Due: 26th of August 2018 at 11:59pm

# COMP 9020 – Assignment 1

Note: In your assignment, how you arrived at your solution is as important (if not more so) than the solution itself and will be assessed accordingly. There may be more than one way to find a solution, and your approach should contain enough detail to justify its correctness. Lecture content can be assumed to be common knowledge.

- 1. (a) Compute gcd(132, 84).
  - (b) Suppose  $a, b \in \mathbb{N}$  are co-prime. What is gcd(a, a + b)?

## Solution:

(a) From the Euclidean algorithm (presented in lectures) we have:

$$\gcd(132, 84) = \gcd(132 - 84, 84)$$

$$= \gcd(48, 84)$$

$$= \gcd(48, 84 - 48)$$

$$= \gcd(48, 36)$$

$$= \gcd(48 - 36, 36)$$

$$= \gcd(12, 36)$$

$$= \gcd(12, 36 - 12)$$

$$= \gcd(12, 24)$$

$$= \gcd(12, 24 - 12)$$

$$= \gcd(12, 12)$$

$$= 12$$

(4 marks)

(b) We have  $a+b\geq a$  and  $\gcd(a,b)=1$ . Therefore, from the Euclidean algorithm we have:

$$gcd(a, a + b) = gcd(a, (a + b) - a) = gcd(a, b) = 1.$$

That is, a and a + b are co-prime. (6 marks)

- 2. For sets A and B, define A \* B to be  $(A \cup B)^c$  (the complement of  $A \cup B$ ).
  - (a) Simplify (A \* B) \* (A \* B). Justify your answer (e.g. using a Venn diagram or some other technique).

- (b) Express  $A^c$  using A and \*. Justify your answer.
- (c) Express  $A \cap B$  using A, B, and \*. Justify your answer.

#### Solution: (A\*B)\*(A\*B)(a) $= ((A * B) \cup (A * B))^c$ (Definition of \*) $=(A*B)^c$ (Idempotence) $=((A\cup B)^c)^c$ (Definition of \*) $= A \cup B$ (Double complement) (3 marks) (b) $=(A\cup A)^c$ (Idempotence) = A \* A(Definition of \*) (3 marks) (c) $A \cap B$ $= ((A^c)^c \cap (B^c)^c)$ (Double complement) $=(A^c \cup B^c)^c$ (De Morgan) $= (A^c) * (B^c)$ (Definition of \*) = (A \* A) \* (B \* B)(from (b)) (4 marks)

- 3. (a) List all possible functions  $f: \{a, b, c\} \rightarrow \{0, 1\}$ 
  - (b) Describe a connection between your answer for (a) and  $Pow(\{a, b, c\})$ .
  - (c) In general, if card(A) = m and card(B) = n, how many:
    - (i) functions are there from A to B?
    - (ii) relations are there between A and B?

## Solution:

- (a) There are eight functions from  $\{a, b, c\}$  to  $\{0, 1\}$ :
  - $f_0: a \mapsto 0, b \mapsto 0, c \mapsto 0$
  - $f_1: a \mapsto 0, b \mapsto 0, c \mapsto 1$
  - $f_2$ :  $a \mapsto 0$ ,  $b \mapsto 1$ ,  $c \mapsto 0$
  - $f_3$ :  $a \mapsto 0$ ,  $b \mapsto 1$ ,  $c \mapsto 1$
  - $f_4$ :  $a \mapsto 1$ ,  $b \mapsto 0$ ,  $c \mapsto 0$
  - $f_5$ :  $a \mapsto 1$ ,  $b \mapsto 0$ ,  $c \mapsto 1$
  - $f_6: a \mapsto 1, b \mapsto 1, c \mapsto 0$
  - $f_7: a \mapsto 1, b \mapsto 1, c \mapsto 1$

(3 marks)

- (b) We observe that the cardinality of  $Pow(\{a, b, c\})$  is equal to the number of functions from  $\{a, b, c\}$  to  $\{0, 1\}$ . Indeed, for each function  $f: \{a, b, c\} \rightarrow \{0, 1\}$  we can associate a unique element of  $Pow(\{a, b, c\})$  given by  $f^{\leftarrow}(1)$ . For example,  $f_0$  corresponds to  $\emptyset$ ;  $f_5$  corresponds to  $\{a, c\}$ . (3 marks)
- (c) In general, if card(A) = m and card(B) = n, there are:
  - (i)  $n^m$  functions from A to B because each of the m elements of A can map to one of n elements of B yielding  $n \times n \times \cdots n = n^m$  possible functions. (2 marks)
  - (ii)  $2^{mn}$  relations between A and B because a relation is a subset of  $A \times B$  and there are  $2^{|A \times B|} = 2^{mn}$  subsets of  $A \times B$ . (2 marks)
- 4. Let  $\Sigma = \{a, b\}$  and  $L = \{w \in \Sigma^* : 3 | \operatorname{length}(w) \}$ .
  - (a) List the elements of  $L^{\leq 3}$  in lexicographic order.

Define  $R \subseteq \Sigma^* \times \Sigma^*$  as follows:  $(w, w') \in R$  if there is a  $v \in \Sigma^*$  such that: either  $wv \in L$  and  $w'v \notin L$ , or  $wv \notin L$  and  $w'v \in L$ . For example  $(a, bbb) \in R$  because for  $v = \lambda$ ,  $av = a \notin L$  and  $bbbv = bbb \in L$ . On the other hand,  $(a, b) \notin R$  because for any  $v \in \Sigma^*$ , length(av) = length(bv); so whenever  $av \in L$ ,  $bv \in L$  and vice-versa.

- (b) Which of the following are elements of R:
  - (i) (abab, baba)?
  - (ii) (ab, abab)?
  - (iii)  $(\lambda, b)$ ?
  - (iv)  $(\lambda, bb)$ ?

(v)  $(\lambda, bbb)$ ?

Now define  $S\subseteq \Sigma^*\times \Sigma^*$  as the complement of R. That is  $(w,w')\in S$  if, and only if,  $(w,w')\notin R$ .

- (b) State a simple rule for determining whether  $(w, w') \in S$ . Hint: consider length(w) length(w')
- (c) Show that S is an equivalence relation. That is, show that S is reflexive, symmetric, and transitive.
- (d) How many equivalence classes does S have?

## Solution:

(a) The elements of  $L^{\leq 3}$  in lexicographic order are:

 $\lambda$ , aaa, aab, aba, abb, baa, bab, bba, bbb

(2 marks)

- (b) We observe that  $(w, w') \in R$  if and only if  $3 \not | \text{length}(w) \text{length}(w')$ .
  - (i) (abab, baba)? No because for all v: length(ababv) = length(babav), so whenever  $ababv \in L$ , we have  $babav \in L$  and vice versa.
  - (ii) (ab, abab)? Yes because for v = a:  $abv = aba \in L$  but  $ababv = ababa \notin L$ .
  - (iii)  $(\lambda, b)$ ? Yes because for  $v = \lambda$ :  $\lambda v = \lambda \in L$  but  $bv = b \notin L$ .
  - (iv)  $(\lambda, bb)$ ? Yes because for  $v = \lambda$ :  $\lambda v = \lambda \in L$  but  $bbv = bb \notin L$ .
  - (v)  $(\lambda, bbb)$ ? No because for all v: length $(\lambda v)$ -length(bbbv) = -3, so whenever  $\lambda v \in L$ , we have  $bbbv \in L$  and vice versa.

(1 mark each)

- (b)  $(w, w') \in S$  if and only if  $3|\operatorname{length}(w) \operatorname{length}(w')$ . (2 marks)
- (c) We need to show reflexivity (R), symmetry (S), and transitivity (T):
  - (R): Since length(w) length(w) = 0 and 3|0 we have that  $(w, w) \in S$  for all  $w \in \Sigma^*$ .

(3 marks)

- (S): Suppose  $(w, w') \in S$ . Then  $3|\operatorname{length}(w) \operatorname{length}(w')$ , i.e.  $\operatorname{length}(w) \operatorname{length}(w') = 3k$  for some  $k \in \mathbb{Z}$ . So  $\operatorname{length}(w') \operatorname{length}(w) = 3k'$  for some  $k' \in \mathbb{Z}$  (namely k' = -k) so  $3|\operatorname{length}(w') \operatorname{length}(w)$ . So  $(w', w) \in S$ .

  (3 marks)
- (T): Suppose  $(w, w') \in S$  and  $(w', w'') \in S$ . Then  $3|\operatorname{length}(w) \operatorname{length}(w')$  and  $3|\operatorname{length}(w') \operatorname{length}(w'')$ . Therefore,  $\operatorname{length}(w) \operatorname{length}(w') = 3k$  and  $\operatorname{length}(w') \operatorname{length}(w'') = 3k'$  for some  $k, k' \in \mathbb{Z}$ . Therefore

$$\begin{aligned} \operatorname{length}(w) &- \operatorname{length}(w'') \\ &= \operatorname{length}(w) - \operatorname{length}(w') + \operatorname{length}(w') - \operatorname{length}(w'') \\ &= 3k + 3k' \\ &= 3(k + k'). \end{aligned}$$

So  $3|\operatorname{length}(w) - \operatorname{length}(w'')$ , and so  $(w, w'') \in S$ .

(3 marks)

(d) S has three equivalence classes:  $[\lambda]$  (i.e. set of all words with length divisible by 3), [a] (set of all words with length 1 more than a multiple of 3), and [aa] (set of all words with length 2 more than a multiple of 3). (2 marks)