

# MATH-329 Nonlinear optimization Homework 3:

## Constrained optimization

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### Part 1 : Projections to cones and stopping criteria in constrained optimization.

#### Question 1

Since  $Q$  is non-empty, Let  $x_0 \in Q$ . Then by definition of the minimizer, we have:

$$\min_{x \in Q} \frac{1}{2} \|x - z\|^2 \leq \frac{1}{2} \|x_0 - z\|^2$$

Inspired by this, let us define the space  $Q' \subseteq Q$  as the intersection of  $Q$  with the closed ball  $\bar{B}(\|x_0 - z\|, z)$  of center  $z$  and radius  $\|x_0 - z\|$ :

$$Q' = Q \cap \bar{B}(\|x_0 - z\|, z)$$

We have  $Q' \neq \emptyset$  since  $x_0 \in Q'$ . Moreover,  $Q'$  is closed since it is the intersection of two closed sets. Finally,  $Q'$  is bounded since it is contained in the closed ball  $\bar{B}(\|x_0 - z\|, z)$ . Therefore,  $Q'$  is compact. By Weierstrass, the function  $f(x) = \frac{1}{2} \|x - z\|^2$  attains its minimum on  $Q'$ . So the set:

$$\text{Proj}_{Q'}(z) = \left\{ x \in Q' : \frac{1}{2} \|x - z\|^2 = \min_{y \in Q'} \frac{1}{2} \|y - z\|^2 \right\}$$

is well-defined and non-empty.

We want to show that  $\text{Proj}_{Q'}(z) = \text{Proj}_Q(z)$ . Let  $x \in Q \setminus Q'$ . Then:

$$\|x - z\| > \|x_0 - z\| \implies$$

$$\frac{1}{2} \|x - z\|^2 > \frac{1}{2} \|x_0 - z\|^2 \geq \min_{y \in Q'} \frac{1}{2} \|y - z\|^2$$

So the minimizer of  $\frac{1}{2} \|y - z\|^2$  on  $Q'$  is also the minimizer of  $\frac{1}{2} \|y - z\|^2$  on  $Q$ . Therefore,  $\text{Proj}_{Q'}(z) = \text{Proj}_Q(z)$ .

#### Question 2

Let  $\mathcal{E} := \mathbb{R}$ ;  $S := \{-1, 1\}$ ;  $z := 0$ . Then clearly  $S$  is closed (union of 2 singletons). Moreover, we have:

$$\frac{1}{2} \|z - s\|^2 = \frac{1}{2} \|s\|^2 = \frac{1}{2} \text{ for all } s \in S \implies$$

$$\text{Proj}_S(z) = \{1, -1\}$$

### Question 3

**Proj<sub>C</sub>(z) = {0} ⇒ z ∈ C<sup>o</sup>:**

Assume Proj<sub>C</sub>(z) = {0}. This implies the point in C closest to z is the origin. So for any x ∈ C:

$$\begin{aligned}\frac{1}{2}\|x - z\|^2 &\geq \frac{1}{2}\|0 - z\|^2 \\ \|x\|^2 - 2\langle x, z \rangle + \|z\|^2 &\geq \|z\|^2 \\ \|x\|^2 - 2\langle x, z \rangle &\geq 0\end{aligned}$$

Now if ⟨x, z⟩ > 0, for some x then we have:

$$\lambda := \frac{\langle x, z \rangle}{\|z\|^2} > 0 \implies \lambda x \in C$$

But then, plugging λx in the inequality above, we get:

$$\begin{aligned}\lambda^2\|x\|^2 - 2\lambda\langle x, z \rangle &\geq 0 \iff \\ \left(\frac{\langle x, z \rangle}{\|z\|^2}\right)^2 \|x\|^2 - 2\frac{\langle x, z \rangle}{\|z\|^2} \langle x, z \rangle &\geq 0 \iff \\ -\frac{\langle x, z \rangle^2}{\|x\|^2} &\geq 0\end{aligned}$$

Which is clearly a contradiction. Therefore, we must have ⟨x, z⟩ ≤ 0 for all x ∈ C. This implies z ∈ C<sup>o</sup>, by definition.

**z ∈ C<sup>o</sup> ⇒ Proj<sub>C</sub>(z) = {0}:**

Assume z ∈ C<sup>o</sup>. We want to show that for all x ∈ C \ {0} we have:

$$\frac{1}{2}\|x - z\|^2 > \frac{1}{2}\|0 - z\|^2$$

By the above calculations this is equivalent to showing that for all x ∈ C \ {0} we have:

$$\|x\|^2 - 2\langle x, z \rangle > 0$$

But this is true since z ∈ C<sup>o</sup>, so ⟨x, z⟩ ≤ 0 for all x ∈ C and x ≠ 0. Therefore, we have:

$$\text{Proj}_C(z) = \{0\}$$

### Question 4

We know from class that x\* ∈ S is a stationary point of f if and only if -∇f(x\*) ∈ (T<sub>x\*</sub>S)<sup>o</sup>. By Question 3, this is equivalent to:

$$\text{Proj}_{T_{x^*}S}(-\nabla f(x^*)) = \{0\}$$

## Question 5

(a)  $v \in \mathbf{Proj}_C(z) \implies \langle v, z - v \rangle = 0$ :

Let:

$$\begin{aligned} g : \mathcal{E} &\rightarrow \mathbb{R} \\ x &\mapsto \frac{1}{2} \|x - z\|^2 \end{aligned}$$

Then  $g$  is differentiable and for  $h \in \mathbb{R}$  and  $v \in \mathcal{E}$  we have:

$$\begin{aligned} g(x + hv) &= \frac{1}{2} \|x + hv - z\|^2 = \frac{1}{2} \|x - z\|^2 + h \langle v, x - z \rangle + \frac{h^2}{2} \|v\|^2 \\ &= g(x) + h \langle v, x - z \rangle + O(h^2) \implies \\ \nabla g(x) &= x - z \end{aligned}$$

If  $v \in \mathbf{Proj}_C(z)$  then  $v$  is a stationary point of  $g$  (as  $v$  is a global minimum of  $g$ ). Therefore, we have:

$$\begin{aligned} \langle \nabla g(v), w \rangle &\geq 0, \quad \forall w \in T_v C \\ \langle v - z, w \rangle &\geq 0, \quad \forall w \in T_v C \end{aligned}$$

It is clear that if we show that  $v, -v \in T_v C$  we are done. But this is true since  $v \in C$  and  $C$  is a cone.

Let  $((1 - \frac{1}{n})v)_{n \in \mathbb{N}^*} \subseteq C$ ,  $((1 + \frac{1}{n})v)_{n \in \mathbb{N}^*} \subseteq C$  then:

$$\lim_{n \rightarrow \infty} \left( (1 - \frac{1}{n})v \right) = v, \quad \lim_{n \rightarrow \infty} \left( (1 + \frac{1}{n})v \right) = v$$

and:

$$\begin{aligned} \lim_{n \rightarrow \infty} \left( \frac{(1 - \frac{1}{n})v - v}{\frac{1}{n}} \right) &= \lim_{n \rightarrow \infty} \left( \frac{-\frac{1}{n}v}{\frac{1}{n}} \right) = -v \\ \lim_{n \rightarrow \infty} \left( \frac{(1 + \frac{1}{n})v - v}{\frac{1}{n}} \right) &= \lim_{n \rightarrow \infty} \left( \frac{\frac{1}{n}v}{\frac{1}{n}} \right) = v \end{aligned}$$

Therefore, by the definition of the tangent cone  $v, -v \in T_v C$ . So we have:

$$\begin{aligned} \langle v - z, v \rangle &\geq 0 \\ \langle v - z, -v \rangle &\geq 0 \implies \\ -\langle v - z, v \rangle &\geq 0 \end{aligned}$$

So we have:

$$\langle v, z - v \rangle = 0$$

(b)  $v_1, v_2 \in \mathbf{Proj}_C(z) \implies \|v_1\| = \|v_2\|$ :

By part (a), we have:

$$\begin{aligned} \langle v_1, z - v_1 \rangle &= 0 \implies \|v_1\|^2 = \langle v_1, v_1 \rangle = \langle v_1, z \rangle \\ \langle v_2, z - v_2 \rangle &= 0 \implies \|v_2\|^2 = \langle v_2, v_2 \rangle = \langle v_2, z \rangle \end{aligned}$$

Since both are minimizers of  $\frac{1}{2} \|x - z\|^2$ , we have:

$$\begin{aligned} \frac{1}{2} \|v_1 - z\|^2 &= \frac{1}{2} \|v_2 - z\|^2 \implies \\ \|v_1\|^2 - 2\langle v_1, z \rangle + \|z\|^2 &= \|v_2\|^2 - 2\langle v_2, z \rangle + \|z\|^2 \implies \\ \|v_1\|^2 &= \|v_2\|^2 \implies \|v_1\| = \|v_2\| \end{aligned}$$

## Question 6

We present an example where the function  $q(x) = \|\text{Proj}_{T_x S}(-\nabla f(x))\|$  is discontinuous on the set  $S$ . Consider the following:

**Function  $f$  and Set  $S$ :**

- Function  $f$ : Define  $f : \mathbb{R} \rightarrow \mathbb{R}$  as  $f(x) = -x^2$ . Its gradient is  $\nabla f(x) = f'(x) = -2x$ .
- Set  $S$ : Define  $S := [0, 1] \subseteq \mathbb{R}$ .

First, let's compute  $T_x S$  for all  $x \in S$ . We have:

- $x \in (1, 0)$ . Then  $x$  is in the interior of  $S$  and  $T_x S = \mathbb{R}$ , by example 7.10. from the lecture notes.
- $x = 0$ . Then 0 is on the boundary of  $S$  and  $T_0 S = [0, +\infty)$ , by example 7.11. from the lecture notes (we can see 0 as the boundary for the half-space  $x \geq 0$  and locally around 0 the two spaces are the same. We conclude by the fact that the tangent cone is a local property, as any sequence  $x_k \rightarrow 0^+$  will be contained in  $(0, 1)$ , for  $k$  big enough).
- $x = 1$ . Then 1 is on the boundary of  $S$  and  $T_1 S = (-\infty, 0]$  (same reasoning as above, but with the half space  $x \leq 1$ ).

Now, let's compute  $\text{Proj}_{T_x S}(-\nabla f(x))$  for all  $x \in S$ . We have:

- $x \in (0, 1)$ . Then  $\text{Proj}_{T_x S}(-\nabla f(x)) = \{-\nabla f(x)\}$ , since  $T_x S = \mathbb{R}$ . So:

$$q(x) = \|2x\| = 2x$$

- $x = 0$ . Then  $\text{Proj}_{T_0 S}(-\nabla f(0)) = \{0\}$ , since  $T_0 S = [0, +\infty)$  and  $-f'(0) = 0$ . So:

$$q(0) = \|0\| = 0$$

- $x = 1$ . Then  $\text{Proj}_{T_1 S}(-\nabla f(1)) = \{0\}$ , since  $T_1 S = (-\infty, 0]$  and  $-f'(1) = 2$ . So:

$$q(1) = \|0\| = 0$$

So:

$$q(x) = \begin{cases} 2x & \text{if } x \in [0, 1) \\ 0 & \text{if } x = 1 \end{cases}$$

Clearly,  $q$  is not continuous at  $x = 1$ .

## Question 7

By example 7.14. from the lecture notes, we have:

$$T_x S = \{v \in \mathbb{R} : \langle v, x \rangle = 0\}, \quad \forall x \in S$$

For all  $x = (x_1, x_2) \in S$ , let  $x^\perp = (x_2, -x_1) \in S$ , then it is clear that:

$$\langle x, x^\perp \rangle = 0 \implies x^\perp \in T_x S$$

Moreover, by a simple argument over the dimensionality of  $T_x S$  and  $\text{span}(x^\perp)$ , we have:

$$T_x S = \text{span}(x^\perp)$$

Since  $T_x S$  is a sub-vector space of  $\mathbb{R}^2$ , of dimension 1 ( $\{x^\perp\}$  is an orthogonal basis), we have:

$$\text{Proj}_{T_x S}(-\nabla f(x)) = \text{Proj}_{\text{span}(x^\perp)}(-\nabla f(x)) = \left\{ \frac{\langle -\nabla f(x), x^\perp \rangle}{\|x^\perp\|^2} x^\perp \right\} = \{-\langle \nabla f(x), x^\perp \rangle x^\perp\}$$

So:

$$q(x) = \| -\langle \nabla f(x), x^\perp \rangle x^\perp \| = \| -\langle \nabla f(x), (x_2, -x_1) \rangle (x_2, -x_1) \| = |\langle \nabla f(x), (x_2, -x_1) \rangle|$$

Which is clearly a continuous map as it is a composition of continuous maps ( $\nabla f$  is continuous by assumption).

### Question 8

Consider  $\mathcal{E} = \mathbb{R}^n$  with a continuously differentiable function  $h : \mathbb{R}^n \rightarrow \mathbb{R}^p$  and  $S = \{x \in \mathbb{R}^n : h(x) = 0\}$ , assuming LICQ holds for all  $x$  in  $S$ .

(a)  $T_x S$ ?

From the lecture notes, we have that if LICQ holds for  $x \in S$  then:

$$T_x S = F_x S = \{v \in \mathcal{E} : \langle \nabla h_i(x), v \rangle = 0; \forall i \in \{1, 2, \dots, p\}\}$$

(b)

Let  $H(x) \in \mathbb{R}^{p \times n}$  such that:

$$H(x) = \begin{pmatrix} \nabla h_1(x)^T \\ \nabla h_2(x)^T \\ \vdots \\ \nabla h_p(x)^T \end{pmatrix}$$

Then we can rewrite  $T_x S$  as:

$$T_x S = \{v \in \mathcal{E} : H(x)v = 0\} = \ker H(x)$$

As  $T_x S$  is a sub-vector space of  $\mathcal{E}$  of dimension  $n - p$  (since all the lines of  $H(x)$  are linearly independent), we know that the projection of  $z \in \mathcal{E}$  exists and is unique.

By SVD, there exist  $V \in \mathbb{R}^{n \times n}$  and  $U \in \mathbb{R}^{p \times p}$  orthogonal matrices and  $D \in \mathbb{R}^{p \times n}$  such that:

$$UH(x)V^T = D$$

$$D = \begin{pmatrix} \sigma_1 & & & \\ & \sigma_2 & & \\ & & \ddots & \\ & & & \sigma_p & & O_{p \times (n-p)} \end{pmatrix}$$

where  $\sigma_1, \dots, \sigma_p$  are the singular values of  $H(x)$ .

Then, since  $U$  and  $V$  are orthogonal, we have:

$$v \in T_x S \iff H(x)v = 0 \iff UH(x)V^T Vv = 0 \iff DVv = 0 \iff Vv \in \ker D$$

and, as  $V$  is orthogonal, we have:

$$\frac{1}{2}\|z - v\|^2 = \frac{1}{2}\|Vz - Vv\|^2$$

So, we have:

$$\text{Proj}_{T_x S}(z) = \text{Proj}_{\ker D}(Vz)$$

## Part 2 : A Frank-Wolfe algorithm

### Question 1

Let  $g : S \rightarrow \mathbb{R}$  be defined by  $g(x) = \langle w, x \rangle$  for a fixed  $w$  and for all  $x \in S$ . Since  $g$  is a continuous function (as the inner product of two vectors in  $\mathbb{R}^n$  is continuous) and since  $S$  is compact, the Weierstrass Extreme Value Theorem guarantees that  $g$  attains its minimum and maximum on  $S$ .

### Question 2

Let  $S = [-1, 1] \times \{0\}$  and  $w = (0, 1)$ . Then, for any  $x \in S$ , we have  $\langle w, x \rangle = 0$ . Thus, every point in  $S$  minimizes the function  $\langle w, x \rangle$ , leading to multiple solutions.

### Question 3

**Why is the Restriction  $0 \leq \eta_k \leq 1$  Important?**

In the Frank-Wolfe algorithm, enforcing  $0 \leq \eta_k \leq 1$  ensures  $x_{k+1}$  remains within the feasible set. Since  $S$  is convex, the convex combination  $(1 - \eta_k)x_k + \eta_k s(x_k)$  lies within  $S$  for any  $\eta_k \in [0, 1]$ . Without this restriction, there's no guarantee that  $x_{k+1}$  stays within  $S$ , possibly violating the optimization problem constraints.

### Question 4

We analyze four key inequalities (B1) to (B4) arising in the Frank-Wolfe algorithm under the assumptions that  $f$  is convex and continuously differentiable, and its gradient  $\nabla f$  is  $L$ -Lipschitz continuous.

#### Inequality Analysis

$$\text{(B1)} \quad f(x_{k+1}) - f(x_k) \leq \nabla f(x_k)^\top (x_{k+1} - x_k) + \frac{L}{2} \|x_{k+1} - x_k\|^2$$

This inequality stems from the first-order Taylor expansion and the Lipschitz continuity of  $\nabla f$ , bounding the error of the linear approximation.

$$\text{(B2)} \quad \leq \eta_k \nabla f(x_k)^\top (s(x_k) - x_k) + \frac{L}{2} \eta_k^2 d_S^2$$

Using the update formula  $x_{k+1} = (1 - \eta_k)x_k + \eta_k s(x_k)$  and the definition of  $d_S$ , the diameter of  $S$ , this inequality bounds the change in  $f$  in terms of the diameter of  $S$  and step size  $\eta_k$ .

$$\text{(B3)} \quad \leq \eta_k \nabla f(x_k)^\top (x^* - x_k) + \frac{L}{2} \eta_k^2 d_S^2$$

Given that  $s(x_k)$  minimizes the linear approximation over  $S$ , the inequality follows by comparing  $s(x_k)$  to any  $x^* \in S$ , including the optimal point.

$$\text{(B4)} \quad \leq \eta_k (f(x^*) - f(x_k)) + \frac{L}{2} \eta_k^2 d_S^2$$

This follows from the convexity of  $f$ , which implies  $f(x^*) - f(x_k) \geq \nabla f(x_k)^\top (x^* - x_k)$ . Substituting this into (B3) yields (B4).

### Question 5

Given  $x_0 \in S$ , let  $x_1$  be produced by the Frank-Wolfe algorithm with  $\eta_0 = \frac{2}{0+2} = 1$ . We show that  $f(x_1) - f(x^*) \leq \frac{L}{2}d_S^2$ , where  $L$  is the Lipschitz constant of  $\nabla f$  and  $d_S$  is the diameter of  $S$ .

#### Proof

1. The update rule for  $x_{k+1}$  in the Frank-Wolfe algorithm is  $x_{k+1} = (1 - \eta_k)x_k + \eta_k s(x_k)$ . For  $k = 0$ , this becomes  $x_1 = s(x_0)$ .
2. By the  $L$ -Lipschitz continuity of  $\nabla f$ , we have

$$f(y) \leq f(x) + \nabla f(x)^\top (y - x) + \frac{L}{2} \|y - x\|^2 \quad \text{for all } x, y \in S.$$

3. Setting  $x = x_1$  and  $y = x^*$ , we get

$$f(x^*) \leq f(x_1) + \nabla f(x_1)^\top (x^* - x_1) + \frac{L}{2} \|x^* - x_1\|^2.$$

4. Rearranging, we obtain

$$f(x_1) - f(x^*) \geq -\nabla f(x_1)^\top (x^* - x_1) - \frac{L}{2} \|x^* - x_1\|^2.$$

5. Since  $x_1$  and  $x^*$  are in  $S$  and  $\|x^* - x_1\|^2 \leq d_S^2$ , we have

$$f(x_1) - f(x^*) \leq \frac{L}{2} d_S^2.$$

Thus, after the first iteration with  $\eta_0 = 1$ , the function value at  $x_1$  is within  $\frac{L}{2}d_S^2$  of the optimal value  $f(x^*)$ .

### Question 6

We prove that for the Frank-Wolfe algorithm with step sizes  $\eta_k = \frac{2}{k+2}$ , the inequality  $f(x_k) - f(x^*) \leq \frac{2Ld_S^2}{k+2}$  holds for all  $k \geq 1$ .

#### Proof by Induction

##### Base Case ( $k = 1$ )

From the previous analysis, we have  $f(x_1) - f(x^*) \leq \frac{Ld_S^2}{2}$ , which satisfies the inequality for  $k = 1$ , as  $\frac{2Ld_S^2}{1+2} = \frac{Ld_S^2}{2}$ .

##### Inductive Step

Assume the inequality holds for some  $k \geq 1$ :

$$f(x_k) - f(x^*) \leq \frac{2Ld_S^2}{k+2}$$

We need to show it holds for  $k + 1$ :

$$f(x_{k+1}) - f(x^*) \leq \frac{2Ld_S^2}{k+3}$$

Using the inequality (B4) from the algorithm:

$$f(x_{k+1}) - f(x_k) \leq \eta_k(f(x^*) - f(x_k)) + \frac{L}{2}\eta_k^2 d_S^2$$

where  $\eta_k = \frac{2}{k+2}$ . Substituting and rearranging gives:

$$\begin{aligned} f(x_{k+1}) - f(x^*) &\leq \left(1 - \frac{2}{k+2}\right) \frac{2Ld_S^2}{k+2} + \frac{L}{2} \left(\frac{2}{k+2}\right)^2 d_S^2 \\ &= \frac{kLd_S^2}{(k+2)^2} + \frac{2Ld_S^2}{(k+2)^2} \\ &= \frac{(k+2)Ld_S^2}{(k+2)^2} \\ &= \frac{Ld_S^2}{k+2} \end{aligned}$$

Using  $\frac{2}{k+2} \leq \frac{2}{k+3}$ , we get:

$$f(x_{k+1}) - f(x^*) \leq \frac{2Ld_S^2}{k+3}$$

Thus, by induction, the inequality holds for all  $k \geq 1$ .

### Question 7

We show that the simplex  $\Delta_n = \{x \in \mathbb{R}^n | x_1, \dots, x_n \geq 0 \text{ and } x_1 + \dots + x_n = 1\}$  is convex, compact, and non-empty.

A set is convex if for any two points in the set, the line segment between them is also in the set. For  $x, y \in \Delta_n$  and  $\lambda \in [0, 1]$ , consider  $z = (1 - \lambda)x + \lambda y$ . Since  $x_i, y_i \geq 0$ , each component  $z_i = (1 - \lambda)x_i + \lambda y_i \geq 0$ . Also,  $\sum_{i=1}^n z_i = (1 - \lambda)\sum_{i=1}^n x_i + \lambda\sum_{i=1}^n y_i = 1$ . Hence,  $z \in \Delta_n$ , proving convexity.

A set is compact if it is closed and bounded.  $\Delta_n$  is closed as it contains all its limit points. It is bounded because for all  $x \in \Delta_n$ ,  $0 \leq x_i \leq 1$  and  $\sum_{i=1}^n x_i = 1$ . Therefore,  $\Delta_n$  is compact.

$\Delta_n$  is non-empty as it contains at least the point  $x = (1, 0, \dots, 0)$ , which satisfies  $x_i \geq 0$  and  $\sum_{i=1}^n x_i = 1$ .

In conclusion, the simplex  $\Delta_n$  is convex, compact, and non-empty.

### Question 8

Given a vector  $w \in \mathbb{R}^n$ , we consider the problem of minimizing  $\langle w, x \rangle$  subject to  $x \in \Delta_n$ , where  $\Delta_n$  is the simplex in  $\mathbb{R}^n$ .

The problem is formulated as:

$$\text{minimize } \langle w, x \rangle \quad \text{subject to } x \in \Delta_n.$$

To minimize  $\langle w, x \rangle$ , we allocate the entire weight to the component of  $x$  corresponding to the smallest component of  $w$ . Let  $i^* = \arg \min_i w_i$ . The minimizing vector  $x$  is such that  $x_{i^*} = 1$  and  $x_i = 0$  for all  $i \neq i^*$ .

The computational complexity of finding this solution is  $O(n)$ , as it requires a linear scan to find the minimum component of the vector  $w$ . The minimizing  $x$  is then directly obtained from the index of this minimum component.



### Question 9

Consider the optimization problem  $\min_{x \in \Delta_n} f(x)$  with  $f(x) = \frac{1}{2} \|Ax - b\|^2$ .

The function  $f$  is continuous and defined on the compact set  $\Delta_n$ ,  
thus by the Weierstrass Extreme Value Theorem,  $f$  attains its minimum on  $\Delta_n$ ,  
ensuring the existence of a solution.

However, the uniqueness of the solution depends on  $A$  and  $b$ .

For instance, if  $A = (1, 0, 0, \dots, 0)$  and  $b < 0$ ,  
then all  $x \in \Delta_n$  such that  $x_1 = 0, x_2 = 0$  minimize  $f(x)$ ,  
indicating that the solution is not necessarily unique.

### Question 10

$$\begin{aligned} f(x) &= \frac{1}{2} \|Ax - b\|^2 = \frac{1}{2} (Ax - b)^\top (Ax - b). \\ f(x + tv) &= \frac{1}{2} \|A(x + tv) - b\|^2 \\ &= \frac{1}{2} (A(x + tv) - b)^\top (A(x + tv) - b) \\ &= \frac{1}{2} ((Ax + Atv - b)^\top (Ax + Atv - b)). \end{aligned}$$

Using a Taylor expansion around  $t = 0$ ,

$$f(x + tv) = f(x) + t \langle u, A^\top (Ax - b) \rangle + O(t^2),$$

where  $u$  = the derivative of  $x + tv$  with respect to  $t$  at  $t = 0$  (which is  $v$ ).

Therefore,  $\nabla f(x) = A^\top (Ax - b)$ .

### Question 11

We analyze the line-search function  $g(\eta) = f((1 - \eta)x + \eta y)$  where  $x, y \in \Delta_n$  and  $f(x) = \frac{1}{2} \|Ax - b\|^2$  to determine the optimal values of  $\eta \in [0, 1]$ .

#### Expression for $g(\eta)$

$$g(\eta) = \frac{1}{2} \|A((1 - \eta)x + \eta y) - b\|^2$$

#### Optimal Value of $\eta$

To find the optimal  $\eta$ , we differentiate  $g(\eta)$  with respect to  $\eta$  and set the derivative to zero:

$$\begin{aligned} g'(\eta) &= \frac{d}{d\eta} \frac{1}{2} \|A((1 - \eta)x + \eta y) - b\|^2 \\ &= (A((1 - \eta)x + \eta y) - b)^\top A(y - x) \end{aligned}$$

Setting  $g'(\eta) = 0$  gives:

$$(A((1 - \eta)x + \eta y) - b)^\top A(y - x) = 0$$

Solving this equation for  $\eta$  gives the optimal value.

## Closed-Form Formula

A closed-form expression for  $\eta$  depends on the specific structure of  $A$ ,  $b$ ,  $x$ , and  $y$ . Without additional assumptions, the exact solution might be complex or not directly obtainable.

## Question 12

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1 function nabla = nabla(x,A,b)
2     nabla = A.'*(A*x-b);
3 end
```

```
1 function sx=linearsubproblem(w)
2     sx=double(w<=min(w));
3     sx=sx./sum(sx);
4 end
```

```
1 function [xbar, gaps]= frank_wolfe(A, x0, b)
2 steps_max=1e5;
3 tol=1e-3;
4
5 xbar=x0;
6 i=0;
7 nabla=nabla(xbar,A,b);
8 sx=linearsubproblem(nabla);
9 gap=dot(nabla,xbar-sx);
10 gaps=[gap];
11 while(i<steps_max&& gap>tol)
12     eta=min( dot(sx-xbar,-nabla)/dot(A*(sx-xbar),A*(sx-xbar)), 1)
13     ;
14     eta=max(eta,0);
15     xbar=(1-eta)*xbar+eta*sx;
16     nabla=nabla(xbar,A,b);
17     sx=linearsubproblem(nabla);
18     gap=dot(nabla,xbar-sx);
19     gaps=[gaps,gap];
20     i=i+1;
21 end
```

## Question 13

```
1 clear
2 load data.mat
3
4
5 x0=abs(randn(size(x)));
6 x0=x0./sum(x0);
7
8 [xbar, gaps] =frank_wolfe(A,x0,b);
9
10 semilogy(gaps);
11
```

```

12 figure;
13 plot_data(x,xbar, A\b);

```

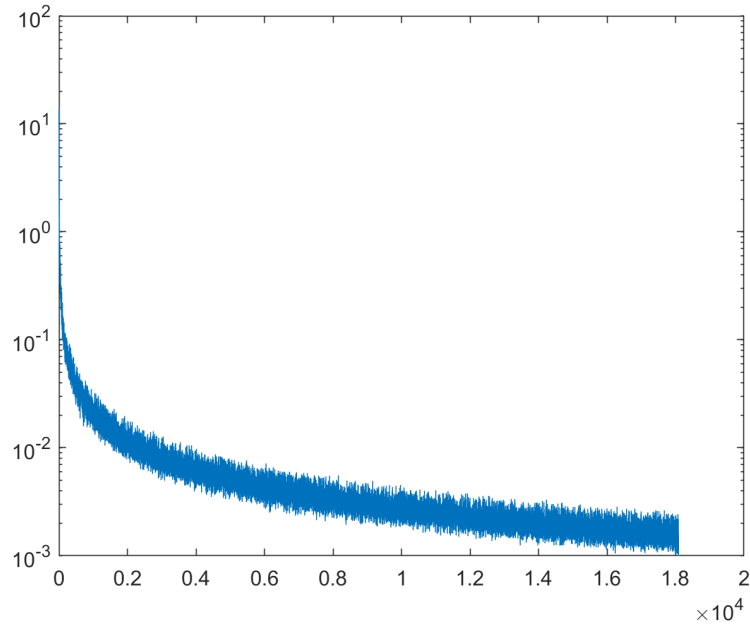


Figure 1: Plot of the Frank-Wolfe gap versus the number of iterations.

## Question 14

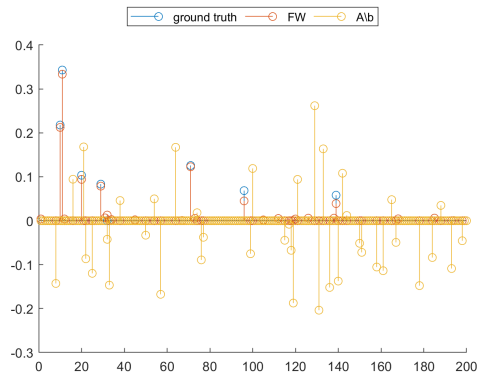


Figure 2: Comparison of the Frank-Wolfe algorithm, the true value and the vector obtained by the Matlab solver.

## Part 3 : KKT conditions and constraint qualifications

### Question 1

Notice that:

$$\begin{aligned} f_i(x) \leq y; \forall i = 1, \dots, N &\iff \\ \max_{i=1, \dots, N} f_i(x) \leq y &\iff \\ f(x) \leq y &\implies \\ \min_x f(x) \leq \min_{(x,y) \in S} y \end{aligned}$$

The other inequality is trivial as  $(x, f(x)) \in S$ . So we can easily conclude that the 2 programs have the same optimal value.

### Question 2

Notice that the constraints for  $S$  are:

$$g_i(x, y) = f_i(x) - y \leq 0; \forall i = 1, \dots, N$$

It is easy to see that:

$$\nabla g_i(x, y) = \begin{pmatrix} \nabla f_i(x) \\ -1 \end{pmatrix}$$

The KKT conditions for  $(x, y) \in S$  are:  $(x, y)$  is a KKT point if there exists  $\lambda \in \mathbb{R}^N$ , with  $\lambda \geq 0$  such that:

$$-(0, \dots, 0, 1) = \sum_{i=1}^N \lambda_i (\nabla f_i(x))^T, -1)$$

and

$$\lambda_i (f_i(x) - y) = 0; \forall i = 1, \dots, N$$

### Question 3

Let  $n = 1$  and:

$$f_1(x) = x, f_2(x) = x^2, f_3(x) = x^3$$

Then:

$$\nabla g_1(x, y) = \begin{pmatrix} 1 \\ -1 \end{pmatrix}, \nabla g_2(x, y) = \begin{pmatrix} 2x \\ -1 \end{pmatrix}, \nabla g_3(x, y) = \begin{pmatrix} 3x^2 \\ -1 \end{pmatrix}$$

For  $(x, y) = (1, 1) \in S$ , it is clear that LICQ doesn't hold as  $\nabla g_1(1, 1), \nabla g_2(1, 1), \nabla g_3(1, 1)$  are linearly dependent as a family of 3 vectors in  $\mathbb{R}^2$  ( $g_i(1, 1) = 0$  for all  $i$ ).

### Question 4

Let  $I(x, y) = \{i \in \{1, \dots, N\} \mid g_i(x, y) = 0\}$ . Then for MFCQ to hold, for all  $(x, y) \in S$  we need to find a point  $(\tilde{x}, \tilde{y}) \in S$  such that:

$$\langle \nabla g_i(x, y), (\tilde{x} - x, \tilde{y} - y) \rangle < 0; \forall i \in I(x)$$

Substituting  $\nabla g_i$ , we need to find  $(\tilde{x}, \tilde{y}) \in S$  such that:

$$\left\langle \begin{pmatrix} \nabla f_i(x) \\ -1 \end{pmatrix}, (\tilde{x} - x, \tilde{y} - y) \right\rangle < 0; \forall i \in I(x)$$

But notice that if we set  $\tilde{x} := x$  and  $\tilde{y} := y + 1$  then:

$$\left\langle \begin{pmatrix} \nabla f_i(x) \\ -1 \end{pmatrix}, (0, 1) \right\rangle = \langle -1, 1 \rangle = -1 < 0; \quad \forall i \in I(x)$$

So MFCQ holds for all  $(x, y) \in S$ .

**(Note:**  $(x, y + 1) \in S$  as  $y + 1 > y \geq f_i(x)$ , for all  $i$ )