


Q1. Propositional Logic: Semantic Tableaux [10]

Determine whether $A \rightarrow B$, $C \rightarrow D$, $B \vee D \rightarrow E$, $\neg E$ logically follows $\neg A \wedge \neg C$ or not.



Q2. Predicate Logic: Translations and Interpretations (10 Marks) [5+5]

In a seaside town, regulations state that every boat must have a life jacket for each passenger to ensure safety. One day, a boat is borrowed without the owner's permission by someone who is not aware that the boat is short on life jackets. When the coast guard stops the boat, they find that there are not enough life jackets for all passengers.

1 Express the following of these statements using quantifiers. Then form the negation of the statement so that no negation is to the left of a quantifier. [5]

(a) Boat users should ensure they have a life jacket for each passenger. [2.5]

$$\textcircled{1} \quad \forall x. \forall y. (b(x) \wedge P(y)) \wedge U(x, y) \rightarrow J(x, y)$$

$$\begin{aligned} & \text{Negation: } \neg (\forall x. \forall y. (b(x) \wedge P(y)) \wedge U(x, y)) \rightarrow J(x, y) \\ & \equiv \exists x. \exists y. \neg ((b(x) \wedge P(y)) \wedge U(x, y)) \rightarrow J(x, y) \end{aligned}$$

(b) No boat user should be penalized for lacking life jackets if they were unaware of the shortage. [2.5]

$$\begin{aligned} \textcircled{1} \quad & \neg (\exists x. \exists y. b(x) \wedge P(y) \wedge U(x, y) \wedge \neg \text{aware}(x) \\ & \quad \wedge \text{Penalized}(x, y)) \\ & \text{Negation: } \exists x. \exists y. b(x) \wedge P(y) \wedge U(x, y) \wedge \neg \text{aware}(x) \\ & \quad \wedge \text{Penalized}(x, y). \end{aligned}$$

2 Let domain $D_x = D_y = \{-2, -1, 0, 1, 2\}$. Explain whether the following statements are true or not. If they are true then define the truth set. [5]

$$\begin{aligned}
 & \text{(a) } \forall x, \exists y \cdot x + y = 0 \\
 & (\exists y, -2 + y = 0) \wedge (\exists y, -1 + y = 0) \wedge (\exists y, 0 + y = 0) \wedge (\exists y, 1 + y = 0) \wedge (\exists y, 2 + y = 0) \\
 & (-2 + 2 = 0) \wedge (-1 + 1 = 0) \wedge (0 + 0 = 0) \wedge (1 + (-1) = 0) \\
 & \wedge (2 + (-2) = 0).
 \end{aligned}$$

True.

$$\begin{aligned}
 & \text{(b) } \exists x, \forall y \cdot x + y = y \\
 & \cancel{\exists x, (x + -2 = y)} \wedge \cancel{(x + -1 = y)} \\
 & \exists x, (x + (-2) = -2) \wedge (x + (-1) = -1) \wedge (x + 0 = 0) \\
 & \wedge (x + 1 = 1) \wedge (x + 2 = 2)
 \end{aligned}$$

True.

$$\begin{aligned}
 & (0 + (-2) = -2) \wedge (0 + (-1) = -1) \wedge (0 + 0 = 0) \\
 & \wedge (0 + 1 = 1) \wedge (0 + 2 = 2).
 \end{aligned}$$

Q3. Predicate Logic: Transformational Proof (10 Marks)

[10]

$$\begin{aligned}
 & \neg \exists x \exists y \cdot Q(x, y) \leftrightarrow Q(y, x) \Leftrightarrow \forall x \forall y \cdot (Q(x, y) \vee Q(y, x)) \wedge \forall y \forall x \cdot \\
 & (\neg Q(y, x) \vee \neg Q(x, y)) \\
 & \text{Show transformational proof of } \neg \exists x \exists y \cdot Q(x, y) \leftrightarrow Q(y, x) \Leftrightarrow \forall x \forall y \cdot (Q(x, y) \vee Q(y, x)) \wedge \forall y \forall x \cdot \\
 & (\neg Q(y, x) \vee \neg Q(x, y)) \\
 & \quad \text{I.H.S.} \\
 & \equiv \forall x \cdot \forall y \cdot \neg (Q(x, y) \rightarrow Q(y, x)) \rightarrow Q(y, x)) \quad \text{DeMorgan's L2x} \\
 & \equiv \forall x \cdot \forall y \cdot \neg (Q(x, y) \rightarrow Q(y, x)) \wedge (Q(y, x) \rightarrow Q(x, y)) \quad \text{Eccrine ce} \\
 & \equiv \forall x \cdot \forall y \cdot \neg ((\neg Q(x, y) \vee Q(y, x)) \wedge (\neg Q(y, x) \vee Q(x, y))) \quad \text{Implic, L2x} \\
 & \equiv \forall x \cdot \forall y \cdot (\neg (\neg Q(x, y) \vee Q(y, x)) \vee \neg (\neg Q(y, x) \vee Q(x, y))) \quad \text{DeMorgan's am 3} \\
 & \equiv \forall x \cdot \forall y \cdot ((\neg \neg Q(x, y) \wedge \neg Q(y, x)) \vee (\neg \neg Q(y, x) \wedge \neg Q(x, y))) \quad \text{DeMorgan's am 3} \\
 & \equiv \forall x \cdot \forall y \cdot ((Q(x, y) \wedge \neg Q(y, x)) \vee (Q(y, x) \wedge \neg Q(x, y))) \quad \text{Negation L2x} \\
 & \equiv \forall x \cdot \forall y \cdot ((Q(x, y) \wedge \neg Q(y, x)) \vee Q(y, x)) \wedge ((Q(x, y) \wedge \neg Q(y, x)) \vee \\
 & \quad \neg Q(x, y)) \quad \text{Dissimil} \\
 & \equiv \forall x \cdot \forall y \cdot ((Q(x, y) \vee Q(y, x)) \wedge (\neg Q(y, x) \vee Q(y, x)) \wedge (Q(x, y) \vee \neg Q(x, y)) \\
 & \quad \wedge (\neg Q(y, x) \vee \neg Q(x, y))) \quad \text{Distributive L2x} \\
 & \equiv \forall x \cdot \forall y \cdot (Q(x, y) \vee Q(y, x)) \wedge T \wedge T \wedge (\neg Q(y, x) \wedge \neg Q(x, y)) \\
 & \quad \text{Excluded Mid}(2x) \\
 & \equiv \forall x \cdot \forall y \cdot ((Q(x, y) \vee Q(y, x)) \wedge (\neg Q(y, x) \wedge \neg Q(x, y))) \quad \text{Simplification I.} \\
 & \equiv \forall x \cdot \forall y \cdot (Q(x, y) \vee Q(y, x)) \wedge \forall x \forall y \cdot (\neg Q(y, x) \wedge \neg Q(x, y)) \\
 & \quad \text{Sessional-2} \quad \text{Spring 2024} \quad \text{Page 5 of 11 Ad-Di, student} \\
 & \equiv \forall x \forall y \cdot (Q(x, y) \vee Q(y, x)) \wedge \forall y \forall x \cdot (\neg Q(y, x) \vee \neg Q(x, y)) \\
 & \quad \text{Swapping variables} \\
 & \quad \text{Proven.}
 \end{aligned}$$

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Q4. Predicate Logic: Natural Deduction Proof Marks) [10]

Proof the following by using natural deduction. you are only allowed to use inference rules.
 $P(x) \rightarrow Q(x) \wedge S(x)$, $\forall x \cdot P(x) \wedge R(x) \vdash_N \forall x \cdot R(x) \wedge S(x)$

$$1. \forall x. P(x) \rightarrow Q(x) \wedge S(x) \quad] \text{ premises}$$

$$2. \forall x. P(x) \wedge R(x).$$

$$3. \exists g. P(g) \rightarrow (Q(g) \wedge S(g)) \quad \forall E, 1$$

$$4. P(g) \rightarrow (Q(g) \wedge S(g)) \quad \forall E, 2$$

$$5. P(g) \wedge R(g) \quad \wedge E, 5$$

$$6. P(g) \quad u, b, \rightarrow E$$

$$7. Q(g) \wedge S(g) \quad \gamma, \wedge E$$

$$8. S(g) \quad s, \wedge E$$

$$9. R(g) \quad b, 9, \wedge I$$

$$10. R(g) \wedge S(g)$$

$$3-10, \forall I.$$

$$11. \forall x. R(x) \wedge S(x)$$