# Blind 75 Coding Questions

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#### Array

- 1. Two Sum
- 2. Best Time to Buy and Sell Stock
- 3. Contains Duplicate
- 4. Product of Array Except Self
- 5. Maximum Subarray
- 6. Maximum Product Subarray
- 7. Find Minimum in Rotated Sorted Array
- 8. Search in Rotated Sorted Array
- 9. 3 Sum
- 10. Container With Most Water

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- 2. Coin Change
- 3. Longest Increasing Subsequence
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- 5. Word Break Problem
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- 7. House Robber
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#### Binary

- 1. Sum of Two Integers
- 2. Number of 1 Bits
- 3. Counting Bits
- 4. Missing Number
- 5. Reverse Bits

#### Graph

- 1. Clone Graph
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- 3. Pacific Atlantic Water Flow
- 4. Number of Islands
- 5. Longest Consecutive Sequence
- 6. Alien Dictionary (Leetcode Premium)
- 7. Graph Valid Tree (Leetcode Premium)
- 8. Number of Connected Components in an Undirected Graph (Leetcode Premium)

#### Interval

- 1. Insert Interval
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- 3. Non-overlapping Intervals
- 4. Meeting Rooms (Leetcode Premium)
- 5. Meeting Rooms II (Leetcode Premium)

#### Matrix

- 1. Set Matrix Zeroes
- 2. Spiral Matrix
- 3. Rotate Image
- 4. Word Search

#### String

- 1. Longest Substring Without Repeating Characters
- 2. Longest Repeating Character Replacement
- 3. Minimum Window Substring
- 4. Valid Anagram
- 5. Group Anagrams
- 6. Valid Parentheses
- 7. Valid Palindrome
- 8. Longest Palindromic Substring
- 9. Palindromic Substrings
- 10. Encode and Decode Strings (Leetcode Premium)

### Heap

- 1. Merge K Sorted Lists
- 2. Top K Frequent Elements
- 3. Find Median from Data Stream

#### LinkedList

- 1. Reverse a Linked List
- 2. Detect Cycle in a Linked List
- 3. Merge Two Sorted Lists
- 4. Merge K Sorted Lists
- 5. Remove Nth Node From End Of List
- 6. Reorder List

#### Tree

- 1. Maximum Depth of Binary Tree
- 2. Same Tree
- 3. Invert/Flip Binary Tree
- 4. Binary Tree Maximum Path Sum
- 5. Binary Tree Level Order Traversal
- 6. Serialize and Deserialize Binary Tree
- 7. Subtree of Another Tree
- 8. Construct Binary Tree from Preorder and Inorder Traversal
- 9. Validate Binary Search Tree
- 10. Kth Smallest Element in a BST
- 11. Lowest Common Ancestor of BST
- 12. Implement Trie (Prefix Tree)
- 13. Add and Search Word
- 14. Word Search II

### 1 Two Sum

Return index of two elements in the array that add upto the target value.

```
class Solution:
    def twoSum(self, nums: List[int], target: int) -> List[int]:
        hash_map = {}

    for idx, num in enumerate(nums):
        res = target - num
        if res not in hash_map:
              hash_map[num] = idx
        else:
        return idx, hash_map[res]
```

## 2 Best Time To Buy and Sell Stock

Return the maximum profit that can be made using the price values provided.

```
class Solution:
    def maxProfit(self, prices: List[int]) -> int:
        profit = 0
        buying_price = float('inf')
        for i in range(0, len(prices)):
            buying_price = min(prices[i], buying_price)
            profit = max(profit, prices[i] - buying_price)
        return profit
```

# 3 Contains Duplicate

Return boolean value if there are duplicate elements in the array.

```
class Solution:
    def containsDuplicate(self, nums: List[int]) -> bool:
        dup = set()
        for i in range(len(nums)):
            if nums[i] not in dup:
                 dup.add(nums[i])
        else:
            return True
    return False
```

## 4 Product of Array Except Self

```
Input: nums = [1,2,3,4] Output: [24,12,8,6]
Input: nums = [-1,1,0,-3,3] Output: [0,0,9,0,0]

class Solution:
    def productExceptSelf(self, nums: List[int]) -> List[int]:
        res = [1] * len(nums)

    prefix = 1
    for i in range(0, len(nums)):
        res[i] *= prefix
        prefix *= nums[i]

    postfix = 1
    for j in range(len(nums)-1, -1, -1):
        res[j] *= postfix
        postfix *= nums[j]
```

## 5 Maximum Subarray

Trick to solving is if the current sum is less than zero we set the current sum to zero.

```
class Solution:
    def maxSubArray(self, nums: List[int]) -> int:
        #Kadane's algorithm repeated work
        max_sum = nums[0]
        curr_sum = 0
        for i in range(0, len(nums)):
            if curr_sum < 0:
                curr_sum = 0
            curr_sum += nums[i]
            max_sum = max(curr_sum , max_sum)
        return max_sum</pre>
```

# 6 Maximum Product Subarray

Technique to solving is to maintain current minimum and current maximum and utilize both to get the maximum answer.

```
class Solution:
   def maxProduct(self, nums: List[int]) -> int:
```

```
#if the max subarray is just one value
result = max(nums)
curr_min = 1
curr_max = 1
for n in nums:
    temp = curr_max
    curr_max = max(curr_max*n, curr_min*n, n)
    curr_min = min(curr_min*n, temp*n, n)
    result = max(result, curr_max)
return result
```

### 7 Container with most water

Two pointer approach and move in the direction of greater height.

```
class Solution:
    def maxArea(self, height: List[int]) -> int:
        right = 0
        left = len(height)-1
        max_area = 0
        while right < left:
            min_height = min(height[right], height[left])
            max_area = max(max_area, min_height * (left-right))
        if height[right] < height[left]:
            right+=1
        else:
            left -= 1
        return max_area</pre>
```

## 8 Find minimum in rotated sorted array

```
class Solution:
    def findMin(self, nums: List[int]) -> int:
        # binary search
    left = 0
    right = len(nums)-1
    while left <= right:
        mid = int(left + (right-left)/2)

    next = (mid+1)%len(nums)
    prev = (mid-1+len(nums))%len(nums)

    if nums[mid] <= nums[prev] and nums[mid] <= nums[next]:
        return nums[mid]
    elif nums[0] <= nums[mid]:</pre>
```

```
\begin{array}{rll} & \text{left} = \text{mid} + 1 \\ & \text{elif nums}[\text{mid}] <= \text{nums}[-1]: \\ & \text{right} = \text{mid} -1 \\ & \text{return nums}[0] \end{array}
```

## 9 Search in rotated sorted array

```
class Solution:
    def search (self, nums: List[int], target: int) -> int:
         def find_min_index(nums):
             left = 0
             right = len(nums)-1
             while left <= right:
                 mid = int(left + (right-left)/2)
                 next = (mid+1)\%len(nums)
                 prev = (mid-1+len(nums))\%len(nums)
                 if nums[mid] <= nums[prev] and nums[mid] <= nums[next]:
                      return mid
                  elif nums[0] \le nums[mid]:
                      left = mid + 1
                  elif nums[mid] \le nums[-1]:
                      right = mid -1
             return -1
         def binary_search(nums, target):
             left = 0
             right = len(nums)-1
             while left <= right:
                 mid = int(left + (right-left)/2)
                 if nums[mid] = target:
                      return mid
                  elif nums[mid] < target:
                      left = mid + 1
                  elif nums[mid] > target:
                      \operatorname{right} \; = \; \operatorname{mid} \; \, -1
```

```
return -1
        ind = find_min_index(nums)
        arr1 = nums[0:ind]
        arr2 = nums[ind:]
        val = binary_search(arr1, target)
        val2 = binary_search(arr2, target)
        return val if val2 == -1 else val2+len(arr1)
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     3 Sum
Sort the array and run two sum II for each unique element in the loop.
class Solution:
    def threeSum(self, nums: List[int]) -> List[List[int]]:
        res = []
        nums = sorted(nums)
        for i in range (0, len (nums)):
            if i > 0 and nums [i] = nums [i-1]:
                 continue
            left = i+1
            right = len(nums)-1
            while left < right:
                summed = nums[left] + nums[right] + nums[i]
                 if summed == 0:
                     res.append((nums[left], nums[right], nums[i]))
                     left += 1
                     while nums[left] = nums[left -1] and left < right:
                         left += 1
                 elif summed < 0:
                     left += 1
                 elif summed > 0:
                     right -=1
```

### 11 Container with most water

return res

```
class Solution: def maxArea(self, height: List[int]) \rightarrow int: left = 0
```

```
right = len(height)-1
max_area = -1
while left < right:
   h = min(height[left], height[right])
   area = (right - left) * h
   max_area = max(max_area, area)
   if height[left] < height[right]:
        left += 1
   elif height[right] < height[left]:
        right -=1
   elif height[right] == height[left]:
        left+=1
return max_area</pre>
```

## 12 Binary - Sum of two integers

We make use of xor a=1, b=0, ans=1. a=1, b=1, ans=0

When we have two one's we have a carry. If a and b are 1, we have a carry. a&b left shift 1

If we dont have a carry we are done.

Code doesn't work in python since the assumption that we make is that the length of integers is 32 bits. However python allows for an infinite length of integers. Java does not.

```
class Solution {
    public int getSum(int a, int b) {
        while( b!= 0){
            int temp = (a &b) <<1;
            a = a ^ b;
            b = temp;
        }
        return a;
    }
}</pre>
```

### 13 Number of one bits

Count the number of 1's in bits.

```
class Solution:  \begin{array}{lll} def \ hammingWeight(self \,,\, n\colon int) \to int\colon \\ total \, = \, 0 \\ while \ n \, != \, 0\colon \\ total \, +\! = \, n \, \%2 \\ n \, = \, n \, >\! > \, 1 \\ return \ total \end{array}
```

## 14 Counting Bits

```
Dynamic Programming Solution

class Solution:
    def countBits(self, n: int) -> List[int]:
        dp = [0]*(n+1)
        offset = 1

    for i in range(1, n+1):
        if offset * 2 == i:
            offset = i

        dp[i] = 1 + dp[i - offset]

    return dp
```

## 15 Missing Number

```
class Solution:
    def missingNumber(self, nums: List[int]) -> int:
        xor1 = 0
        xor2 = nums[0]
        for i in range(1, len(nums)+1):
            xor1 ^= i
        for j in range(1, len(nums)):
            xor2 ^= nums[j]

    return xor1 ^ xor2
```

## 16 Reverse Bits

```
class Solution:
    def reverseBits(self, n: int) \rightarrow int:
        res = 0
        for i in range(0, 32):
        res = res << 1
        if n % 2 == 1:
            res += 1

        n = n >> 1
        return res
```

# 17 DP - Climbing stairs

You are climbing a staircase. It takes n steps to reach the top.

Each time you can either climb 1 or 2 steps. In how many distinct ways can you climb to the top?

```
class Solution:
    def __init__(self):
        self.map = {}

    def climbStairs(self, n: int) -> int:
        if n == 0:
            return 0
        if n == 1:
            return 1
        if n == 2:
            return 2

    if n in self.map:
        return self.map[n]

    self.map[n] = self.climbStairs(n-1) + self.climbStairs(n-2)
    return self.map[n]
```

## 18 Coin Change

```
class Solution:
    def coinChange(self, coins: List[int], amount: int) -> int:
        memo = \{\}
        def unbounded(coins, i, target):
            key = (i, target)
            if key in memo:
                return memo[key]
            if target = 0:
                return 0
            if i \le 0 and target > 0:
                return float ('inf')
            if coins[i-1] > target:
                memo[key] = unbounded(coins, i-1, target)
                return memo[key]
            memo[key] = min(1+unbounded(coins, i, target-coins[i-1]),
                            unbounded (coins, i-1, target))
```

```
return memo[key]
res = unbounded(coins, len(coins), amount)
if res == float('inf'):
    return -1
return res
```

# 19 Longest Increasing Subsequence

Recursive memoized solution.

```
Best to explore using nested loops as well.
```

```
class Solution:
    def lengthOfLIS(self, nums: List[int]) -> int:
        memo = \{\}
        def lcs(nums, prev, current):
            key = (prev, current)
            if key in memo:
                return memo[key]
            if len(nums) == current:
                return 0
            c1 = 0
            if (prev = -1) or (nums[prev] < nums[current]):
                c1 = 1 + lcs (nums, current, current+1)
            c2 = lcs(nums, prev, current+1)
            memo[key] = max(c1, c2)
            return memo[key]
        return lcs(nums, -1, 0)
```

# 20 Longest Common Subsequence

```
class Solution:
   def longestCommonSubsequence(self, text1: str, text2: str) -> int:
```

```
memo = \{\}
          def lcs(s1, s2, i, j):
               key = (i, j)
               if key in memo:
                    return memo[key]
               if i \le 0 or j \le 0:
                    return 0
               if s1[i-1] = s2[j-1]:
                   memo[key] = 1 + lcs(s1, s2, i-1, j-1)
                    return memo[key]
              memo [\text{key}] = \max(\text{lcs}(\text{s1}, \text{s2}, \text{i}, \text{j-1}), \text{lcs}(\text{s1}, \text{s2}, \text{i-1}, \text{j}))
               return memo[key]
          return lcs(text1, text2, len(text1), len(text2))
21
      Word Break
class Solution:
     \label{eq:condition} def \ wordBreak(\, self \, , \ s \colon \ str \, , \ wordDict \colon \ List[\, str \, ] \, ) \ -\!\!\!> \ bool \colon
          dp = [False] * (len(s)+1)
          dp[len(s)] = True
          for i in range (len (s), -1, -1):
               for w in wordDict:
                    if i+len(w) \le len(s) and s[i : i + len(w)] = w:
                         dp[i] = dp[i+len(w)]
                    if dp[i]:
                         break
          return dp[0]
Recursive DP
class Solution:
     def wordBreak(self, s: str, wordDict: List[str]) -> bool:
         memo = \{\}
          def wb(s, wordDict):
               if s in memo:
                    return memo[s]
```

```
if len(s) == 0:
    memo[""] = True
    return True

for word in wordDict:
    if s.startswith(word) and len(s) >= len(word):
        if wb(s[len(word):], wordDict):
            return True

memo[s] = False

return False

return wb(s, wordDict)
```

### 22 Combination Sum

Given an array of distinct integers candidates and a target integer target, return a list of all unique combinations of candidates where the chosen numbers sum to target. You may return the combinations in any order.

Time Complexity =  $2\hat{T}$  (Where T is the target value, as the height of the decision tree is determined by the target value.)

```
class Solution:
    def combinationSum(self, candidates: List[int], target: int) -> List[List[int]]:
        res = []

    def dfs(i, curr, total):
        if total == target:
            res.append(curr.copy())
            return

        if i >= len(candidates) or total > target:
            return

        curr.append(candidates[i])
        dfs(i, curr, total+candidates[i])
        curr.pop()
        dfs(i+1, curr, total)

        return

    dfs(0, [], 0)
    return res
```

### 23 House Robbers

```
Find the recurrence, code up the recursion.
```

```
Recurrence = max(arr[0] + dp(arr[2:]), dp(arr[1:]))
class Solution:
    def rob(self, nums: List[int]) -> int:
        # recurrence
        #rob = max(nums[0] + rob(nums[2:]), rob(nums[1:]))
        memo = {}

    def dp(nums, n):
        if n >= len(nums):
            return 0

        if n in memo:
            return memo[n]

        memo[n] = max(nums[n] + dp(nums, n+2), dp(nums, n+1))
        return memo[n]

    return dp(nums, 0)
```

### 24 House Robber II

This time the first and last house are connected so we split the array into two parts.

```
We run our house robber I function on arr[:-1] and arr[1:]
```

```
class Solution:
    def rob(self, nums: List[int]) -> int:
        if len(nums) == 1:
            return nums[0]

    memo = {}

    def dp(nums, n):
        if n >= len(nums):
            return 0
        if n in memo:
            return memo[n]

    memo[n] = max(nums[n] + dp(nums, n+2), dp(nums, n+1))
    return memo[n]
```

```
arr1 = dp(nums[:-1], 0)
memo = {}
arr2 = dp(nums[1:], 0)
max_ret = max(arr1, arr2)
return max_ret
```