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## 1 牛客计算几何

```
#include <bits/stdc++.h>
using namespace std;
// using point t=long long;
using point t=long double; //全局数据类型
constexpr point t eps=1e-8;
constexpr point t INF=numeric limits < point t >:: max();
constexpr long double PI=3.14159265358979323841;
// 点与向量
template < typename T > struct point
{
    T x, y;
    bool operator == (const point &a) const {return (abs (x-
       a.x) \le eps \&\& abs(y-a.y) \le eps);
    bool operator < (const point &a) const {if (abs(x-a.x))
       \leq eps) return y\leq a.y-eps; return x\leq a.x-eps;
    bool operator > (const point &a) const {return !(*this <
       a \mid | *this == a);
    point operator+(const point &a) const {return {x+a.x,
       y+a.y;
    point operator - (const point &a) const {return {x-a.x,
       y-a.y;
    point operator -() const {return \{-x,-y\};}
    point operator*(const T k) const {return {k*x,k*y};}
    point operator/(const T k) const {return \{x/k,y/k\};}
    T operator*(const point &a) const {return x*a.x+y*a.y
       ;} // 点积
    T operator \(^(\const \point \&a) \const \{return x \const \y - y \cdot a \cdot x\}\)
       ;} // 叉积,注意优先级
    int toleft (const point &a) const {const auto t=(*this
       )^a; return (t>eps)-(t<-eps);} // to-left 测试
    T len2() const {return (*this)*(*this);} // 向量长度
       的平方
    T dis2 (const point &a) const {return (a-(*this)).len2
       ();} // 两点距离的平方
    // 涉及浮点数
    long double len() const {return sqrtl(len2());} //
       向量长度
    long double dis(const point &a) const {return sqrtl(
       dis2(a));} // 两点距离
```

```
long double ang(const point &a) const {return acosl(
       \max(-1.01, \min(1.01, ((*this)*a)/(len()*a.len()))))
       ;} // 向量夹角
    point rot(const long double rad) const {return {x*cos
       (rad)-y*sin(rad),x*sin(rad)+y*cos(rad)};} // 逆时
       针旋转(给定角度)
    point rot(const long double cosr, const long double
       sinr) const {return {x*cosr-y*sinr, x*sinr+y*cosr
       };  // 逆时针旋转(给定角度的正弦与余弦)
};
using Point=point < point t >;
// 极角排序
struct argcmp
    bool operator()(const Point &a, const Point &b) const
        const auto quad = [](const Point &a)
            if (a.y<-eps) return 1;</pre>
            if (a.y>eps) return 4;
            if (a.x<-eps) return 5;</pre>
            if (a.x>eps) return 3;
            return 2;
        };
       const int qa=quad(a),qb=quad(b);
        if (qa!=qb) return qa < qb;
       const auto t=a^b;
        // if (abs(t) \le eps) return a*a \le b*b - eps; // 不同
           长度的向量需要分开
       return t>eps;
    }
};
// 直线
template < typename T > struct line
    point <T> p, v; // p 为直线上一点, v 为方向向量
    bool operator == (const line &a) const {return v.toleft
       (a.v) == 0 \&\& v.toleft(p-a.p) == 0;
    int toleft(const point <T> &a) const {return v.toleft(
       a-p);} // to-left 测试
   bool operator <(const line &a) const // 半平面交算法
       定义的排序
```

```
{
       if (abs(v^a.v) \le eps \&\& v*a.v = -eps) return toleft
          (a.p) == -1;
       return argcmp()(v,a.v);
   // 涉及浮点数
   point <T > inter(const line &a) const {return p+v*((a.v
       ^(p-a.p))/(v^a.v));} // 直线交点
   long double dis(const point <T> &a) const {return abs(
      v^(a-p))/v.len();} // 点到直线距离
   point <T> proj(const point <T> &a) const {return p+v*((
      v*(a-p))/(v*v));} // 点在直线上的投影
};
using Line=line < point t >;
// 线段
template < typename T > struct segment
   point < T > a, b;
   bool operator < (const segment &s) const {return
      make_pair(a,b)<make_pair(s.a,s.b);}
   // 判定性函数建议在整数域使用
   // 判断点是否在线段上
   // -1 点在线段端点 | 0 点不在线段上 | 1 点严格在线段
   int is on (const point <T> &p) const
       if (p==a \mid p==b) return -1;
       return (p-a). to left (p-b)==0 && (p-a)*(p-b)<-eps;
   // 判断线段直线是否相交
   // -1 直线经过线段端点 | 0 线段和直线不相交 | 1 线段
       和直线严格相交
   int is_inter(const line <T> &1) const
       if (1.toleft(a)==0 || 1.toleft(b)==0) return -1;
       return 1. toleft(a)!=1. toleft(b);
   // 判断两线段是否相交
```

```
// -1 在某一线段端点处相交 | 0 两线段不相交 | 1 两线
       段严格相交
   int is inter(const segment < T > &s) const
       if (is_on(s.a) || is_on(s.b) || s.is_on(a) || s.
           is_on(b)) return -1;
       const line T > 1 \{a, b-a\}, 1s \{s.a, s.b-s.a\};
       return 1. toleft(s.a)*1. toleft(s.b)==-1 && 1s.
           toleft(a)*ls.toleft(b)==-1;
   }
   // 点到线段距离
   long double dis(const point <T> &p) const
       if ((p-a)*(b-a) < -eps \mid | (p-b)*(a-b) < -eps) return
           min(p.dis(a),p.dis(b));
       const line T > 1\{a, b-a\};
       return 1. dis(p);
   }
   // 两线段间距离
   long double dis(const segment <T > &s) const
       if (is_inter(s)) return 0;
       return min(\{dis(s.a), dis(s.b), s.dis(a), s.dis(b)\})
};
using Segment=segment<point t>;
// 多边形
template < typename T > struct polygon
{
   vector < point <T>> p; // 以逆时针顺序存储
   size t nxt(const size t i) const {return i==p.size()
       -1?0:i+1;
    size t pre(const size t i) const {return i==0?p. size
       ()-1:i-1;
   // 回转数
   // 返回值第一项表示点是否在多边形边上
   // 对于狭义多边形, 回转数为 0 表示点在多边形外, 否则
       点在多边形内
   pair < bool, int > winding (const point < T > &a) const
```

```
{
        int cnt=0;
        for (size_t = 0; i \le p. size(); i++)
            const point <T> u=p[i], v=p[nxt(i)];
            if (abs((a-u)^{(a-v)}) \le eps \& (a-u)*(a-v) \le eps
               ) return \{true, 0\};
            if (abs(u.y-v.y) <= eps) continue;
            const Line uv = \{u, v-u\};
            if (u.y \le v.y - eps \&\& uv.toleft(a) \le 0) continue;
            if (u.y>v.y+eps \&\& uv.toleft(a)>=0) continue;
            if (u.y < a.y-eps \&\& v.y >= a.y-eps) cnt++;
            if (u.y>=a.y-eps && v.y<a.y-eps) cnt--;
        return {false, cnt};
    }
    // 多边形面积的两倍
   // 可用于判断点的存储顺序是顺时针或逆时针
   T area() const
       T sum=0;
        )];
        return sum;
    }
    // 多边形的周长
   long double circ() const
        long double sum=0;
        for (size t i=0; i < p. size(); i++) sum+=p[i]. dis(p[
           nxt(i)]);
        return sum;
    }
};
using Polygon=polygon<point_t >;
//凸多边形
template < typename T > struct convex: polygon < T >
    // 闵可夫斯基和
    convex operator + (const convex &c) const
    {
        const auto &p=this->p;
```

```
vector < Segment > e1 (p. size ()), e2 (c.p. size ()), edge (
        p. size()+c.p. size());
    vector < point <T>> res; res.reserve(p.size()+c.p.
        size());
    const auto cmp=[](const Segment &u, const Segment
        &v) {return argcmp()(u.b-u.a,v.b-v.a);};
    for (size_t i=0; i \le p. size(); i++) e1[i]=\{p[i], p[
        this \rightarrow nxt(i);
    for (size t i=0; i < c.p. size(); i++) e2[i] = \{c.p[i], c.p[i], c.p[i]\}
        .p[c.nxt(i)]};
    rotate(el.begin(), min element(el.begin(), el.end()
        , cmp), e1. end());
    rotate(e2.begin(), min_element(e2.begin(), e2.end()
        , cmp), e2. end());
    merge(e1.begin(),e1.end(),e2.begin(),e2.end(),
        edge.begin(),cmp);
    const auto check = [](const vector < point <T>> &res,
        const point <T> &u)
    {
         const auto back1=res.back(), back2=*prev(res.
            end(),2);
         return (back1-back2). toleft (u-back1)==0 && (
            back1-back2)*(u-back1) \ge -eps;
    };
    auto u=e1[0].a+e2[0].a;
    for (const auto &v:edge)
         while (res.size()>1 && check(res,u)) res.
            pop back();
         res.push_back(u);
        u=u+v \cdot b-v \cdot a;
    if (res.size()>1 && check(res, res[0])) res.
        pop_back();
    return {res};
}
// 旋转卡壳
// 例: 凸多边形的直径的平方
T rotcaliper() const
    const auto &p = this -> p;
    if (p.size()==1) return 0;
    if (p. size() == 2) return p[0]. dis2(p[1]);
    const auto area = [](const point < T > &u, const point <
        T > &v, const point < T > &w) \{return (w-u)^(w-v); \};
```

```
T ans=0;
    for (size_t i=0, j=1; i < p. size(); i++)
        const auto nxti=this->nxt(i);
        ans=max({ans,p[j].dis2(p[i]),p[j].dis2(p[nxti])}
        while (area(p[this \rightarrow nxt(j)], p[i], p[nxti]) > =
            area(p[j],p[i],p[nxti]))
            j = this \rightarrow nxt(j);
            ans=max({ans,p[j].dis2(p[i]),p[j].dis2(p[
                }
    return ans;
}
// 判断点是否在凸多边形内
// 复杂度 O(logn)
// -1 点在多边形边上 | 0 点在多边形外 | 1 点在多边形
int is_in(const point <T> &a) const
    const auto &p=this \rightarrow p;
    if (p. size() == 1) return a == p[0]? -1:0;
    if (p.size()==2) return segmentT>\{p[0],p[1]\}.
       is on (a)? -1:0;
    if (a==p[0]) return -1;
    if ((p[1]-p[0]) \cdot toleft(a-p[0]) == -1 \mid | (p.back()-p)
       [0]).toleft(a-p[0]) == 1) return 0;
    const auto cmp=[&](const point <T> &u, const point <
       T > &v) { return (u-p[0]) . toleft (v-p[0]) ==1; };
    const size_t i=lower_bound(p.begin()+1,p.end(),a,
       cmp)-p.begin();
    if (i==1) return segment T>\{p[0], p[i]\}. is_on(a)
        ? -1:0;
    if (i==p.size()-1 \&\& segment<T>\{p[0],p[i]\}.is on(
       a)) return -1;
    if (segment < T > {p[i-1],p[i]}.is_on(a)) return -1;
    return (p[i]-p[i-1]). toleft (a-p[i-1])>0;
}
// 凸多边形关于某一方向的极点
// 复杂度 O(logn)
// 参考资料: https://codeforces.com/blog/entry/48868
```

```
const
        const auto &p=this \rightarrowp;
        const auto check=[&](const size_t i){return dir(p
            [i]).toleft(p[this -> nxt(i)] - p[i])>=0;};
        const auto dir0=dir(p[0]); const auto check0=
            check(0);
        if (! \operatorname{check} 0 \&\& \operatorname{check}(p.\operatorname{size}()-1)) return 0;
        const auto cmp=[&](const point <T> &v)
            const size t vi=&v-p.data();
            if (vi==0) return 1;
            const auto checkv=check(vi);
            const auto t = dir0. toleft (v-p[0]);
            if (vi==1 \&\& checkv==check0 \&\& t==0) return
                1;
            return checkv^(checkv==check0 && t <=0);
        };
        return partition point (p. begin (), p. end (), cmp)-p.
            begin();
    }
    // 过凸多边形外一点求凸多边形的切线,返回切点下标
    // 复杂度 O(logn)
    // 必须保证点在多边形外
    pair < size t , size t > tangent (const point <T > &a) const
        const size t i=extreme([&](const point <T> &u){
            return u-a;});
        const size t j=extreme([&](const point <T> &u){
            return a-u;});
        return \{i,j\};
    }
    // 求平行于给定直线的凸多边形的切线, 返回切点下标
    // 复杂度 O(logn)
    pair < size_t > tangent(const line <T> &a) const
        const size_t i=extreme([\&](...){return a.v;});
        const size_t j=extreme([&](...){return -a.v;});
        return \{i, j\};
    }
};
using Convex=convex<point t>;
```

template < typename F > size\_t extreme (const F & dir)

```
// 圆
struct Circle
    Point c;
   long double r;
   bool operator == (const Circle &a) const {return c==a.c
       && abs(r-a.r) \le eps;
   long double circ() const {return 2*PI*r;}
                                              // 周长
   long double area() const {return PI*r*r;}
                                             // 面积
   // 点与圆的关系
   // -1 圆上 | 0 圆外 | 1 圆内
   int is in (const Point &p) const {const long double d=
       p.dis(c); return abs(d-r) \le eps?-1:d \le r-eps;
   // 直线与圆关系
   // 0 相离 | 1 相切 | 2 相交
   int relation (const Line &1) const
       const long double d=1.dis(c);
       if (d>r+eps) return 0;
       if (abs(d-r) \le eps) return 1;
       return 2;
   }
    // 圆与圆关系
    // -1 相同 | 0 相离 | 1 外切 | 2 相交 | 3 内切 | 4 内
   int relation (const Circle &a) const
       if (*this==a) return -1;
       const long double d=c.dis(a.c);
       if (d>r+a.r+eps) return 0;
       if (abs(d-r-a.r) \le eps) return 1;
        if (abs(d-abs(r-a.r)) \le eps) return 3;
       if (d < abs(r-a.r)-eps) return 4;
       return 2;
   }
   // 直线与圆的交点
   vector < Point > inter(const Line &1) const
       const long double d=1.dis(c);
       const Point p=1.proj(c);
```

```
const int t=relation(1);
    if (t==0) return vector <Point >();
    if (t==1) return vector Point > \{p\};
    const long double k = sqrt(r*r-d*d);
    return vector < Point > {p-(1.v/1.v.len())*k,p+(1.v/1
       .v.len())*k};
}
// 圆与圆交点
vector < Point > inter (const Circle &a) const
    const long double d=c.dis(a.c);
    const int t=relation(a);
    if (t==-1 | | t==0 | | t==4) return vector < Point > ()
    Point e=a.c-c; e=e/e.len()*r;
    if (t==1 | | t==3)
        if (r*r+d*d-a.r*a.r>=-eps) return vector <
            Point > \{c+e\};
        return vector <Point > {c-e};
    const long double costh = (r*r+d*d-a.r*a.r)/(2*r*d)
        , sinth = sqrt(1-costh*costh);
    return vector < Point > {c+e.rot(costh, -sinth), c+e.
       rot(costh, sinth)};
}
// 圆与圆交面积
long double inter_area(const Circle &a) const
{
    const long double d=c.dis(a.c);
    const int t=relation(a);
    if (t==-1) return area();
    if (t < 2) return 0;
    if (t>2) return min(area(), a. area());
    const long double costh1 = (r*r+d*d-a.r*a.r)/(2*r*d
       ), costh2 = (a.r*a.r+d*d-r*r)/(2*a.r*d);
    const long double sinth1=sqrt(1-costh1*costh1),
       sinth2 = sqrt(1-costh2*costh2);
    const long double th1=acos(costh1),th2=acos(
       costh2);
    return r*r*(th1-costh1*sinth1)+a.r*a.r*(th2-
       costh2*sinth2);
}
```

```
// 过圆外一点圆的切线
vector < Line > tangent (const Point &a) const
    const int t=is_in(a);
    if (t==1) return vector <Line >();
    if (t == -1)
        const Point v = \{-(a-c).y, (a-c).x\};
        return vector <Line > {{a, v}};
    Point e=a-c; e=e/e.len()*r;
    const long double costh=r/c.dis(a), sinth=sqrt(1-
        costh*costh);
    const Point t1=c+e. rot (costh, -sinth), t2=c+e. rot (
        costh, sinth);
    return vector <Line > \{a, t1-a\}, \{a, t2-a\}\};
}
// 两圆的公切线
vector < Line > tangent (const Circle & a) const
    const int t=relation(a);
    vector <Line> lines;
    if (t==-1 \mid | t==4) return lines;
    if (t==1 | | t==3)
        const Point p=inter(a)[0], v = \{-(a.c-c).y, (a.c-c)\}
            c).x};
        lines.push_back({p,v});
    const long double d=c.dis(a.c);
    const Point e=(a.c-c)/(a.c-c).len();
    if (t \le 2)
    {
        const long double costh = (r-a.r)/d, sinth = sqrt
            (1-\cosh*\cosh);
        const Point d1=e.rot(costh, -sinth), d2=e.rot(
            costh, sinth);
        const Point u1=c+d1*r, u2=c+d2*r, v1=a.c+d1*a.r
            , v2=a.c+d2*a.r;
        lines.push_back({u1,v1-u1}); lines.push_back
            (\{u2, v2-u2\});
    if (t==0)
```

```
(1-\cosh*\cosh);
             const Point d1=e.rot(costh, -sinth), d2=e.rot(
                 costh, sinth);
             const Point u1=c+d1*r, u2=c+d2*r, v1=a.c-d1*a.r
                 , v2=a.c-d2*a.r;
             lines.push_back({u1,v1-u1}); lines.push_back
                 (\{u2, v2-u2\});
        return lines;
    }
    // 圆的反演
    tuple < int, Circle, Line > inverse (const Line &1) const
         const Circle null c = \{\{0.0, 0.0\}, 0.0\};
        const Line null_1 = \{ \{0.0, 0.0\}, \{0.0, 0.0\} \};
         if (1. toleft(c) == 0) return \{2, null_c, 1\};
        const Point v=1. toleft(c)==1? Point \{1.v.y, -1.v.x\}:
            Point \{-1.v.y, 1.v.x\};
        const long double d=r*r/l.dis(c);
        const Point p=c+v/v.len()*d;
        return \{1,\{(c+p)/2,d/2\},null_1\};
    }
    tuple < int, Circle, Line > inverse (const Circle & a) const
        const Circle null_c = \{\{0.0, 0.0\}, 0.0\};
        const Line null 1 = \{\{0.0, 0.0\}, \{0.0, 0.0\}\};
        const Point v=a.c-c;
         if (a.is in (c) == -1)
             const long double d=r*r/(a.r+a.r);
             const Point p=c+v/v.len()*d;
             return \{2, null_c, \{p, \{-v, y, v, x\}\}\};
        if (c==a.c) return \{1,\{c,r*r/a.r\},null\ 1\};
        const long double d1=r*r/(c.dis(a.c)-a.r), d2=r*r
            /(c.dis(a.c)+a.r);
        const Point p=c+v/v.len()*d1, q=c+v/v.len()*d2;
         return \{1,\{(p+q)/2,p.dis(q)/2\},null_1\};
    }
};
// 圆与多边形面积交
```

const long double costh = (r+a.r)/d, sinth = sqrt

```
long double area_inter(const Circle &circ, const Polygon &
   poly)
{
    const auto cal = [](const Circle & circ, const Point & a,
       const Point &b)
        if ((a-circ.c).toleft(b-circ.c)==0) return 0.01;
        const auto ina=circ.is in(a), inb=circ.is in(b);
        const Line ab = \{a, b-a\};
        if (ina && inb) return ((a-circ.c)^{(b-circ.c)})/2;
        if (ina &&!inb)
             const auto t=circ.inter(ab);
             const Point p=t. size () ==1?t[0]:t[1];
             const long double ans = ((a-circ.c)^{(p-circ.c)})
             const long double th = (p-circ.c).ang(b-circ.c)
             const long double d=circ.r*circ.r*th/2;
             if ((a-circ.c).toleft(b-circ.c)==1) return
                ans+d;
             return ans-d;
        if (!ina && inb)
             const Point p=circ.inter(ab)[0];
             const long double ans = ((p-circ.c) \land (b-circ.c))
                /2;
             const long double th = (a - circ.c).ang(p - circ.c)
             const long double d=circ.r*circ.r*th/2;
             if ((a-circ.c).toleft(b-circ.c)==1) return
                ans+d;
             return ans-d;
        const auto p=circ.inter(ab);
        if (p. size() == 2 \&\& Segment\{a,b\}. dis(circ.c) \le circ
            .r+eps)
        {
             const long double ans = ((p[0] - circ.c)^{(p[1] - circ.c)})
                circ.c))/2;
             const long double th1 = (a-circ.c).ang(p[0]-
                circ.c), th2 = (b-circ.c). ang (p[1]-circ.c);
             const long double d1=circ.r*circ.r*th1/2,d2=
                circ.r*circ.r*th2/2:
```

```
if ((a-circ.c).toleft(b-circ.c)==1) return
                ans+d1+d2;
            return ans-d1-d2;
        const long double th = (a - circ.c).ang(b - circ.c);
        if ((a-circ.c).toleft(b-circ.c)==1) return circ.r
           *circ.r*th/2;
        return -circ.r*circ.r*th/2;
    };
    long double ans = 0;
    for (size t i=0; i < poly.p.size(); i++)
        const Point a=poly.p[i],b=poly.p[poly.nxt(i)];
        ans += cal(circ, a, b);
    return ans;
}
// 点集的凸包
// Andrew 算法, 复杂度 O(nlogn)
Convex convexhull (vector < Point > p)
    vector < Point > st;
    if (p.empty()) return Convex{st};
    sort (p. begin (), p. end ());
    const auto check = [](const vector < Point > &st, const
       Point &u)
    {
        const auto back1=st.back(),back2=*prev(st.end()
        return (back1-back2).toleft(u-back1) <=0;
    };
    for (const Point &u:p)
        while (st.size()>1 && check(st,u)) st.pop_back();
        st.push back(u);
    size t k=st.size();
    p.pop_back(); reverse(p.begin(),p.end());
    for (const Point &u:p)
        while (st.size()>k && check(st,u)) st.pop back();
        st.push back(u);
    st.pop back();
```

```
}
// 半平面交
// 排序增量法, 复杂度 O(nlogn)
// 输入与返回值都是用直线表示的半平面集合
vector < Line > halfinter (vector < Line > 1, const point_t lim
         =1e9)
{
            const auto check = [] (const Line &a, const Line &b, const
                         Line &c) { return a. toleft (b. inter (c)) \leq 0; };
            // 无精度误差的方法,但注意取值范围会扩大到三次方
           /*const auto check=[](const Line &a, const Line &b,
                      const Line &c)
            {
                        const Point p=a.v*(b.v^c.v), q=b.p*(b.v^c.v)+b.v*(
                                  c \cdot v^{(b \cdot p-c \cdot p)} - a \cdot p *(b \cdot v^{c} \cdot v);
                        return p. toleft(q) < 0;
            }; */
            1. push back(\{\{-\lim_{n \to \infty}, \{0, -1\}\}\}); 1. push back(\{\{0, -\lim_{n \to \infty}, \{0, -1\}\}\})
                       \},\{1,0\}\});
            1. push_back(\{\{lim, 0\}, \{0, 1\}\}); 1. push_back(\{\{0, lim\}\}); 1. push_back(\{\{0, lim\}\}\}); 1. push_back(\{\{0, lim\}\}); 1. push_back(\{\{0, lim\}\}\}); 1. push_back(\{\{0, lim\}\}); 1. push_back(\{0, lim\}\}); 1. push_back(\{0, lim\}\}); 1. push_back(\{0, lim\}]); 1. push_back(\{0, lim\}\}); 1. push_back(\{0, li
                       \},\{-1,0\}\});
            sort(1.begin(),1.end());
            deque < Line > q;
            for (size t i=0; i<1.size(); i++)
                         if (i>0 \&\& 1[i-1].v.toleft(1[i].v)==0 \&\& 1[i-1].v
                                  *1[i].v>eps) continue;
                        while (q. size ()>1 && check(1[i],q.back(),q[q. size
                                   () -2])) q.pop back();
                        while (q.size()>1 && check(1[i],q[0],q[1])) q.
                                   pop front();
                         if (!q.empty() && q.back().v.toleft(l[i].v)<=0)</pre>
                                   return vector <Line >();
                        q.push_back(1[i]);
            while (q. size()>1 \&\& check(q[0],q.back(),q[q.size()
                       -2])) q.pop_back();
            while (q. size()>1 \&\& check(q.back(),q[0],q[1])) q.
                      pop_front();
            return vector < Line > (q. begin(), q. end());
}
// 点集形成的最小最大三角形
// 极角序扫描线, 复杂度 O(n^2 logn)
```

return Convex{st};

```
// 最大三角形问题可以使用凸包与旋转卡壳做到 O(n^2)
pair < point_t > minmax_triangle(const vector < Point >
    &vec)
    if (\text{vec.size}() \le 2) return \{0,0\};
    vector < pair < int , int >> evt;
    evt.reserve(vec.size()*vec.size());
    point t maxans=0, minans=INF;
    for (size t i=0; i < vec. size(); i++)
        for (size t j=0; j < vec. size(); j++)
             if (i==j) continue;
             if (\text{vec}[i] == \text{vec}[i]) minans = 0;
             else evt.push back({i,j});
    }
    sort(evt.begin(),evt.end(),[&](const pair < int, int > &u
        , const pair < int, int > &v)
    {
        const Point du=vec[u.second]-vec[u.first], dv=vec[
            v.second]-vec[v.first];
         return argcmp()(\{du.y,-du.x\},\{dv.y,-dv.x\});
    });
    vector < size t > vx(vec.size()), pos(vec.size());
    for (size t = 0; i < vec. size(); i + +) vx[i] = i;
    sort(vx.begin(), vx.end(), [\&](int x, int y) \{return vec[
        x \le vec[y]; \});
    for (size \ t \ i=0; i < vx. size(); i++) \ pos[vx[i]]=i;
    for (auto [u,v]:evt)
    {
        const size t i=pos[u], j=pos[v];
        const size t = \min(i, j), r = \max(i, j);
        const Point vecu=vec[u], vecv=vec[v];
         if (1>0) minans=min(minans, abs((vec[vx[1-1]]-vecu
            )^{(vec[vx[1-1]]-vecv))};
         if (r < vx \cdot size() - 1) minans=min(minans, abs((vec[vx[
            r+1]]-vecu)^(vec[vx[r+1]]-vecv));
        maxans=max(\{maxans, abs((vec[vx[0]]-vecu)^(vec[vx[vecu]))\}
            [0]] - vecv)), abs ((vec[vx.back()]-vecu)^(vec[vx.
            back()]-vecv))});
         if (i \le j) swap(vx[i], vx[j]), pos[u]=j, pos[v]=i;
    return {minans, maxans};
}
```

```
// 平面最近点对
// 扫描线, 复杂度 O(nlogn)
point_t closest_pair(vector < Point > points)
    sort(points.begin(),points.end());
    const auto cmpy=[](const Point &a, const Point &b){if
       (abs(a.y-b.y) \le eps) return a.x \le b.x-eps; return a.y
       <b.y-eps;};
    multiset < Point, decltype(cmpy) > s {cmpy};
    point t ans=INF;
    for (size t i=0, l=0; i < points. size(); i++)
        const point_t sqans=sqrtl(ans)+1; // 整数情况
        // const point t sqans=sqrtl(ans)+1; // 浮点数情
            况
        while (1 \le i \&\& points[i].x-points[1].x \ge sqans) s.
            erase (s. find (points [1++]));
        for (auto it=s.lower_bound(Point{-INF, points[i].y
            -sqans }); it !=s . end () &&it ->y-points [i]. y \le sqans
            ; it++)
        {
             ans=min(ans, points[i].dis2(*it));
        s.insert(points[i]);
    return ans;
}
// 判断多条线段是否有交点
// 扫描线, 复杂度 O(nlogn)
bool segs_inter(const vector < Segment > & segs)
{
    if (segs.empty()) return false;
    using seq_t=tuple < point_t , int , Segment >;
    const auto seqcmp = [](const seq_t &u, const seq_t &v)
    {
        const auto [u0, u1, u2]=u;
        \quad \textbf{const} \quad \textbf{auto} \quad [\ v0\ ,v1\ ,v2\ ] = v\ ;
        if (abs(u0-v0) \le eps) return make_pair(u1,u2) <
            make_pair(v1, v2);
        return u0<v0-eps;
    };
    vector < seq_t > seq;
    for (auto seg:segs)
    {
        if (seg.a.x>seg.b.x+eps) swap(seg.a,seg.b);
```

```
seq.push_back({seg.a.x,0,seg});
        seq.push_back({seg.b.x,1,seg});
    }
    sort(seq.begin(),seq.end(),seqcmp);
    point_t x_now;
    auto cmp=[&](const Segment &u, const Segment &v)
        if (abs(u.a.x-u.b.x) \le eps || abs(v.a.x-v.b.x) \le
           eps) return u.a.y<v.a.y-eps;
        return ((x \text{ now-u.a.x})*(u.b.y-u.a.y)+u.a.y*(u.b.x-u.b.y-u.a.y)
           u.a.x))*(v.b.x-v.a.x) < ((x_now-v.a.x)*(v.b.y-v.a.x))
           a.y)+v.a.y*(v.b.x-v.a.x))*(u.b.x-u.a.x)-eps;
    };
    multiset < Segment, decltype (cmp) > s {cmp};
    for (const auto [x,o,seg]:seq)
        x_now=x;
        const auto it=s.lower_bound(seg);
        if (o==0)
        {
            if (it!=s.end() && seg.is_inter(*it)) return
                true;
            if (it!=s.begin() && seg.is_inter(*prev(it)))
                 return true;
            s.insert(seg);
        }
        else
        {
            if (next(it)!=s.end() && it!=s.begin() && (*
                prev(it)).is_inter(*next(it))) return true
            s.erase(it);
        }
    return false;
}
// 多边形面积并
// 轮廓积分,复杂度 O(n^2 log n), n为边数
// ans[i] 表示被至少覆盖了 i+1 次的区域的面积
vector < long double > area_union (const vector < Polygon > &
   polys)
{
    const size t siz=polys.size();
    vector < vector < pair < Point , Point >>> segs(siz);
```

```
const auto check = [](const Point &u, const Segment &e) {
   return !((u \le a.a \&\& u \le a.b) || (u \ge a.a \&\& u \ge a.b)); ; ;
auto cut_edge=[&](const Segment &e, const size_t i)
    const Line le {e.a, e.b-e.a};
    vector < pair < Point , int >> evt;
    evt.push back(\{e.a,0\}); evt.push back(\{e.b,0\});
    for (size t j=0; j < polys. size(); j++)
    {
         if (i==j) continue;
         const auto &pj=polys[j];
         for (size_t k=0; k < pj.p. size(); k++)
             const Segment s = \{pj.p[k], pj.p[pj.nxt(k)]\}
             if (le.toleft(s.a)==0 \&\& le.toleft(s.b)
                 ==0)
                  evt.push_back({s.a,0});
                  evt.push_back({s.b,0});
             else if (s.is_inter(le))
                  const Line ls\{s.a, s.b-s.a\};
                  const Point u=le.inter(ls);
                  if (le.toleft(s.a)<0 \&\& le.toleft(s.b)
                     ) \ge 0 evt.push_back({u,-1});
                  else if (le.toleft(s.a) >= 0 \&\& le.
                      toleft(s.b)<0) evt.push back({u
                     ,1));
             }
         }
    sort(evt.begin(),evt.end());
    if (e.a>e.b) reverse(evt.begin(),evt.end());
    int sum=0;
    for (size \ t \ i=0; i < evt. size(); i++)
         sum += evt[i].second;
         const Point u=evt[i]. first , v=evt[i+1]. first;
         if (!(u==v) \&\& check(u,e) \&\& check(v,e)) segs
            [sum].push back({u,v});
         if (v==e.b) break;
    }
};
```

```
for (size_t i=0; i < polys.size(); i++)
        const auto &pi=polys[i];
        for (size_t k=0; k< pi.p. size(); k++)
             const Segment ei={pi.p[k],pi.p[pi.nxt(k)]};
             cut edge(ei,i);
    }
    vector < long double > ans(siz);
    for (size \ t \ i=0; i < siz; i++)
    {
        long double sum=0;
        sort(segs[i].begin(), segs[i].end());
        int cnt=0;
        for (size_t j=0; j < segs[i]. size(); j++)
             if (j>0 \&\& segs[i][j]==segs[i][j-1]) segs[i]
                +(++cnt)].push_back(segs[i][j]);
             else cnt=0, sum+=segs[i][j]. first segs[i][j].
                second;
        ans [i] = sum/2;
    return ans;
}
// 圆面积并
// 轮廓积分, 复杂度 O(n^2 logn)
// ans[i] 表示被至少覆盖了 i+1 次的区域的面积
vector < long double > area union (const vector < Circle > &
   circs)
{
    const size t siz=circs.size();
    using \ arc\_t = tuple < Point \ , long \ double \ , long \ double \ , long
       double >;
    vector < vector < arc_t >> arcs(siz);
    const auto eq = [](const arc_t &u, const arc_t &v)
    {
        const auto [u1, u2, u3, u4]=u;
        const auto [v1, v2, v3, v4] = v;
        return u1 == v1 && abs(u2-v2) \le eps && abs(u3-v3) \le eps
            eps && abs (u4-v4) \le eps;
    };
```

```
auto cut_circ = [&](const Circle &ci, const size_t i)
    vector < pair < long double, int >> evt;
    evt.push_back({-PI,0}); evt.push_back({PI,0});
    int init = 0;
    for (size_t j=0; j < circs. size(); j++)
        if (i==j) continue;
        const Circle &cj=circs[j];
        if (ci.r < cj.r - eps \&\& ci.relation(cj) >= 3) init
            ++;
        const auto inters=ci.inter(cj);
        if (inters.size()==1) evt.push_back({atan21((
            inters [0] - ci.c).y, (inters [0] - ci.c).x), (0);
        if (inters.size()==2)
        {
             const Point dl=inters[0]-ci.c, dr=inters
                [1] - ci.c;
             long double argl=atan21(dl.y,dl.x), argr=
                 atan21 (dr.y, dr.x);
             if (abs(argl+PI) \le eps) argl=PI;
             if (abs(argr+PI) <= eps) argr=PI;</pre>
             if (argl>argr+eps)
             {
                 evt.push back({argl,1}); evt.
                     push back(\{PI, -1\});
                 evt.push back({-PI,1}); evt.push back
                     (\{argr, -1\});
             }
             else
             {
                 evt.push back({argl,1});
                 evt.push back({argr,-1});
             }
        }
    }
    sort(evt.begin(),evt.end());
    int sum=init;
    for (size_t i=0; i < evt. size(); i++)
        sum+=evt[i].second;
        if (abs(evt[i].first-evt[i+1].first)>eps)
            arcs[sum].push back({ci.c,ci.r,evt[i].
            first, evt[i+1]. first });
        if (abs(evt[i+1].first-PI) \le eps) break;
```

```
};
    const auto oint = [](const arc_t & arc)
        const auto [cc, cr, l, r] = arc;
        if (abs(r-1-PI-PI) \le eps) return 2.01*PI*cr*cr;
        return cr*cr*(r-1)+cc.x*cr*(sin(r)-sin(1))-cc.y*
            cr*(cos(r)-cos(1));
    };
    for (size_t i=0; i < circs. size(); i++)
        const auto &ci=circs[i];
        cut circ(ci,i);
    vector < long double > ans(siz);
    for (size_t = 0; i < siz; i++)
        long double sum=0;
        sort(arcs[i].begin(),arcs[i].end());
        int cnt=0;
        for (size_t j=0; j < arcs[i]. size(); j++)
             if (j>0 && eq(arcs[i][j],arcs[i][j-1])) arcs[
                i+(++cnt)].push_back(arcs[i][j]);
             else cnt=0, sum+=oint(arcs[i][j]);
        ans [i] = sum/2;
    return ans;
}
```