

Linked Lists

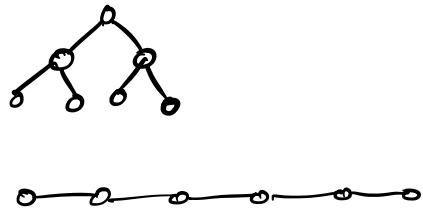
May 13, 2022

AGENDA:

- What are Linked Lists ?
- Why Linked Lists ?
- Creating a Linked List
- Basic operations
 - Traversal
 - Insertion
 - Deletion

Arrays

- * Contiguous.
- * Homogeneous.
- * Sequential / Linear



TC of accessing an element
let s^{th} element. $\rightarrow \frac{O(1)}{P}$

TC of insertion / deletion of element in array. $\rightarrow O(N)$

Linked List

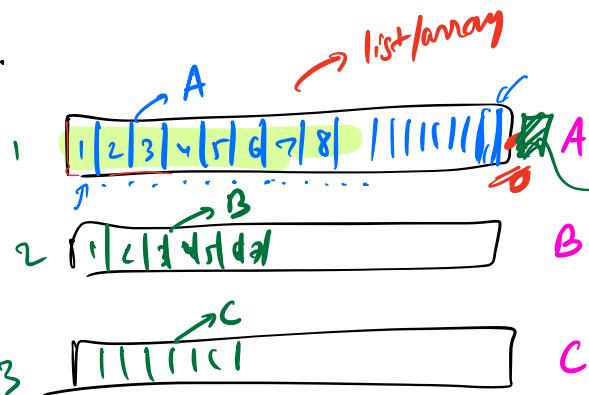
- * Linear / sequential (Same as arrays)



- * Not contiguous.

{ TC of accessing an element $\rightarrow O(N)$
** TC of insertion / deletion $\rightarrow O(1)$.

Library.



Paste a sticky note.

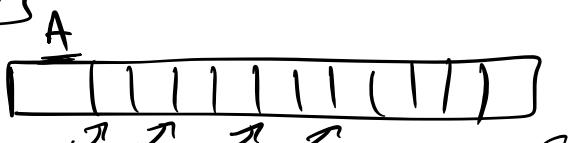
There is A2
bookshelf,
and it is located
at the Hall 2
of library.



so shelves.

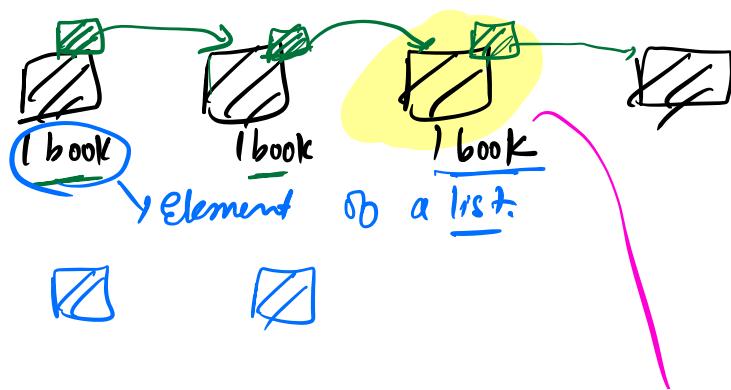


Linked list

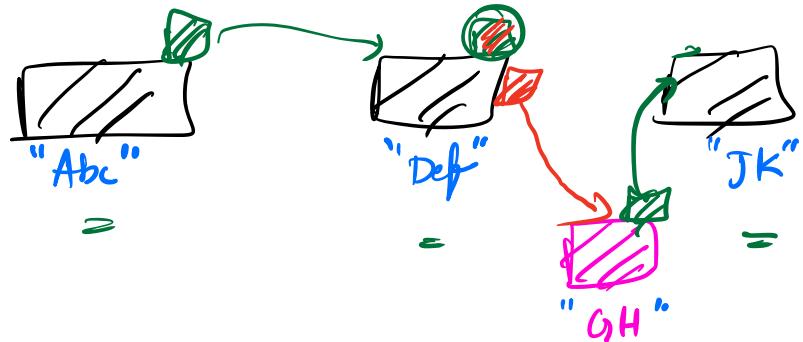


that got filled up.

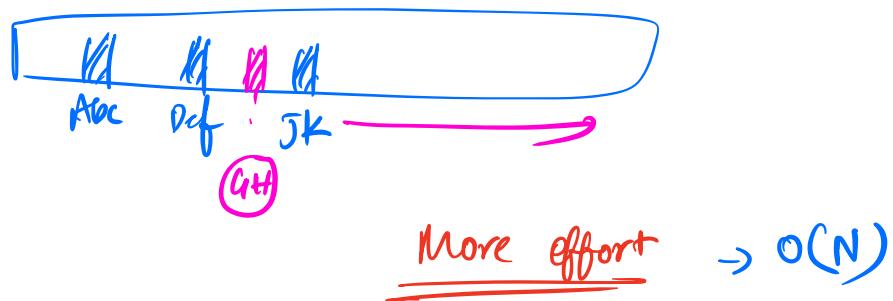
$O(N)$ → shift elements in an array.



① BOOK
Sticky note
 ↗ direction
Address.



Very min effort. $O(1)$ time.



linked list

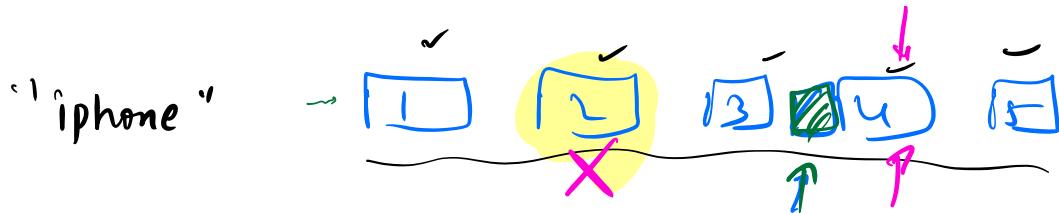
does not mean
linking different list.

means

linking elements in
form of a list.

Why linked list?

Search engine



"Sears"

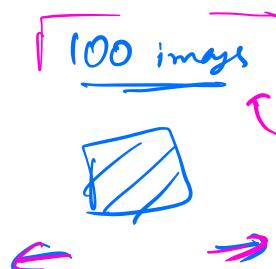


In insertion / deletion
needs to be efficient

Random Access with $O(1)$ times
is not needed.

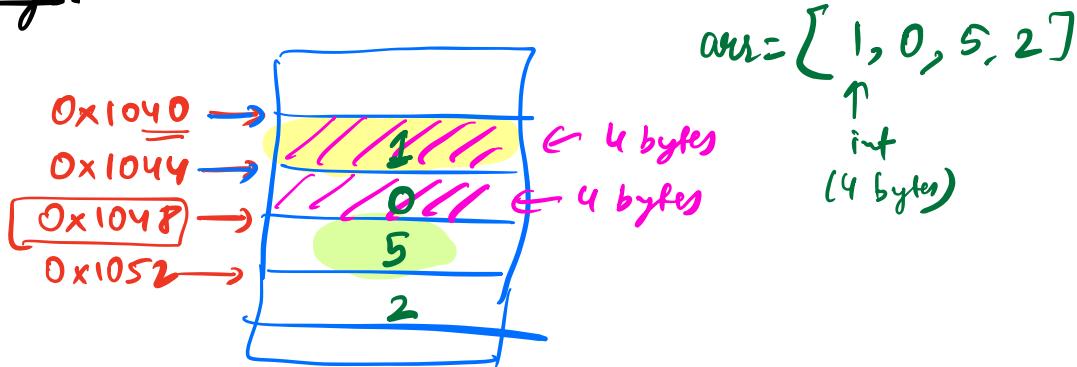
Access. $O(N)$ is justified.

** Image Viewer



What happens inside the memory?

Arrays.



$$arr[2] = ?$$

`arr[0]` will be stored at `0x1040`
`arr[2]` " " `0x1048`



$B + W \rightarrow$ 1st element
 $B + 2W \rightarrow$ 2nd element
 $B + 3W \rightarrow$ 3rd element.

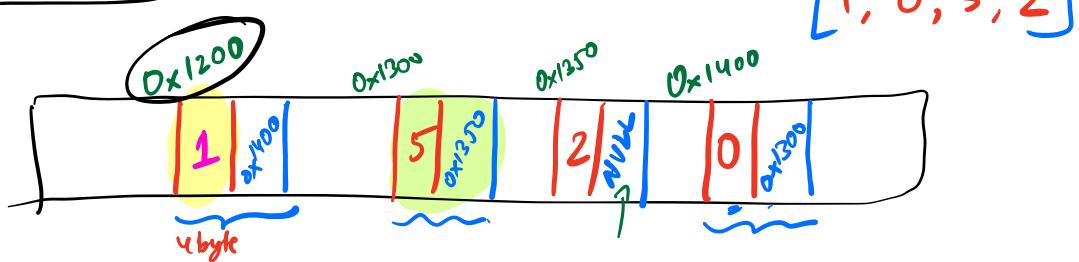
Given the address of memory,
TC to get the value stored in that
memory?

$\downarrow O(1)$ in case of RAM

RAM \rightarrow Random Access Memory

You can randomly access any memory location in constant time.

linked Lists.



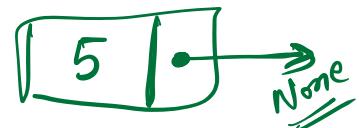
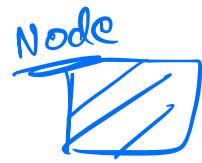
→ 0x1200 ?
 ↑
 head

arg = [1, 0, 5, 2]
↓
variable itself
takes care of
starting address

Node
(Val
Address of next node)

class Node :

def __init__(self, val):
 self.data = val
 self.next = None

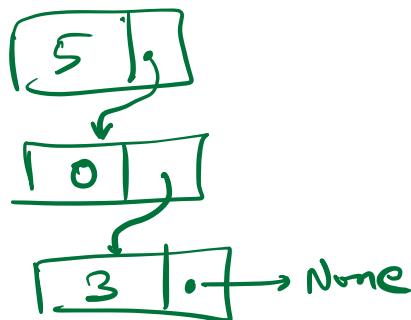


Create 3 Nodes and link all of it.

[5, 0, 3]

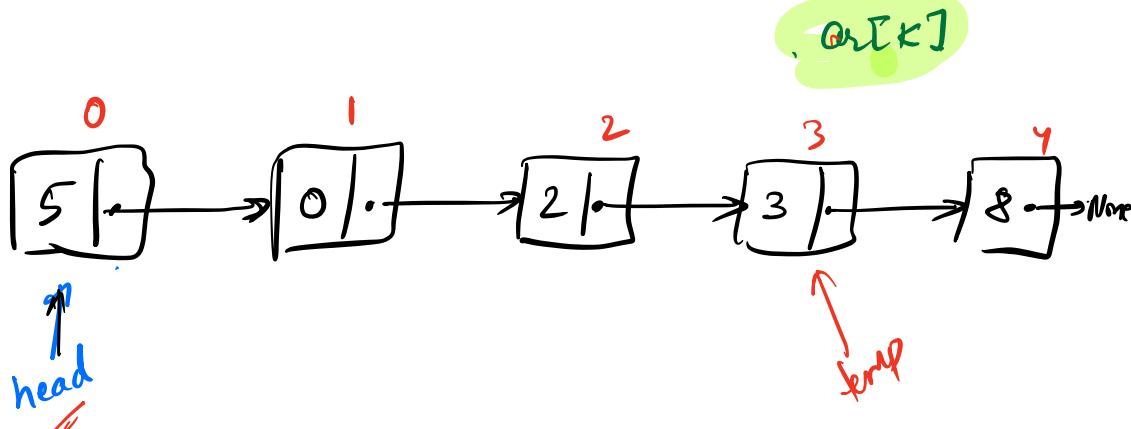
head = Node(5)
second-node = Node(0)
third-node = Node(3)

head.next = second-node
second-node.next = third-node
head = second .



* You will be given the 'head' node.

Q. Given a LL, (given the head), find the Kth element.



temp = temp.next

def findKthNode (head , k):

temp = head
for i in range(k):
 temp = temp.next

return temp.data

None.next

None has no attribute named 'next'.

Base cases

K = 10

if head is of a linked list:

```
def findKthNode(head, k):
```

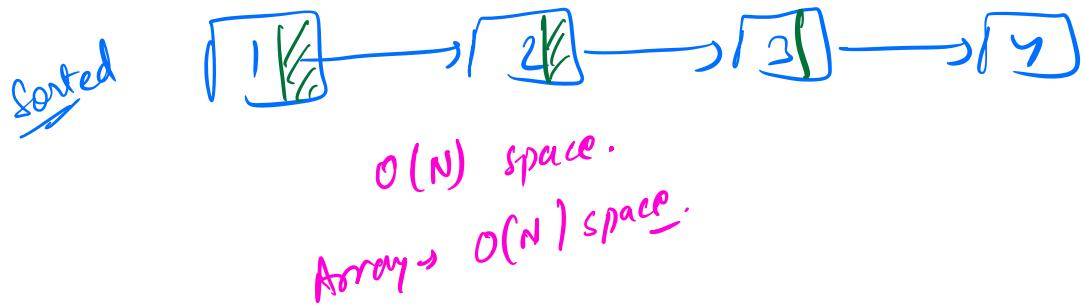
```
temp = head
for i in range(k):
    if temp.next == None:
        return None
    .
    .
    .
temp = temp.next
if temp == None:
    return None
return temp.data
```

TC: $O(N)$
SC: $O(1)$

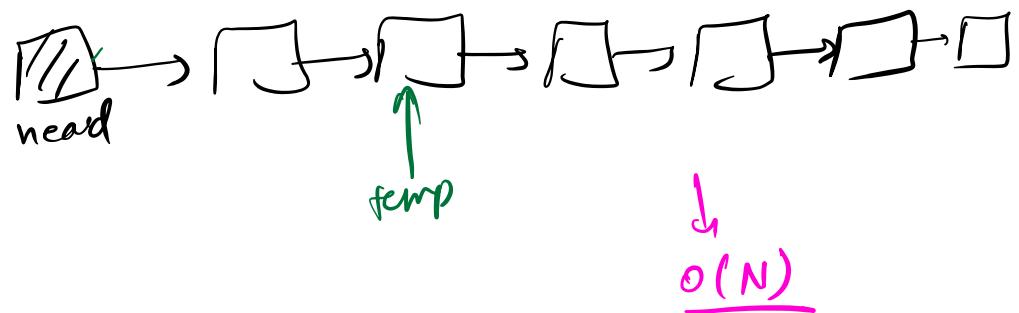
```
print(head) ←
print(type(head))
↓
class Node
```

Break till 10:22

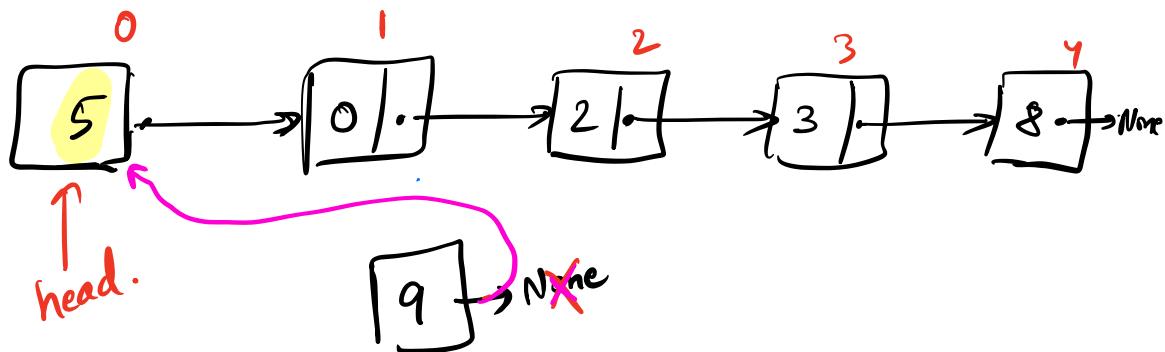
```
→ () {} [] ←
↓
imbuilt data structure
```



Find the length of linked list.



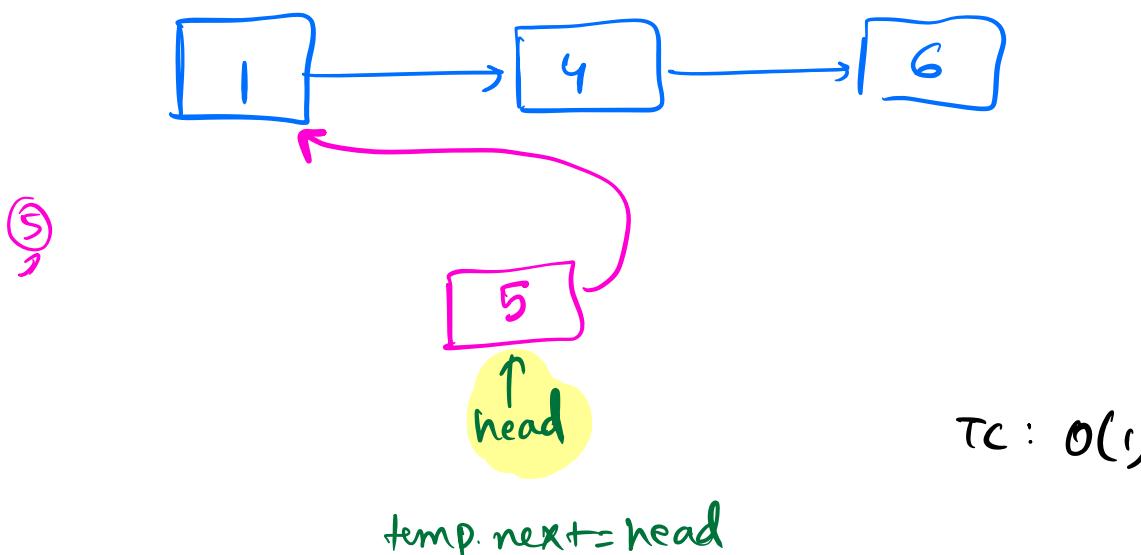
Insertion in a linked list.



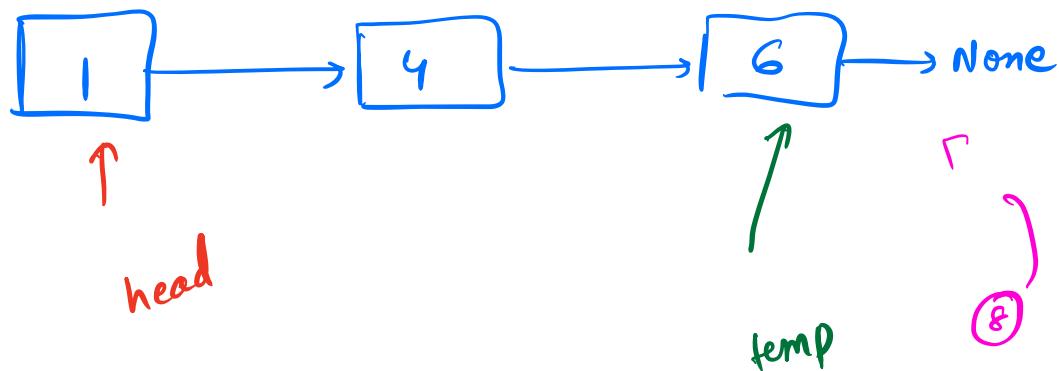
* Insertion at the beginning.
a value.

```
def insert_at_beginning( head, val ) :
```

```
    temp = Node( val )
    temp. next = head
    head = temp
```



* Insert at end.



def insert-at-end(head, val) :

x = Node(8) ←



temp = head

if head is None:

head = x return

while (temp.next is not None) :

temp = temp.next

temp.next = x

x.next = None × × ← Not needed.

Edge-case

What if the list is empty?

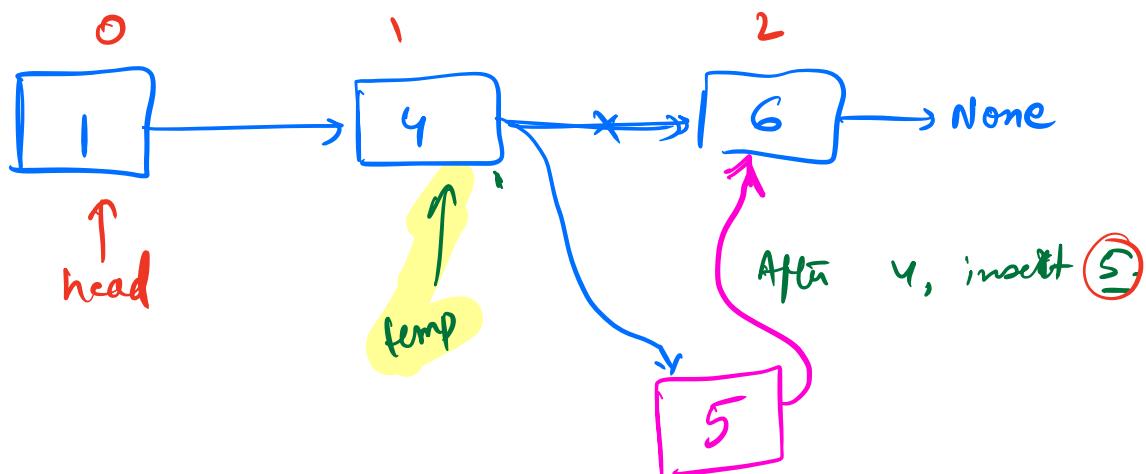
head = None.

}

\downarrow
TC: $O(N)$

tail. node $\rightarrow O(1)$ time

** Insertion at the middle.



$\left\{ \begin{array}{l} x = \text{Node}(5) \\ x.\text{next} = \text{temp}.\text{next} \\ \text{temp}.\text{next} = x \end{array} \right.$

def insert-at-middle(val, K, head):

$0 \leq K < N$

// Insert after
'K'. (0-indexed)

→ temp = find Kth Node(head, K)

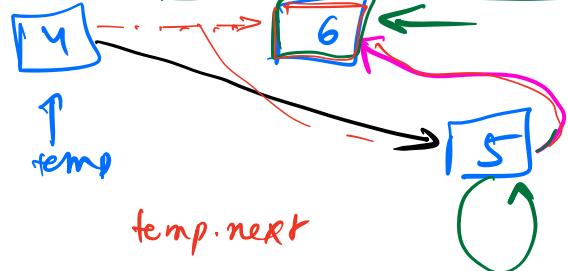
if temp == None:

insert-at-start(head, val)

$x = \text{Node}(5)$
 $x.\text{next} = \text{temp.next}$
 $\text{temp.next} = x$

(on)

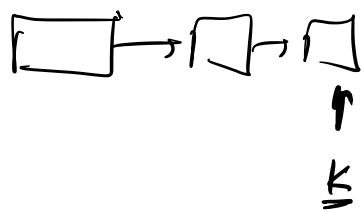
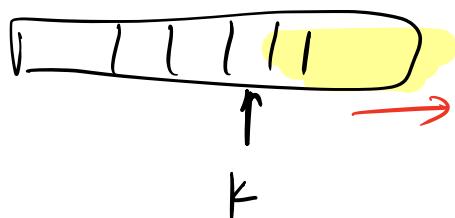
Wrong \times $x = \text{Node}(5)$ \times
 $\text{temp.next} = x$ \times
 $x.\text{next} = \text{temp.next}$



TC: $O(N)$

SC: $O(1)$

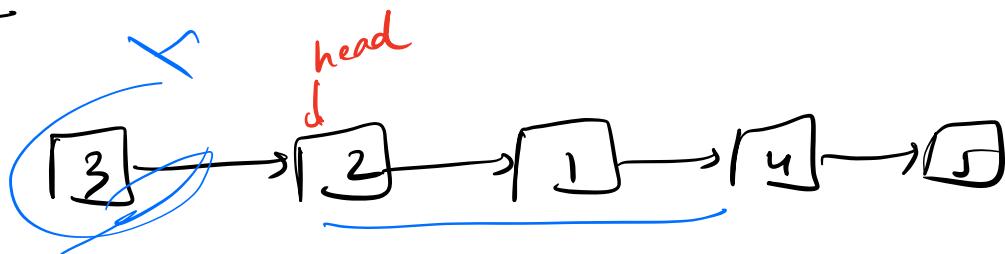
$O(1)$



2

0(1)

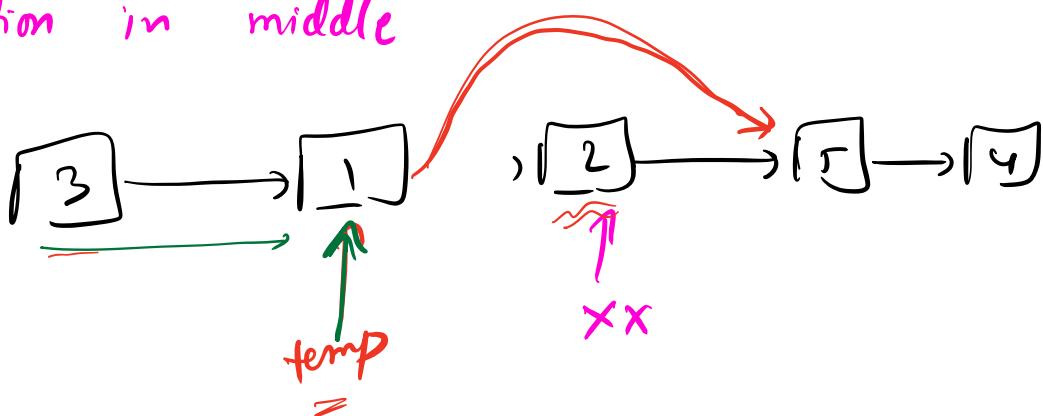
Deletion



- ① Delete the start node .

`head = head.next`

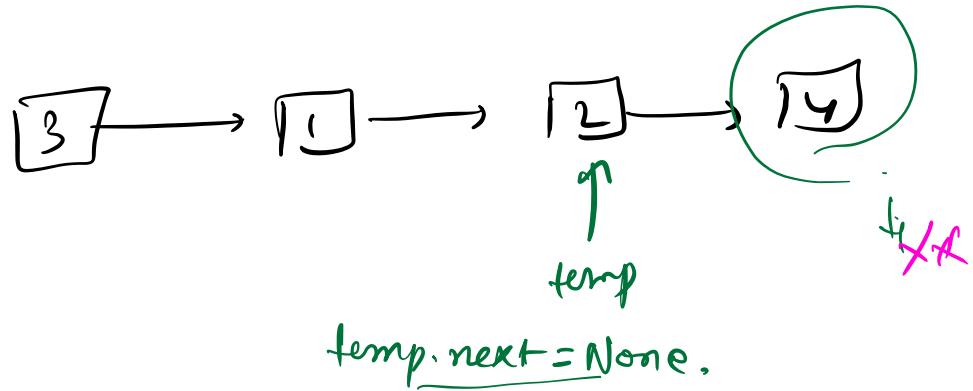
- ## ② Deletion in middle



`temp->next = temp->next->next`

Handle edge case.

- ③ Delete the last element.



while (temp.next.next is not None)

↑
Reach to the second last
node.

HW

Try implementing all of these
on your own.

Handle all the edge cases.