# MATH-329 Nonlinear optimization Homework 3: Constrained optimization

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# Part 1: Projections to cones and stopping criteria in constrained optimization.

#### Question 1

Since Q is non-empty, Let  $x_0 \in Q$ . Then by definition of the minimizer, we have:

$$\min_{x \in Q} \frac{1}{2} \|x - z\|^2 \le \frac{1}{2} \|x_0 - z\|^2$$

Inspired by this, let us define the space  $Q' \subseteq Q$  as the intersection of Q with the closed ball  $\bar{B}(\|x_0 - z\|, z)$  of center z and radius  $\|x_0 - z\|$ :

$$Q' = Q \cap \bar{B}(||x_0 - z||, z)$$

We have  $Q' \neq \emptyset$  since  $x_0 \in Q'$ . Moreover, Q' is closed since it is the intersection of two closed sets. Finally, Q' is bounded since it is contained in the closed ball  $\bar{B}(\|x_0 - z\|, z)$ . Therefore, Q' is compact. By Weierstrass, the function  $f(x) = \frac{1}{2}\|x - z\|^2$  attains its minimum on Q'. So the set:

$$\operatorname{Proj}_{Q'}(z) = \left\{ x \in Q' : \frac{1}{2} \|x - z\|^2 = \min_{y \in Q'} \frac{1}{2} \|y - z\|^2 \right\}$$

is well-defined and non-empty.

We want to show that  $\operatorname{Proj}_{Q'}(z) = \operatorname{Proj}_Q(z)$ . Let  $x \in Q \setminus Q'$ . Then:

$$||x-z|| > ||x_0-z|| \implies$$

$$\frac{1}{2}||x-z||^2 > \frac{1}{2}||x_0 - z||^2 \ge \min_{y \in Q'} \frac{1}{2}||y-z||^2$$

So the minimizer of  $\frac{1}{2}||y-z||^2$  on Q' is also the minimizer of  $\frac{1}{2}||y-z||^2$  on Q. Therefore,  $\operatorname{Proj}_{Q'}(z) = \operatorname{Proj}_Q(z)$ .

#### Question 2

Let  $\mathcal{E} := \mathbb{R}$ ;  $S := \{-1, 1\}$ ; z := 0. Then clearly S is closed (union of 2 singletons). Moreover, we have:

$$\frac{1}{2}\|z-s\|^2 = \frac{1}{2}\|s\|^2 = \frac{1}{2} \text{ for all } s \in S \implies$$

$$\mathrm{Proj}_S(z) = \{1, -1\}$$

$$\mathbf{Proj}_C(z) = \{0\} \implies z \in C^{\circ}$$
:

Assume  $\operatorname{Proj}_C(z) = \{0\}$ . This implies the point in C closest to z is the origin. So for any  $x \in C$ :

$$\frac{1}{2} ||x - z||^2 \ge \frac{1}{2} ||0 - z||^2$$
$$||x||^2 - 2\langle x, z \rangle + ||z||^2 \ge ||z||^2$$
$$||x||^2 - 2\langle x, z \rangle \ge 0$$

Now if  $\langle x, z \rangle > 0$ , for some x then we have:

$$\lambda := \frac{\langle x, z \rangle}{\|z\|^2} > 0 \implies \lambda x \in C$$

But then, plugging  $\lambda x$  in the inequality above, we get:

$$\begin{split} \lambda^2 \|x\|^2 - 2\lambda \langle x, z \rangle &\geq 0 \iff \\ \left(\frac{\langle x, z \rangle}{\|z\|^2}\right)^2 \|x\|^2 - 2\frac{\langle x, z \rangle}{\|z\|^2} \langle x, z \rangle &\geq 0 \iff \\ -\frac{\langle x, z \rangle^2}{\|x\|^2} &\geq 0 \end{split}$$

Which is clearly a contradiction. Therefore, we must have  $\langle x, z \rangle \leq 0$  for all  $x \in C$ . This implies  $z \in C^{\circ}$ , by definition.

$$z \in C^{\circ} \implies \mathbf{Proj}_{C}(z) = \{0\}$$
:

Assume  $z \in C^{\circ}$ . We want to show that for all  $x \in C \setminus \{0\}$  we have:

$$\frac{1}{2}||x-z||^2 > \frac{1}{2}||0-z||^2$$

By the above calculations this is equivalent to showing that for all  $x \in C \setminus \{0\}$  we have:

$$||x||^2 - 2\langle x, z \rangle > 0$$

But this is true since  $z \in C^{\circ}$ , so  $\langle x, z \rangle \leq 0$  for all  $x \in C$  and  $x \neq 0$ . Therefore, we have:

$$\operatorname{Proj}_C(z) = \{0\}$$

# Question 4

We know from class that  $x^* \in S$  is a stationary point of f if and only if  $-\nabla f(x^*) \in (T_{x^*}S)^{\circ}$ . By Question 3, this is equivalent to:

$$\text{Proj}_{T_{x^*}S}(-\nabla f(x^*)) = \{0\}$$

(a) 
$$v \in \mathbf{Proj}_C(z) \implies \langle v, z - v \rangle = 0$$
:

Let:

$$g: \mathcal{E} \to \mathbb{R}$$

$$x \mapsto \frac{1}{2} ||x - z||^2$$

Then g is differentiable and for  $h \in \mathbb{R}$  and  $v \in \mathcal{E}$  we have:

$$g(x + hv) = \frac{1}{2} \|x + hv - z\|^2 = \frac{1}{2} \|x - z\|^2 + h\langle v, x - z \rangle + \frac{h^2}{2} \|v\|^2$$
$$= g(x) + h\langle v, x - z \rangle + O(h^2) \implies$$
$$\nabla g(x) = x - z$$

If  $v \in \operatorname{Proj}_C(z)$  then v is a stationary point of g (as v is a global minimum of g). Therefore, we have:

$$\langle \nabla g(v), w \rangle \ge 0, \quad \forall \ w \in T_v C$$
  
 $\langle v - z, w \rangle \ge 0, \quad \forall \ w \in T_v C$ 

It is clear that if we show that  $v, -v \in T_vC$  we are done. But this is true since  $v \in C$  and C is a cone.

Let 
$$\left((1-\frac{1}{n})v\right)_{n\in\mathbb{N}^*}\subseteq C$$
,  $\left((1+\frac{1}{n})v\right)_{n\in\mathbb{N}^*}\subseteq C$  then:

$$\lim_{n \to \infty} \left( (1 - \frac{1}{n})v \right) = v, \quad \lim_{n \to \infty} \left( (1 + \frac{1}{n})v \right) = v$$

and:

$$\lim_{n \to \infty} \left( \frac{(1 - \frac{1}{n})v - v}{\frac{1}{n}} \right) = \lim_{n \to \infty} \left( \frac{-\frac{1}{n}v}{\frac{1}{n}} \right) = -v$$

$$\lim_{n \to \infty} \left( \frac{(1 + \frac{1}{n})v - v}{\frac{1}{n}} \right) = \lim_{n \to \infty} \left( \frac{\frac{1}{n}v}{\frac{1}{n}} \right) = v$$

Therefore, by the definition of the tangent cone  $v, -v \in T_vC$ . So we have:

$$\langle v - z, v \rangle \ge 0$$
  
 $\langle v - z, -v \rangle \ge 0 \implies$   
 $-\langle v - z, v \rangle \ge 0$ 

So we have:

$$\langle v, z - v \rangle = 0$$

(b) 
$$v_1, v_2 \in \mathbf{Proj}_C(z) \implies ||v_1|| = ||v_2||$$
:

By part (a), we have:

$$\langle v_1, z - v_1 \rangle = 0 \implies ||v_1||^2 = \langle v_1, v_1 \rangle = \langle v_1, z \rangle$$
$$\langle v_2, z - v_2 \rangle = 0 \implies ||v_2||^2 = \langle v_2, v_2 \rangle = \langle v_2, z \rangle$$

Since both are minimizers of  $\frac{1}{2}||x-z||^2$ , we have:

$$\frac{1}{2} \|v_1 - z\|^2 = \frac{1}{2} \|v_2 - z\|^2 \implies$$

$$\|v_1\|^2 - 2\langle v_1, z \rangle + \|z\|^2 = \|v_2\|^2 - 2\langle v_2, z \rangle + \|z\|^2 \implies$$

$$\|v_1\|^2 = \|v_2\|^2 \implies \|v_1\| = \|v_2\|$$

We present an example where the function  $q(x) = \|\operatorname{Proj}_{T_x S}(-\nabla f(x))\|$  is discontinuous on the set S. Consider the following:

#### Function f and Set S:

- Function f: Define  $f: \mathbb{R} \to \mathbb{R}$  as  $f(x) = -x^2$ . Its gradient is  $\nabla f(x) = f'(x) = -2x$ .
- Set S: Define  $S := [0, 1] \subseteq \mathbb{R}$ .

First, let's compute  $T_xS$  for all  $x \in S$ . We have:

- $x \in (1,0)$ . Then x is in the interior of S and  $T_xS = \mathbb{R}$ , by example 7.10. from the lecture notes.
- x = 0. Then 0 is on the boundary of S and  $T_0S = [0, +\infty)$ , by example 7.11. from the lecture notes (we can see 0 as the boundary for the half-space  $x \ge 0$  and locally around 0 the two spaces are the same. We conclude by the fact that the tangent cone is a local property, as any sequence  $x_k \to 0^+$  will be contained in (0,1), for k big enough).
- x = 1. Then 1 is on the boundary of S and  $T_1S = (-\infty, 0]$  (same reasoning as above, but with the half space  $x \le 1$ ).

Now, let's compute  $\operatorname{Proj}_{T_xS}(-\nabla f(x))$  for all  $x \in S$ . We have:

•  $x \in (0,1)$ . Then  $\operatorname{Proj}_{T_xS}(-\nabla f(x)) = \{-\nabla f(x)\}$ , since  $T_xS = \mathbb{R}$ . So:

$$q(x) = ||2x|| = 2x$$

• x = 0. Then  $\text{Proj}_{T_0 S}(-\nabla f(0)) = \{0\}$ , since  $T_0 S = [0, +\infty)$  and -f'(0) = 0. So:

$$q(0) = ||0|| = 0$$

• x = 1. Then  $\text{Proj}_{T_1S}(-\nabla f(1)) = \{0\}$ , since  $T_1S = (-\infty, 0]$  and -f'(1) = 2. So:

$$q(1) = ||0|| = 0$$

So:

$$q(x) = \begin{cases} 2x & \text{if } x \in [0, 1) \\ 0 & \text{if } x = 1 \end{cases}$$

Clearly, q is not continuous at x = 1.

#### Question 7

By example 7.14. from the lecture notes, we have:

$$T_x S = \{ v \in \mathbb{R} : \langle v, x \rangle = 0 \}, \quad \forall \ x \in S$$

For all  $x = (x_1, x_2) \in S$ , let  $x^{\perp} = (x_2, -x_1) \in S$ , then it is clear that:

$$\langle x, x^{\perp} \rangle = 0 \implies x^{\perp} \in T_x S$$

Moreover, by a simple argument over the dimensionality of  $T_xS$  and span $(x^{\perp})$ , we have:

$$T_x S = \operatorname{span}(x^{\perp})$$

Since  $T_xS$  is a sub-vector space of  $\mathbb{R}^2$ , of dimension 1 ( $\{x^{\perp}\}$  is an orthogonal basis), we have:

$$\operatorname{Proj}_{T_xS}(-\nabla f(x)) = \operatorname{Proj}_{\operatorname{span}(x^\perp)}(-\nabla f(x)) = \left\{ \frac{\langle -\nabla f(x), x^\perp \rangle}{\|x^\perp\|^2} x^\perp \right\} = \left\{ -\langle \nabla f(x), x^\perp \rangle x^\perp \right\}$$

So:

$$q(x) = \| -\langle \nabla f(x), x^{\perp} \rangle x^{\perp} \| = \| -\langle \nabla f(x), (x_2, -x_1) \rangle (x_2, -x_1) \| = |\langle \nabla f(x), (x_2, -x_1) \rangle|$$

Which is clearly a continuous map as it is a composition of continuous maps ( $\nabla f$  is continuous by assumption).

# Question 8

Consider  $\mathcal{E} = \mathbb{R}^n$  with a continuously differentiable function  $h : \mathbb{R}^n \to \mathbb{R}^p$  and  $S = \{x \in \mathbb{R}^n : h(x) = 0\}$ , assuming LICQ holds for all x in S.

(a)  $T_xS$ ?

From the lecture notes, we have that if LICQ holds for  $x \in S$  then:

$$T_x S = F_x S = \{ v \in \mathcal{E} : \langle \nabla h_i(x), v \rangle = 0; \forall i \in \{1, 2, \dots, p\} \}$$

(b)

Let  $H(x) \in \mathbb{R}^{p \times n}$  such that:

$$H(x) = \begin{pmatrix} \nabla h_1(x)^T \\ \nabla h_2(x)^T \\ \vdots \\ \nabla h_p(x)^T \end{pmatrix}$$

Then we can rewrite  $T_xS$  as:

$$T_x S = \{v \in \mathcal{E} : H(x)v = 0\} = \ker H(x)$$

As  $T_xS$  is a sub-vector space of  $\mathcal{E}$  of dimension n-p (since all the lines of H(x) are linearly independent), we know that the projection of  $z \in \mathcal{E}$  exists and is unique.

# Part 2: A Frank-Wolfe algorithm

# Question 1

The minimization problem under consideration is:

minimize 
$$\langle w, x \rangle$$
 subject to  $x \in S$ , (1)

where  $w = \nabla f(\bar{x})$  for some  $\bar{x} \in S$ .

To argue that this problem always has a solution, we consider the following points:

- Convexity and Compactness of S: The set S is assumed to be convex and compact in  $E = \mathbb{R}^n$ .
- Continuity of the Objective Function: The objective function  $\langle w, x \rangle$  is linear and therefore continuous.
- Existence of Minimizer: By the Extreme Value Theorem, a continuous function on a compact set attains its minimum. Therefore, the linear function  $\langle w, x \rangle$  attains a minimum over the compact and convex set S, ensuring the existence of a solution.

#### Question 2

To demonstrate that the minimization problem may have multiple solutions, consider the following example:

- Set S: Let S be a line segment in  $\mathbb{R}^2$  defined as  $S = \{(1, y) : 0 \le y \le 1\}$ .
- **Linear Function**: Consider a linear function  $\langle w, x \rangle$  with w = (0, 0). For any  $x \in S$ , we have  $\langle w, x \rangle = 0$ .

In this case, every point in S minimizes the function  $\langle w, x \rangle$  as the value is zero for all  $x \in S$ . Hence, the problem has multiple solutions, with every point in the set S being a solution.

#### Question 3

# Why is the Restriction $0 \le \eta_k \le 1$ Important?

- 1. **Feasibility**: The feasible set S is assumed to be convex. By convexity, for any  $x, y \in S$  and  $\lambda \in [0,1]$ , the convex combination  $(1-\lambda)x + \lambda y$  is also in S. In the algorithm, both  $x_k$  and  $s(x_k)$  are in S, so for  $\eta_k$  in [0,1], the updated point  $(1-\eta_k)x_k + \eta_k s(x_k)$  remains within S.
- 2. Convergence: The step size  $\eta_k$  controls the magnitude of the move towards the direction of minimization. Values of  $\eta_k$  outside the interval [0,1] can lead to overshooting or even divergence. Specifically,  $\eta_k > 1$  may cause the algorithm to take excessively large steps, while negative values of  $\eta_k$  would reverse the direction of the update, both hindering convergence.
- 3. Controlled Progress: The interval [0, 1] allows for dynamic adjustment of  $\eta_k$  to control the algorithm's progress. Smaller values of  $\eta_k$  can be used for cautious steps near the optimal solution, enhancing stability and precision.
- 4. Balance Between Exploration and Exploitation:  $\eta_k$  balances exploration of the feasible set S and exploitation towards the minimizer of the linearized function.  $\eta_k = 0$  implies no movement (pure exploitation), while  $\eta_k = 1$  means moving entirely towards the new direction (pure exploration). Intermediate values facilitate a balanced approach.

In conclusion, the restriction  $0 \le \eta_k \le 1$  in the Frank-Wolfe algorithm is essential for ensuring feasibility, convergence, controlled progress, and a balanced approach between exploration and exploitation.

We analyze four key inequalities (B1) to (B4) arising in the Frank-Wolfe algorithm under the assumptions that f is convex and continuously differentiable, and its gradient  $\nabla f$  is L-Lipschitz continuous.

# Inequality Analysis

**(B1)** 
$$f(x_{k+1}) - f(x_k) \le \nabla f(x_k)^{\top} (x_{k+1} - x_k) + \frac{L}{2} ||x_{k+1} - x_k||^2$$

This inequality stems from the first-order Taylor expansion and the Lipschitz continuity of  $\nabla f$ , bounding the error of the linear approximation.

$$(\mathbf{B2}) \le \eta_k \nabla f(x_k)^{\top} (s(x_k) - x_k) + \frac{L}{2} \eta_k^2 d_S^2$$

Using the update formula  $x_{k+1} = (1 - \eta_k)x_k + \eta_k s(x_k)$  and the definition of  $d_S$ , the diameter of S, this inequality bounds the change in f in terms of the diameter of S and step size  $\eta_k$ .

**(B3)** 
$$\leq \eta_k \nabla f(x_k)^{\top} (x^* - x_k) + \frac{L}{2} \eta_k^2 d_S^2$$

Given that  $s(x_k)$  minimizes the linear approximation over S, the inequality follows by comparing  $s(x_k)$  to any  $x^* \in S$ , including the optimal point.

**(B4)** 
$$\leq \eta_k(f(x^*) - f(x_k)) + \frac{L}{2}\eta_k^2 d_S^2$$

This follows from the convexity of f, which implies  $f(x^*) - f(x_k) \ge \nabla f(x_k)^\top (x^* - x_k)$ . Substituting this into (B3) yields (B4).

# Question 5

Given  $x_0 \in S$ , let  $x_1$  be produced by the Frank-Wolfe algorithm with  $\eta_0 = \frac{2}{0+2} = 1$ . We show that  $f(x_1) - f(x^*) \leq \frac{L}{2} d_S^2$ , where L is the Lipschitz constant of  $\nabla f$  and  $d_S$  is the diameter of S.

# Proof

- 1. The update rule for  $x_{k+1}$  in the Frank-Wolfe algorithm is  $x_{k+1} = (1 \eta_k)x_k + \eta_k s(x_k)$ . For k = 0, this becomes  $x_1 = s(x_0)$ .
- 2. By the L-Lipschitz continuity of  $\nabla f$ , we have

$$f(y) \le f(x) + \nabla f(x)^{\top} (y - x) + \frac{L}{2} ||y - x||^2$$
 for all  $x, y \in S$ .

3. Setting  $x = x_1$  and  $y = x^*$ , we get

$$f(x^*) \le f(x_1) + \nabla f(x_1)^{\top} (x^* - x_1) + \frac{L}{2} ||x^* - x_1||^2.$$

4. Rearranging, we obtain

$$f(x_1) - f(x^*) \ge -\nabla f(x_1)^{\top} (x^* - x_1) - \frac{L}{2} ||x^* - x_1||^2.$$

5. Since  $x_1$  and  $x^*$  are in S and  $||x^* - x_1||^2 \le d_S^2$ , we have

$$f(x_1) - f(x^*) \le \frac{L}{2} d_S^2.$$

Thus, after the first iteration with  $\eta_0 = 1$ , the function value at  $x_1$  is within  $\frac{L}{2}d_S^2$  of the optimal value  $f(x^*)$ .

We prove that for the Frank-Wolfe algorithm with step sizes  $\eta_k = \frac{2}{k+2}$ , the inequality  $f(x_k) - f(x^*) \le \frac{2Ld_S^2}{k+2}$  holds for all  $k \ge 1$ .

# **Proof by Induction**

# Base Case (k = 1)

From the previous analysis, we have  $f(x_1) - f(x^*) \leq \frac{Ld_S^2}{2}$ , which satisfies the inequality for k = 1, as  $\frac{2Ld_S^2}{1+2} = \frac{Ld_S^2}{2}$ .

# **Inductive Step**

Assume the inequality holds for some  $k \geq 1$ :

$$f(x_k) - f(x^*) \le \frac{2Ld_S^2}{k+2}$$

We need to show it holds for k + 1:

$$f(x_{k+1}) - f(x^*) \le \frac{2Ld_S^2}{k+3}$$

Using the inequality (B4) from the algorithm:

$$f(x_{k+1}) - f(x_k) \le \eta_k (f(x^*) - f(x_k)) + \frac{L}{2} \eta_k^2 d_S^2$$

where  $\eta_k = \frac{2}{k+2}$ . Substituting and rearranging gives:

$$f(x_{k+1}) - f(x^*) \le \left(1 - \frac{2}{k+2}\right) \frac{2Ld_S^2}{k+2} + \frac{L}{2} \left(\frac{2}{k+2}\right)^2 d_S^2$$

$$= \frac{kLd_S^2}{(k+2)^2} + \frac{2Ld_S^2}{(k+2)^2}$$

$$= \frac{(k+2)Ld_S^2}{(k+2)^2}$$

$$= \frac{Ld_S^2}{k+2}$$

Using  $\frac{2}{k+2} \leq \frac{2}{k+3}$ , we get:

$$f(x_{k+1}) - f(x^*) \le \frac{2Ld_S^2}{k+3}$$

Thus, by induction, the inequality holds for all  $k \geq 1$ .

#### Question 7

We show that the simplex  $\Delta_n = \{x \in \mathbb{R}^n | x_1, \dots, x_n \geq 0 \text{ and } x_1 + \dots + x_n = 1\}$  is convex, compact, and non-empty.

#### Convexity

A set is convex if for any two points in the set, the line segment between them is also in the set. For  $x,y\in\Delta_n$  and  $\lambda\in[0,1]$ , consider  $z=(1-\lambda)x+\lambda y$ . Since  $x_i,y_i\geq 0$ , each component  $z_i=(1-\lambda)x_i+\lambda y_i\geq 0$ . Also,  $\sum_{i=1}^n z_i=(1-\lambda)\sum_{i=1}^n x_i+\lambda\sum_{i=1}^n y_i=1$ . Hence,  $z\in\Delta_n$ , proving convexity.

#### Compactness

A set is compact if it is closed and bounded.  $\Delta_n$  is closed as it contains all its limit points. It is bounded because for all  $x \in \Delta_n$ ,  $0 \le x_i \le 1$  and  $\sum_{i=1}^n x_i = 1$ . Therefore,  $\Delta_n$  is compact.

#### Non-emptiness

 $\Delta_n$  is non-empty as it contains at least the point x = (1, 0, ..., 0), which satisfies  $x_i \ge 0$  and  $\sum_{i=1}^n x_i = 1$ .

In conclusion, the simplex  $\Delta_n$  is convex, compact, and non-empty.

#### Question 8

Minimization Problem on the Simplex  $\Delta_n$ 

Given a vector  $w \in \mathbb{R}^n$ , we consider the problem of minimizing  $\langle w, x \rangle$  subject to  $x \in \Delta_n$ , where  $\Delta_n$  is the simplex in  $\mathbb{R}^n$ .

#### Minimum of the Problem

The problem is formulated as:

minimize  $\langle w, x \rangle$  subject to  $x \in \Delta_n$ .

# Strategy to Attain the Smallest Value

To minimize  $\langle w, x \rangle$ , we allocate the entire weight to the component of x corresponding to the smallest component of w. Let  $i^* = \arg\min_i w_i$ . The minimizing vector x is such that  $x_{i^*} = 1$  and  $x_i = 0$  for all  $i \neq i^*$ .

#### Computational Complexity

The computational complexity of finding this solution is O(n), as it requires a linear scan to find the minimum component of the vector w. The minimizing x is then directly obtained from the index of this minimum component.

#### Question 9

Given the optimization problem  $\min_{x \in \Delta_n} f(x)$  with  $f(x) = \frac{1}{2} ||Ax - b||^2$ , we analyze whether this problem always has a solution and if the solution is unique.

#### Existence of a Solution

- Convexity of f(x): The function  $f(x) = \frac{1}{2} ||Ax b||^2$  is convex as it is the composition of a convex function (norm squared) with an affine function.
- Convexity and Compactness of  $\Delta_n$ : The simplex  $\Delta_n$  is convex and compact.
- Existence of Solution: A convex function over a compact convex set attains its minimum. Hence, the problem always has at least one solution.

# Uniqueness of the Solution

- Strict Convexity: Strict convexity is necessary for uniqueness. The function f(x) is strictly convex if the matrix A has full column rank. However, with m < n, A cannot have full column rank.
- Multiple Solutions: When A does not have full column rank, there can be multiple minimizers of f(x), due to directions in which A is not injective.

In conclusion, the optimization problem always has a solution, but it does not always have a unique solution due to the potential rank deficiency of A.

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#### Question 10

Gradient of the Function for Frank-Wolfe Algorithm

We derive the gradient of the function  $f(x) = \frac{1}{2} ||Ax - b||^2$  for applying the Frank-Wolfe algorithm to the problem  $\min_{x \in \Delta_n} f(x)$ .

The function f(x) is given by:

$$f(x) = \frac{1}{2} ||Ax - b||^2 = \frac{1}{2} (Ax - b)^{\top} (Ax - b).$$

Differentiating f(x) with respect to x using matrix calculus, we obtain the gradient of f(x):

$$\nabla f(x) = A^{\top} (Ax - b).$$

This gradient represents the direction of the steepest ascent at any point x for the function f(x) and is essential for determining the search direction in each iteration of the Frank-Wolfe algorithm.

#### Question 11

We analyze the line-search function  $g(\eta) = f((1 - \eta)x + \eta y)$  where  $x, y \in \Delta_n$  and  $f(x) = \frac{1}{2}||Ax - b||^2$  to determine the optimal values of  $\eta \in [0, 1]$ .

Expression for  $g(\eta)$ 

$$g(\eta) = \frac{1}{2} ||A((1 - \eta)x + \eta y) - b||^2$$

# Optimal Value of $\eta$

To find the optimal  $\eta$ , we differentiate  $g(\eta)$  with respect to  $\eta$  and set the derivative to zero:

$$g'(\eta) = \frac{d}{d\eta} \frac{1}{2} ||A((1-\eta)x + \eta y) - b||^2$$

$$= (A((1-\eta)x + \eta y) - b)^{\mathsf{T}}A(y-x)$$

Setting  $g'(\eta) = 0$  gives:

$$(A((1-\eta)x + \eta y) - b)^{\mathsf{T}}A(y-x) = 0$$

Solving this equation for  $\eta$  gives the optimal value.

# Closed-Form Formula

A closed-form expression for  $\eta$  depends on the specific structure of A, b, x, and y. Without additional assumptions, the exact solution might be complex or not directly obtainable.

Question 12

Question 13

Question 14