pattern matching
- Fortran and Python
  - Primitive type with assignment and several operations
- Java
  - Primitive via the `String` class
- Perl, JavaScript, Ruby, AWK and PHP
  - Provide built-in pattern matching, using regular expressions

# User-Defined Ordinal Types

- An ordinal type is one in which the range of possible values can be easily associated with the set of positive integers
- Examples of primitive ordinal types in Java
  - `integer`
  - `char`
  - `boolean`

# Enumeration Types

- All possible values, which are named constants, are provided in the definition

- C# example

```
enum days {mon, tue, wed, thu, fri, sat, sun};
```

- Design issues

- Is an enumeration constant allowed to appear in more than one type definition, and if so, how is the type of an occurrence of that constant checked?
- Are enumeration values coerced to integer?
- Any other type coerced to an enumeration type?

# Evaluation of Enumerated Type

- Aid to readability, e.g., no need to code a color as a number
- Aid to reliability, e.g., compiler can check:
  - operations (don't allow colors to be added)
  - No enumeration variable can be assigned a value outside its defined range
  - Ada, C#, and Java 5.0 provide different support for enumeration than C++: enumeration type variables in these languages are not coerced into integer types



# Subrange Types

- An ordered contiguous subsequence of an ordinal type
  - Example: 12..18 is a subrange of integer type

- Ada's design

```
type Days is (mon, tue, wed, thu, fri, sat, sun);
```

```
subtype Weekdays is Days range mon..fri;
```

```
subtype Index is Integer range 1..100;
```

```
Day1: Days;
```

```
Day2: Weekday;
```

```
Day2 := Day1;
```

# Subrange Evaluation

- Aid to readability
  - Make it clear to the readers that variables of subrange can store only certain range of values
- Reliability
  - Assigning a value to a subrange variable that is outside the specified range is detected as an error

# Implementation of User-Defined Ordinal Types

- Enumeration types are implemented as integers
- Subrange types are implemented like the parent types with code inserted (by the compiler) to restrict assignments to subrange variables

# Array Types

- An array is a homogeneous aggregate of data elements in which an individual element is identified by its position in the aggregate, relative to the first element.

# Array Design Issues

- What types are legal for subscripts?
- Are subscripting expressions in element references range checked?
- When are subscript ranges bound?
- When does allocation take place?
- Are ragged or rectangular multidimensional arrays allowed, or both?
- What is the maximum number of subscripts?
- Can array objects be initialized?
- Are any kind of slices supported?

# Array Indexing

- *Indexing* (or subscripting) is a mapping from indices to elements  
array\_name (index\_value\_list)  $\mapsto$  an element
- **Index Syntax**
  - Fortran and Ada use parentheses
    - Ada explicitly uses parentheses to show uniformity between array references and function calls because both are *mappings*
  - Most other languages use square brackets

# Arrays Index (Subscript) Types

- FORTRAN, C: integer only
- Ada: integer or enumeration (includes Boolean and char)
- Java: integer types only
- Index range checking
  - C, C++, Perl, and Fortran do not specify range checking
  - Java, ML, C# specify range checking
  - In Ada, the default is to require range checking, but it can be turned off

# Subscript Binding and Array Categories

- *Static*: subscript ranges are statically bound and storage allocation is static (before run-time)
  - Advantage: efficiency (no dynamic allocation)
- *Fixed stack-dynamic*: subscript ranges are statically bound, but the allocation is done at declaration time
  - Advantage: space efficiency



# Subscript Binding and Array Categories (continued)

- *Stack-dynamic*: subscript ranges are dynamically bound and the storage allocation is dynamic (done at run-time)
  - Advantage: flexibility (the size of an array need not be known until the array is to be used)
- *Fixed heap-dynamic*: similar to fixed stack-dynamic: storage binding is dynamic but fixed after allocation (i.e., binding is done when requested and storage is allocated from heap, not stack)

# Subscript Binding and Array Categories (continued)

- Heap-dynamic: binding of subscript ranges and storage allocation is dynamic and can change any number of times
  - Advantage: flexibility (arrays can grow or shrink during program execution)

# Subscript Binding and Array Categories (continued)

- C and C++ arrays that include `static` modifier are static
- C and C++ arrays without `static` modifier are fixed stack-dynamic
- C and C++ provide fixed heap-dynamic arrays
- C# includes a second array class `ArrayList` that provides fixed heap-dynamic
- Perl, JavaScript, Python, and Ruby support heap-dynamic arrays

# Array Initialization

- Some language allow initialization at the time of storage allocation

- C, C++, Java, C# example

```
int list [] = {4, 5, 7, 83};
```

- Character strings in C and C++

```
char name [] = "freddie";
```

- Arrays of strings in C and C++

```
char *names [] = {"Bob", "Jake", "Joe"};
```

- Java initialization of String objects

```
String[] names = {"Bob", "Jake", "Joe"};
```

# Array Initialization

- C-based languages

- `int list [] = {1, 3, 5, 7};`
- `char *names [] = {"Mike", "Fred", "Mary Lou"};`

- Ada

- `List : array (1..5) of Integer :=  
 (1 => 17, 3 => 34, others => 0);`

- Python

- List comprehensions

```
list = [x ** 2 for x in range(12) if x % 3 == 0]
```

```
puts [0, 9, 36, 81] in list
```

# Heterogeneous Arrays

- A *heterogeneous array* is one in which the elements need not be of the same type
- Supported by Perl, Python, JavaScript, and Ruby

# Arrays Operations

- APL provides the most powerful array processing operations for vectors and matrixes as well as unary operators (for example, to reverse column elements)
- Ada allows array assignment but also catenation
- Python has array assignments, but they are only reference changes. Python also supports array catenation and element membership operations
- Ruby also provides array catenation
- Fortran provides *elemental* operations because they are between pairs of array elements
  - For example, + operator between two arrays results in an array of the sums of the element pairs of the two arrays

# Rectangular and Jagged Arrays

- A rectangular array is a multi-dimensional array in which all of the rows have the same number of elements and all columns have the same number of elements
- A jagged matrix has rows with varying number of elements
  - Possible when multi-dimensional arrays actually appear as arrays of arrays
- C, C++, and Java support jagged arrays
- Fortran, Ada, and C# support rectangular arrays (C# also supports jagged arrays)



# Slices

- A slice is some substructure of an array; nothing more than a referencing mechanism
- Slices are only useful in languages that have array operations

# Slice Examples

- Python

```
vector = [2, 4, 6, 8, 10, 12, 14, 16]
```

```
mat = [[1, 2, 3], [4, 5, 6], [7, 8, 9]]
```

`vector [3:6]` is a three-element array

`mat[0][0:2]` is the first and second element of the first row of `mat`

- Ruby supports slices with the `slice` method

`list.slice(2, 2)` returns the third and fourth elements of `list`

# Implementation of Arrays

- Access function maps subscript expressions to an address in the array
- Access function for single-dimensioned arrays:  
$$\text{address}(\text{list}[k]) = \text{address}(\text{list}[\text{lower\_bound}]) + ((k - \text{lower\_bound}) * \text{element\_size})$$

	0	1	...	$j-1$	$j$	...	$n-1$
0							
1							
$\vdots$							
$i-1$							
$i$					⊗		
$\vdots$							
$m-1$							

# Accessing Multi-dimensional Arrays

- Two common ways:
  - Row major order (by rows) – used in most languages
  - Column major order (by columns) – used in Fortran
  - A compile-time descriptor for a multidimensional array

Multidimensioned array
Element type
Index type
Number of dimensions
Index range 0
⋮
Index range $n - 1$
Address

# Locating an Element in a Multi-dimensioned Array

- General format

Location ( $a[l,j]$ ) = address of  $a$  [ $\text{row\_lb}, \text{col\_lb}$ ] + ((( $l - \text{row\_lb}$ ) \*  $n$ ) + ( $j - \text{col\_lb}$ )) \*  $\text{element\_size}$

	1	2	...	$j-1$	$j$	...	$n$
1							
2							
$\vdots$							
$i-1$							
$i$					⊗		
$\vdots$							
$m$							

# Compile-Time Descriptors

Array
Element type
Index type
Index lower bound
Index upper bound
Address

Single-dimensioned array

Multidimensioned array
Element type
Index type
Number of dimensions
Index range 1
⋮
Index range $n$
Address

Multidimensional array

# Associative Arrays

- An *associative array* is an unordered collection of data elements that are indexed by an equal number of values called *keys*
  - User-defined keys must be stored
- Design issues:
  - What is the form of references to elements?
  - Is the size static or dynamic?
- Built-in type in Perl, Python, Ruby, ...

# Associative Arrays in Perl

- Names begin with % ; literals are delimited by parentheses

```
%hi_temps = ("Mon" => 77, "Tue" => 79, "Wed" => 65, ...);
```

- Subscripting is done using braces and keys

```
$hi_temps{"Wed"} = 83;
```

- Elements can be removed with **delete**

```
delete $hi_temps{"Tue"};
```



# Related Implementations in Java/C++

- C++ map or Java HashMap
- These are parameterized using `<>` and type names, so the type of the key and value are fixed
- Advantage: compile time type checking

# Record Types

- A *record* is a possibly heterogeneous aggregate of data elements in which the individual elements are identified by names
- Design issues:
  - What is the syntactic form of references to the field?
  - Are elliptical references allowed

# Definition of Records in COBOL

- COBOL uses level numbers to show nested records; others use recursive definition

```
01 EMP-REC.
```

```
    02 EMP-NAME.
```

```
        05 FIRST PIC X(20).
```

```
        05 MID    PIC X(10).
```

```
        05 LAST   PIC X(20).
```

```
    02 HOURLY-RATE PIC 99V99.
```

# Related Implementations in C/C++

- **struct:** a collection of elements of different types

```
struct employeeType {  
    int id;  
    char name[25];  
    int age;  
    float salary;  
    char dept;  
};  
struct employeeType employee;  
...  
employee.age = 45;
```

# References to Records

- Record field references

## 1. COBOL

`field_name OF record_name_1 OF ... OF record_name_n`

## 2. Others (dot notation)

`record_name_1.record_name_2. ... record_name_n.field_name`

- Fully qualified references must include all record names
- Elliptical references allow leaving out record names as long as the reference is unambiguous, for example in COBOL

`FIRST`, `FIRST OF EMP-NAME`, and `FIRST of EMP-REC` are elliptical references to the employee's first name

# Operations on Records

- Assignment is very common if the types are identical
- Ada allows record comparison
- COBOL provides `MOVE CORRESPONDING`
  - Copies a field of the source record to the corresponding field in the target record

# Evaluation and Comparison to Arrays

- Records are used when collection of data values is heterogeneous
- Access to array elements is much slower than access to record fields, because subscripts are dynamic (field names are static)
- Dynamic subscripts could be used with record field access, but it would disallow type checking and it would be much slower

# Tuple Types

- A tuple is a data type that is similar to a record, except that the elements are not named
- Used in Python, ML, and F# to allow functions to return multiple values
  - Python
    - Closely related to its lists, but immutable
    - Create with a tuple literal

```
myTuple = (3, 5.8, 'apple')
```

Referenced with subscripts

Catenation with `+` and deleted with `del`



# List Types

- Lists in LISP and Scheme are delimited by parentheses and use no commas

(A B C D) and (A (B C) D)

- Data and code have the same form

As data, (A B C) is literally what it is

As code, (A B C) is the function A applied to the parameters B and C

- The interpreter needs to know which a list is, so if it is data, we quote it with an apostrophe

' (A B C) is data

# List Types (continued)

- List Operations in Scheme

- CAR returns the first element of its list parameter

(CAR ' (A B C) ) returns A

- CDR returns the remainder of its list parameter after the first element has been removed

(CDR ' (A B C) ) returns (B C)

- CONS puts its first parameter into its second parameter, a list, to make a new list

(CONS 'A (B C) ) returns (A B C)

- LIST returns a new list of its parameters

(LIST 'A 'B ' (C D) ) returns (A B (C D) )

# List Types (continued)

- Python Lists

- The list data type also serves as Python's arrays
- Unlike Scheme, Common LISP, ML, and F#, Python's lists are mutable
- Elements can be of any type
- Create a list with an assignment

```
myList = [3, 5.8, "grape"]
```

# List Types (continued)

- Python Lists (continued)

- List elements are referenced with subscripting, with indices beginning at zero

`x = myList[1]`    Sets `x` to 5.8

- List elements can be deleted with `del`

`del myList[1]`

- List Comprehensions – derived from set notation

`[x * x for x in range(6) if x % 3 == 0]`

`range(12)` creates `[0, 1, 2, 3, 4, 5, 6]`

Constructed list: `[0, 9, 36]`

# List Types (continued)

- C++ supports lists through standard templates like vector and list
- Both C# and Java supports lists through their generic heap-dynamic collection classes, `List` and `ArrayList`, respectively
- Types are fixed and are named in `<>` at declaration time

# Union Types

- A *union* is a type whose variables are allowed to store different type values at different times during execution
- Design issues
  - Should type checking be required?
  - Should unions be embedded in records?

# Discriminated vs. Free Unions

- Fortran, C, and C++ provide union constructs in which there is no language support for type checking; the union in these languages is called *free union*
- Type checking of unions require that each union include a type indicator called a *discriminant*
  - Supported by Ada

# C/C++ union

```
union u {  
    char    addr[4];  
    short   pieces[2];  
    long    full;  
} ipaddr;
```

```
union spacesave {  
    char    c;  
    int     i;  
    float   f;  
} ss;
```



# Evaluation of Unions

- Free unions are unsafe
  - Do not allow type checking
- Java and C# do not support unions
  - Reflective of growing concerns for safety in programming language
- Ada's discriminated unions are safe

# User Defined Types

- A class in Java or C++
- Extensions of classes using generics/templates
- Java Interfaces

# Pointer and Reference Types

- A *pointer* type variable has a range of values that consists of memory addresses and a special value, *nil*
- Provide the power of indirect addressing
- Provide a way to manage dynamic memory
- A pointer can be used to access a location in the area where storage is dynamically created (the *heap*)

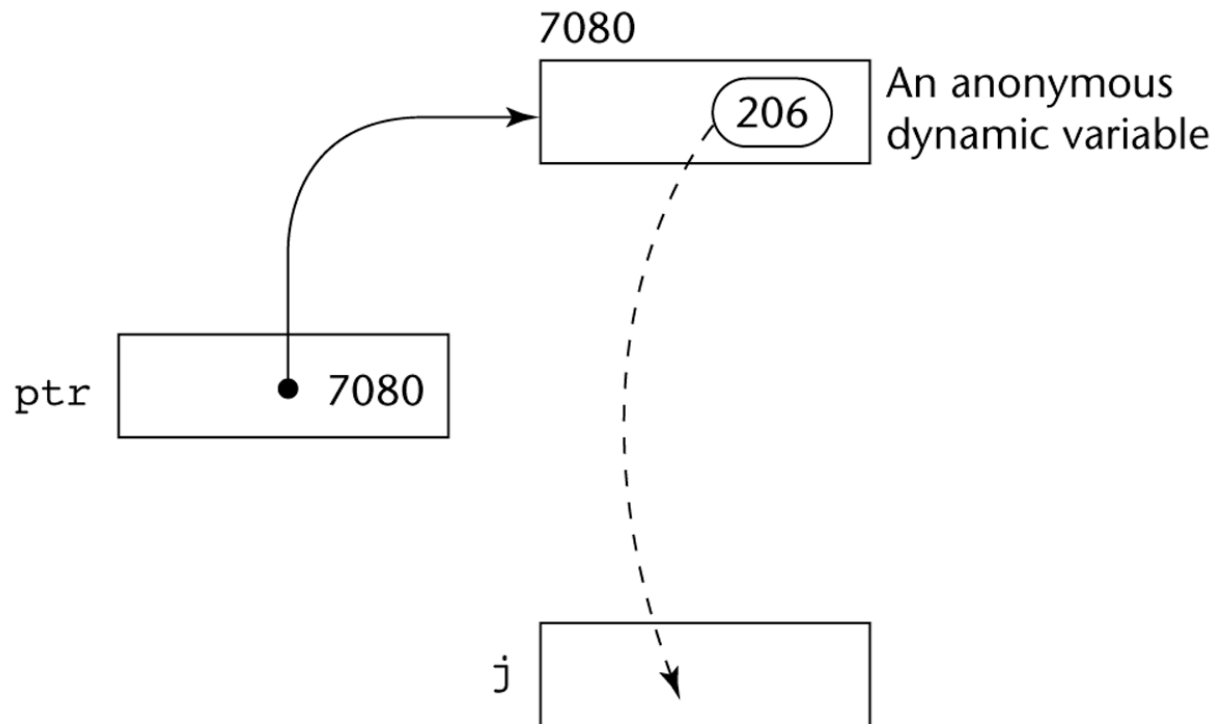
# Design Issues of Pointers

- What are the scope of and lifetime of a pointer variable?
- What is the lifetime of a heap-dynamic variable?
- Are pointers restricted as to the type of value to which they can point?
- Are pointers used for dynamic storage management, indirect addressing, or both?
- Should the language support pointer types, reference types, or both?

# Pointer Operations

- Two fundamental operations: assignment and dereferencing
- Assignment is used to set a pointer variable's value to some useful address
- Dereferencing yields the value stored at the location represented by the pointer's value
  - Dereferencing can be explicit or implicit
  - C++ uses an explicit operation via `*`  
`j = *ptr`  
sets `j` to the value located at `ptr`

# Pointer Assignment Illustrated



The assignment operation  $j = *ptr$

# Problems with Pointers

- Dangling pointers (dangerous)
  - A pointer points to a heap-dynamic variable that has been deallocated
- Lost heap-dynamic variable
  - An allocated heap-dynamic variable that is no longer accessible to the user program (often called *garbage*)
    - Pointer `p1` is set to point to a newly created heap-dynamic variable
    - Pointer `p1` is later set to point to another newly created heap-dynamic variable
    - The process of losing heap-dynamic variables is called *memory leakage*

# Pointers in C and C++

- Extremely flexible but must be used with care
- Pointers can point at any variable regardless of when or where it was allocated
- Used for dynamic storage management and addressing
- Pointer arithmetic is possible
- Explicit dereferencing and address-of operators
- Domain type need not be fixed (`void *`)
  - `void *` can point to any type and can be type checked (cannot be de-referenced)



# Pointer Arithmetic in C and C++

```
float stuff[100];  
float *p;  
p = stuff;
```

`*(p+5)` is equivalent to `stuff[5]` and `p[5]`

`*(p+i)` is equivalent to `stuff[i]` and `p[i]`

# Reference Types

- C++ includes a special kind of pointer type called a *reference type* that is used primarily for formal parameters
  - Advantages of both pass-by-reference and pass-by-value
- Java extends C++'s reference variables and allows them to replace pointers entirely
  - References are references to objects, rather than being addresses
- C# includes both the references of Java and the pointers of C++

# Evaluation of Pointers

- Dangling pointers and dangling objects are problems as is heap management
- Pointers are like `goto`'s--they widen the range of cells that can be accessed by a variable
- Pointers or references are necessary for dynamic data structures--so we can't design a language without them

# Heap Management

- A very complex run-time process
- Single-size cells vs. variable-size cells
- Two approaches to reclaim garbage
  - Reference counters (*eager approach*): reclamation is gradual
  - Mark-sweep (*lazy approach*): reclamation occurs when the list of variable space becomes empty

# Reference Counter

- Reference counters: maintain a counter in every cell that store the number of pointers currently pointing at the cell
  - *Disadvantages*: space required, execution time required, complications for cells connected circularly
  - *Advantage*: it is intrinsically incremental, so significant delays in the application execution are avoided

# Mark-Sweep

- The run-time system allocates storage cells as requested and disconnects pointers from cells as necessary; mark-sweep then begins
  - Every heap cell has an extra bit used by collection algorithm
  - All cells initially set to garbage
  - All pointers traced into heap, and reachable cells marked as not garbage
  - All garbage cells returned to list of available cells
  - Disadvantages: in its original form, it was done too infrequently. When done, it caused significant delays in application execution. Contemporary mark-sweep algorithms avoid this by doing it more often—called incremental mark-sweep

# Variable-Size Cells

- All the difficulties of single-size cells plus more
- Required by most programming languages
- If mark-sweep is used, additional problems occur
  - The initial setting of the indicators of all cells in the heap is difficult
  - The marking process is nontrivial
  - Maintaining the list of available space is another source of overhead

# Type Checking

- Generalize the concept of operands and operators to include subprograms and assignments
- *Type checking* is the activity of ensuring that the operands of an operator are of compatible types
- A *compatible type* is one that is either legal for the operator, or is allowed under language rules to be implicitly converted, by compiler- generated code, to a legal type
  - This automatic conversion is called a *coercion*.
- A *type error* is the application of an operator to an operand of an inappropriate type



# Type Checking (continued)

- If all type bindings are static, nearly all type checking can be static
- If type bindings are dynamic, type checking must be dynamic
- A programming language is *strongly typed* if type errors are always detected
- Advantage of strong typing: allows the detection of the misuses of variables that result in type errors

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